

Teaching Logic Programming at the Budapest University of Technology

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Overview of the talk

- The Budapest University of Technology and Economics
- The Declarative Programming course
- Elective LP courses
- Other educational activities
- Conclusions

The Budapest University of Technology and Economics

- BUTE — Budapest University of Technology and Economics:
 - Established in 1782
 - Initially: Institutum Geometrico-Hydrotechnicum
 - within the Faculty of Liberal Arts, University of Buda
 - Present:
 - Seven faculties, >110 departments and institutes
 - About 1000 full-time lecturers, several hundred research staff
 - Approx. 10,000 students, of which 10% are foreigners
 - BUTE issues about 70% of Hungary's engineering diplomas
- Faculty of Electrical Engineering and Informatics — the home of computer-related education
 - About 200 full time teaching staff
 - Approx. 3,000 students, majoring in Electric Engineering and Technical Informatics (roughly equivalent to Computer Science)
- Five year (10 semester) education leading to an MSc/MEng degree
- More info at: <http://www.bute.hu/en/>

Logic Programming education at BUTE

- Students exposed to LP-related topics (and the relevant subjects):
 - Majors in Technical Informatics at Faculty of Electrical Eng. and Informatics (approx. 400/year)
 - Mathematical Logic — semester 3, compulsory course
 - Declarative Programming (DP) — semester 4, compulsory course
 - Highly Efficient Logic Programming (HELP) — elective course
 - Selected Topics from Logic Programming (STLP) — elective student seminar
 - AI laboratory — semester 7 (specialisation Intelligent Systems)
 - Majors in Mathematics, at Faculty of Natural Sciences, with specialisation in Algebra and applications (approx. 10/year)
 - Logic Programming — LP (semester 7, compulsory course, same as the logic programming part of DP)
 - Students at some other faculties (e.g. Mechanical Engineering)
 - various courses on AI and Programming
- The author is involved in the DP (LP), HELP, STLP courses — this is the focus of the present talk

The Declarative Programming course

- Course topics
 - Functional programming, SML — Péter Hanák
 - Logic programming, Prolog — Péter Szeredi
- Course data
 - Base position: Semester 4 (spring), cross semester: 5 (autumn)
 - Students enrolled:
 - spring 2004: 473
 - autumn 2004: 162
 - Two 2*45 minute lectures/week, 14 weeks/semester
 - No laboratory exercises :-(
 - Presented — 20 times so far — since 1994
(until 1999, under the title Programming Paradigms)
- Preceding courses
 - Programming languages: C, C++, Java
 - Mathematical Logic

Global course structure

- Major lecture blocks

| | |
|------------------------|--|
| Lecture 1: | Introducing the declarative programming paradigm |
| Lectures 2-8: | Logic Programming, part 1 |
| Lectures 9-15: | Functional Programming, part 1 |
| Lectures 16-21: | Logic Programming, part 2 |
| Lectures 22-27: | Functional Programming, part 2 |
| Lecture 28: | Summary, outlook |

- Links between the functional and logic programming parts

- Common concepts and techniques, e.g. tail-recursion, accumulators, construction and decomposition using pattern-matching
- Common major assignment, occasionally common minor assignments

