



Automated Roundtrip Engineering in Static Aspect Weaving and Transformations

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Static Aspect Weaving

... and other program transformations

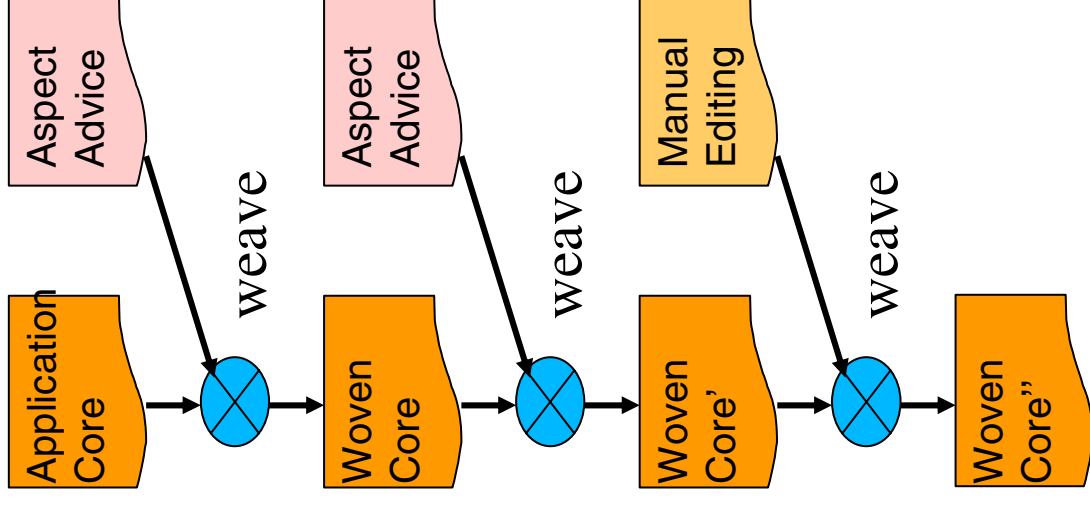
- Static Aspect Weaving
- C macro expansion
- Rule-based source-to-source transformation
- Interactive (manual) editing

→ Sequences of (syntactic) program transformations (*weaving steps*)

ARE (Automatic Roundtrip Engineering) Problem:

How to trace back and consistently commit (manual) changes in transformed code to the right source of origin?

- WovenCode → Core source
- WovenCode → Aspect source



Formalisation: Concrete Syntax Tree (CST)



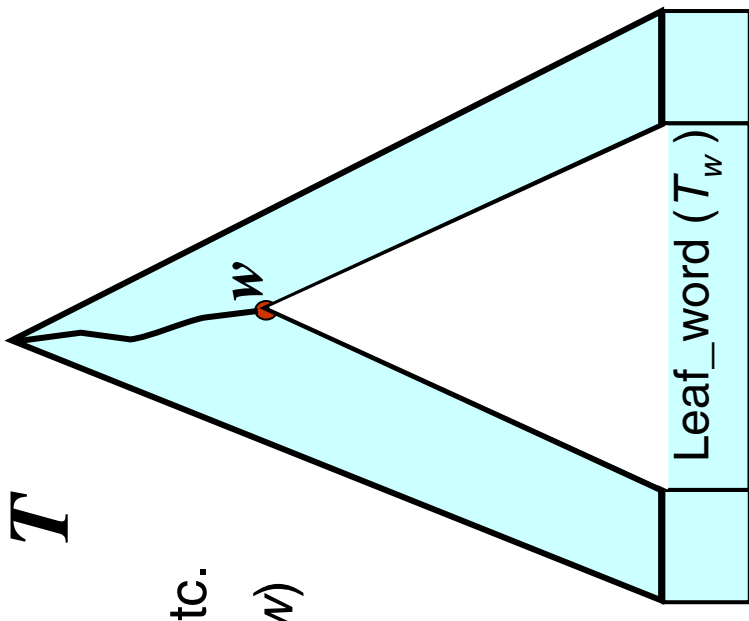
CST: Parse tree / Derivation tree

- Including whitespaces, semicolons, comments etc.
- Inner nodes w annotated with nonterminal: $NT(w)$
- Leaf node b annotated with terminal: $NT(b)$
- Inner node w with children $v_1, \dots, v_k$:

Production $NT(w) \rightarrow NT(v_1) \dots NT(v_k)$

Syntactic Type of CST:

$NT(\text{root})$



Generic CST / tree fragment: CST with "open" leaves
(= non-expanded nonterminals)

