YouTube Popularity Dynamics and Third-party Authentication

Niklas Carlsson

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Keynote at the 10th IEEE Workshop on Network Measurements (IEEE WNM @LCN), Nov. 2016



YouTube Popularity Dynamics, Edge Caching, Third-party Authentication, and Interactive Videos Streaming

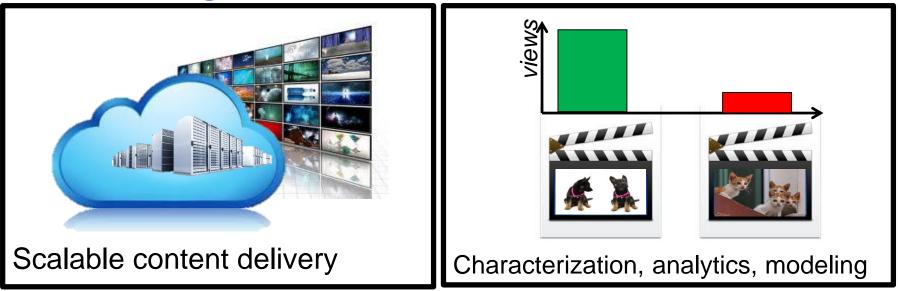
"some topics I am very excited to talk about" ...

The work here was in collaboration ...

- Including with students (alphabetic order):
 - Youmna Borghol (NICTA, Australia)
 - Vengatanathan Krishnamoorthi (Linköping University, Sweden)
 - Siddharth Mitra (IIT Dehli, India)
 - Anna Vapen (Linköping University, Sweden)
- ... and non-student collaborators (alphabetic order):
 - Martin Arlitt (HP Labs, USA, and University of Calgary, Canada)
 - György Dan (KTH, Sweden)
 - Derek Eager (University of Saskatchewan, Canada)
 - Anirban Mahanti (NICTA, Australia)
 - Nahid Shahmehri (Linköping University, Sweden)

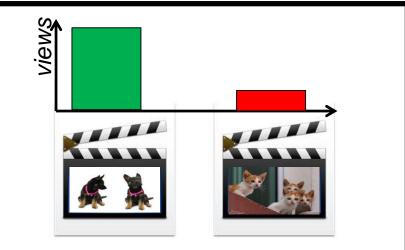




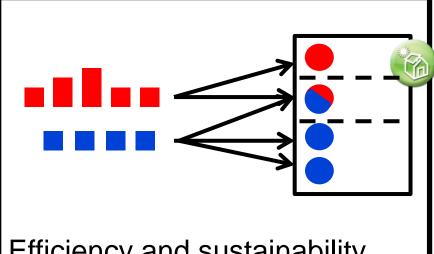




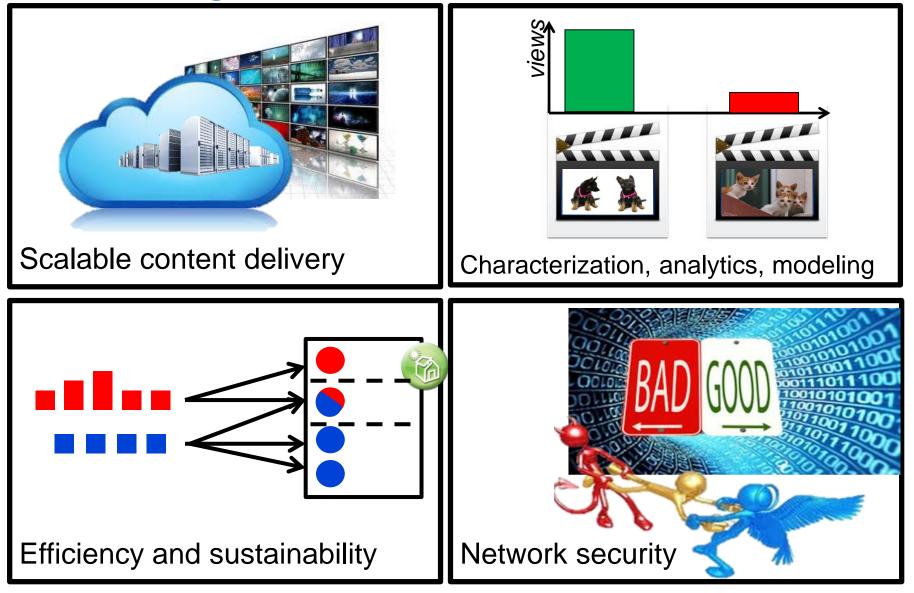
Scalable content delivery

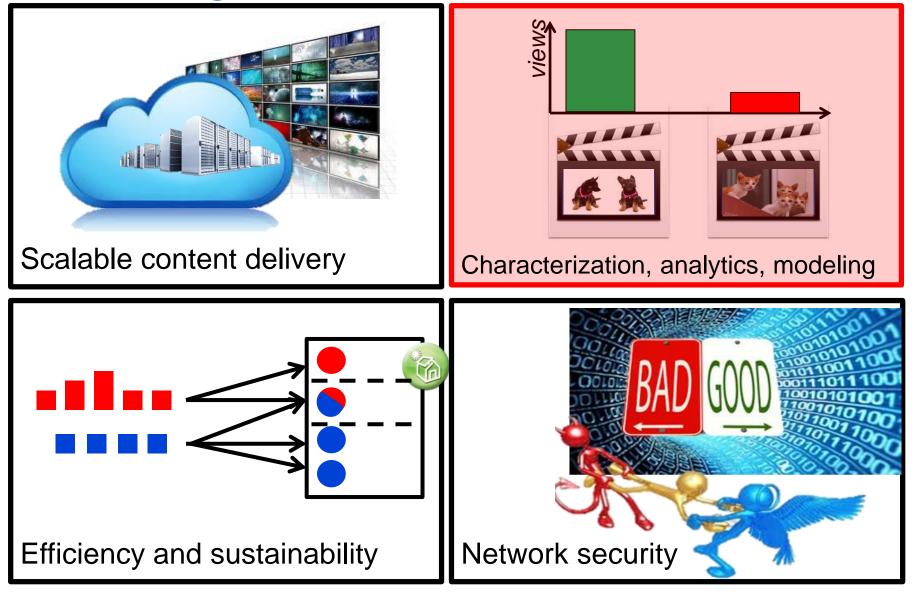


Characterization, analytics, modeling



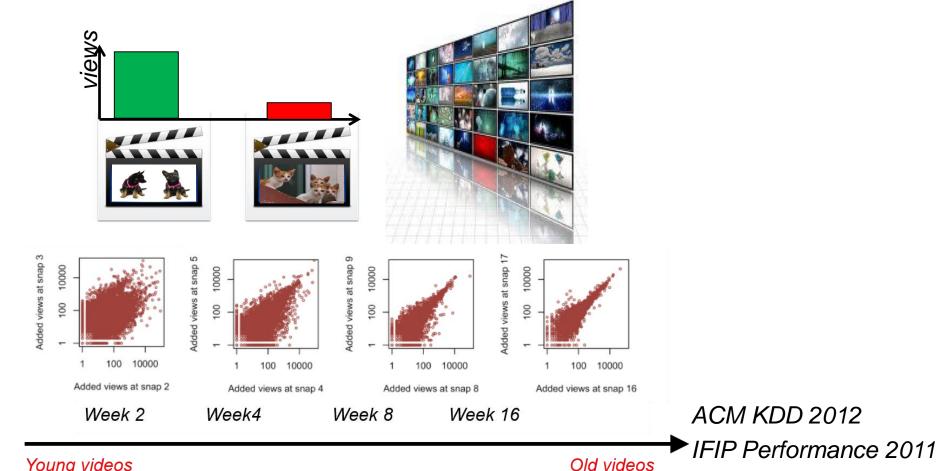
Efficiency and sustainability





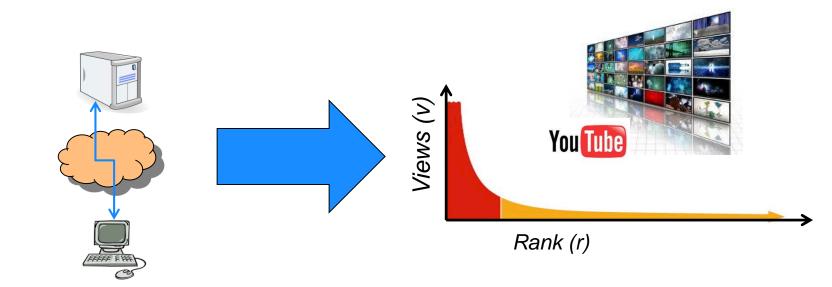
In this talk ...

... model+understand popularity ...



Young videos

... popularity dynamics and caching ...

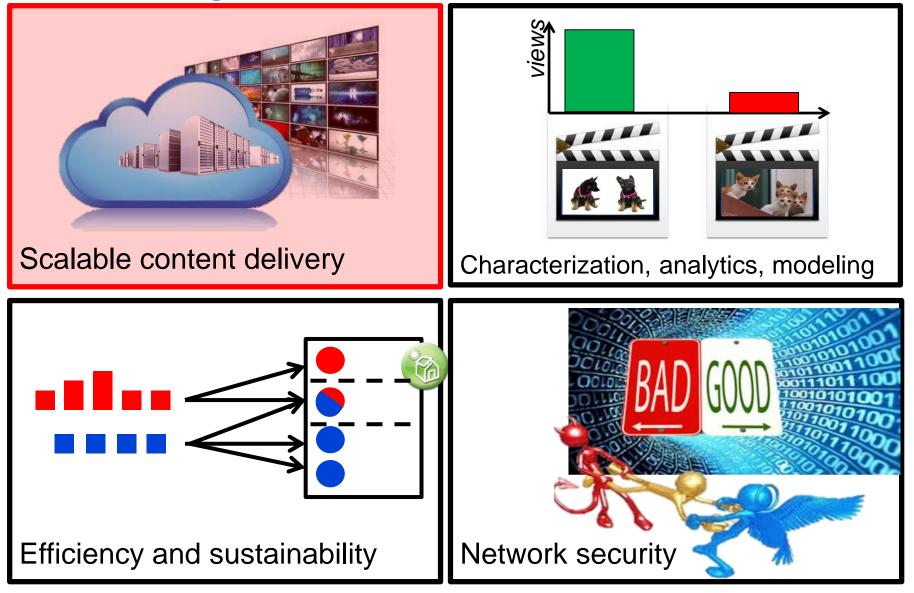


IEEE TPDS 2016/2017 * (*accepted last month)

... third-party authentication ...



IEEE IC 2016 IFIP SEC 2015 PAM 2014



... innovative new streaming media ...



ACM MM 2015 ACM MM 2014 ACM CCR 2013

So let's start ...



Video streaming landscape



Video streaming landscape



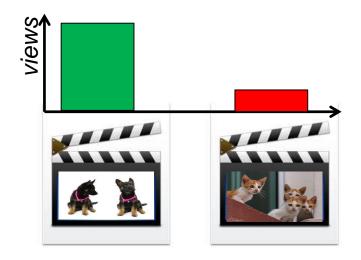
Video streaming landscape



Motivation

- Streaming services contribute to over 60% of the global Internet traffic currently
- By 2020, this share is expected to be over 80%
- Systems need to be well understood, scalable, and efficient to match growth projections





The Untold Story of the Clones: Content-agnostic Factors that Impact YouTube Video Popularity

Proc. ACM SIGKDD 2012.

Characterizing and Modeling Popularity of User-generated Videos *Proc. IFIP PERFORMANCE* 2011.



 Video dissemination (e.g., YouTube) can have widespread impacts on opinions, thoughts, and cultures

Motivation

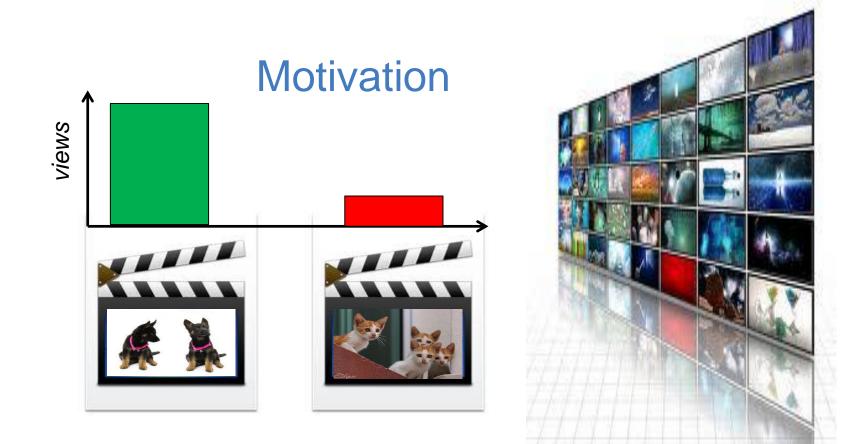




 Not all videos will reach the same popularity and have the same impact



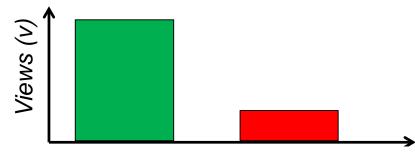
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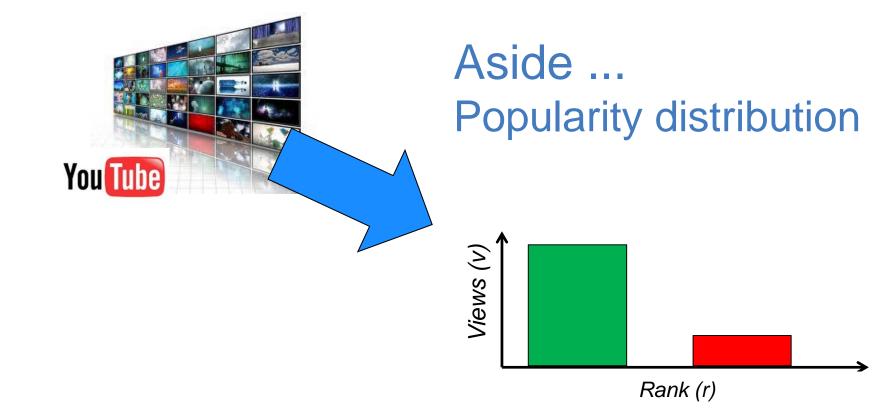
Aside ... Popularity distribution

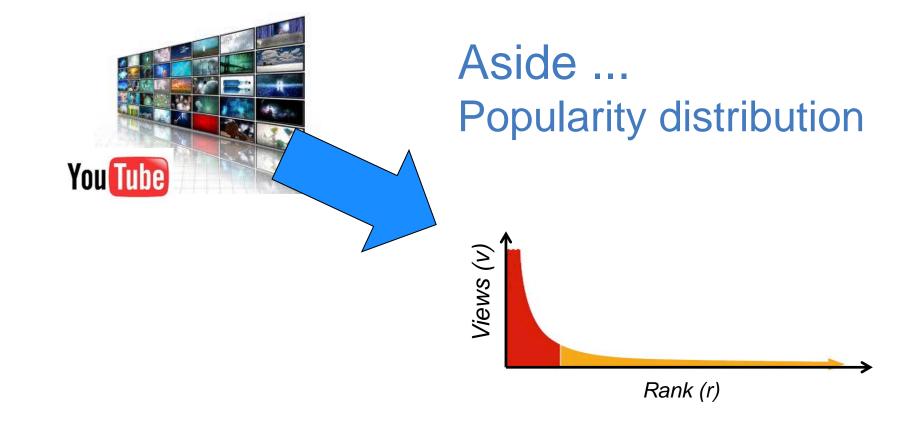


Rank (r)