Topics in LP — part 1

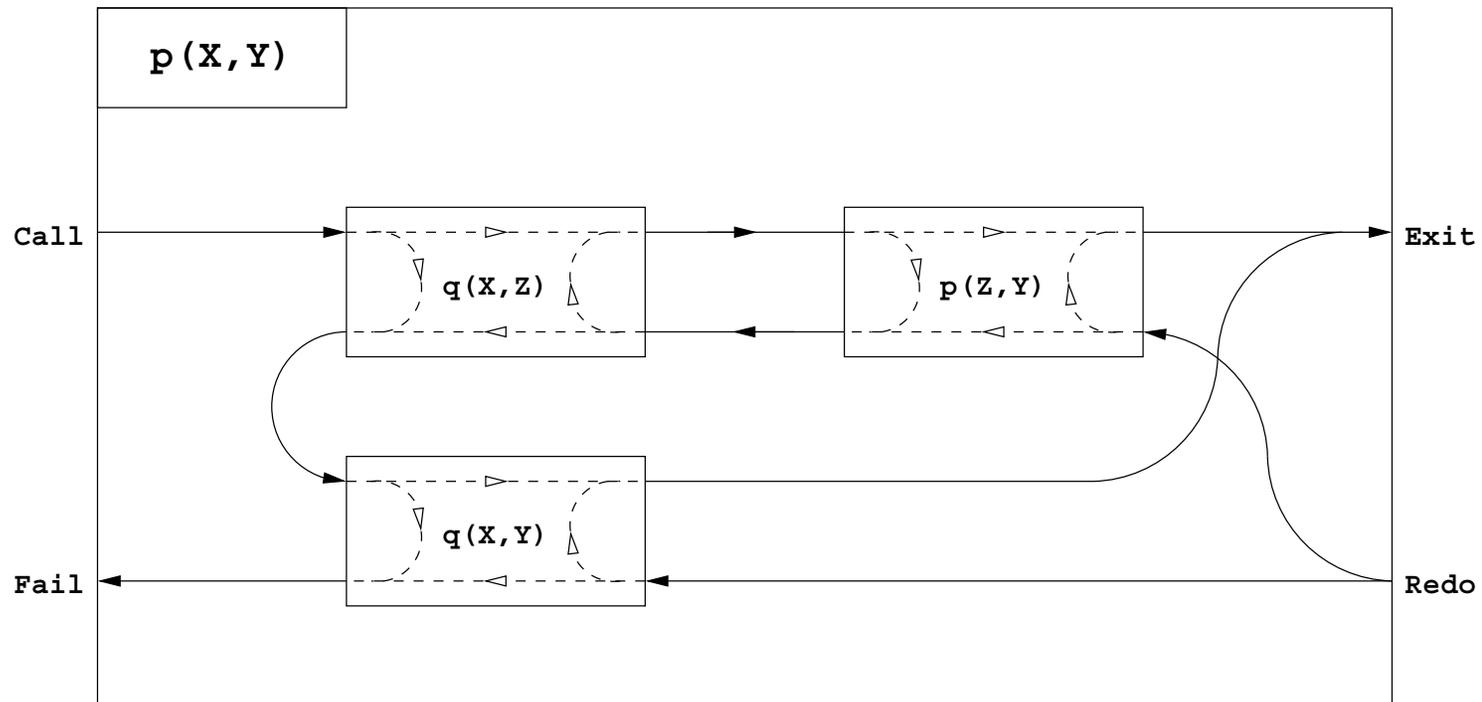
| | |
|------------|---|
| Lecture 1. | Introducing the declarative programming paradigm. A very simple declarative subset of the C language (Cékla = beet-root in Hungarian :-). |
| Lecture 2. | Introductory Prolog examples (family relations, summing numbers in binary trees), Prolog as a subset of logic, declarative semantics. |
| Lecture 3. | Procedural semantics of Prolog, execution models (goal-reduction and procedure-box models). |
| Lecture 4. | Data structures, unification, the logic variable. |
| Lecture 5. | Operators, disjunction, negation, if-then-else. |
| Lecture 6. | Lists, basic list handling library predicates. |
| Lecture 7. | Example: finding paths in a graph, using various representations. |
| Lecture 8. | Prolog syntax summary. |