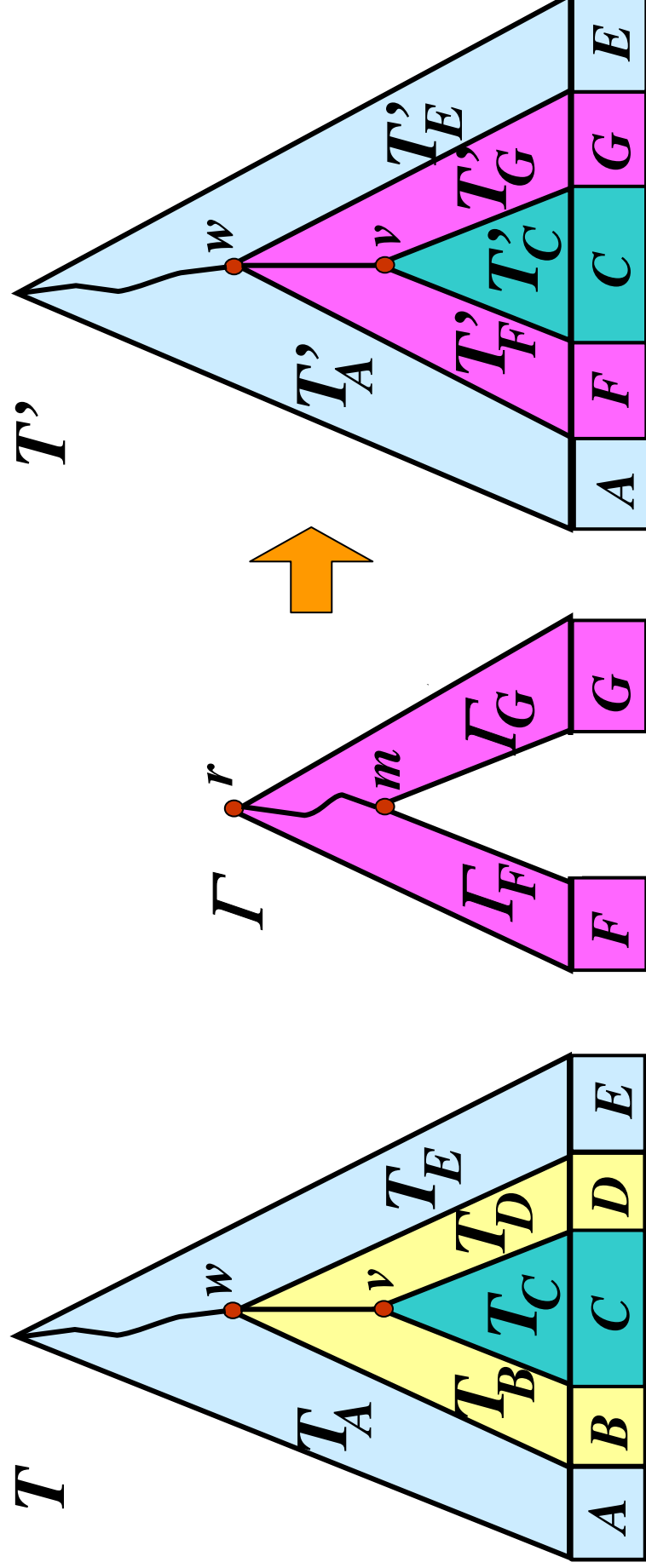
Syntactic substitutability of T_w
only with tree (fragment) of type compatible to $NT(w)$

Formalisation (2)

Double-Grafting Transformation on CSTs



■ (1,1)-Double-Grafting-Transformation:

