Quality-adaptive Prefetching for Interactive Branched Video using HTTP-based Adaptive Streaming

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LIU EXPANDING REALITY

too sad too violent

too sad too violent too scary

 $\bullet \bullet \bullet$

too sad too violent too scary

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too sad too violent too scary

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YouTube	Interlude	and even books!
tourube	interiude	













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 - Playback stalls and quality fluctuations

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 - Multiple encodings of each fragment (defined in manifest file)
| Chunk1 | Chunk2 | Chunk3 | Chunk4 | Chunk5 |
|--------|--------------------------------------|--|--|--|
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850 Kb/s	Chunk1	Chunk2	Chunk3	Chunk4	Chunk5
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 - Goal: Seamless playback even if user decision at last possible moment

- We develop a simple analytic model which allows us to define the prefetching problem as an optimization problem
 - Maximizes expected playback quality while avoiding stalls
- Based on our findings, we design optimized policies that determine:
 - 1. When different chunks should be downloaded
 - 2. What quality level should be selected for each of these chunks
 - 3. How to manage playback buffers and (multiple) TCP connections such as to ensure smooth playback experience without excessive workahead (buffering)
- The design and implementation of the framework
- Experimental evaluation of our policies, which provide insights into the importance of careful adaptive policies

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Problem: Maximize quality, given playback deadlines and bandwidth conditions



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Objective function

maximize playback quality



Objective function

maximize
$$\sum_{i=1}^{n_e} q_i l_i + \sum_{i=n_e+1}^{n_e+|\mathcal{E}^b|} q_i l_i$$












































































current segment











- Playback deadlines
 - for seamless playback without stalls



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 - Current segment: e.g., 2 and 3



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• Playback deadlines

Download

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- First chunks next segment: e.g., 4, 7, and 10 $t_i^c \le t_i^d = \tau + \sum_{j=1}^{n_e} l_j$ if $n_e < i \le n_e + |\mathcal{E}^b|$ Time at which branch point is reached

Download completion times



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• Download times t_i^c , rate estimations, and parallel connections

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- At download completion
 - Decide number of chunks to download next
 - Decide quality level of chunks
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 - Maximize expected weighted playback
- Exponential number of candidate schedules
- Our optimized policies restrict the number of candidate schedules to consider
 - Policies differ in number of candidate schedules and how aggressive they are (in choosing qualities)

Policy	Connections	Schedules considered	Objective
All schedules	1≤c _i ≤C ^{max}	Q ^M , where M=n _e + $ \xi_b $ -m	-
Optimized non- increasing quality	1≤c _i ≤C ^{max}	(M+Q-1 Q-1	$\sum^{n_e} \alpha I + \sum^{n_e + \xi_b } \alpha I$
Optimized maintainable quality	1≤c _i ≤C ^{max}	Q	<u>Z</u> <i>Y i i i i n</i> _e +1

- Total number of schedules: Q^M
- Optimized non-increasing quality:
 - Constraint: Qualities of consecutive chunks are non-increasing
- Optimized maintainable quality:
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- Single connection: baseline comparing to policies which do not use multiple connections
- Greedy bandwidth: bandwidth aggressive as opposed to aggressive quality choices
- Naïve: benchmark to regular branched video players



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Greedy bandwidth	1≤c _i ≤C ^{max}		$\sum_{i=j}^{j+m} q_i l_i$

- Single connection: baseline comparing to policies which do not use multiple connections
- Greedy bandwidth: bandwidth aggressive as opposed to aggressive quality choices
- Naïve: benchmark to regular branched video players



Test Scenario

2 3 4 5 Worst case scenario always pick the last segment at last possible moment

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Test Scenario Worst case scenario always pick the last segment at last possible moment

- Default scenario:
 - Chunks per segment: 5
 - Branches per branch point: 4
 - Branch points: 3



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- Default scenario:
 - Chunks per segment: 5
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- Results are averages over 30 experiments



Naïve policy: does not perform prefetching

- Stalls at every branch point
- Note: High playback rate is misleading on its own



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Optimized maintainable quality provides best tradeoff



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- Much lower stall probability
- Tradeoff is somewhat lower playback rate



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• Optimized non-increasing quality is more aggressive

- Higher playback rate
- More stalls


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- Optimized non-increasing quality is more aggressive
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- Single connection does not use parallel connections
 - Good (slightly higher) playback rate
 - Much more stalls



- Single connection does not use parallel connections
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- Single connection does not use parallel connections
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- Greedy bandwidth aggressively grabs bandwidth
 - Lower playback rate
 - More stalls



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Impact of Round-trip Times (RTTs)



- Quality decreases with larger RTTs
 - Playback rate decrease with RTT
 - Stall probability increase with RTT

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Impact of Segment Lengths





- Quality increases with more chunks per segment
- Very many stalls if segments are too short



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- Stalls frequent when too many branch options
 - Single connection struggles the most



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- Stalls frequent when too many branch options
 - Single connection struggles the most

Impact of Competing Flows





- Player adapts playback rate based on competing traffic
- Parallel connection polices see increased benefits when competing traffic
 - E.g., Single connection policy has much more stalls when competing flows



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- Most HAS players perform ON-OFF switching based on two buffer thresholds: T_{min} and T_{max}
- If buffer > T_{min} - Start playback
- If buffer > T_{max} Suspend download
- If buffer $< T_{min}$ Resume download



• Most HAS players perform ON-OFF switching based on two buffer thresholds: T_{min} and T_{max}

ON

3

OFF

ON

OFF

ON

OFF

ON

9

10

Time_nd

- If buffer > T_{min} Time, - Start playback Video
 - If buffer > T_{max}
 Suspend download
 - If buffer $< T_{min}$
 - Resume download

buffer > T_{min}

Download

Playback

- Most HAS players perform ON-OFF switching based on two buffer thresholds: T_{min} and T_{max}
- If buffer > T_{min} Time_nd Time - Start playback Video ON OFF OFF ON OFF ON ON • If buffer > T_{max} Suspend download Download 9 10 3 • If buffer $< T_{min}$ Playback Resume download

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- If buffer > T_{min} Time_nd Time - Start playback Video OFF ON OFF ON OFF • If buffer > T_{max} Suspend download Download 6 9 10 • If buffer $< T_{min}$ Playback Resume download buffer > T_{max} buffer > T_{min} buffer > T_{min}

 How to handle workahead when video contains branches?


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Example Buffer occupancy (s) Time (s)

- Designed and implemented branched video player that achieve seamless streaming without playback interruptions
- Designed optimized policies that maximize playback quality while ensuring sufficient workahead to avoid stalls
- Evaluation shows that solution effectively adapt quality levels and number of parallel connections so as to provide best possible video quality, given current conditions

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Software: http://www.ida.liu.se/~nikca/mm14.html







Quality-adaptive Prefetching for Interactive Branched Video using HTTP-based Adaptive Streaming





Vengatanathan Krishnamoorthi, Niklas Carlsson, Derek Eager, Anirban Mahanti, Nahid Shahmehri

Software: http://www.ida.liu.se/~nikca/mm14.html

www.liu.se