SAMPLE SLIDE: Procedure-box model example

$p(X, Y) \text{ :- } q(X, Z), p(Z, Y) .$

$p(X, Y) \text{ :- } q(X, Y) .$

$q(1, 2) . q(2, 3) , q(2, 4) .$



SAMPLE SLIDE: Box model — OO view, the „next solution” method of p/2

```

bool p::next ()
{ switch(clno) {
  case 0:
    // entry point for the Call port
    clno = 1;
    // enter clause 1:
    qaptr = new q(x, &z);
    // create a new instance of subgoal q(X,Z)
    redo11:
    if(!qaptr->next()) {
      // if q(X,Z) fails
      delete qaptr;
      // destroy it,
      goto cl2;
      // and continue with clause 2 of p/2
    }
    pptr = new p(z, py);
    // otherwise, create a new instance of subgoal p(Z,Y)
  case 1:
    // (enter here for Redo port if clno==1)
    /* redo12: */
    if(!pptr->next()) {
      // if p(Z,Y) fails
      delete pptr;
      // destroy it,
      goto redo11;
      // and continue at redo port of q(X,Z)
    }
    return true;
    // otherwise, exit via the Exit port
  cl2:
    clno = 2;
    // enter clause 2:
    qbpPtr = new q(x, py);
    // create a new instance of subgoal q(X,Y)
  case 2:
    // (enter here for Redo port if clno==1)
    /* redo21: */
    if(!qbpPtr->next()) {
      // if q(X,Y) fails
      delete qbpPtr;
      // destroy it,
      return false;
      // and exit via the Fail port
    }
    return true;
    // otherwise, exit via the Exit port
} }

```

p(X,Y) :- q(X,Z), p(Z,Y).

p(X,Y) :- q(X,Y).

Topics in LP — part 2

| | |
|-------------|---|
| Lecture 16. | Pruning the search space. Control predicates. |
| Lecture 17. | Determinism, indexing, tail-recursion, accumulators. |
| Lecture 18. | Rewriting imperative programs to Prolog, collecting and enumerating solutions. |
| Lecture 19. | Meta-logical built-in predicates. |
| Lecture 20. | Modularity, meta-predicates, meta-programming, dynamic predicates. |
| Lecture 21. | Definite clause grammars, „traditional” built-in predicates. |
| Lecture 28. | Brief outlook on LP extensions (external interfaces, coroutining, constraints). |

SAMPLE SLIDE: Meta-logical predicates, implementing term-precedes

```
% T1 precedes T2-t in standard order. (Equivalent to T1 @< T2, except for
% variables, these are ordered by their first occurrence in T1-T2.)
```

```
precedes(T1, T2) :-
    \+ \+ (numbervars(T1-T2, 0, _), prec(T1, T2)).
```

```
% class(+T, -C): Term T belongs to term class C.
```

```
class(T, C) :-
    (   T='$VAR'(_) -> C=0           % variable
    ;   float(T)     -> C=1         % float
    ;   integer(T)   -> C=2         % integer
    ;   atom(T)      -> C=3         % atom
    ;   compound(T) -> C=4         % compound
    ).
```

```
% Numbevar'd term T1 precedes term T2-t.
```

```
prec(T1, T2) :-
    class(T1, C1), class(T2, C2),
    (   C1 == C2 ->
        (   C1 == 1 -> T1 < T2   % floating point numbers
        ;   C1 == 2 -> T1 < T2   % integers
        ;   struct_prec(T1, T2)  % variables, atoms, compounds
        )
    ;   C1 < C2           % different term classes
    ).
```

SAMPLE SLIDE: implementing term-precedes (contd.)

```

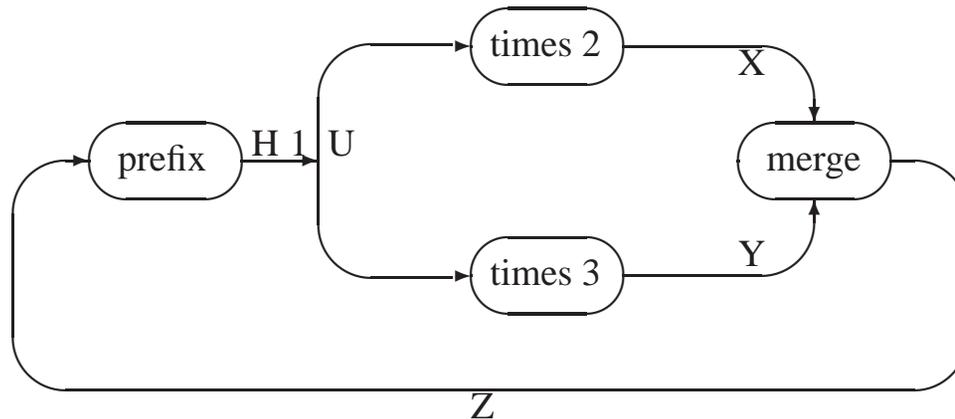
% S1 precedes S2 (S1 and S2 are callable terms, i.e. atoms or compounds).
struct_prec(S1, S2) :-
    functor(S1, F1, N1), functor(S2, F2, N2),
    ( N1 < N2 -> true
    ; N1 == N2,
      ( F1 == F2 -> args_prec(1, N1, S1, S2)
      ; atom_prec(F1, F2)
      )
    ).

% For the first argument position i, N0=<i=<N, for which
% S1[i] differs from S2[i], S1[i] precedes S2[i] (and such i exists).
args_prec(N0, N, S1, S2) :-
    N0 =< N,
    arg(N0, S1, A1), arg(N0, S2, A2),
    ( A1 = A2 -> N1 is N0+1, args_prec(N1, N, S1, S2)
    ; prec(A1, A2)
    ).

% Atom A1 precedes atom A2.
atom_prec(A1, A2) :-
    atom_codes(A1, C1), atom_codes(A2, C2), struct_prec(C1, C2).