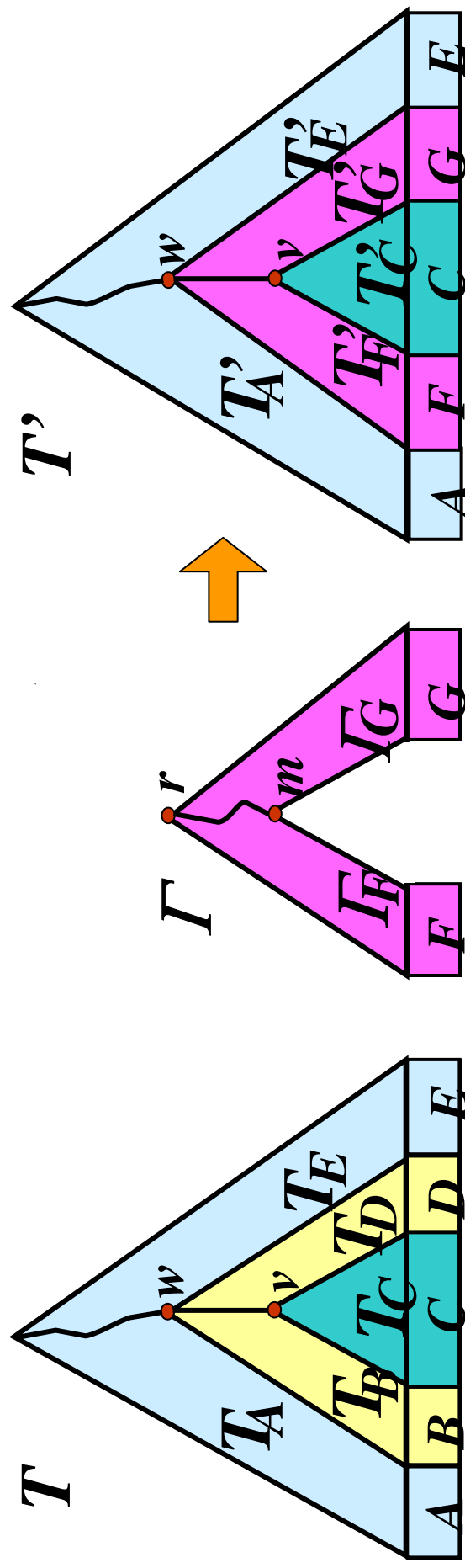


Notation: $T \stackrel{I}{\vdash} T'$

Syntactically correct
if $NT(w) = NT(r)$
and $NT(v) = NT(m)$

Example 1: Static Aspect Weaving

- Static Aspects (*before, after, around* join point)
(e.g. early version of Aspect-J)
- Application of an *advice* at a *join point (shadow)*
is a (1,1)-Double-Grafting-Transformation
- With missing *proceed()*-pseudocall (= no *m-node*):
(1,0)-Double-Grafting Transformation



AOP Example



File Edit Window

Core:

```
1 public class BankTransaction {
2
3   public static void main(
4     String[] args){
5     BankAccount ac =
6       new BankAccount(500, 400);
7     ac.credit(200);
8     ac.credit(100);
9   }
10 }
11
12
```

Aspect:

```
public aspect secureTrans {

  pointcut checkings(BankAccount bk):
    execution(*
      BankAccount.credit(float))
    && target(bk);

  before(BankAccount bk) : checkings(bk) {
    assert bk.getBalance() > 250;
  }
  after(BankAccount bk) returning:
    checkings(bk) {
    assert bk.getBalance() > 0;
  }
}
```

Woven code:

```
5 BankAccount ac =
6   new BankAccount(500, 400);
7   assert ac.getBalance() > 250;
8   ac.credit(200);
9   assert ac.getBalance() > 0;
10  assert ac.getBalance() > 250;
11  ac.credit(200);
12  assert ac.getBalance() > 0;
```

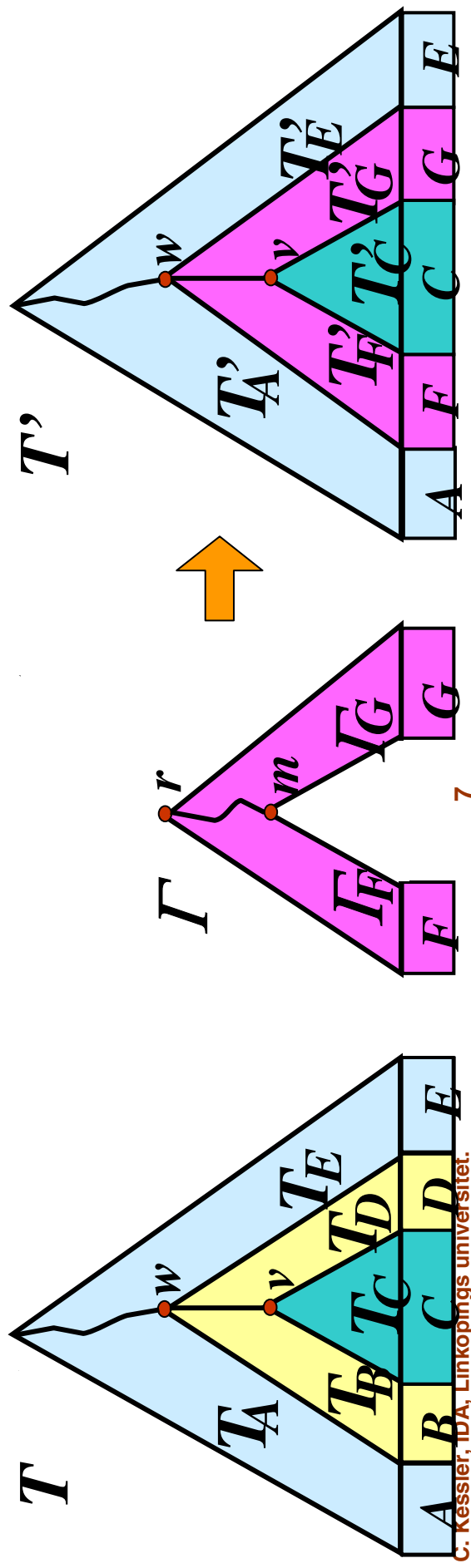
The woven code is in Aspect-J not visible for the programmer! (Weaving is done at bytecode level ...)

