Clock Synchronization in Ethernet



Example network





GPS to ECU4

- Common that systems implement GPS corrections to bring down GPS error from ~10 meters to sub-meter precision, sometimes even centimeter level precision
- Task scheduling and variable Ethernet communication latency can cause an end-to-end delay to consumer application of 100-200 ms
- At 100 km/h, this leads to multiple meters of error
- <u>Solution:</u> GPS application time stamps data before transmitting, the consumer application can make temporal corrections without loss of GPS accuracy



Why do we need to synchronize clocks?

- Time stamp data and events
- Synchronous data acquisition
- Sensor fusion
- Synchronous actuation
 - Collision-imminent steering and braking
 - presentation of audio and video
- Big-data analytics without loss of temporal information
- Establish causal relationships that led to a failure
- Synchronized task execution
- Scheduling of packets on the network (e.g., TDMA)



IEEE Std 802.1AS-2011

- gPTP Generalized Precision Time Protocol
 - Grand Master (GM)
 - Time Slave
 - Time Relay (i.e., Ethernet switches)
- BMCA Best Master Clock selection Algorithm
 - "GM capable" components
 - Distributed algorithm to elect GM
 - Re-elects new GM if current GM disappears



IEEE Std 802.1AS-2011

- gPTP (precision time protocol; profile of IEEE 1588 for full-duplex Ethernet and WiFi)
 - Time sync frames are transmitted by a Grand Master (GM)
 - Time sync frames are time stamped on ingress and egress
 - Time sync frames are modified with corrections based on time measurements and propagated through the network to time slaves
 - Time slaves adjust their time based on received time sync frames
 - Link delay is measured at each port
 - Syntonization: Rate ratio is measured by each endpoint and switch (except GM), because clocks may have different frequencies.
- BMCA (Best Master Clock selection Algorithm)
 - Election process based on Announce frames to select the Grand Master and configure port roles, thereby establishing a path to all time slaves
 - Priorities in Announce frames are configurable



802.1AS at a glance

- Best Master Clock selection Algorithm (BMCA) to build spanning tree
- Alternative: static configuration





802.1AS: Three components



- Mostly not used
- Best clock is fixed
- Unless we have a plug & play network, or we need redundancy



Latency measurement

- Rarely in use (due to fixed wire lengths)
- Latency well known and preprogrammed

Distribution of time



> Always used

IEEE 802.1AS frame formats

- Reserved multicast address: 01-80-C2-00-00E
- Reserved EtherType: 0x88F7
- 7 frame types
- Event frames are time stamped on ingress and egress

gPTP frame type	Function	Class	Value in header
Sync	Time sync	Event	0x0
Follow_Up	Time Sync	General	0x8
Pdelay_Req	Link Delay	Event	0x2
Pdelay_Resp	Link Delay	Event	0x3
Pdelay_Resp_Follow_Up	Link Delay	General	0xA
Announce	BMCA	General	0xB
Signaling	Power Saving (e.g., request reduced frequency in delay measurements)	General	0xC



Rest of this module

- Go through selected parts of Kevin Stanton's tutorial: <u>http://ieee802.org/1/files/public/docs2014/as-</u> <u>kbstanton-8021AS-tutorial-0714-v01.pdf</u>
- Application on an example
- Brief overview of recent amendments to the standard
 - Multiple time domains
 - Multiple synchronization paths
 - Multiple grand masters



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- All other switches and microcontrollers will send Announce messages with priority 255 (i.e., they are slaves; not grand master capable)
- Highest priority Announce message is propagated by switches on all Ra ports for which asCapable == TRUE (inferior Announce messages are discarded)
 - All switches will configure port roles based on reception of Announce, making ECU0 switch the GM
 - Microcontrollers can participate as slaves in BMCA or be statically configured to be slaves
 - If a link fails, switches will reconfigure port roles (Announce messages are sent continuously in BMCA)
 - If the current GM fails, the next best GM capable device will • become GM and the switches will reconfigure port roles (Announce messages are sent continuously)







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2. Port roles established and gPTP operation establishes time sync (a sync tree is created with the ECU0 Ethernet switch as the root)





Example network





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gPTP Sync and FollowUp in operation (ECU0 switch to Radar portion of the sync tree)





Recommended configurations

- Sync and FollowUp frequency: 8 times per second
- PDelay_Req frequency: 3 times in each direction per link per second
- Initial link delay parameters: calibrated (e.g., in automotive, constant wiring length, no plug and play)
- Frequency of Announce frames: 1 per second
- Static priorities and clock qualities defined for each Grand Master capable device.
- asCapable is set to true/false depending on the desired part of the network to participate in 802.1AS



Sync status

- 802.1AS does not define any mechanism to detect errors in the way slaves synchronize to the master
 - Can be done in the application layer by having each slave transmitting their synchronized time



Without BMCA (static port role configuration)









Other ways to implement redundancy

- Ongoing standardization project (mature, likely to be published in 2018)
- Drafts are available
- Redundant sync paths
- Redundant grand masters (hot standby)



Redundant sync paths

- Two redundant synchronization trees from a single GM, with each synchronization tree in > a different gPTP domain
- Time is transported on multiple paths)

INKÖPING

Second sync tree (only) to react faster to sync losses >



Grand Master redundancy

- > One primary GM and one hot-standby GM (which are separated in two gPTP domains)
- > Hot-standby operating: the secondary GM has to be synchronized to the primary GM
- (preselected) second GM will seamless take over the time transport when something happens with first GM
 - > No time needed to select & announce new GM
 - > Reduces grandmaster change over time



GM and synchronization path redundancy

- > Each GM establishes two sync trees, resulting in a total of four sync trees that are separated in four gPTP domains
 - > Preselected 2nd GM
 - > Preconfigured redundant paths





Summary

- There are many applications that need global notion of time
 - GPS, Sensor fusion, Audio and video, Data logging and postprocessing
- 802.1AS provides 1-2 microsecond precision synchronization; network service to the application
 - Election of grand master and establish sync tree
 - Measure delays and rate differences
 - Propagate time frames through sync tree
- Several solutions exist to provide redundant synchronization