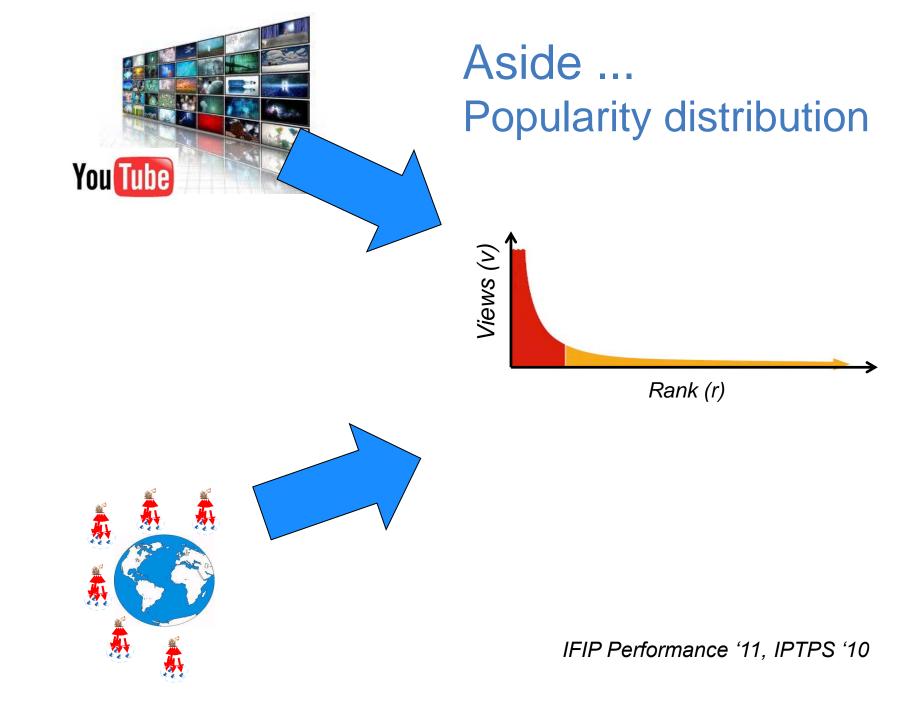


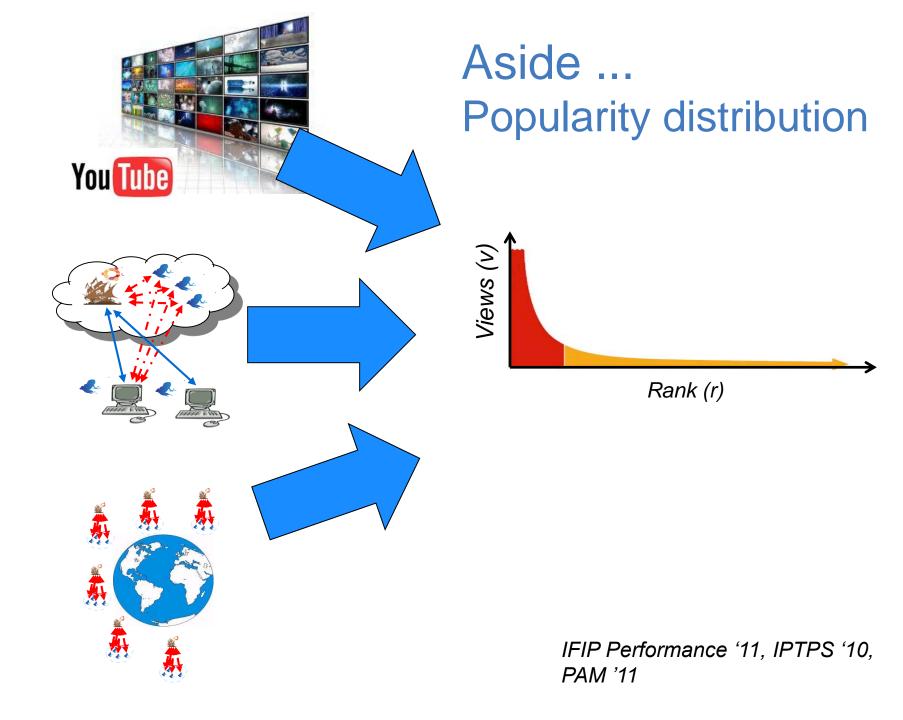


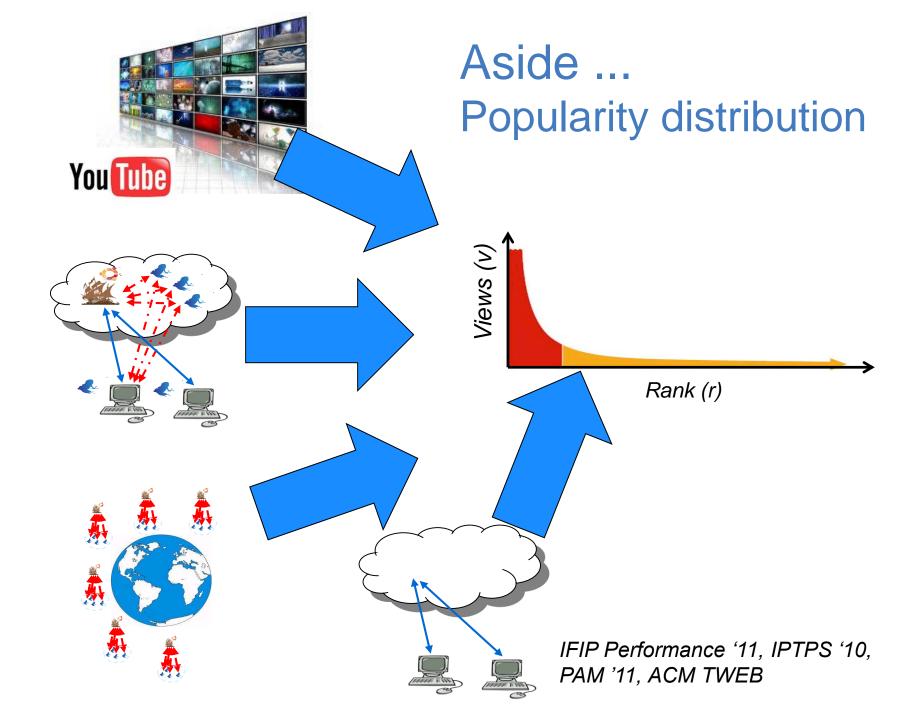


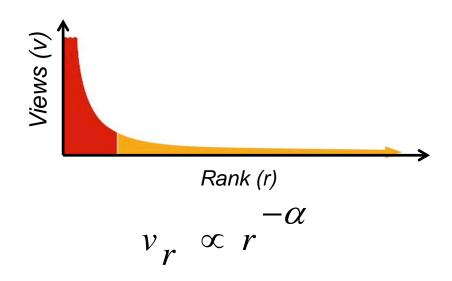


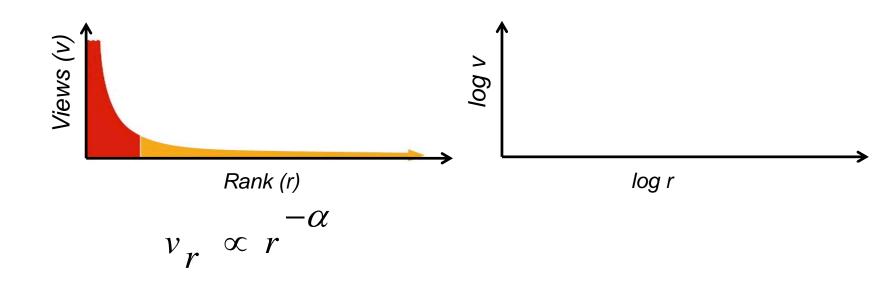
IFIP Performance '11

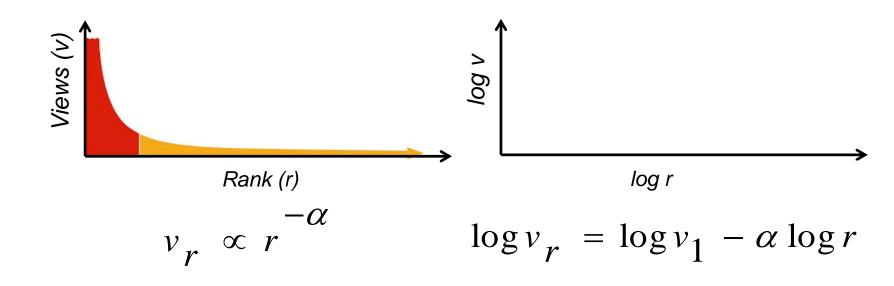


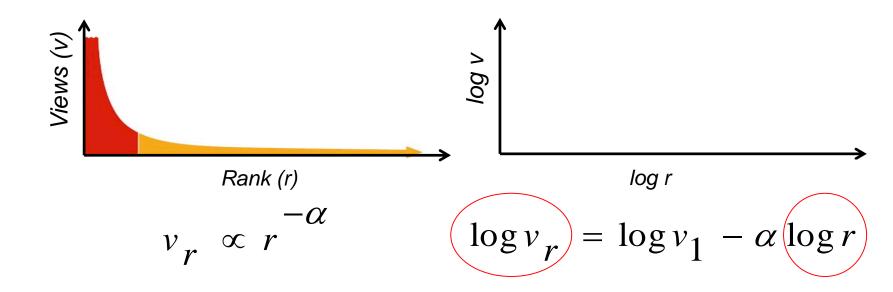


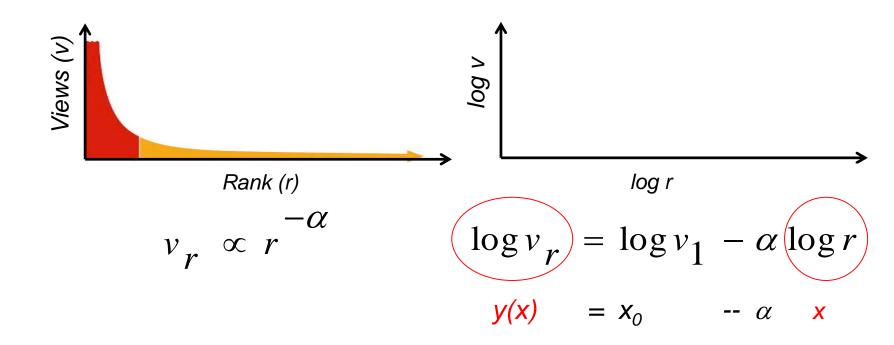


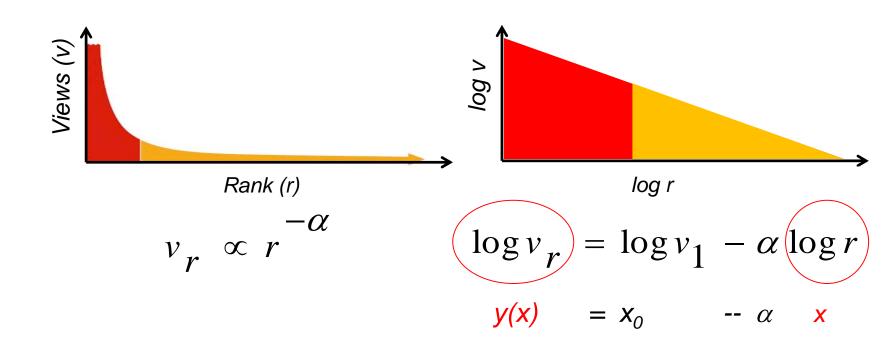


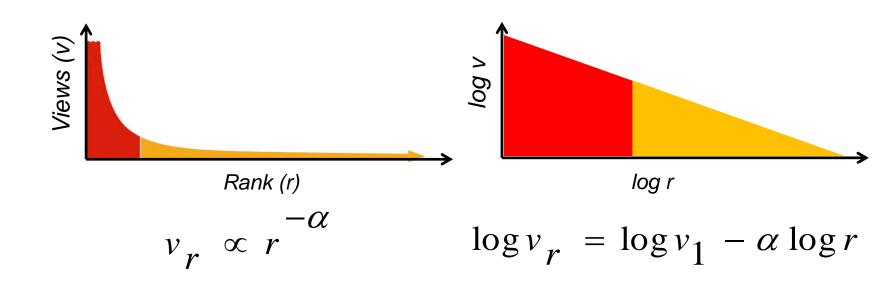


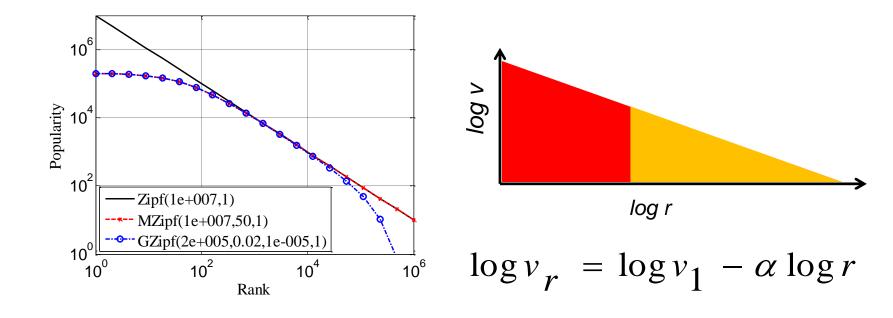


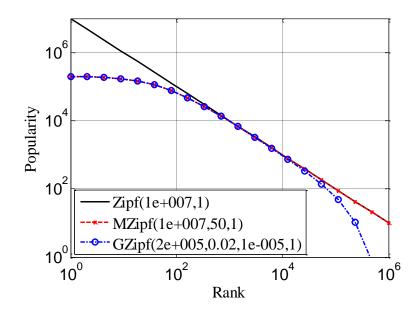




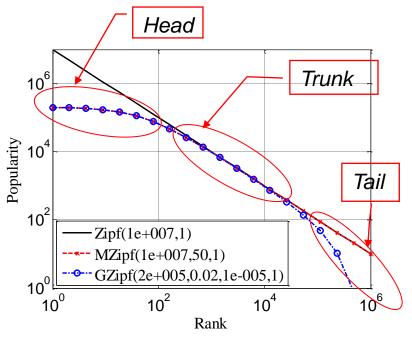








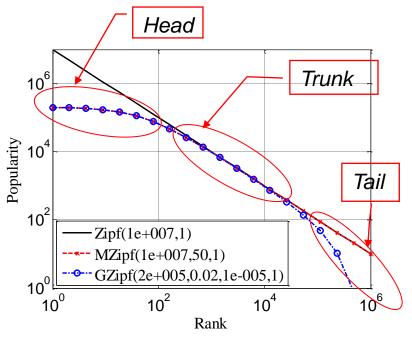
E.g., ACM TWEB, PAM '11 IFIP Performance '11, IPTPS '10



- Popularity distribution statistics and models
 - Across services (impact on system design)
 - Lifetime vs current
 - Over different time period (churn)
 - Different sampling methods

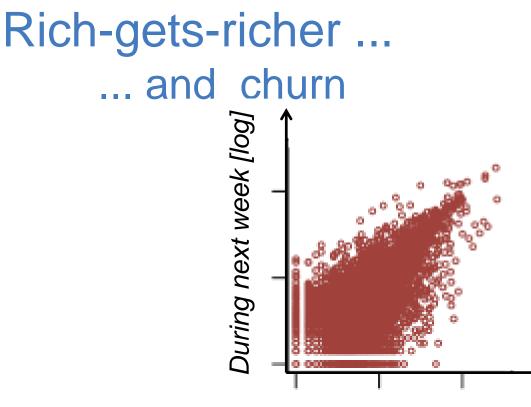
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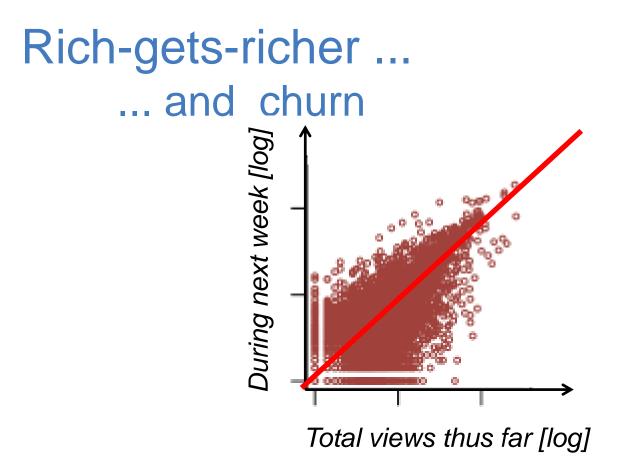
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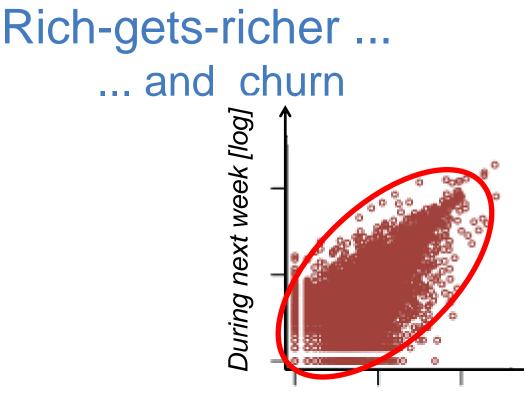


Total views thus far [log]

E.g., IFIP Performance '11



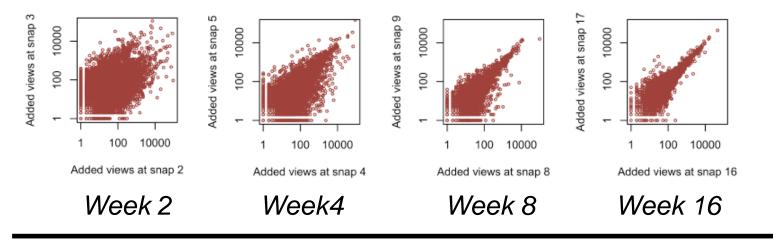
The more views a video has, the more views it is likely to get in the future



Views during week [log]

- The more views a video has, the more views it is likely to get in the future
- The relative popularity of the individual videos are highly non-stationary

Rich-gets-richer and churn



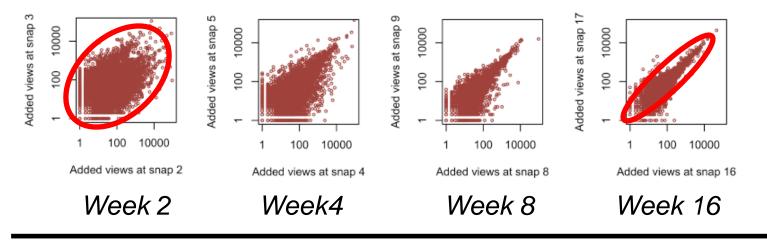
Young videos

Old videos

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E.g., IFIP Performance '11

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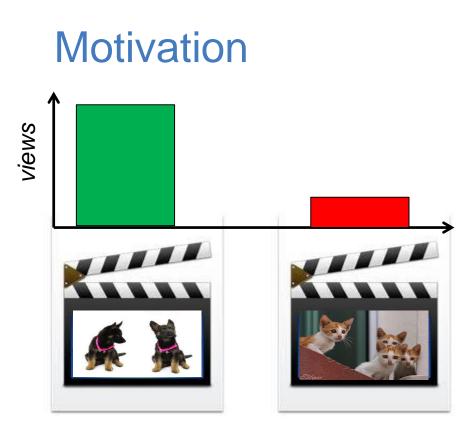