Example 2: Editing

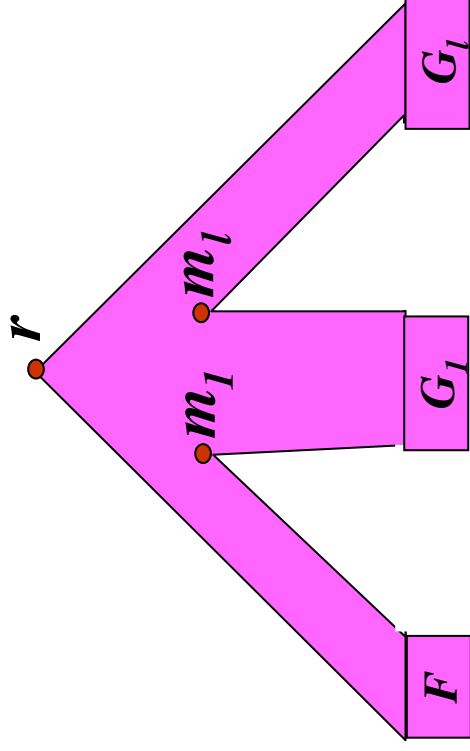
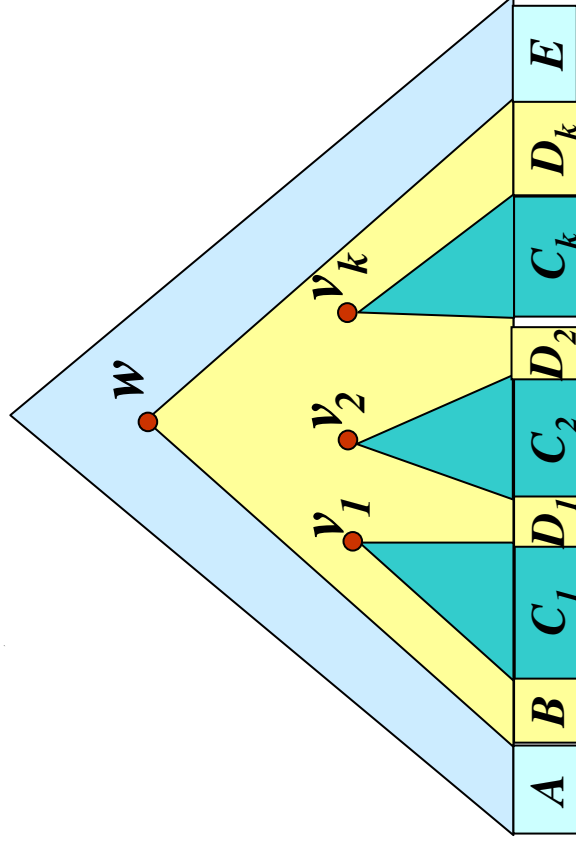
- Insert: (1,1)-Double-Grafting-Transformation
- Delete, replace: (1,0)-Double-Grafting-Transformation

Requirements:

- Program again syntactically correct after editing
- Editor coupled with CST search structure –
Clicking / Cursor on program point locates CST-node



