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Abstract Interpretation

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Outline

- Summary of Data Flow Analysis (previouse lecture)
- Problems left open
- Abstract interpretation idea

Complete Partial Order (CPO)

- Partially ordered sets $(U,\sqsubseteq)$ over a universe U
  - Smallest element $\bot \in U$
  - Partial order relation $\sqsubseteq$
- Ascending chain $C=[c_1,c_2,\ldots] \subseteq U$
  - Smallest element c_1
  - $c_{i-1} \sqsubseteq c_i$
  - Maybe finite or countable: constructor for next element $c_i = next([c_1,c_2,\ldots,c_{i-1}])$
- Unique largest element s of chain $C=[c_1,c_2,\ldots]$
  - $c_i \sqsubseteq s$ (larger than all chain elements c_i)
  - s called supremum $s = \bigsqcup(C)$
- Ascending chain property: any (may be countable) ascending chain $C \subseteq U$ has an element c_i with
  - i is finite and
  - for all elements $c_{<i} \sqsubseteq c_i$ and
  - for all elements $c_{>i} = c_i$,

  then $c_i = \bigsqcup(C)$
- Example: $(\wp^{\mathbb{N}}, \subseteq)$ and $C=[\varnothing,\{1\},\{1,2\},\{1,2,3\},\ldots]$, $c_i = \{\max(c_{i-1})+1\} \cup c_{i-1}$ $s = \cup(C) = \mathbb{N}$ but ascending chain property does not hold

CPOs and Lattices

- Lattice $L = (U,\bigsqcup,\bigsqcap)$
  - any two elements a, b of U have
    - an infimum $\bigsqcap(a, b)$ - unique largest smaller of a, b
    - a supremum $\bigsqcup(a, b)$ - unique smallest bigger of a, b
  - unique smallest element $\bot$ (bottom)
  - unique largest element $\top$ (top)
- A lattice $L = (U,\bigsqcup,\bigsqcap)$ defines two kinds of CPOs $(U,\sqsubseteq)$
  - upwards:
    - $a \sqsubseteq b \Leftrightarrow a \bigsqcup b = b$, smallest $\bot$,
    - finite heights $\Rightarrow$ ascending chain property holds ($c_i = \top$)
  - downwards:
    - $b \sqsubseteq a \Leftrightarrow a \bigsqcup b = b$ ($\Leftrightarrow a \bigsqcap b = a$), smallest $\top$,
    - finite heights $\Rightarrow$ ascending chain property holds ($c_i = \bot$)

Special lattices of importance

- Boolean Lattice over $U=\{true, false\}$
  - $\bot = true$, $\top = false$, $true \sqsubseteq false$, $\bigsqcup(a,b)=a \vee b$, $\bigsqcap(a,b)=a \wedge b$
  - Finite heights
- Generalization: Bit Vector Lattice over $U = \{true, false\}^n$
  - Finite heights if n is finite
- Power Set Lattice $\wp^S$ over S (set of all subsets of a set S)
  - $\bot = \varnothing$, $\top = S$, $\sqsubseteq = \subseteq$, $\bigsqcup(a,b)=a \cup b$, $\bigsqcap(a,b)=a \cap b$ or the dual lattice
  - $\bot = S$, $\top = \varnothing$, $\sqsubseteq = \supseteq$, $\bigsqcup(a,b)=a \cap b$, $\bigsqcap(a,b)=a \cup b$
  - Finite heights if S is finite

Functions on CPOs

- Functions $f: U \to U'$ (if not indicated otherwise, we assume $U = U'$)
- f monotone: $x \sqsubseteq y \Rightarrow f(x) \sqsubseteq f(y)$ with $x, y \in U$
- f continuous: $f(\bigsqcup C) = \bigsqcup f(C)$ with $f(C) = f([x_1,x_2,\ldots]) = [f(x_1),f(x_2),\ldots]$

- f continuous $\Rightarrow f$ monotone,

- f monotone $\wedge$ U is finite $\Rightarrow f$ continuous
- f monotone $\wedge$ $(U,\sqsubseteq)$ a CPO with ascending chain property $\Rightarrow f$ continuous
- f monotone $\wedge$ $(U,\bigsqcup,\bigsqcap)$ a lattice with finite heights $\Rightarrow f$ continuous

Example

- $U = \mathcal{P}^{\mathbb{N}}$ (set of all subsets of Natural numbers N)
- Define:
  - $f(u) = \varnothing \Leftrightarrow u \in U$ finite
  - $f(u) = \mathbb{N} \Leftrightarrow u \in U$ infinite
- f is monotone $u \subseteq u' \Rightarrow f(u) \subseteq f(u')$, e.g.,
  - $\varnothing \subseteq \{0\} \subseteq \{0,1\} \subseteq \ldots \Rightarrow f(\varnothing) \subseteq f(\{0\}) \subseteq f(\{0,1\}) \subseteq \ldots = \varnothing \subseteq \varnothing \subseteq \varnothing \subseteq \ldots$
- f is not continuous $f(\cup C) \neq \cup f(C)$, e.g.:
  - $C = [\varnothing, \{0\}, \{0,1\}, \ldots]$
  - $f(C) = [f(\varnothing), f(\{0\}), f(\{0,1\}), \ldots] = [\varnothing, \varnothing, \varnothing, \ldots]$
  - $\cup f(C) = \cup [f(\varnothing), f(\{0\}), f(\{0,1\}), \ldots] = \cup [\varnothing, \varnothing, \varnothing, \ldots] = \varnothing$
  - $\cup C = \cup [\varnothing, \{0\}, \{0,1\}, \ldots] = \mathbb{N}$
  - $f(\cup C) = f(\cup [\varnothing, \{0\}, \{0,1\}, \ldots]) = f(\mathbb{N}) = \mathbb{N}$
- Note: Power Set Lattice $\mathcal{P}^{\mathbb{N}} = (\mathbb{N}, \cup, \cap)$ is not of finite heights