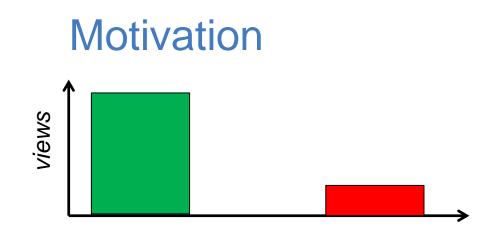
Young videos

Old videos

- The more views a video has, the more views it is likely to get in the future
- The relative popularity of the individual videos are highly non-stationary
- Some long-term popularity

E.g., IFIP Performance '11





• Some popularity differences due to content differences

Motivation

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- But also because of other "content-agnostic" factors
 - The latter factors are of considerable interest but it has been difficult to accurately study them

Motivation

- Some popularity differences due to content differences
- But also because of other "content-agnostic" factors
 - The latter factors are of considerable interest but it has been difficult to accurately study them

In general, existing works **do not** take content differences into account ... (e.g., large number of rich-gets-richer studies)

 Develop and apply a methodology that is able to accurately assess, both qualitatively and quantitatively, the impacts of various content-agnostic factors on video popularity

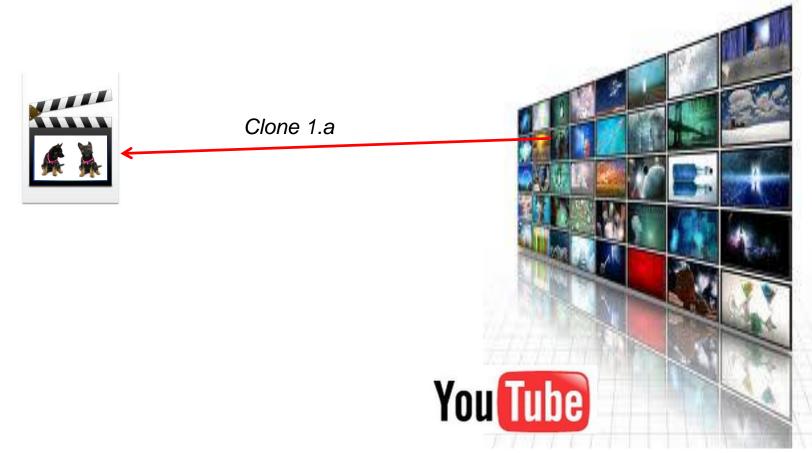
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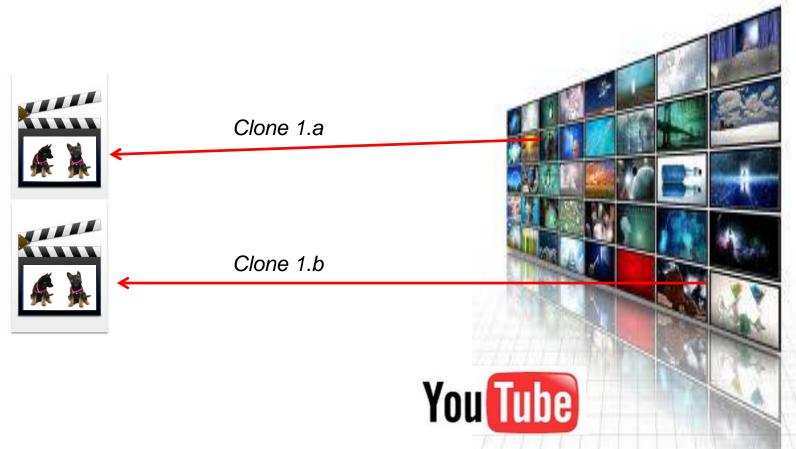
- Clones
 - Videos that have "identical" content (e.g., same audio and video track)



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- Clones
 - Videos that have "identical" content
- Clone set
 - Set of videos that have "identical" content •

You Tube



Clone set 1

- Clones
 - Videos that have "identical" content
- Clone set
 - Set of videos that have "identical" content





You Tube

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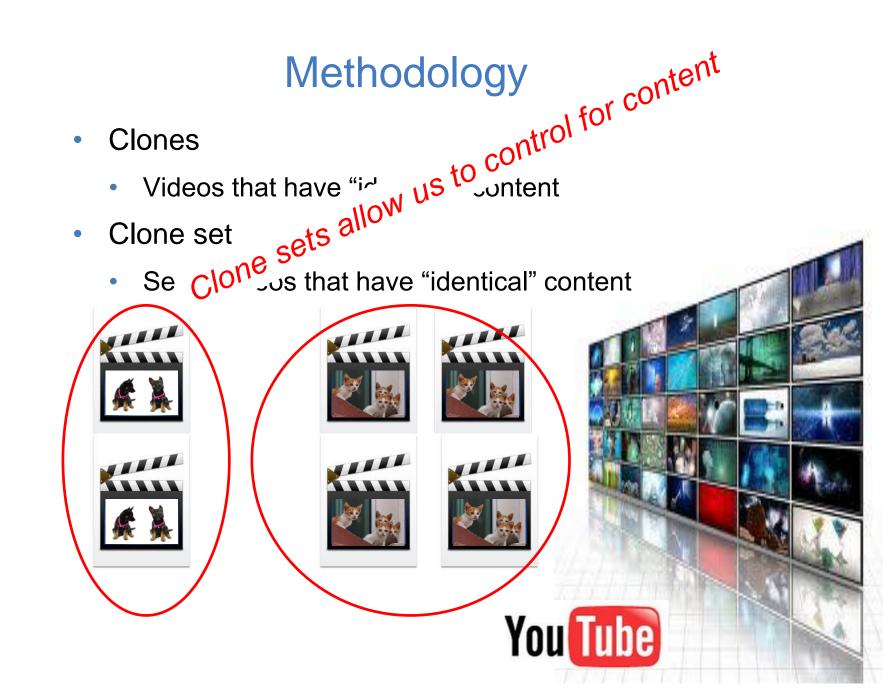




You Tube

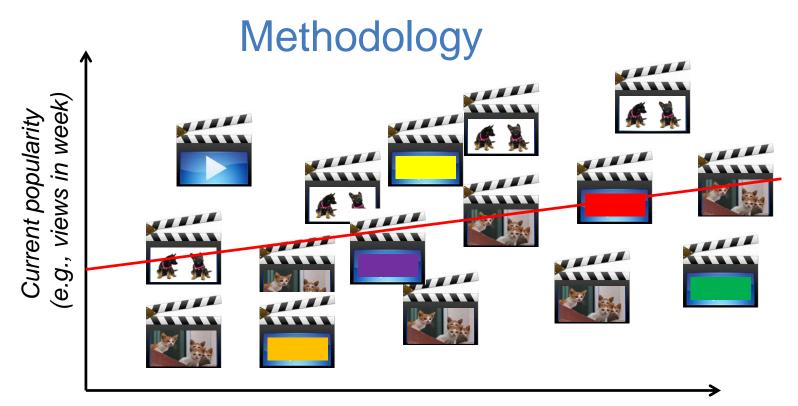
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- Analyze how different factors impact the current popularity while accounting for differences in content
 - 1) Baseline: Aggregate video statistics (ignoring clone identity)
 - 2) Individual clone set statistics
 - 3) Content-based statistics



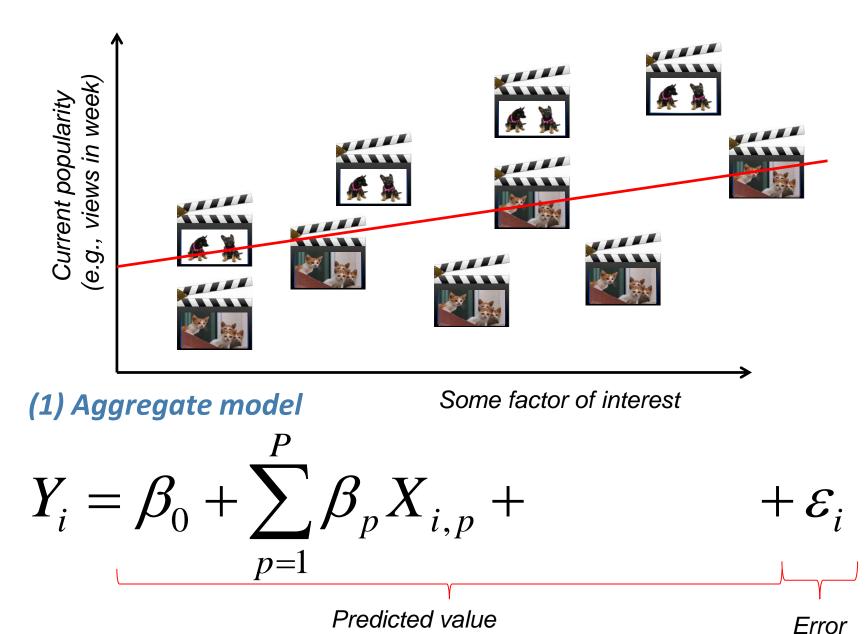
Some factor of interest



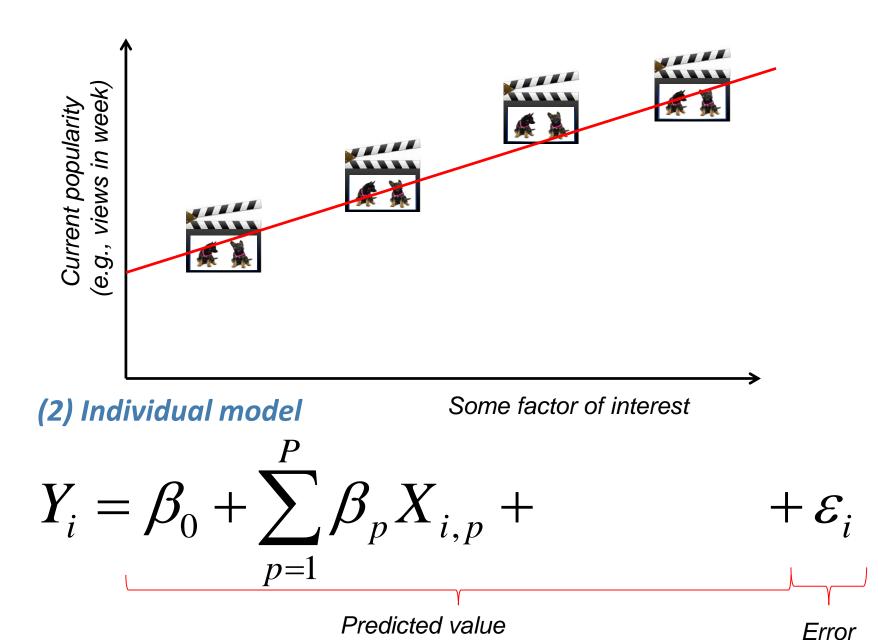
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Focus on clone sets

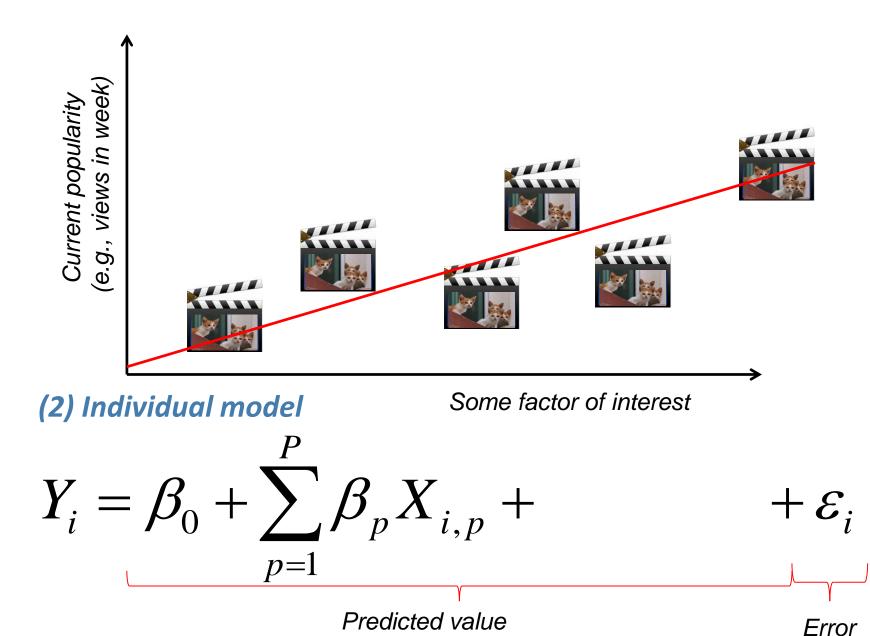
Methodology: (1) Aggregate model



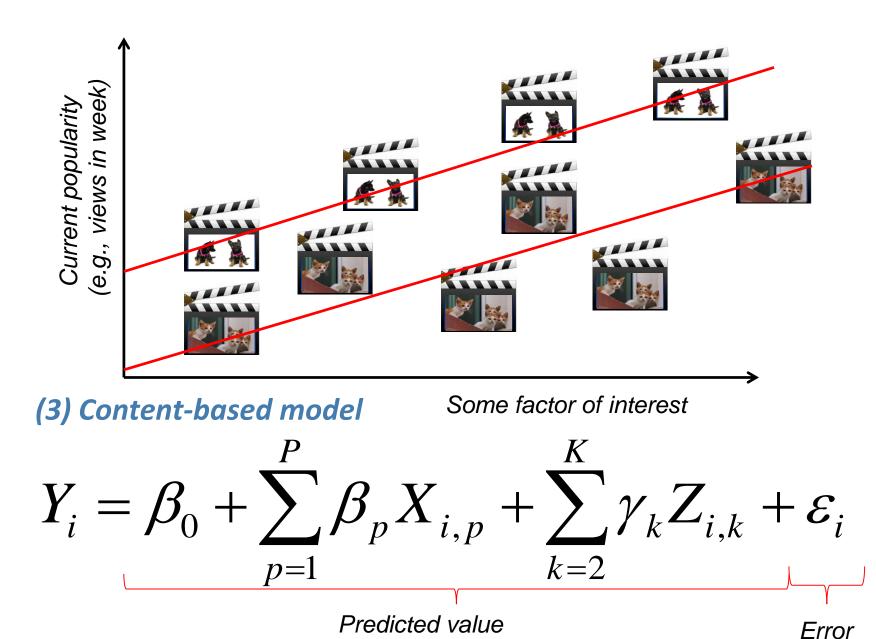
Methodology: (2) Individual model



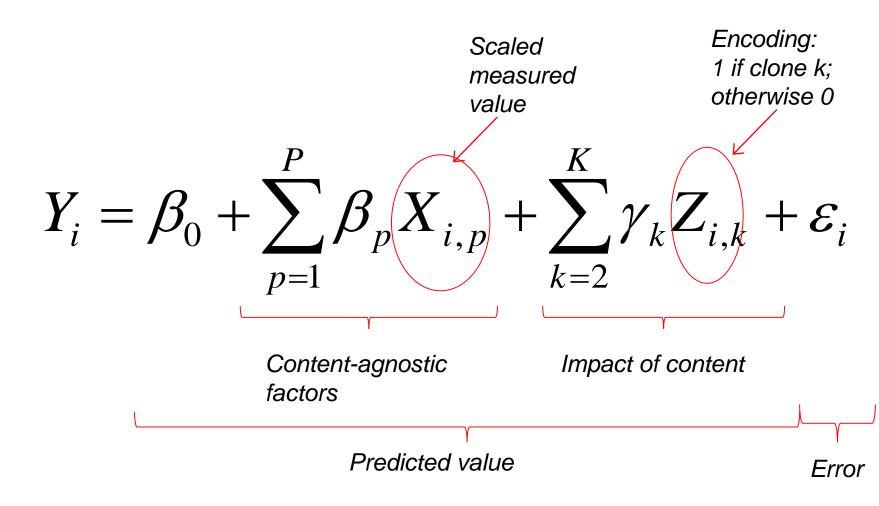
Methodology: (2) Individual model



Methodology: (3) Content-based model



Methodology: (3) Content-based model



Data collection

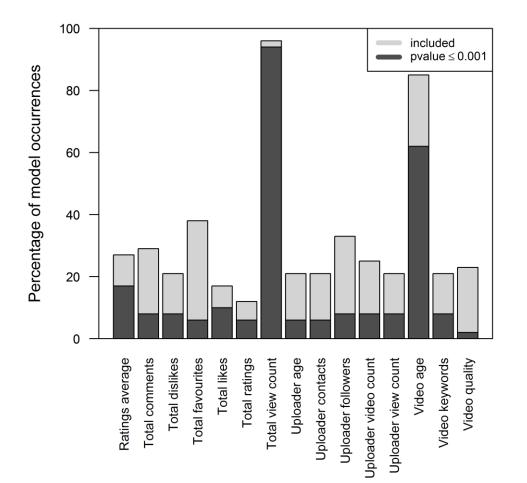
- Identified large set of clone sets
 - 48 clone sets with 17 94 videos per clone set (median = 29.5)
 - 1,761 clones in total
- Collect statistics for these sets (API + HTML scraping)
 - Video statistics (2 snapshots \Rightarrow lifetime + weekly rate statistics)
 - Historical view count (100 snapshots since upload)
 - Influential events (and view counts associated with these)

Analysis approach

- Example question: Which content-agnostic factors most influence the current video popularity, as measured by the view count over a week?
- Use standard statistical tools
 - E.g., PCA; correlation and collinearity analysis; multi-linear regression with variable selection; hypothesis testing
- Linearity assumptions validated using range of tests and techniques
 - Some variables needed transformations
 - Others where very weak predictors on their own (but in some cases important when combined with others!!)