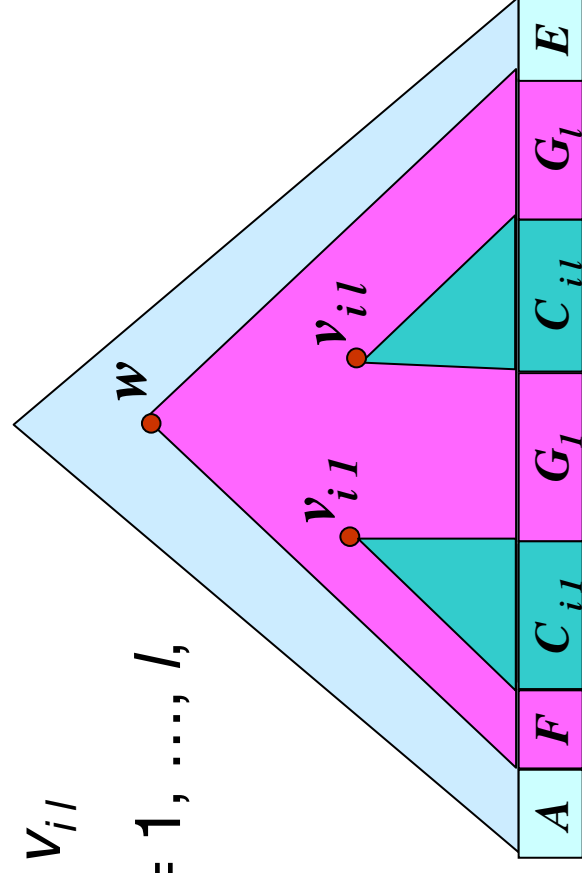
(k, l) -Double-Grafting-Transformation



+ Mapping: $m_1 \rightarrow v_{i1}, \dots, m_k \rightarrow v_{i1}$

where v_{ij} in $\{v_1, \dots, v_k\}$ for $j = 1, \dots, l$,

and $\text{NT}(m_j) = \text{NT}(v_{ij})$



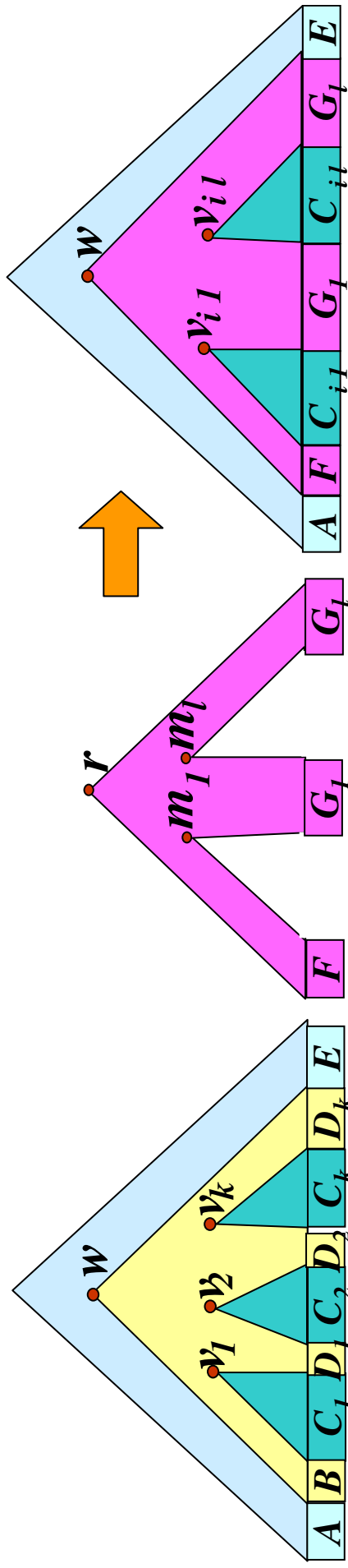
Example 3: C - Preprocessor

Macro Expansion

```
#define MAX( X, Y ) ( (X) > (Y)? X : Y )
```

```
... = MAX( u+1, 23 );
```

