## Assignment 4

for Discrete Structures 2

- 4.1 Explain what is wrong in the presented erroneous inductive proof on page 2.
- 4.2 Let  $(A, \leq), (B \sqsubseteq)$  be complete lattices, and  $f: A \to B$  be a monotonic function<sup>1</sup>. Prove that

$$f(\bigvee C) \sqsupseteq \bigsqcup \{ f(x) \mid x \in C \}$$

for any subset  $C \subseteq A$ .

 $\bigvee$  and  $\bigsqcup$  stand for the lub's, respectively, in  $(A, \leq)$  and in  $(B \sqsubseteq)$ .

[This problem is not related to induction, but rather to lattices and ccpo's. But a special case of this property is used in the lecture notes, as a lemma in the context of transfinite induction.]

4.3 Assume that  $f: A \to A$  is a monotonic function on a complete lattice  $(A, \leq)$ . Prove that  $f^{\beta} \leq f^{\alpha}$  for any ordinals  $\beta < \alpha$ .

Make it clear which induction principle you apply. You may use the fact (proved in the lecture note) that  $f^{\alpha} \leq f^{\alpha+1}$ . Note that it may be not necessary to refer to the inductive assumption in one of the cases within your proof.

<sup>&</sup>lt;sup>1</sup>I.e., if  $x \leq y$  then  $f(x) \sqsubseteq f(y)$ .

## An erroneous inductive proof

Let M be an infinite set, f a function on M, and  $L \subseteq M$ . Property P(L)says that f is constant on L.  $P(L): \text{ If } x, y \in L \text{ then } f(x) = f(y)$ We "prove" P(L) for any finite non empty L. Base case, |L| = 1:  $x, y \in L \implies x = y \implies f(x) = f(y)$ . Inductive step: Assume P(L') holds for any |L'| = n. Let  $|L| = n + 1, x, y \in L$ . 1.  $x = y \implies f(x) = f(y)$ . 2.  $x \neq y$   $L \setminus \{y\}$   $L: \underbrace{x \cdots z \cdots y}_{L \setminus \{x\}}$   $|L \setminus \{x\}| = n = |L \setminus \{y\}|,$   $f(x) = f(z) \text{ for each } z \in L \setminus \{y\},$   $f(y) = f(z) \text{ for each } z \in L \setminus \{x\}.$ Thus f(x) = f(z) = f(y).

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Hence f(x) = f(y) for any  $x, y \in L$ .

What is wrong?