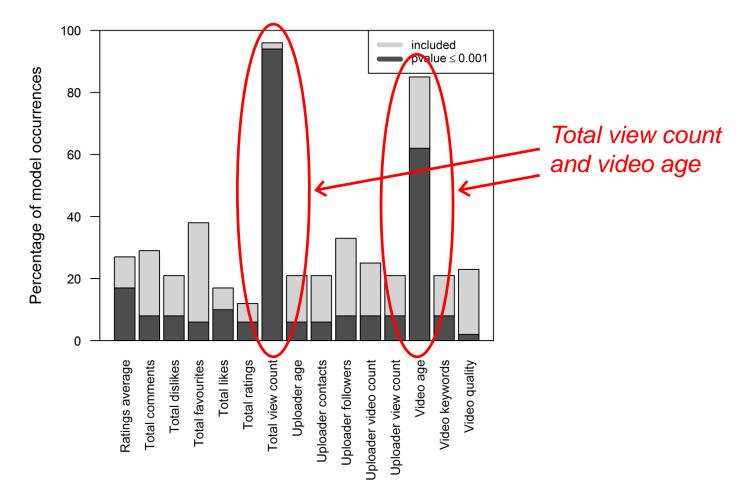
Which factors matter?

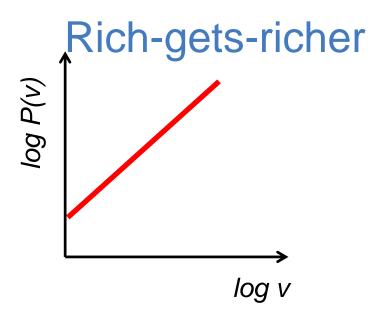
 Using multi-linear regression with variable reduction (e.g., best subset with Mallow's Cp)

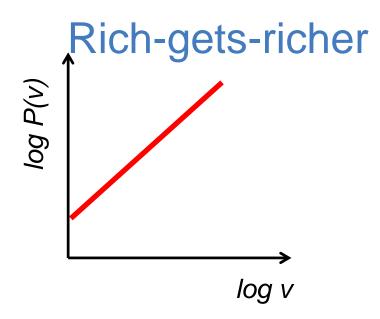


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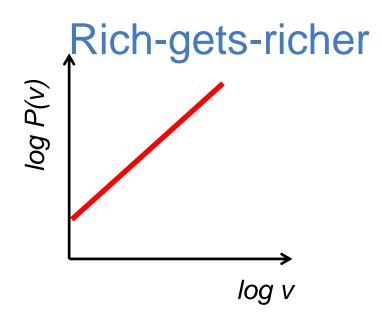
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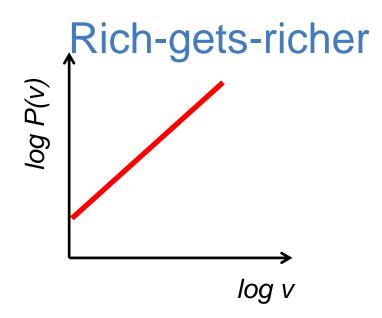




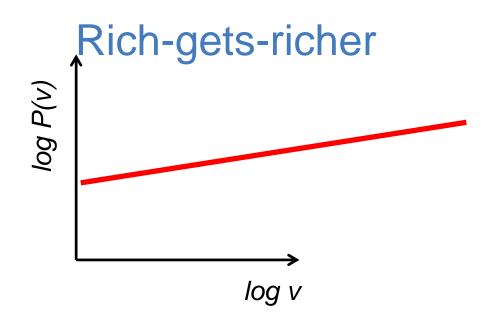
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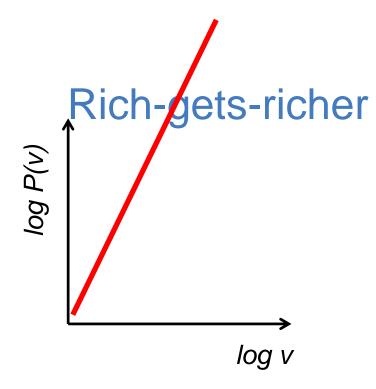
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 - Sub-linear: α < 1 (the rich may get richer, but at a slower rate)
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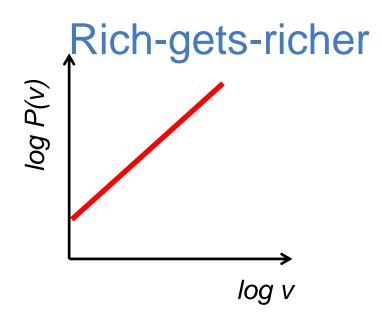
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	α	σ	90%	95%		H ₀ : α≥1	
Individual							
Content-based							
Aggregate							

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Content-based	1.003	-0.014	0.98-1.027	0.976-1.031	0.81	0.59	0.4
Aggregate	0.932	-0.016	0.906-0.958	0.901-0.963	REJECT	REJECT	1

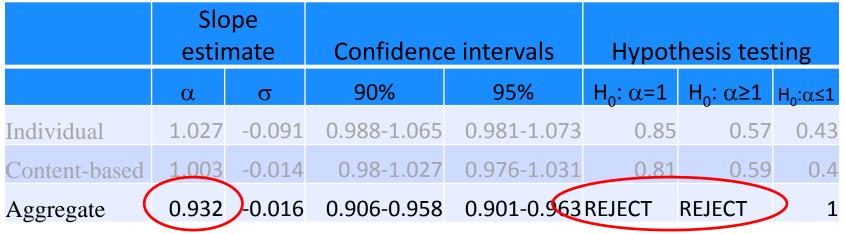
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- If not accounting for content, sub-linear preferential attachment

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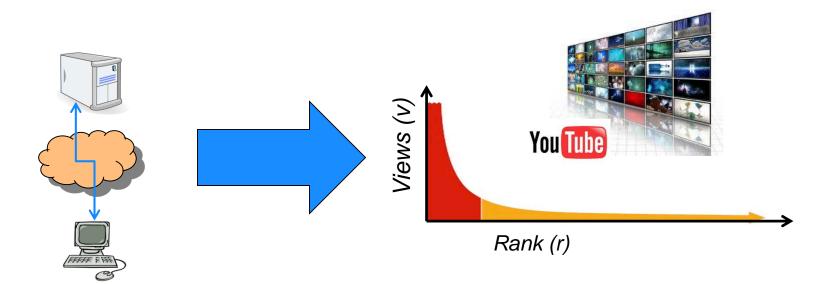
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Clone lessons ... (ACM SIGKDD 2012)

- Develop and apply a clone set methodology
 - Accurately assess (both qualitatively and quantitatively) the impacts of various content-agnostic factors on video popularity
- When controlling for video content, we observe a strong linear ``rich-get-richer'' behavior
 - Except for very young videos, the total number of previous views the most important factor; video age second most important
- Our findings also confirm that inaccurate conclusions can be reached when not controlling for video content



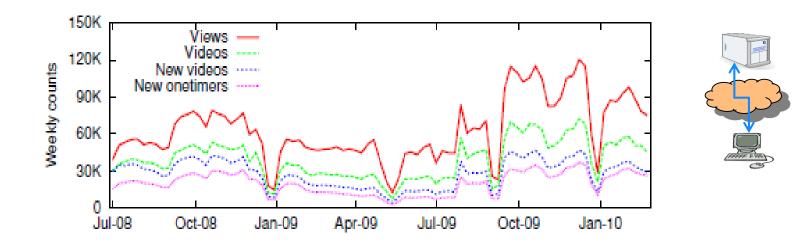
Ephemeral Content Popularity at the Edge and Implications for On-Demand Caching

IEEE Transactions on Parallel and Distributed Systems (IEEE TPDS), 2016.

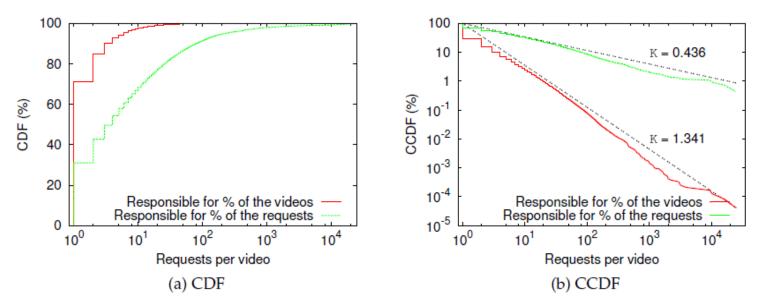
Motivation and observations



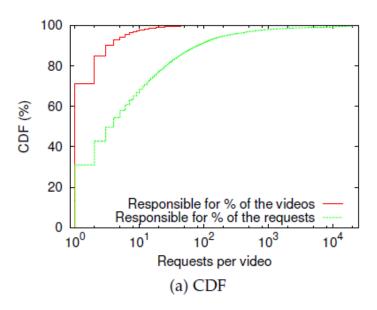
- Ephemeral content popularity seen with many content delivery applications
 - At edge this results in many "one timers" (a.k.a. "one hit wonders")
 - Makes indiscriminate on-demand caching highly inefficient, since many items added to the cache will not be requested again



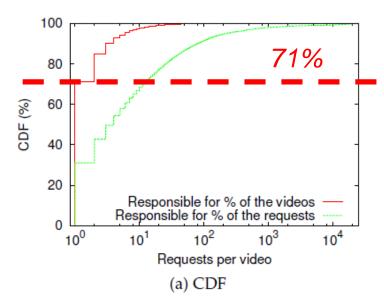
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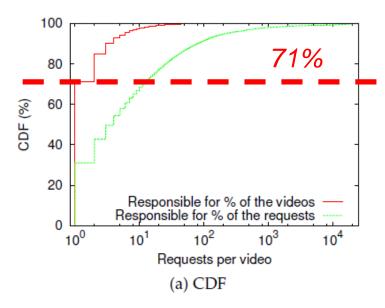
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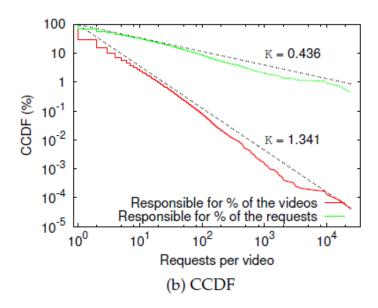
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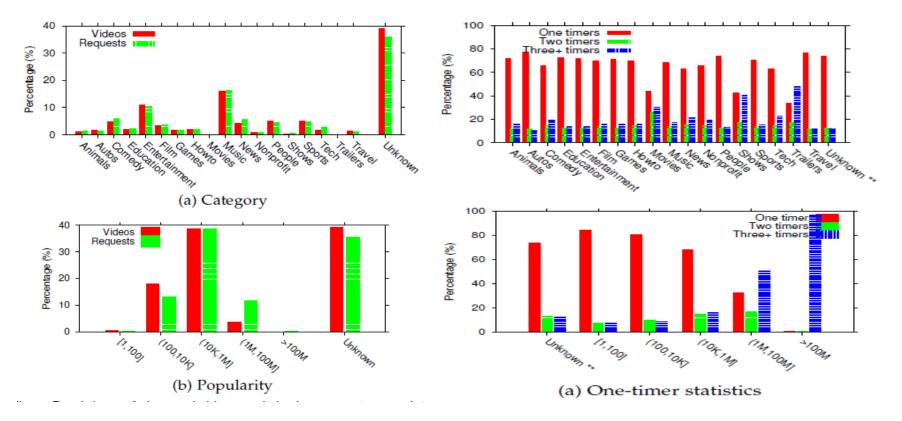


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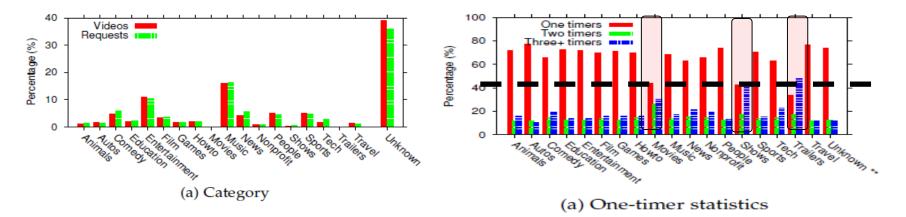
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 - 71% of the requested videos are "one-timers"
 - Demonstrate the need for selective caching policies
 - Popularity follow power law (and Zipf)

Characterizing of "one timers"



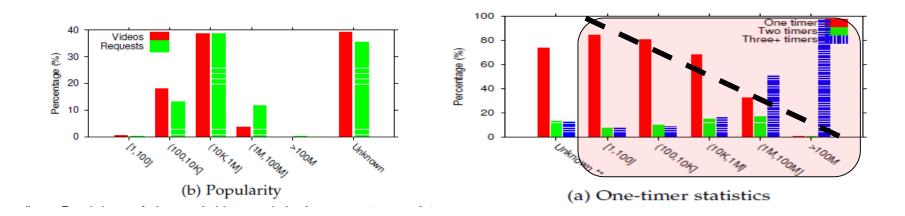
 Using meta data about these videos, we take a closer look at one-timers and other videos receiving few views

Characterizing of "one timers"



- Using meta data about these videos, we take a closer look at one-timers and other videos receiving few views
 - Fewer one-timers among movies, shows, and trailers
 - Strong (negative) correlation between global popularity and one-timers

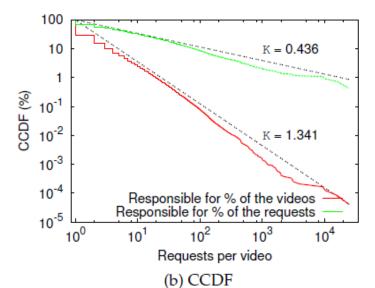
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	Scale parameter estimation				
Time period	Video basis	Request basis			
1 week	3.31 ± 0.02	2.036 ± 0.005			
1 month	2.958 ± 0.006	1.773 ± 0.002			
2 months	2.818 ± 0.004	1.693 ± 0.001			
6 months	2.625 ± 0.002	1.5635 ± 0.0005			
1 year	2.452 ± 0.001	1.4690 ± 0.0003			
All	2.341 ± 0.001	1.4359 ± 0.0003			

Trace duration dependence in power-law fitting.