is a (2,4)-Double-Grafting-Transformation

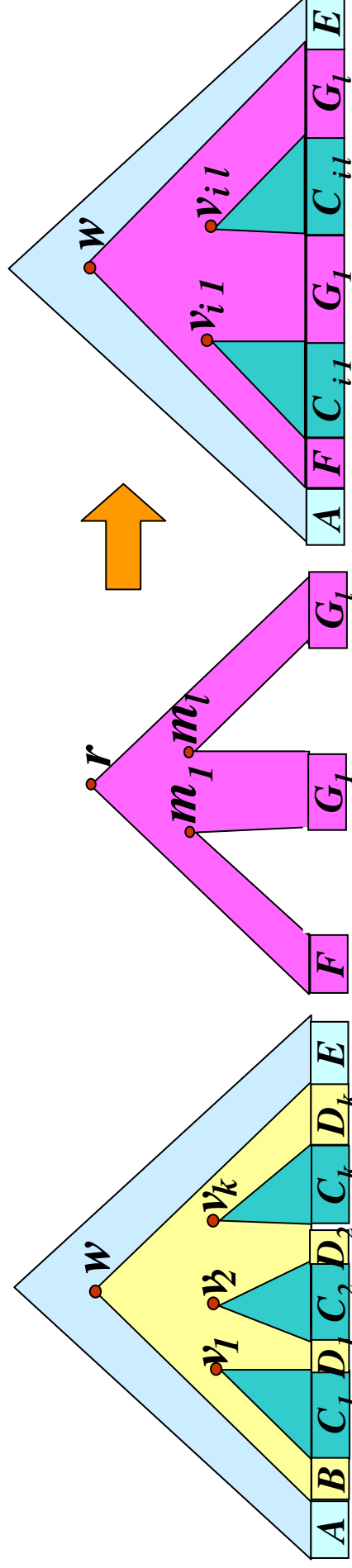


Example 4: Rule-based Source-to-Source Transformation Systems

- Generic Transformation Rule



is a (2,1)-Double-Grafting-Transformation



Static Aspect Weaving

... and other program transformations

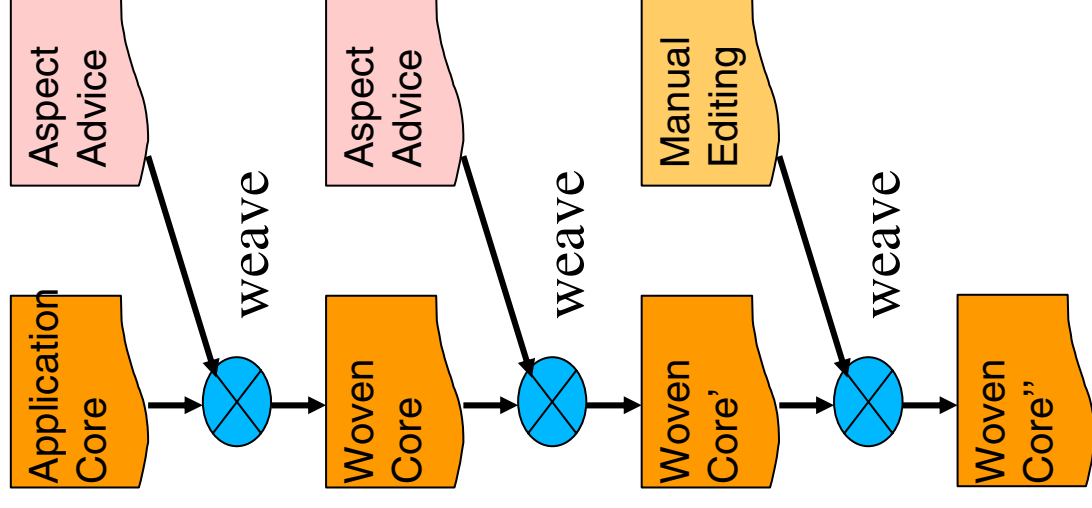
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ARE Example



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            > 250;
    }

    after(BankAccount bk) returning:
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    }
}
```

ARE Notice:

This modification applies to aspect
secureTrans, advice **before** (line 8).
Commit update to aspect source?

Commit

Ignore

Commit

Ignore