```

SAMPLE SLIDE: Coroutinging: Simplified Hamming problem (partial code)



```

% List H contains the first N numbers whose all prime factors are 2 or 3.
hamming(N, H) :-

```

```

    U = [1|H], times(U, 2, X), times(U, 3, Y),
    merge(X, Y, Z), prefix(N, Z, H).

```

```

% times(X, M, Z): Elements of list Z are those of X multiplied by M.

```

```

:- block times(-, ?, ?).

```

```

times([A|X], M, Z) :- B is M*A, Z = [B|U], times(X, M, U).

```

```

times([], _, []).

```

```

(...).

```

Assignments

- How to teach a programming language without laboratory exercises?
 - Assignments: multiple minor and a single major assignment (common to LP and FP)
 - Interactive exercising tool (in ETS, see later)
 - All these are non-compulsory
- Minor assignments
 - 4-6 assignments altogether in LP and FP
 - Fairly simple problems, < 20 lines of code needed
 - Submission within two weeks
- Major assignment
 - A single task to be solved in Prolog and/or SML
 - A scalable problem, most often a logic puzzle
 - Even a simple generate-and-test solution gets some credit
 - Best solutions compete in a „ladder contest” for extra credit
 - Submission in 4-6 weeks

SAMPLE: Minor assignment 1

- Write in Cékla a function `palindrome (n)`
 - `palindrome (n)` returns the smallest integer $b > 1$
 - such that n written in base b is a palindrome.
- Main characteristics of the Cékla C subset:
 - integer type only
 - functions, `if`, and `return` statements
 - `+`, `-`, `*`, `/`, and `%` arithmetic operators, the six comparison operators

The solution of minor assignment 1

```
/* The reverse of number 'n' in base 'base' appended to 'revn' */
int reverse(int n, int base, int revn) {
    if (n == 0)
        return revn;
    else {
        int last_digit = n % base;
        return reverse(n/base, base, revn*base+last_digit);
    }
}
```

```
/* The smallest integer 'b' >= 'base', so that 'n' written
   in base 'b' is a palindrome */
int palindrome1(int n, int base){
    if (reverse(n, base, 0) == n)
        return base;
    else return palindrome1(n, base+1);
}
```

```
/* The smallest integer 'b', so that 'n' written in base 'b' is a palindrome */
int palindrome(int n) {
    return palindrome1(n,2);
}
```

Major assignment in spring 2002: the Magic Spiral puzzle

- The puzzle:
 - A square board of $n * n$ fields is given.
 - Place integer numbers, chosen from the range $[1..m]$, $m \leq n$, on the board, so that:
 1. in each row and each column all integers in $[1, m]$ occur exactly once (and so there are $n - m$ empty fields);
 2. along the spiral starting from the top left corner, the integers follow the pattern $1, 2, \dots, m, 1, 2, \dots, m, \dots$ (number m is called the period of the spiral).
 - Initially, some numbers are already placed on the board.
- An example puzzle and its solution:

| | | | | | | |
|---|--|--|---|--|--|--|
| | | | | | | |
| | | | 4 | | | |
| 1 | | | | | | |
| | | | | | | |
| | | | | | | |
| | | | | | | |
| | | | | | | |

| | | | | | | |
|---|---|---|---|---|---|---|
| | | 1 | 2 | 3 | | 4 |
| 2 | | 3 | 4 | 1 | | |
| 1 | 3 | 4 | | | 2 | |
| 4 | 2 | | | | 3 | 1 |
| 3 | 1 | | | | 4 | 2 |
| | 4 | | 3 | 2 | 1 | |
| | | 2 | 1 | 4 | | 3 |