Motivated by our power-law characterization and fittings, we use a Zipf model

$$P(\text{cache miss}) = 1 - \frac{\sum_{i=k}^{\infty} i^{-\alpha}(i-k)}{\sum_{i=1}^{\infty} i^{-\alpha+1}},$$
$$P(\text{cache insertion}) = \frac{\sum_{i=k}^{\infty} i^{-\alpha}}{\sum_{i=1}^{\infty} i^{-\alpha+1}}.$$

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 - Cache on kth request policy
 - Lower bound "oracle" policies
 - Exact knowledge (exact number of views)

$$P(\text{cache insertion}) = \frac{\sum_{i=k+1}^{\infty} i^{-\alpha}}{\sum_{i=1}^{\infty} i^{-\alpha+1}},$$
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 - Oracle with limited knowledge
 - Binary knowledge (above or below X views)
 - Knows total views, if more than X
 - Knows total views, if less than X

$$P(\text{cache miss}) = 1 - \frac{\sum_{i=k}^{\infty} i^{-\alpha} (i-1) - (k-1) \sum_{i=k}^{X-1} i^{-\alpha}}{\sum_{i=1}^{\infty} i^{-\alpha+1}}.$$
(9)

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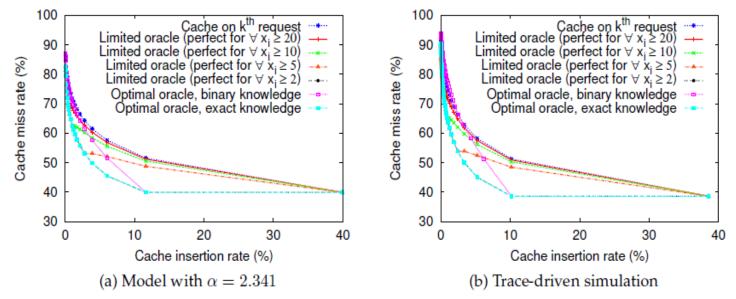
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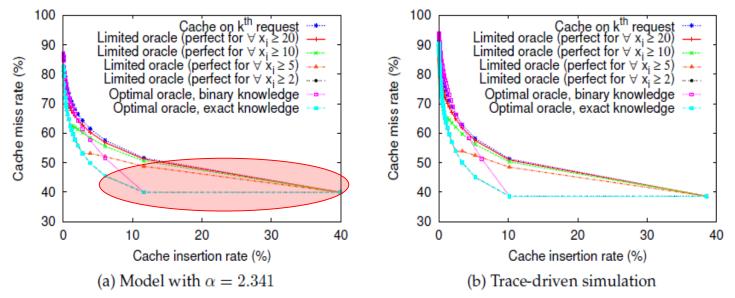
$$P(\text{cache miss}) = 1 - \frac{\sum_{i=k}^{\infty} i^{-\alpha}(i-1)}{\sum_{i=1}^{\infty} i^{-\alpha+1}}, \quad if X \ge k$$

$$P(\text{cache miss}) = 1 - \frac{\sum_{i=k}^{\infty} i^{-\alpha}(i-k)}{\sum_{i=1}^{\infty} i^{-\alpha+1}}, \quad otherwise.$$

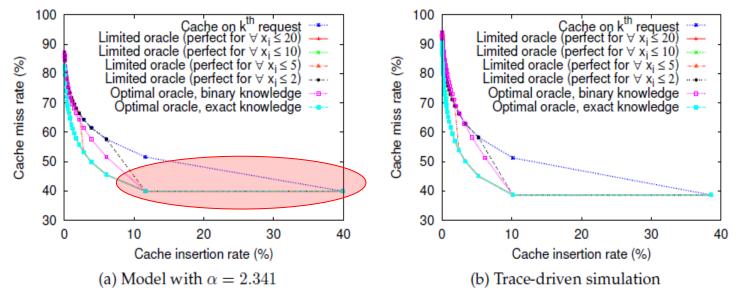
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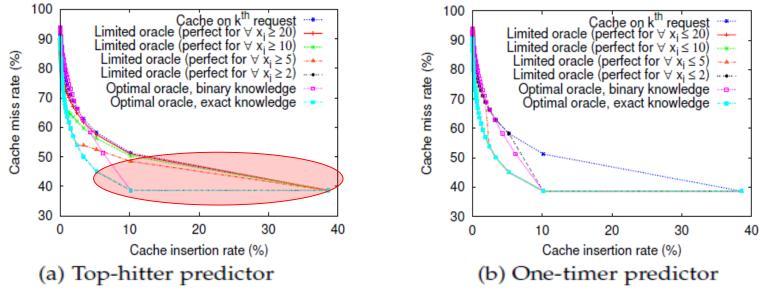
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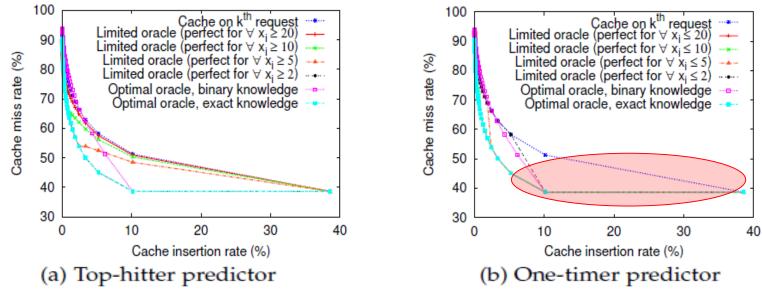


- Evaluation using both model and traces
 - Similar results
- Limited knowledge
 - Noticeable gap if only knows total for videos with more than X
 - Smaller gap if can predict one-timers (and ones with few views)



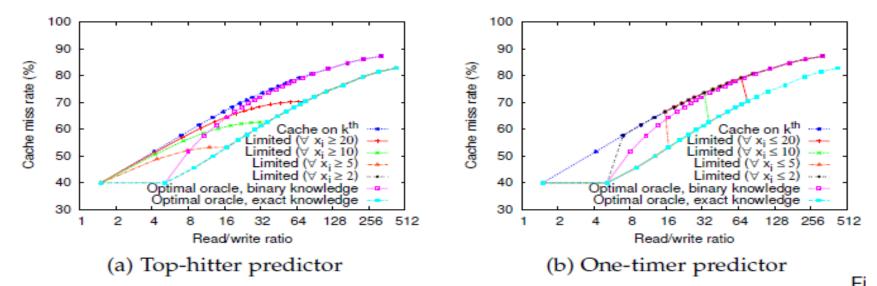
Gap suggest room for improvement

Fi



- Gap suggest room for improvement
 - One-timer prediction may close the gap

Fi



- Gap suggest room for improvement
 - One-timer prediction may close the gap
- Also looked at SSD scenario
 - Read/write ratio vs cache miss rate

Leverage biases in the probabilities that a request will be a one-timer

- Leverage biases in the probabilities that a request will be a one-timer
 - Characterized the one-timers and their request patterns (see paper)

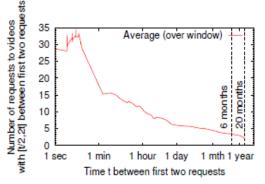


Fig. 15. The expected number of requests $E[x_i|t]$ for videos whose first inter-request time is at least t/2 and at most 2t, plotted as a function of the logarithmic mid-point t.

E.g., Inter-request time dependence ...

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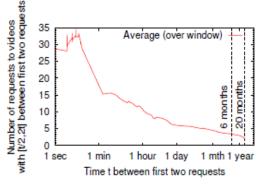


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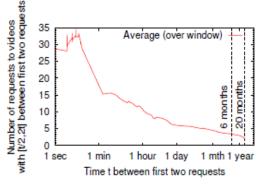


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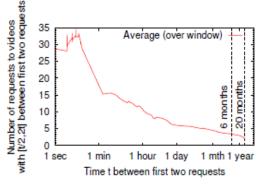


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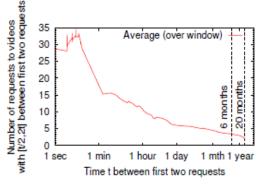


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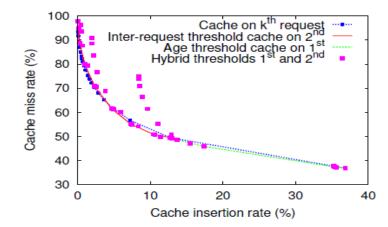


Fig. 22. Summary of the cache performance tradeoffs of the two threshold-based policies, for all videos with known age.

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 - Trace-driven analysis

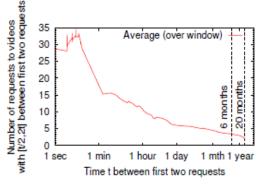


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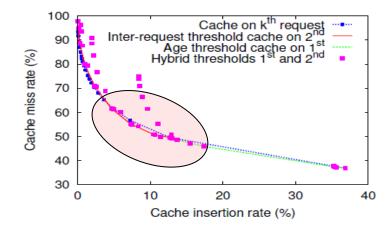


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 - Some small improvements (but still a large gap ...)

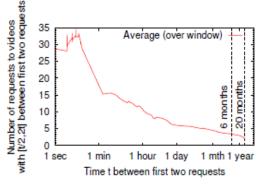


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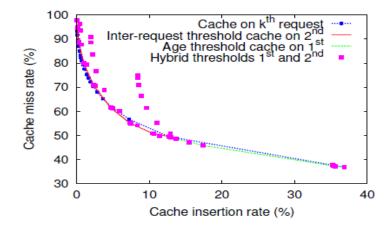


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 - Trace-driven analysis
 - Model to give delimiting insights for case when accurate prediction only possible for a subset of videos

Lessons for edge caching (TPDS paper)

- Collected and analyzed a longitudinal edge dataset
 - All YouTube video accesses over a 20-month period
 - Most videos receive few view (e.g., 71% one-timers)
 - Requests per video accurately modelled using power-law distribution
- Use novel workload model and trace-driven simulations to study the performance of alternative edge caching policies
 - Cache on kth request found able to greatly reduce the cache insertion rate, at the cost of relatively modest increases in cache miss rate
- Assess the potential room for improvements through use of content characteristics
 - Oracles suggest there is room for substantial improvements
 - However, would require the prediction of the number of future requests to the content items that are the least popular
 - This problem is both difficult and not well explored, as most research has focused on predicting the most popular contents ...



A Look at the Third-Party Identity Management Landscape

IEEE Internet Computing, 2016.

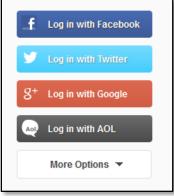
Information Sharing and User Privacy in the Third-party Identity Management Landscape

Proc. IFIP SEC 2015

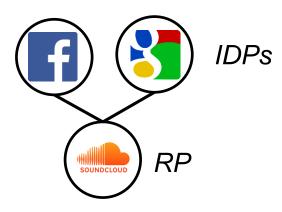
Third-party Identity Management Usage on the Web, Proc

Proc. PAM 2014

Third-party Web Authentication

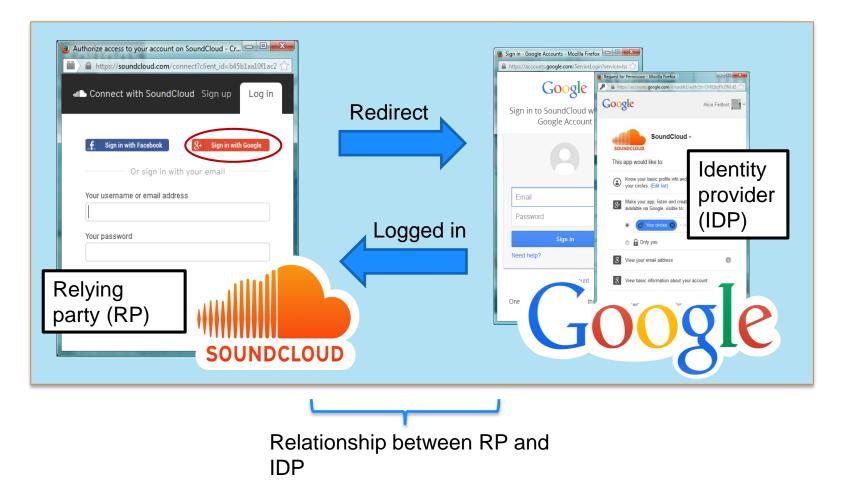


IDPs



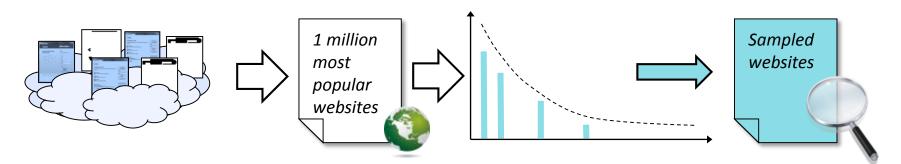
- Use an existing **IDP** (identity provider) account to access an **RP** (relying party)
- Fewer logins
 - Stronger authentication can be used
- Information sharing between websites
 - Privacy leaks!

Third-party Authentication Scenario



Large-scale Crawling

- Popularity-based logarithmic sampling
 - 80,000 points uniformly on a logarithmic range
 - Pareto-like distribution
 - Capturing data from different popularity segments



Large-scale Crawling

- Selenium-based crawling and relationship identification
- Able to process Web 2.0 sites with interactive elements
- Low number of false positives
- Validation with semi-manual classification and text-matching

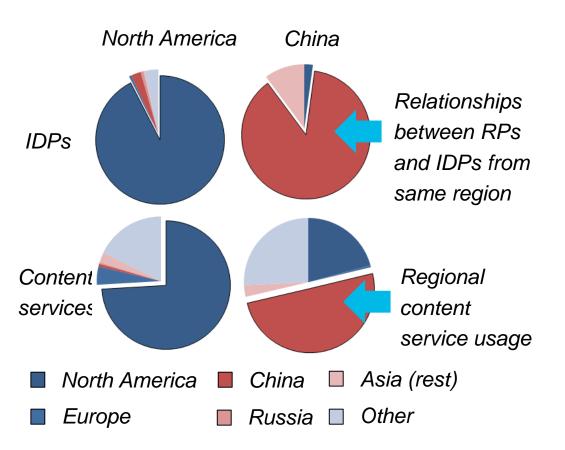


IDPs vs Content Delivery Services

Content providers:

Import images, scripts etc. from other sites (third-party content providers)

IDPs are much more popular sites than content providers.



Service-based Analysis

Likely to be RPs



News, file sharing, info

Likely to be IDPs



Social/portal

Using IDPs from the social/portal category



File sharing, info

Early adopters, using several IDPs



Video, tech

Using IDPs from their own category



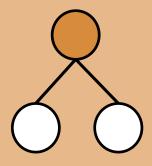
Commerce, tech

Not RPs or IDPs



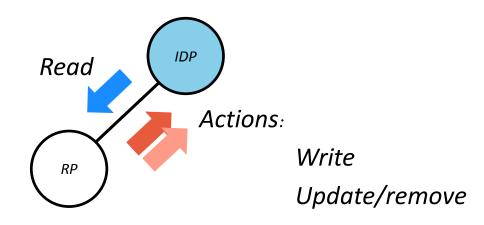
Ads, CDN

Third-parties and Privacy Risks



App Rights and Information Flows

Request for Permission - Mozilla Firefox	- • ×
A https://accounts.google.com/o/oauth2/auth?zt=ChRQI	ozFIcDN1d2: 🏠
Google Alice F	edtest
SoundCloud ~ This app would like to:	
Know your basic profile info and list of people in your circles. (Edit list)	1
Make your app, listen and creative activity available via Google, visible to:	0
Only you	
8 View your email address	0
8 View basic information about your account	0
SoundCloud and Google will use this information in accordance respective terms of service and privacy policies.	with their ccept
App rights examp	le



- Data being sent
- Risks related to
 - Data types
 - Combinations of types

Our Studies on Privacy Risks

- Categorization app-rights data
 - Manual study on the top 200 most popular websites
- Targeted login tests

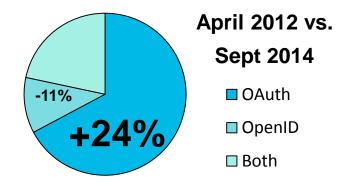
- Longitudinal analysis of privacy risks
 - 200 websites over three years



Privacy risks

Protocol Selection

- OpenID
 - Authentication protocol
 - Decreasing in popularity
- OAuth
 - RP may write/update info on IDP
 - Rich user data is shared
 - Increasingly popular



IDP Selection

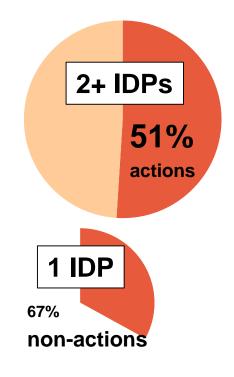
- Top 200 April 2012: 69 RPs and 180 relationships
- Same sites, April 2015: **+15** RPs and **+33** relationships
- Many pairs and triples of popular IDPs
 - 75% of these RPs are selecting all their IDPs from the top 5 most popular IDPs





Risk Types

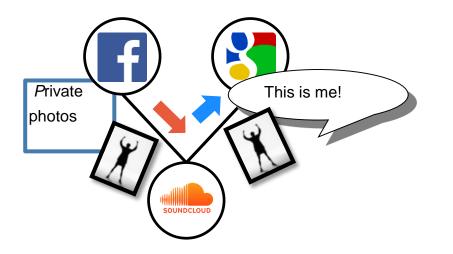
- Only a few relationships in the most privacy preserving category
- 2+ IDPs: More than half are using actions
 - Dangerous when having several IDPs
 - Potential multi-IDP leakage





News and file sharing RPs: most frequent users of actions

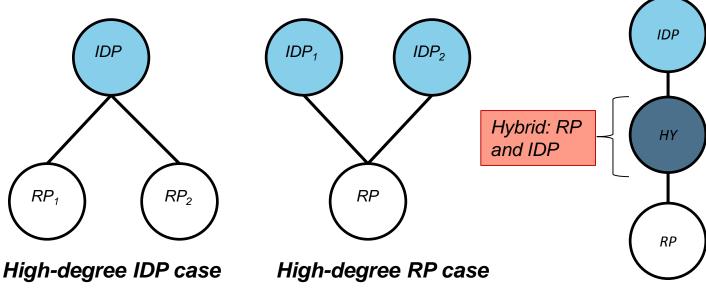
Multi-account Information Risks



Connecting several IDPs to an RP

- Cross account leakage
- Unwanted combinations of conflicting information
- RPs handle multi-IDP usage badly

Structures in the RP-IDP Landscape



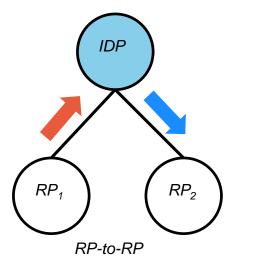
- IDP having many RPs •
- Top IDPs •

- RP having many IDPs •
- Specialized IDPs

Hybrid case

Hybrids are both RP and IDP

RP-to-RP Leakage Example



RP-to-RP leaks	February 2014		April 2015	
IDP	All	Severe	All	Severe
Facebook	645	150	473	66
Twitter	110	110	110	110
Google	91	0	91	0

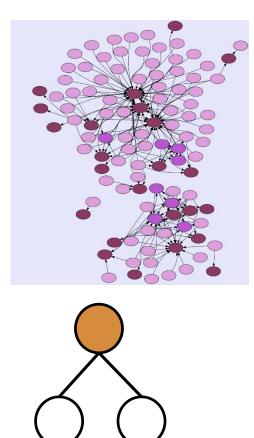
Dataset with 44 RPs using Facebook, 14 using Twitter and 12 using Google

- Potential RP-to-RP leaks
 - Data posted to IDP from RP1
 - Data read from IDP to RP2

Contributions and Findings

- Large-scale RP-IDP study + methodology
 - Categorization of RP-IDP relationships
- Longitudinal changes in the RP-IDP landscape
 - Protocol analysis
 - Privacy risks and information sharing

• Simple web authentication often lack in user privacy





Quality-adaptive Prefetching for Interactive Branched Video using HTTP-based Adaptive Streaming

Proc. ACM Multimedia 2014.