```
7 assert ac.getBalance()
8 ac.credit(200);
9 assert ac.getBalance() > 0;
10 assert ac.getBalance() > 250;
11 ac.credit(200);
12 assert ac.getBalance() > 0;
```

Problem:

Find Origin of a Program Point

- Given: Weaving sequence (History)

$$T_0 \vdash^{I_1} T_1 \vdash^{I_2} T_2 \vdash^{I_3} T_3 \dots \vdash^{I_n} T_n$$

- Given: Program point coordinate w.r.t. T_n
(row r , column c) in Editor
or, linear index p in Lexeme Sequence (leaf word)
- Compute: Source-File + Coordinate of p in Origin of p
either in the core T_0
or in one of the I_i , $i = 1 \dots n$
- Method: Store relative coordinate transformation to each
Double-Grafting-Transformation in the weaving sequence

→ Transposition tree

Transposition Tree Construction (2)

Forward-Transposition:

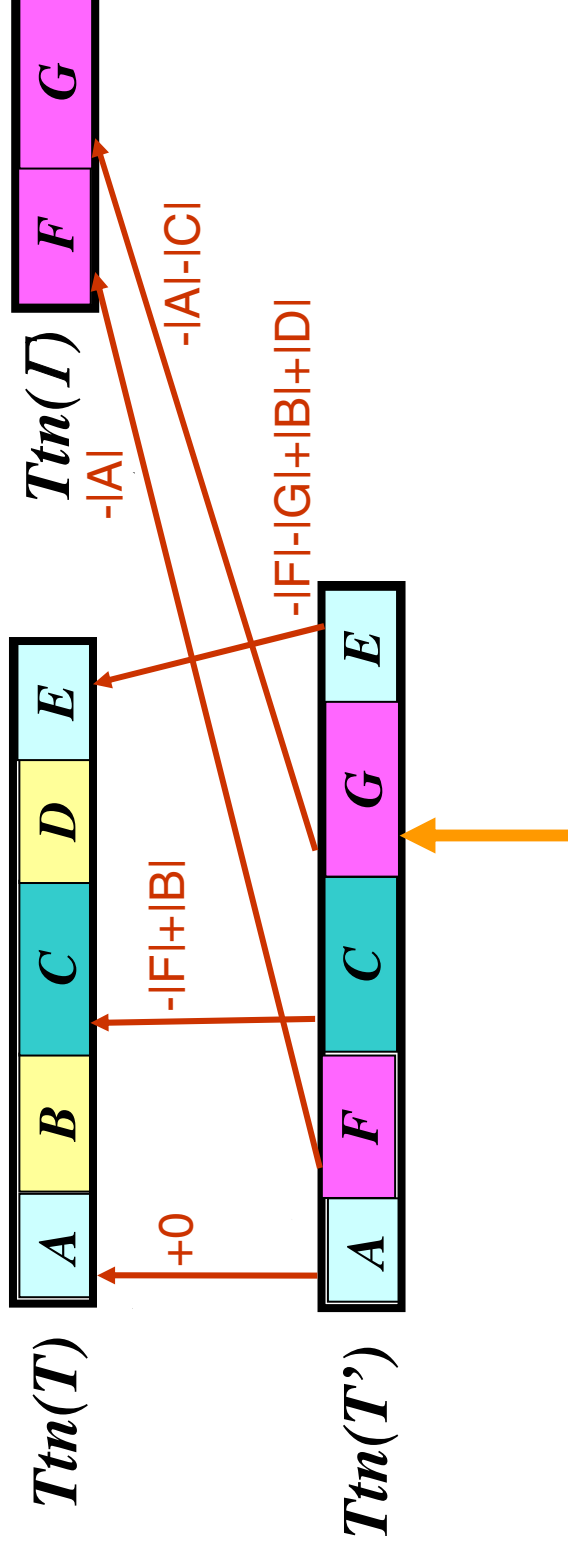
$$fsrc_T(v) = \begin{cases} fsrc_T(v) & \text{if } v \text{ in } T_A, T_C, T_E \\ fsrc_T(v) & \text{if } v \text{ in } T_F, T_G \\ \text{undefined,} & \text{if } v \text{ in } T_B, T_D \end{cases}$$

$$fpos_T(v) = \begin{cases} fpos_T(v) & \text{if } v \text{ in } T_A \\ fpos_T(v) + |A| & \text{if } v \text{ in } T_F \\ fpos_T(v) + |A| - |B| + |F| & \text{if } v \text{ in } T_C \\ fpos_T(v) + |A| - |B| + |F| + |C| & \text{if } v \text{ in } T_G \\ fpos_T(v) + |A| - |B| + |F| + |C| - |D| + |G| & \text{if } v \text{ in } T_E \end{cases}$$

Transposition Tree Construction (3)

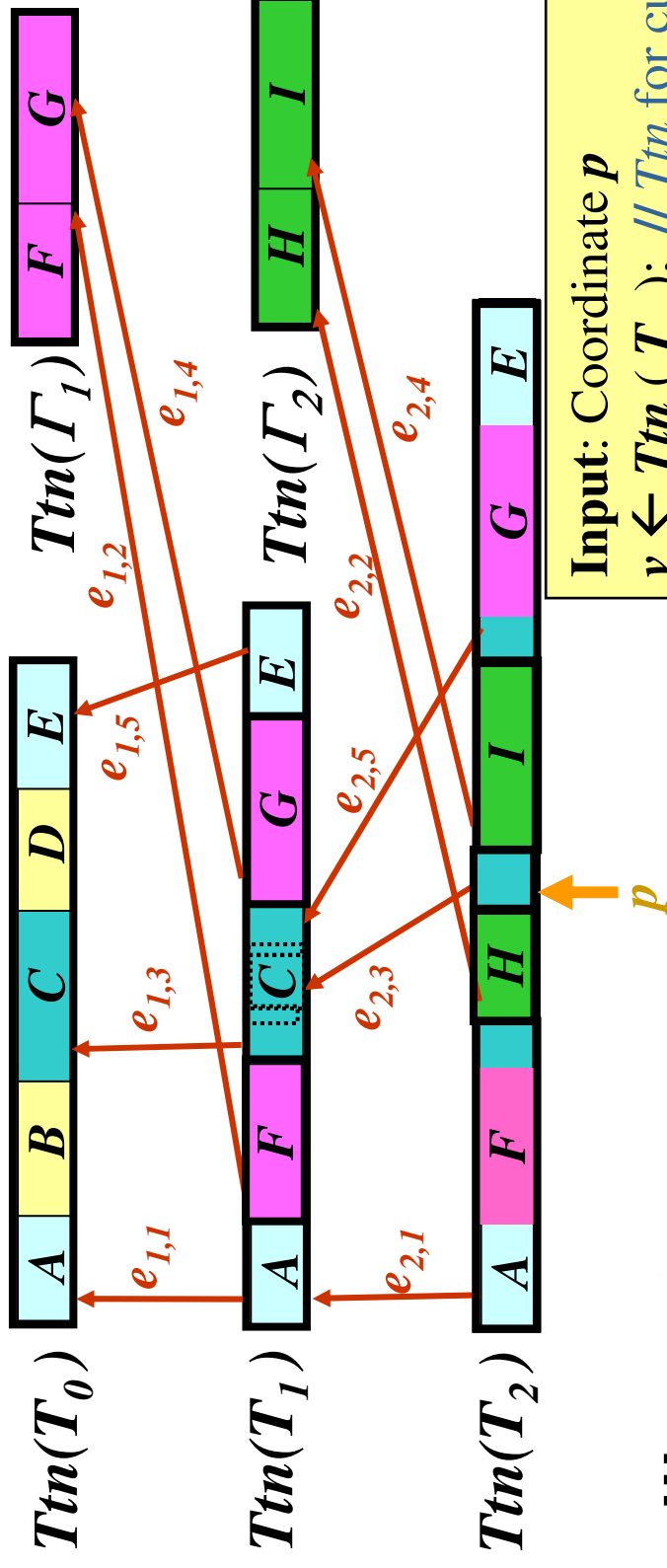
- All coordinates in the same segment ($A, B, \dots$) are mapped (transposed) in the same way
- Backward transposition stored blockwise

→ Nodes / Edges in Transposition tree



Transposition Tree Construction (4)

- Incremental construction yields tree of depth $\leq n$



...

Data structure needs $O(n)$ space
(Constant depends on k, l)

Search backwards in time $O(n)$

Input: Coordinate p

$v \leftarrow Ttn(T_n)$; // Ttn for current CST

for t from n downto 1 do

find interval i of p in v (Binary search)

$p \leftarrow \text{Backwardtransp}[e_{t,i}](p)$;