Major assignment in spring 2003: the Snake puzzle

- The puzzle:
 - A rectangular board of $n * m$ fields is given.
 - Mark certain fields on the board as belonging to the *snake*:
 - The snake consists of a sequence of neighbouring fields (i.e. fields sharing a side).
 - The position of the snake head (the first field) and of the tail (last field) is given.
 - The snake can touch itself only diagonally.
 - Certain fields of the table contain a number.
 - The snake cannot pass through fields containing a number.
 - If a field contains number c , then among the eight side and diagonal neighbours of this field there are exactly c which contain a snake piece. (cf. the minesweeper game).
- An example puzzle and its solution:

| | | | | |
|----------|--|----------|----------|--|
| H | | T | | |
| | | 5 | | |
| 4 | | | | |
| | | | | |
| | | | 3 | |

| | | | | |
|----------|---|----------|----------|---|
| ● | | ● | ● | |
| ● | ● | 5 | ● | ● |
| 4 | ● | | | ● |
| | ● | ● | ● | ● |
| | | | 3 | |

Major assignment in spring 2004: the Clouds puzzle

- The puzzle:

- A board of $n * m$ fields is given. Mark certain fields on the board as belonging to a *cloud*:
 - Clouds occupy an area of rectangular shape, their width and height is at least two units.
 - No clouds touch each other, not even diagonally.
- Supplied information:
 - the size of the board;
 - the number of cloudy fields in certain rows/columns of the board (-1 means unknown);
 - the presence (+) or absence (-) of a cloud at certain fields of the board.

- An example puzzle and its solution:

```

+---+---+---+---+---+
|   |   |   |   |   | 4
+---+---+---+---+---+
|   |   |   |   | - | 4
+---+---+---+---+---+
|   |   |   |   |   | 0
+---+---+---+---+---+
|   | + |   |   |   | -1
+---+---+---+---+---+
|   |   |   |   |   | 4
+---+---+---+---+---+
  4   4  -1  4   2

```

```

+---+---+---+---+---+
| # | # | # | # | . | 4
+---+---+---+---+---+
| # | # | # | # | - | 4
+---+---+---+---+---+
| . | . | . | . | . | 0
+---+---+---+---+---+
| # | + | . | # | # | -1
+---+---+---+---+---+
| # | # | . | # | # | 4
+---+---+---+---+---+
  4   4  -1  4   2

```

Tools used in teaching Declarative Programming

- The Electronic Teaching aSsistant (ETS) is a Web based facility providing
 - access to a database of students of the course, and their results,
 - assignment submission and automatic testing,
 - facilities for student exercising.
- The RDBG reduction debugger visualisation tool
 - step-by-step building of the Prolog search tree
 - skipping, unleashing, breakpoints
- The MATCH plagiarism-detection tool
 - detects similar assignment solutions
 - comparing the call graphs of the programs
 - with front-ends for both Prolog and SML
- Other tools
 - The Cékla interpreter
 - The `xdvipresent` package for slides

Student exercising with ETS

- Each exercise belongs to a scheme and a topic and has a difficulty:
 - The **scheme** determines the presentation and processing **format**, e.g. the student is expected to predict what happens when a given goal is run
 - The **topic** corresponds to the **part of the material** taught, e.g. list processing
 - The **difficulty** can be easy, medium, or hard.
- Prolog schemes
 - Execution: decide on the result of executing a given goal
 - Success/failure/error: deterministic goal, can fail or raise an exception. A single variable assignment is asked for.
 - Multiple variables: supply the substitutions of all variables of a deterministic goal (typically unification).
 - All solutions: enumerate all solutions of a non-deterministic goal, in proper order.
 - Canonical form: type in the canonical form of a Prolog term (cf. `write_canonical`).
 - Programming: write a Prolog program following a given specification.