Empowering the Creative User: Personalized HTTP-based Adaptive Streaming of Multi-path Nonlinear Video *Proc. ACM FhMN@SIGCOMM 2013.* (Also in *ACM CCR*). **Best paper award**

Bandwidth-aware Prefetching for Proactive Multi-video Preloading and Improved HAS Performance *Proc. ACM Multimedia 2015.*

Motivation

- Content personalization and personalized streaming
 - Regular web content is dynamic and personalized, while videos have remained largely unchanged
 - Viewer's tastes vary significantly
 - Personalized streaming is relatively unexplored and several interesting questions remain open





too sad too violent

too sad too violent too scary

• • •

too sad too violent too scary

• • •

... or where we may have wanted our favorite character to make a different choice...

too sad too violent too scary

•••

... or where we may have wanted our favorite character to make a different choice...





	Weekday	Shower 🔶 Est 🔮	Work 🔶	Lunch	Backhome 🔶	
Beginday 🕂		Dressup		Skiplunch 🔶		
		Moresleep 🕒		Watch TV		Dinner
	Holiday	Wakeup 🔶	Cooklunch	Visit friend +	Goout	

Allow user to selects between multiple storylines or alternative endings

	Weekday	Shower 🔹	Work 😍	Lunch	Backhome 🔶	
Beginday 💽		Moresteep	Cooklunch	Skip lunch 🔶 Watch TV 🔶		Dinner 🕒 Sleep 🗭
	Holiday 💽	Wakeup	Cooklunch	Visit friend	Goout	

Allow user to selects between multiple storylines or alternative endings

Weekday	0	Shower Eat	0	Work 💽		Lunch	0					
		Dressup	0	Work		Skip lunch	0	Backhome	0			
		Moresleep	0			Watch TV	0			Dinner	Sleep	•
Holiday	0	Wakeup	0	Cooklunch 🔶	Z	Visit friend	0	Goout	0			

Allow user to selects between multiple storylines or alternative endings

Weekday	0	4	Shower	0	Work	0	Lunch	0	Backhome	•			
			Dressup	0			Skip lunch	0	Backhome	0			
			Moresleep	0			Watch TV	0			Dinner	Sleep	0
Holiday	0				Cooklunch	0	Visit friend	0	Goout	0			
		X	Wakeup	0			Visit meno	0					

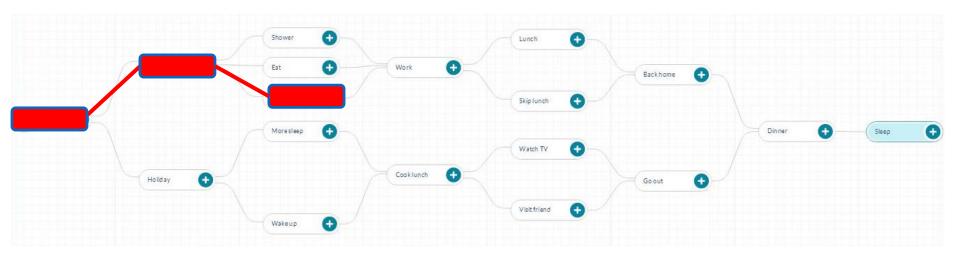
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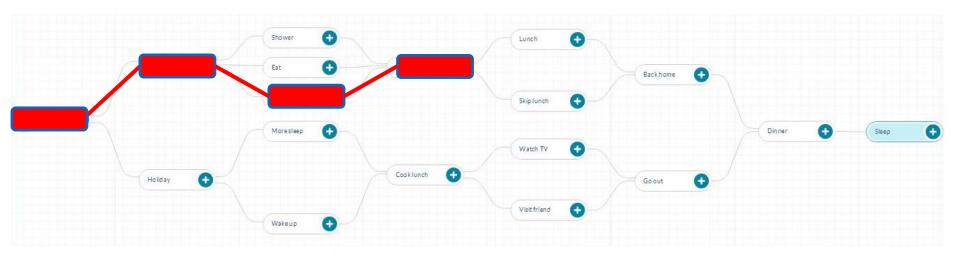
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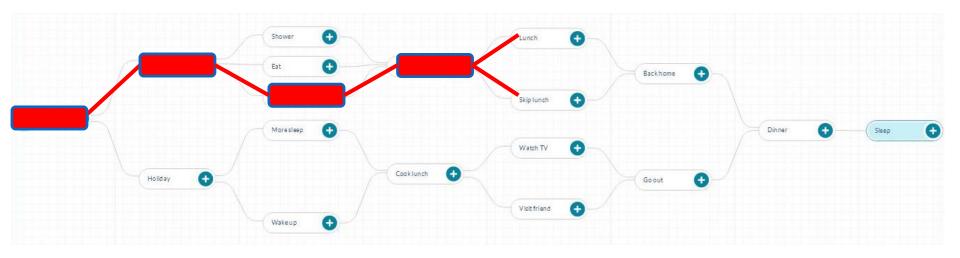
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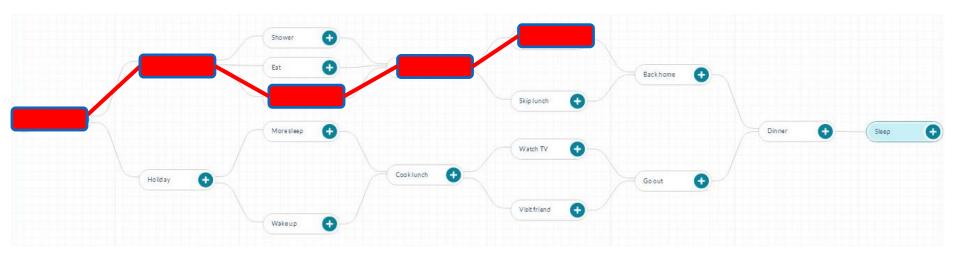
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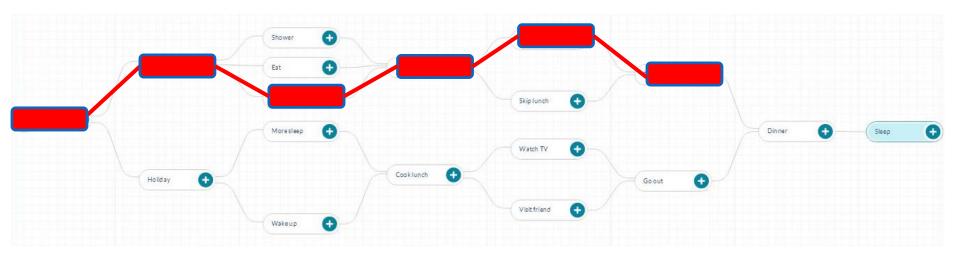
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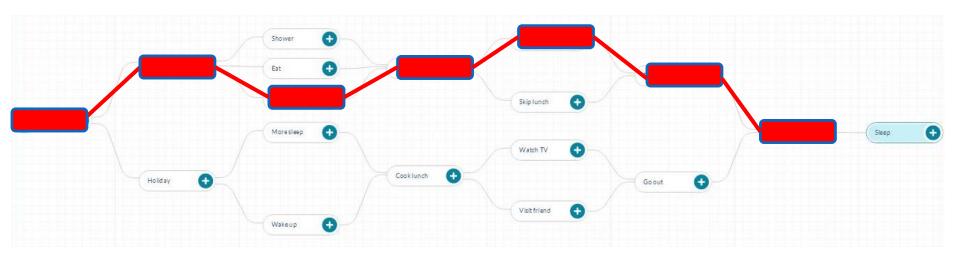
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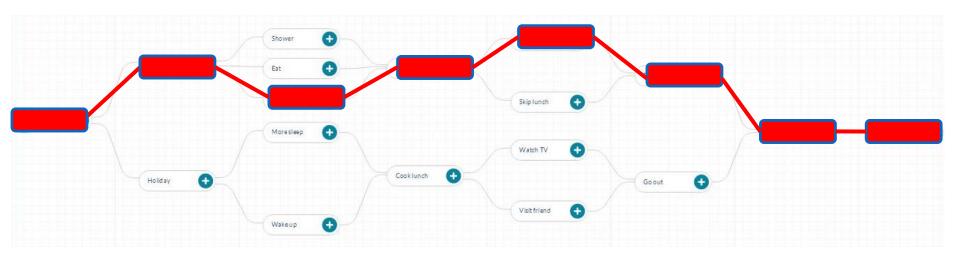
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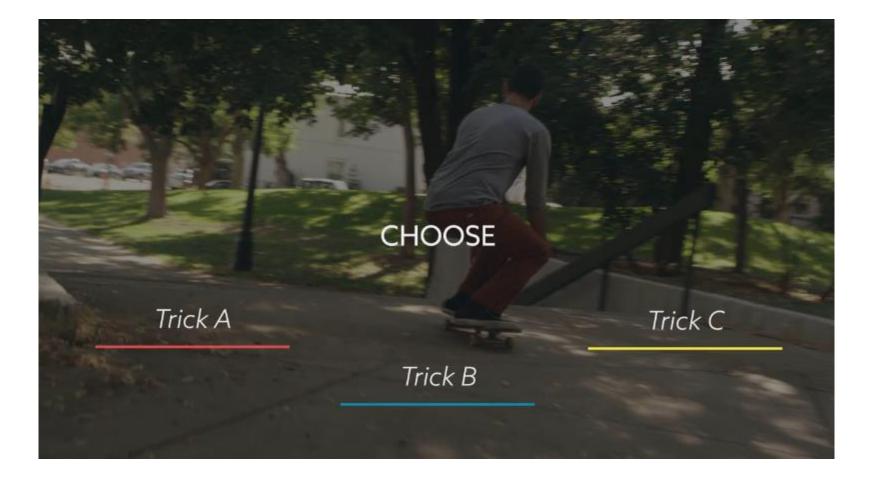


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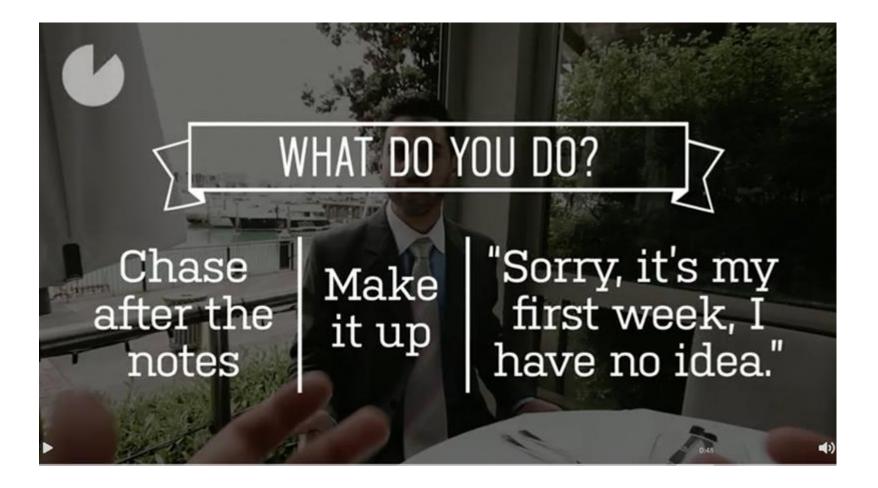


Allow user to selects between multiple storylines or alternative endings

Video personalization through user interaction



Video personalization through user interaction

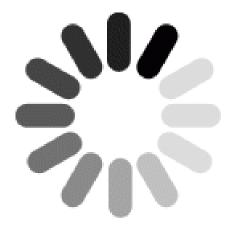


We have solved ...

The problem of providing seamless playback in the presence of multiple branch options

We have solved ...

The problem of providing seamless playback in the presence of multiple branch options

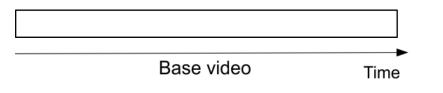


We have solved ...

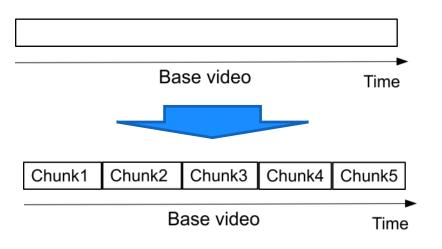
The problem of providing seamless playback in the presence of multiple branch options

- HTTP-based Adaptive Streaming
- Path and quality-aware prefetching



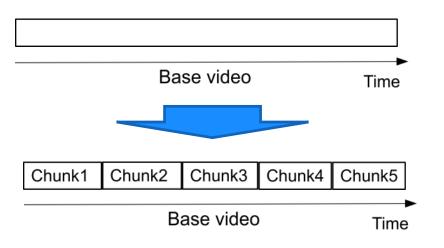


- HTTP-based streaming
 - Video is split into chunks
 - •
 - - •
 - •



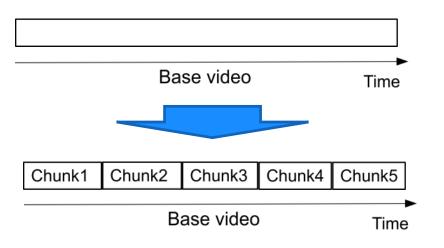
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 - - •

 - •
 - •

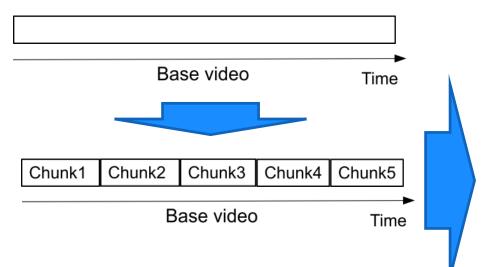


- HTTP-based streaming
 - Video is split into chunks
 - Easy firewall traversal and caching

•

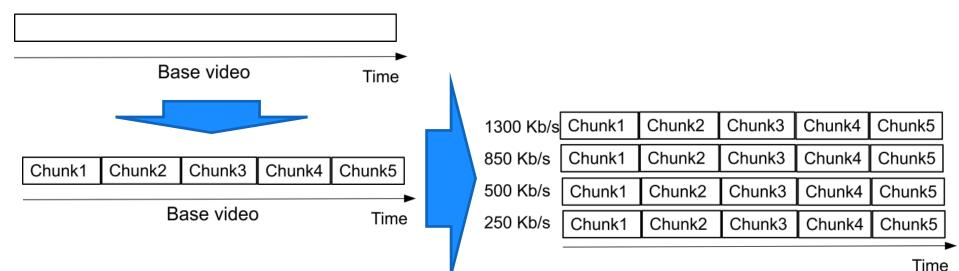


- HTTP-based streaming
 - Video is split into chunks
 - Easy firewall traversal and caching
 - Easy support for interactive VoD



HTTP-based streaming

- Video is split into chunks
- Easy firewall traversal and caching
- Easy support for interactive VoD
- HTTP-based adaptive streaming



HTTP-based streaming

- Video is split into chunks
- Easy firewall traversal and caching
- Easy support for interactive VoD
- HTTP-based adaptive streaming
 - Multiple encodings of each chunk (defined in manifest file)

1300 Kb/s	Chunk1	Chunk2	Chunk3	Chunk4	Chunk5
850 Kb/s	Chunk1	Chunk2	Chunk3	Chunk4	Chunk5
500 Kb/s	Chunk1	Chunk2	Chunk3	Chunk4	Chunk5
250 Kb/s	Chunk1	Chunk2	Chunk3	Chunk4	Chunk5

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HTTP-based Adaptive Streaming (HAS)

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- HTTP-based streaming
 - Video is split into chunks
 - Easy firewall traversal and caching
 - Easy support for interactive VoD
- HTTP-based adaptive streaming
 - Multiple encodings of each chunk (defined in manifest file)
 - Clients adapt quality encoding based on buffer/network conditions

Time

HTTP-based Adaptive Streaming (HAS)

1300 Kb/s Chunk1 Chunk2 Chunk3 Chunk4 Chunk5 Chunk1 Chunk2 850 Kb/s Chunk3 Chunk4 Chunk5 Chunk3 500 Kb/s Chunk1 Chunk2 Chunk4 Chunk5 250 Kb/s Chunk2 Chunk5 Chunk1 Chunk3 Chunk4

- HTTP-based streaming
 - Video is split into chunks
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 - Easy support for interactive VoD
- HTTP-based adaptive streaming
 - Multiple encodings of each chunk (defined in manifest file)
 - Clients adapt quality encoding based on buffer/network conditions

Time

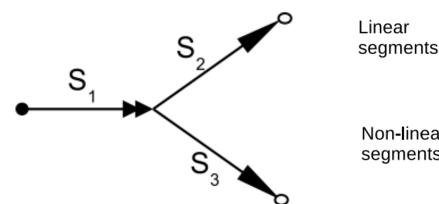
HTTP-based Adaptive Streaming (HAS)

1300 Kb/s Chunk1 Chunk2 Chunk3 Chunk4 Chunk5 Chunk1 Chunk2 850 Kb/s Chunk3 Chunk4 Chunk5 500 Kb/s Chunk1 Chunk2 Chunk3 Chunk4 Chunk5 250 Kb/s Chunk2 Chunk5 Chunk1 Chunk3 Chunk4