$v \leftarrow \text{target}(e_{t,i})$;

if v is Ttn of an aspect I break

od

Propagating Changes (1)

■ Weaving sequenz can always be *replayed*

- Undo Γ_n by replaying $\Gamma_1, \dots, \Gamma_{n-1}$ on T_0

At a change (e.g. edit):

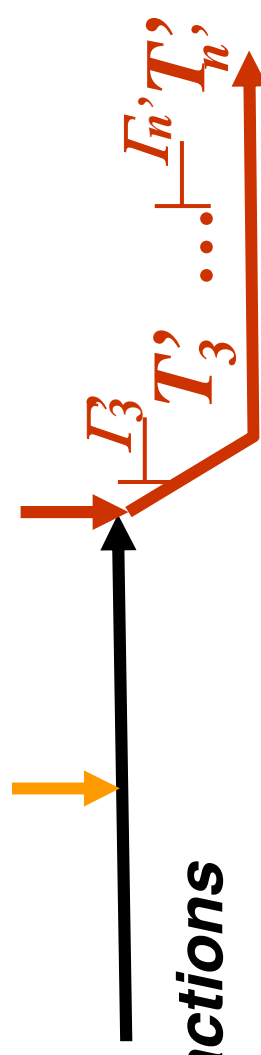
- Origin in core $T_0 \rightarrow$ 2 Options (User dialog):
 - Store and commit edit as local patch Γ_{n+1} of T_n
 - Patch edit directly in the core (as Γ_0), then replay $\Gamma_1, \dots, \Gamma_n$
- Origin in Aspect-Advice $\rightarrow$ 3 Options (User dialog):
 - Store and commit edit as local patch Γ_{n+1}
 - Patch edit directly in Advice (as Γ_0), then replay $\Gamma_1, \dots, \Gamma_n$
 - Commit after Γ_t for some $0 < t < n$, then replay $\Gamma_{t+1}, \dots, \Gamma_n$

Propagating Changes (2)

Problems:

- Not all T_i are (semantically) correct programs
- Paradox (cf. "Back to the future II")
 - Change can lead to disruptive changes in the replay – code segments created "later" may disappear, new ones appear...

$$\frac{T_0 \vdash^I \quad T_1 \vdash^{E_1} \quad T_2 \vdash^{E_2} \quad T_3 \vdash^{E_3} \quad \dots \quad T_n \vdash^{E_n}}{\quad}$$



Remedy: *Weaving Transactions*

- Encapsulate inconsistent intermediate configurations
- Limit scope of propagating changes, undo option

Application?



- Planned: Interactive Parallelisation Tool
 - Rule-based automatic transformations
 - Semiautomatic transformations (Skeleton expansion)
 - Manual parallelisation (forall, sh, \$ivdep, \$distribute ...)
 - Sequential application core remains intact
 - Co-Evolution of application core and woven-in parallelisation code

Summary

- (k, l) -Double-Grafting-Transformation
 - Models static aspect weaving, Macro expansion, Edits ...
- Store weaving sequence
 - with transposition tree
 - find origin of a program point
 - Replay
 - User dialog for consistent change propagation
- ☺ Need not know reverse transformation to each step
- ☺ Space-efficient data structure

References



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