Student exercising with ETS — examples

- Scheme: „Canonical form”, topic: lists/operators

```
[ [], [] ]      - (1, 2)      1 - - 1
```

- Scheme: „Success/failure/error”, topic: unification

```
| ?- [X, 1 | X] = [_, _].
| ?- [X, a, X] = [1, 2, a, 1, 2].
```

- Scheme: „Multiple variables”, topic: unification

```
| ?- g(1+2+3, [a, b]) = g(X+Y, [U|V]).
| ?- h([H, G], H*G) = h([Q/1|R], P/Q*3).
```

- Scheme: „Success/failure/error”, topic: list processing and control predicates

```
| ?- length(X, 1), member(a, X).
| ?- member(X, [1, 2, 3]), !, X < 3.
```

- Scheme: „Multiple solutions”, topic: list processing:

| | |
|----------|---|
| Program: | <code>app([X L1], L2, [X L3]) :- app(L1, L2, L3). app([], L, L).</code> |
| Goal: | <code> ?- app(L, [a _], [a,b,a,b,a]).</code> |

Elective courses — Highly Efficient Logic Programming

- 14 lectures, 2*45 minutes each, so far held 8 times since 1997
- Course topics
 - Mercury — efficiency through streamlining and cleaning the LP language (2 lectures)
 - CLP(\mathcal{X}) — efficiency through additional reasoning capabilities (12 lectures)
 - the CLP(B), CLP(R/Q), CLP(FD) and CHR libraries of SICStus Prolog
 - major assignment usually shared with the preceding DP course
- Typical course layout (1 unit = 45 minutes)

| | Topic | Time | Slides |
|----|--|----------|------------|
| 1. | Prolog extensions relevant for CLP | 3 units | 15 slides |
| 2. | The CLP(\mathcal{X}) scheme and CLP(R/Q) | 3 units | 19 slides |
| 3. | CLP(B) | 2 units | 8 slides |
| 4. | CLP(FD) | 14 units | 109 slides |
| 5. | Constraint Handling Rules | 2 units | 11 slides |
| 6. | Mercury | 4 units | 23 slides |
| | Σ | 28 units | 185 slides |

- For more details see Szeredi's paper in LNAI 3010

Other elective courses

- Selected Topics from Logic Programming — student seminar
 - 14 lectures, 90 minutes each, so far held twice, in 2001 and 2003
 - I hold the first few lectures, normally on Prolog implementation
 - Subsequent (45 or 90 minute) lectures held by students, topics:
 - theory of LP
 - parallel logic programming,
 - object-oriented, graphical, and web-related extensions,
 - abstract interpretation,
 - tracing,
 - overviews of concrete Prolog implementations
- The Semantic Web and Ontologies (by Gergely Lukácsy and Péter Szeredi)
 - Web technologies, RDF, Description Logics, etc.
 - Link to LP: several students chose to write the assignment, a tableaux based reasoner for a Description Logic, in Prolog

Other LP-related activities — Directed Projects

- Semesters 8 and 9, weekly load equivalent to 6*45 minutes course
- Some Directed Projects I led:
 - Applying CLP(FD) for scheduling plastic moulding machines — Tamás Benkő, 1997
 - Interfacing Prolog to Corba — Gábor Gesztesi and Gábor Marosi, 1997–98
 - Debugging CLP(FD) programs — Dávid Hanák and Tamás Szeredi, 2000
(this led to the development of a new SICStus library, see the poster and the talk at WELP)
 - Using CHR for reasoning in Description Logics — Bence Szász, 2002
 - A Prolog based RDF reasoning system — Gergely Lukácsy, 2002
 - A Prolog-Java interface using sockets — Péter Biener, 2003
- Typically Directed Projects serve as the basis for the MSc Thesis
 - All above projects, except for the Corba one, led to an MSc Thesis