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Time

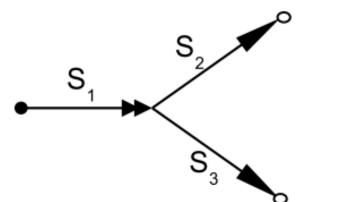
HAS-based interactive branched video



Linear	1	2	3	4	5	6	7	8	9	
segments	Se	gmer	nt 1	Se	gmei	ment 3				
Non-linear										1
segments	1	2	7	4	8	6	3	8	9	
	Segment 1			Se	Segment 2			Segment 3		

- Branched video and branch points
 - The video can include branch points, with multiple branch choices
 - User selects which segment to play back next

HAS-based interactive branched video

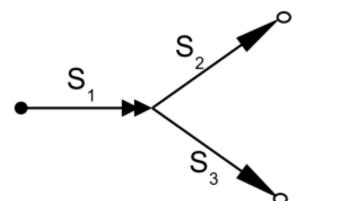


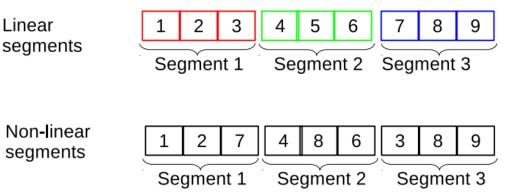
Linear segments	1	2	3	4	5	6	7	8	9
	Seç	Segment 1 Segment 2 Segment 3							
Non-linear	1	2	7		0	6	2	0	
segments		2	1	4	8	6	3	8	9
	Se	Segment 1			Segment 2			Segment 3	

- Branched video and branch points
 - The video can include branch points, with multiple branch choices
 - User selects which segment to play back next
- Our solution: Combine branched video and HAS

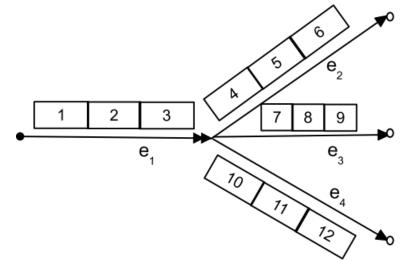
HAS-based interactive branched video

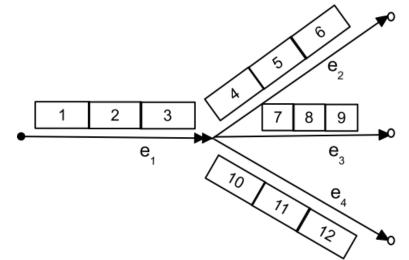
Linear



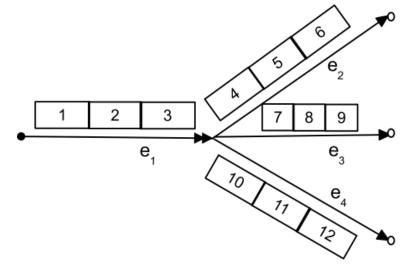


- Branched video and branch points
 - The video can include branch points, with multiple branch choices
 - User selects which segment to play back next
- Our solution: Combine branched video and HAS
- Goal: Seamless playback even if user decision at last possible moment



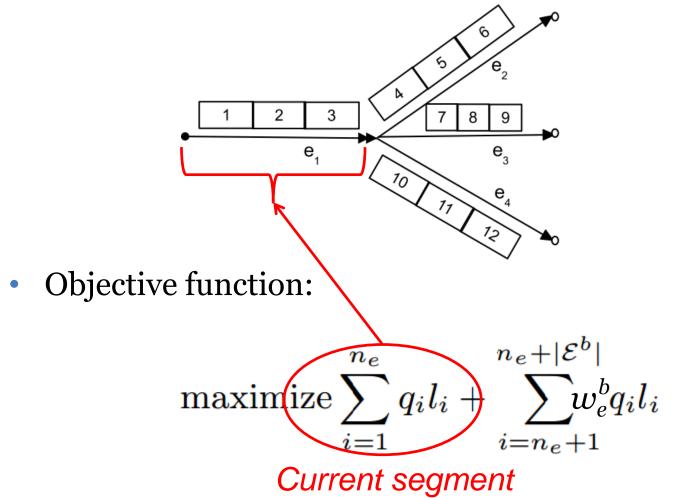


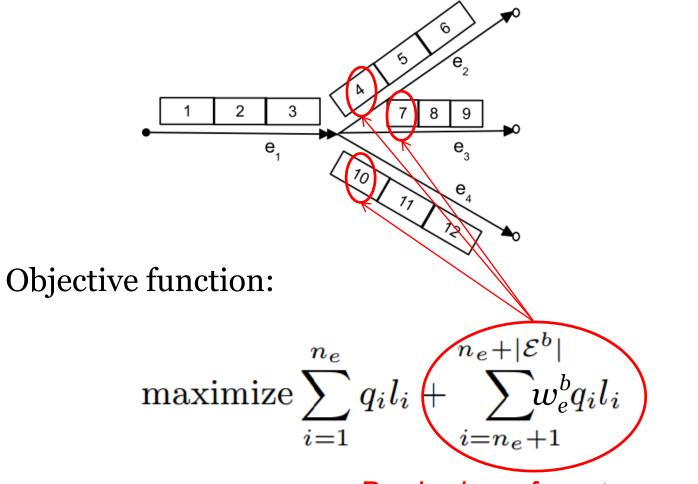
• Problem: Maximize quality, given playback deadlines and bandwidth conditions



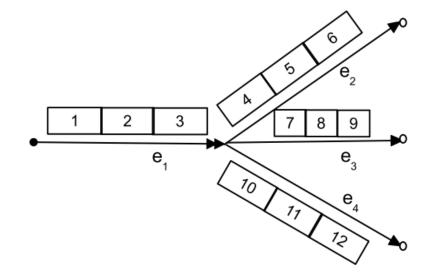
• Objective function:

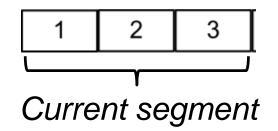
maximize playback quality

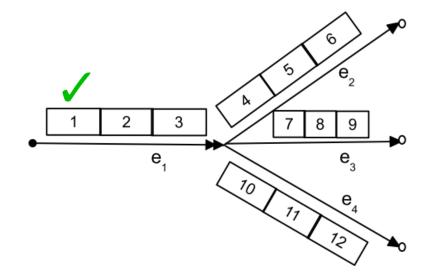


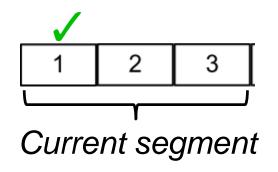


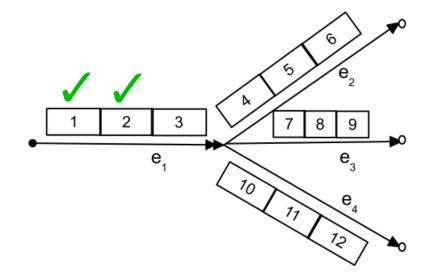
Beginning of next segment

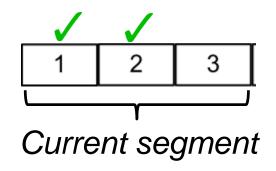


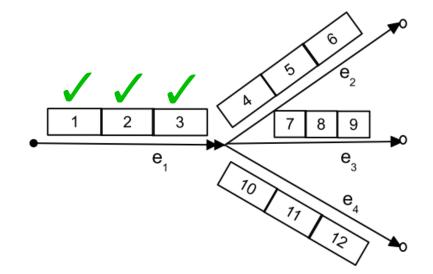


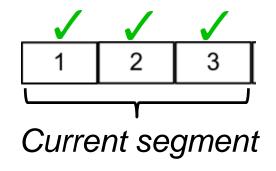


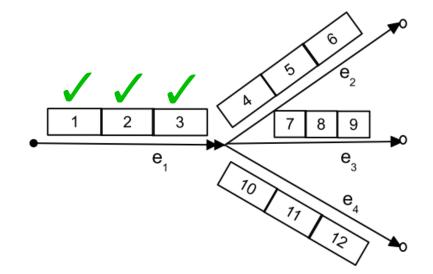


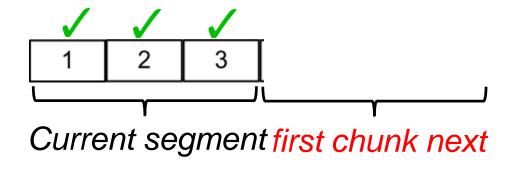


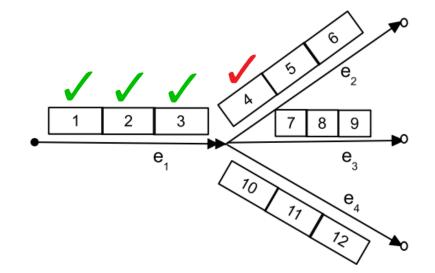


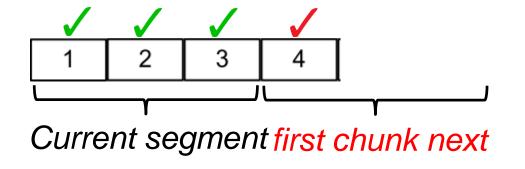


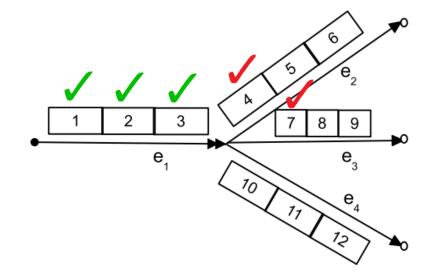


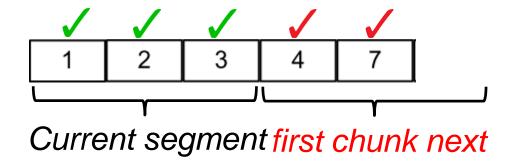


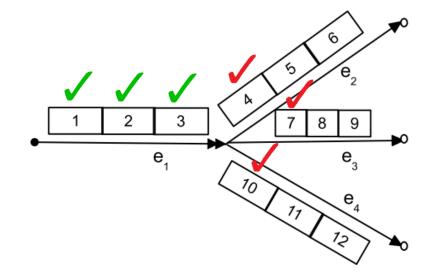


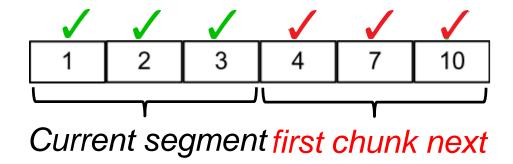


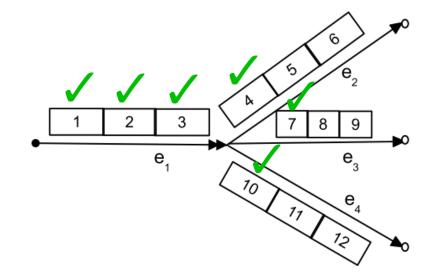


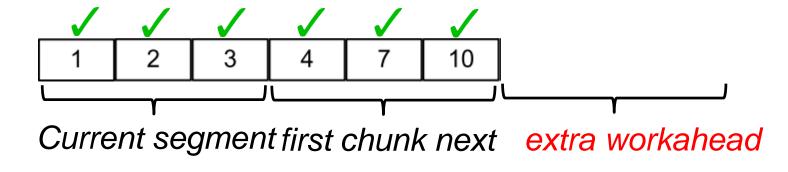


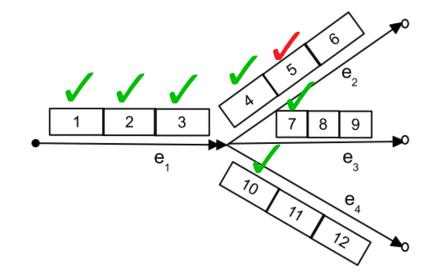


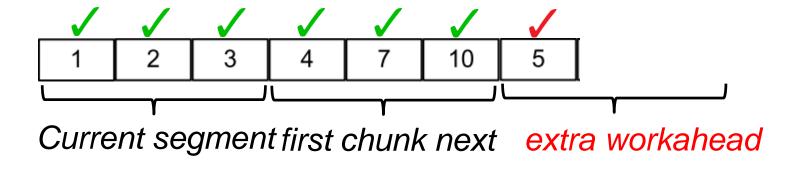


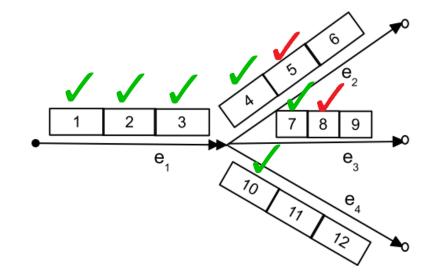


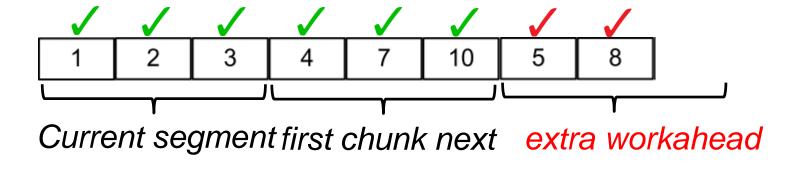


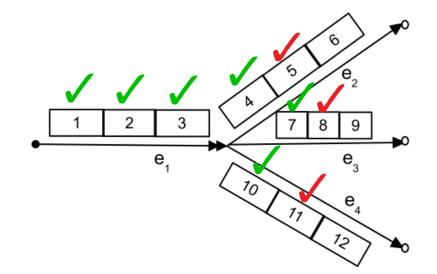


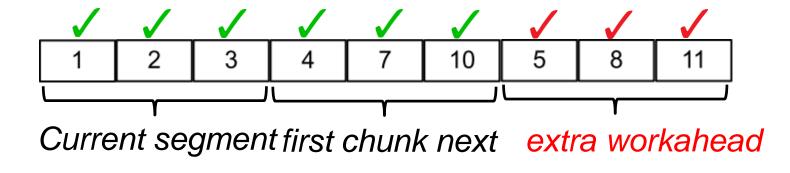


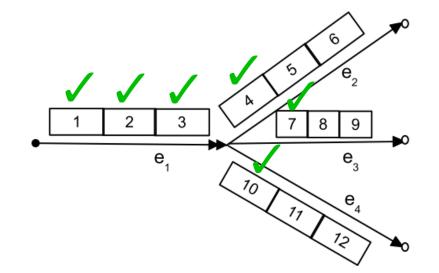


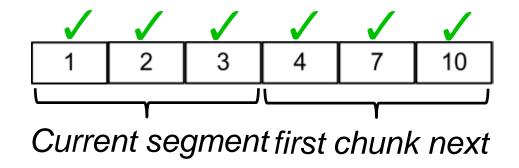


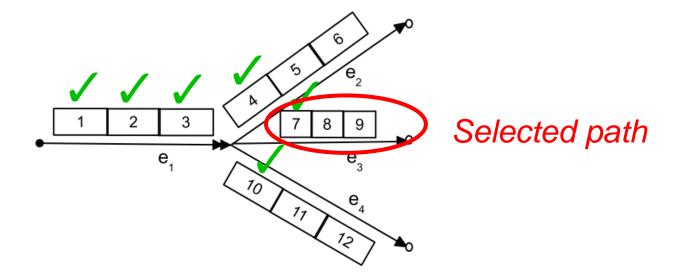


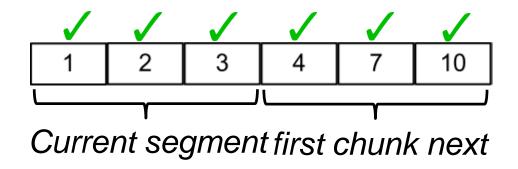


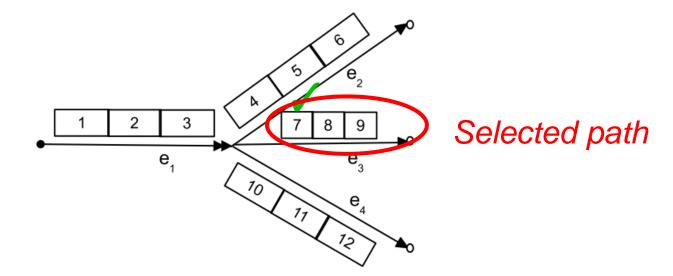


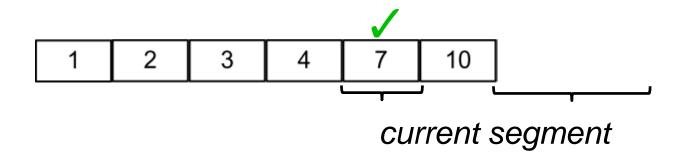


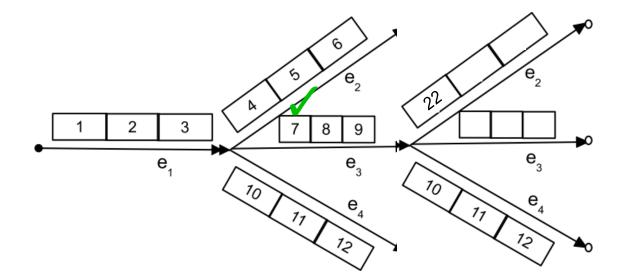


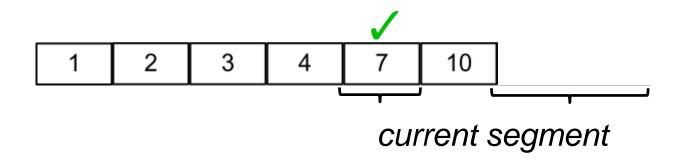


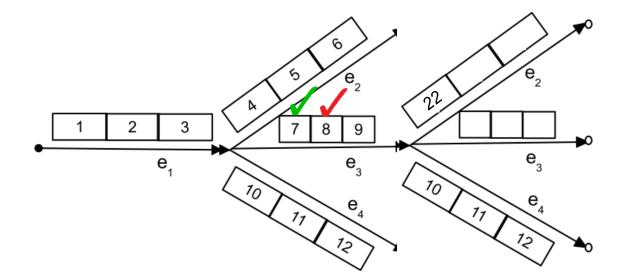


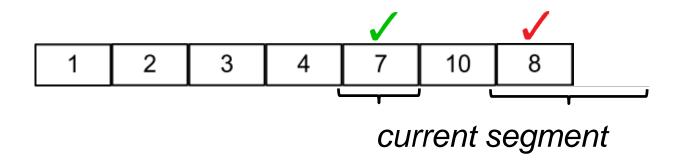


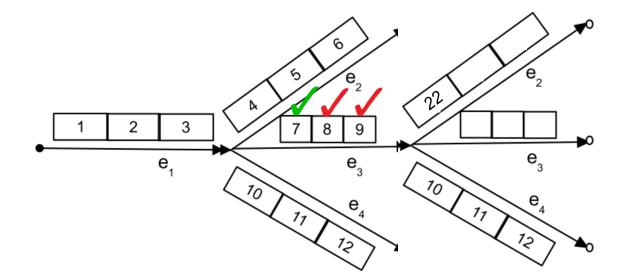


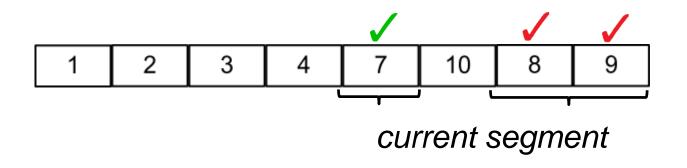


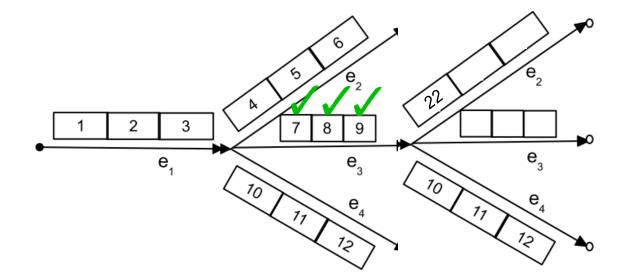


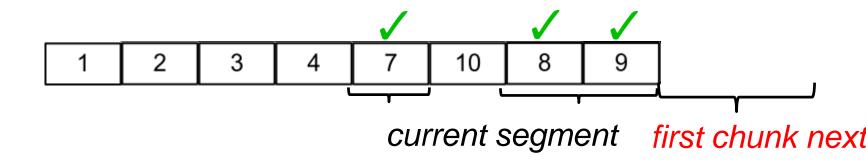


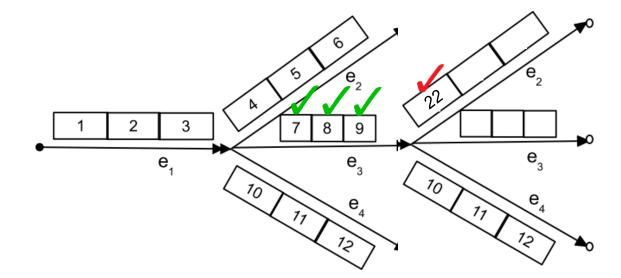


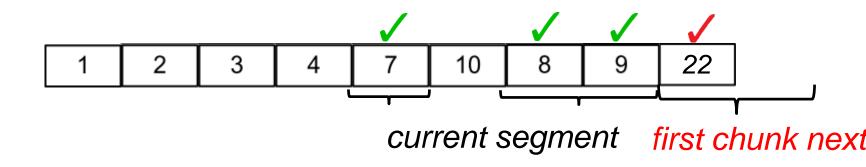


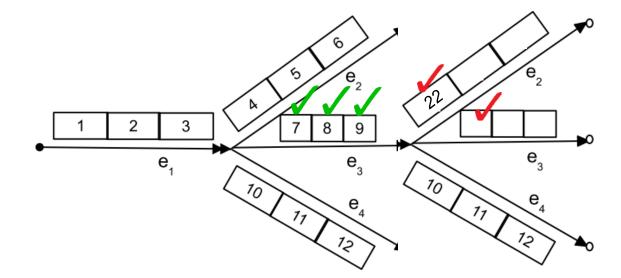


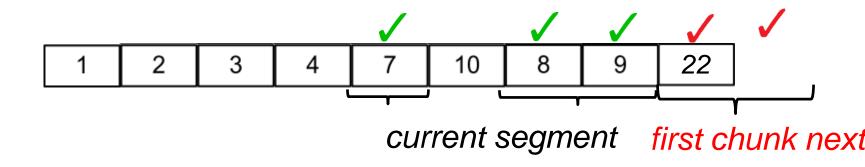


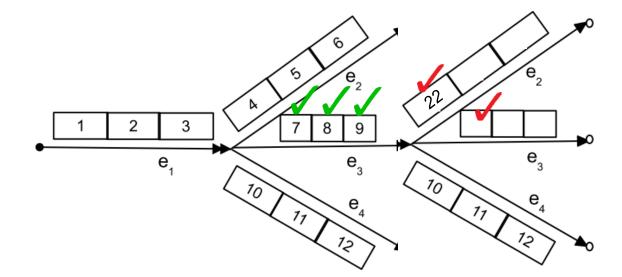


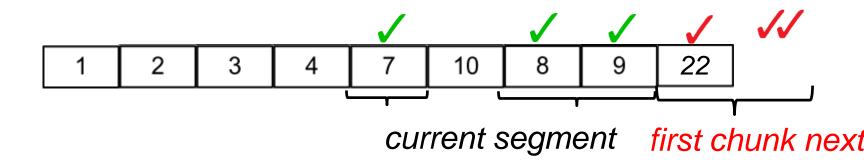


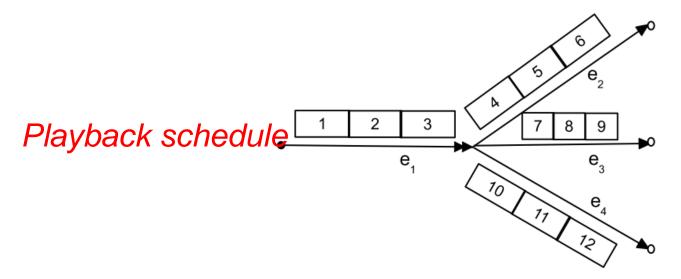




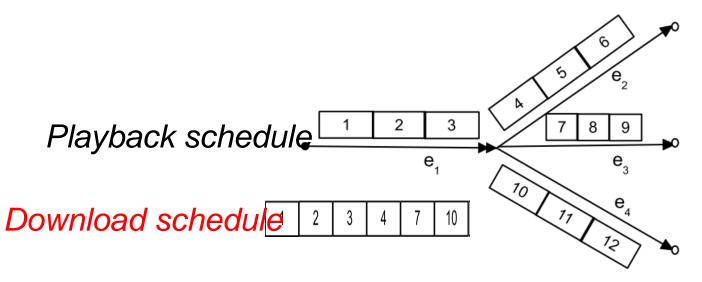




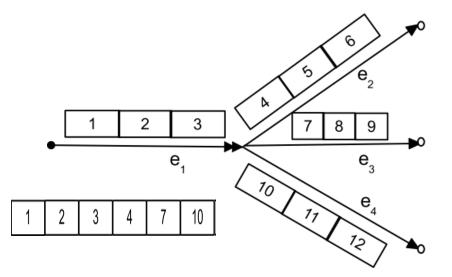




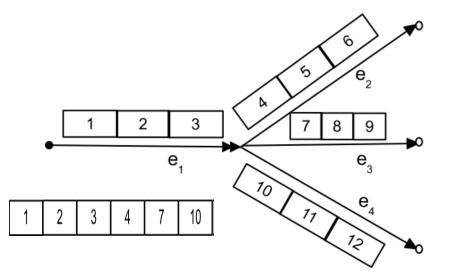
- Playback deadlines
 - for seamless playback without stalls



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 - for seamless playback without stalls

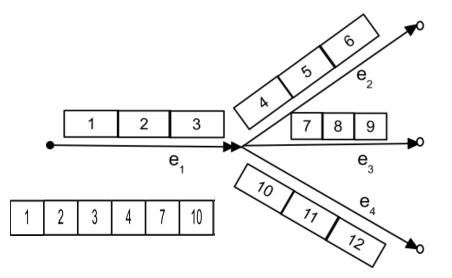


- Playback deadlines
 - for seamless playback without stalls
 - Current segment: e.g., 2 and 3



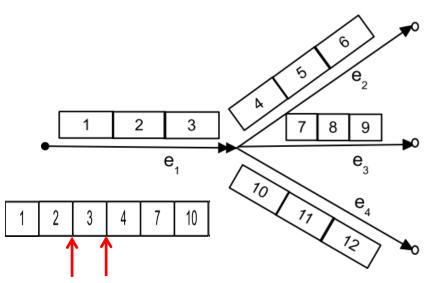
- Playback deadlines
 - for seamless playback without stalls
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$$t_i^c \le t_i^d = \tau + \sum_{j=1}^{i-1} l_j, \text{ if } 1 \le i \le n_e$$



- Playback deadlines
 - for seamless playback without stalls
 - Current segment: e.g., 2 and 3

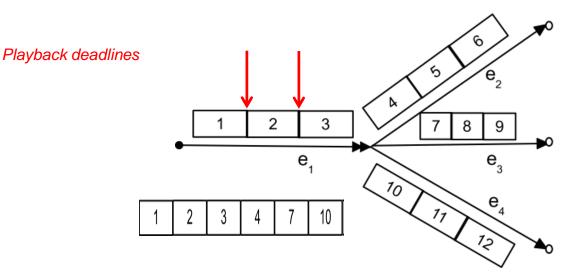
$$\underbrace{t_i^c}_{i} \leq t_i^d = \tau + \sum_{j=1}^{i-1} l_j, \quad \text{if } 1 \leq i \leq n_e$$
 Download completion time



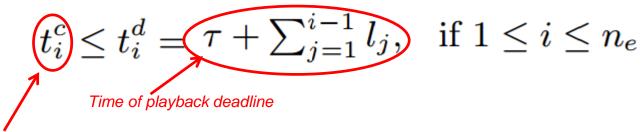
Download completion times

- Playback deadlines
 - for seamless playback without stalls
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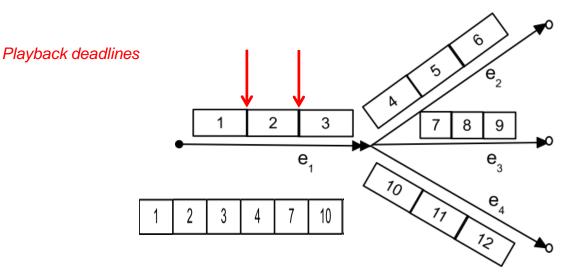
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 Download completion time



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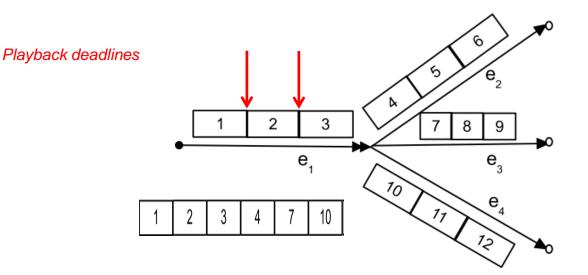


Download completion time



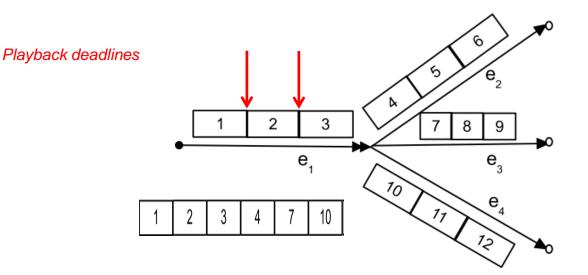
- Playback deadlines
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$$t_i^c \leq t_i^d = \tau + \sum_{j=1}^{i-1} l_j, \quad \text{if } 1 \leq i \leq n_e$$

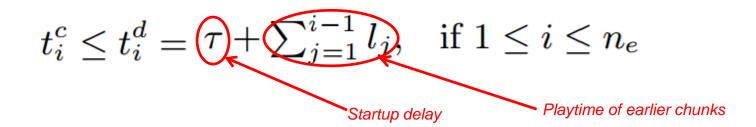


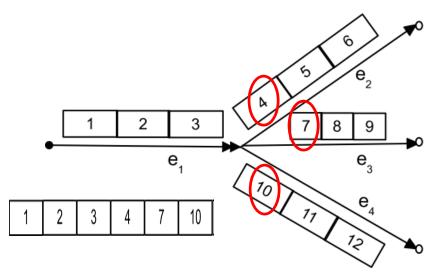
- Playback deadlines
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 - Current segment: e.g., 2 and 3

$$t_i^c \le t_i^d = \underbrace{\tau}_{j=1}^{i-1} l_j, \quad \text{if } 1 \le i \le n_e$$

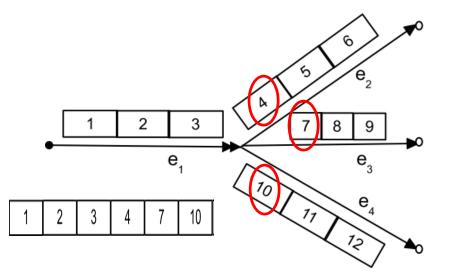


- Playback deadlines
 - for seamless playback without stalls
 - Current segment: e.g., 2 and 3



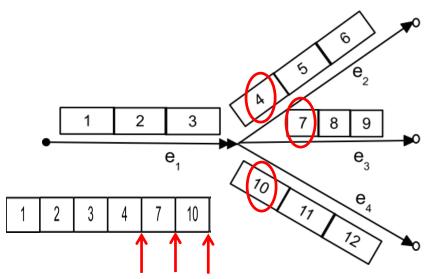


- Playback deadlines
 - for seamless playback without stalls
 - First chunks next segment: e.g., 4, 7, and 10



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 - for seamless playback without stalls
 - First chunks next segment: e.g., 4, 7, and 10

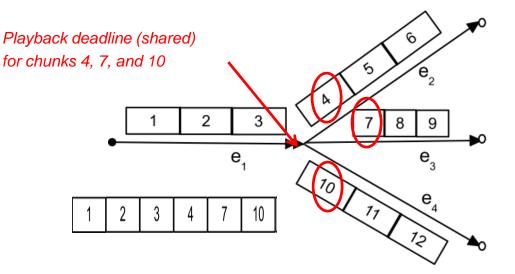
$$t_i^c \le t_i^d = \tau + \sum_{j=1}^{n_e} l_j, \text{ if } n_e < i \le n_e + |\mathcal{E}^b|$$



Download completion times

- Playback deadlines
 - for seamless playback without stalls
 - First chunks next segment: e.g., 4, 7, and 10

$$t_i^c \leq t_i^d = \tau + \sum_{j=1}^{n_e} l_j, \quad \text{if } n_e < i \leq n_e + |\mathcal{E}^b|$$



- Playback deadlines
 - for seamless playback without stalls
 - First chunks next segment: e.g., 4, 7, and 10

$$t_i^c \leq t_i^d = \tau + \sum_{j=1}^{n_e} l_j$$
, if $n_e < i \leq n_e + |\mathcal{E}^b|$
Time at which branch point is reached

Download completion times

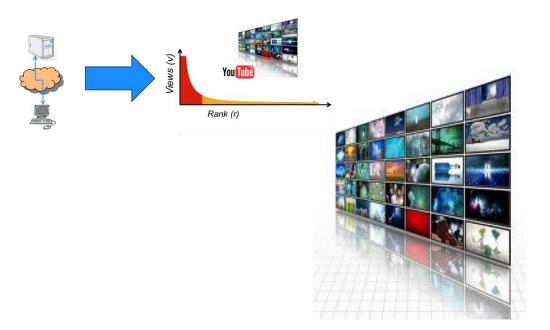
Interactive Branched Video Contributions

- Designed and implemented branched video player that achieve seamless streaming without playback interruptions
- Designed optimized policies that maximize playback quality while ensuring sufficient workahead to avoid stalls
- Evaluation shows that solution effectively adapt quality levels and number of parallel connections so as to provide best possible video quality, given current conditions

• Extensions, generalizations, and variations include "multi-file prefetching for impatient users" [*Proc. ACM Multimedia 2015*]







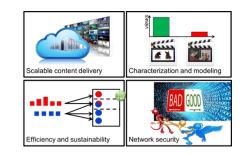














Thanks for listening!