Other LP-related activities — MSc Theses

- MSc Theses I supervised (further to Theses resulting from Directed Projects):
 - Implementation of a constraint reasoning system — Tamás Rozmán, 1997
 - Conversion of document description languages — Zsolt Lente, 2000
 - Knowledge-based tools for information integration — Attila Fokt, 2000
 - Computer Support for Declarative Programming Courses — Dávid Hanák, 2001
 - Verification of object-oriented models using constraints — Péter Tarján, 2001
 - Transforming object-constraints to logic — Károly Opor, 2002
 - Logic-based methods for planning queries on heterogeneous data sources — Lukács Tamás Berki, 2003

Other LP-related activities — Student Research Projects

- Student Research Projects — a long tradition at Hungarian Universities
 - Yearly, faculty-level Student Conference (TDK in Hungarian):
 - 40-60 page paper
 - 20 minute presentation
 - best papers get 1st, 2nd and 3rd prizes
 - Biennial National Student Conferences
- Student Conference in 2003 at the Faculty of Electrical Engineering and Informatics at BUTE:
 - 118 presentations,
 - involving approx. 170 students

Other LP-related activities — Student Research Projects (contd.)

- Student Conference Papers related to LP:
 - Solving a stock exchange allocation problem (using CLP), Dániel Varró, 1998, I. Prize.
 - Conversion of SGML languages, Zsolt Lente, 1999, II. Prize.
 - Comparison of source program structures, Gergely Lukácsy, 2000, I. Prize, Rector's Special Prize; I. Prize at the National Student Conference, 2001
 - Efficient access of an object-oriented database from logic programs, Ambrus Wagner, 2000
 - A Web-based student exercising system for teaching programming languages, András György Békés, Lukács Tamás Berki, 2001
 - Intelligent querying and reasoning on the Web, Gergely Lukácsy, 2002, I. Prize; Special Prize of the Hungarian W3C Office at the National Student Conference, 2003
 - Visualisation of Prolog program execution, Tamás Nepusz, 2003, II. Prize
 - Using abstract interpretation in SICStus Prolog, Balázs Leitem, 2003, III. Prize

Discussion

- Good points
 - Each year about 400 students of BUTE get acquainted with logic programming
 - Logic programming attracts the attention of talented students
 - programming style previously unknown to most students
 - the course shows some problems where the new paradigm can be used
 - the major assignment poses a challenge, as it would be quite difficult to solve using imperative languages
 - the ladder contest adds the thrill of peer-to-peer competition
 - Most talented students get further involved with LP:
 - learn constraint logic programming in the HELP course (top 10%)
 - explore further areas of logic programming in the STLP seminar (top 2–3%)
 - work on logic programming as part of his/her Directed Project, Student Research Project, or Thesis (top 1%)
 - Student evaluation of the LP courses is fairly positive

Discussion (contd.)

- Debatable points
 - Practice vs. theory of LP
 - the courses focus more on the practical programming side than on the theoretical side of LP
 - this seems to fit the students' interest better
 - theoretical foundations get discussed in the STLP seminar
 - Language(s) taught
 - a single FP+LP language (such as Oz, or Mercury) would avoid a lot of student confusion (mostly syntactic)
 - Prolog and SML have a much larger user/application base, and so will be more likely of use to the graduate
 - switching to a new language requires major investment on the lecturers' part
 - consequently, we stay with the present setup, and try to link the LP and FP topics as much as possible

Future Work

- Electronic Teaching aSsistant
 - stand-alone exercising tool need for students without Internet access
 - set of exercising should be made larger
 - extension of the administrative side to support multiple courses and multiple semesters
- The RDBG execution visualisation tool — several extensions underway
- The Match plagiarism-detection tool
 - front-ends for new languages, such as Cékla
- ...

Acknowledgements

- Péter Hanák, exactly 10 years ago, invited me to teach declarative programming at BUTE.
- Enthusiastic student helpers (we had about 50 of these!) did a lot to make the task of teaching LP and FP possible
- Several students carried the major burden of development of various utilities and tools:
 - András György Békés
 - Tamás Benkő
 - Lukács Tamás Berki
 - Dávid Hanák
 - Gergely Lukácsy
 - Tamás Nepusz
 - Tamás Rozmán