Fixed Point Theorem (Knaster-Tarski)

Fixed point of a function: X with $f(X) = X$

For CPO $(U, \sqsubseteq)$ and monotone functions f: $U \to U$
- Minimum (or least or smallest) fixed point X exists
- X is unique

For CPO $(U, \sqsubseteq)$ with smallest element $\perp$ and continuous functions f: $U \to U$
- Minimum fixed point $X = \bigsqcup f^n(\perp)$
- X iteratively computable

CPO $(U, \sqsubseteq)$ fulfills ascending chain property $\Rightarrow X$ is computable effectively
Special cases:
- $(U, \sqsubseteq)$ with U finite,
- $(U, \sqsubseteq)$ defined by a finite heights lattice.

Monotone DFA Framework

- Solution of a set of DFA equations is a fix point computation
- Contribution of a computation A of kind K (*Alloc*, *Add, Load, Store, Call* …) is modeled by monotone transfer function
  - f_K: $U \to U$,
  - Set F of transfer functions is closed under composition and (obviously) the composed functions are monotone as well
- Contribution of predecessors Pre of A is modeled by supremum $\bigsqcup$ of predecessor properties (successor $Succ$, resp., for backward problems)
- Existence of the smallest fix point X is guaranteed, if domain U of analysis values $P(A)$ completely partially ordered $(U, \sqsubseteq)$
- It is efficiently computable if $(U, \sqsubseteq)$ additionally fulfills the ascending chain property

Monotone DFA Framework (cont'd)

- Monotone DFA Framework: $(U, \sqsubseteq, F, \iota)$
  - $(U, \sqsubseteq)$ a CPO of analysis values fulfilling the ascending chain property
  - $F = \{f_K: U \to U, f_{K'}: U \to U, \ldots\}$ set of monotone transfer functions (closed under composition, analysis problem specific)
  - $\iota \in U$ initial value (analysis problem-specific)
- Analysis instance of a Monotone DFA Framework is given by a graph G
  - $G = (N, E, n^1)$ data flow graph of a specific program, with
  - the start node $n^1 \in N$
- $((N \times U \times U)^{|N|}, \sqsubseteq_{\text{Vector}})$ defines a CPO:
  - Let $a = (i, x_{in}, x_{out})$, $b = (j, y_{in}, y_{out})$, $a, b \in (N \times U \times U)$
    $a \sqsubseteq_{\text{Triple}} b \Leftrightarrow i = j \wedge x_{in} \sqsubseteq y_{in} \wedge x_{out} \sqsubseteq y_{out}$
  - Let $m = [a^1, a^2, \ldots, a^{|N|}]$, $n = [b^1, b^2, \ldots, b^{|N|}]$, $m, n \in (N \times U \times U)^{|N|}$
    $m \sqsubseteq_{\text{Vector}} n \Leftrightarrow a^1 \sqsubseteq_{\text{Triple}} b^1 \wedge a^2 \sqsubseteq_{\text{Triple}} b^2 \wedge \ldots \wedge a^{|N|} \sqsubseteq_{\text{Triple}} b^{|N|}$
  - Smallest element is vector $[(n^1, \iota, \perp), (n^2, \perp, \perp), \ldots, (n^{|N|}, \perp, \perp)]$

Monotone DFA Framework (cont'd)

- Data flow equations define monotone functions in $(N \times U \times U, \sqsubseteq)$:
  $P_{in}(A) = \bigsqcup_{X \in Pre(A)} (P_{out}(X))$
  $P_{out}(A) = f_{Kind(A)}(P_{in}(A))$ with $f_{Kind(A)} \in F$ transfer function of A
- Smallest fix point of this system of equations is efficiently computable since
  - $(N \times U \times U)$ and hence $(N \times U \times U)^{|N|}$ completely partially ordered and fulfill the ascending chain property
  - System of equations defines monotone function in $(N \times U \times U)^{|N|}$
- Data flow analysis algorithm:
  - Start with the smallest element: $[(n^1, \iota, \perp), (n^2, \perp, \perp), \ldots, (n^{|N|}, \perp, \perp)]$
  - Apply equations in any (fair) order
  - Until no $P_{in}(A)$ nor $P_{out}(A)$ changes

4 DFA Equations Schemata

- forward and must: $P_{in}(A) = \bigsqcup_{X \in Pre(A)} P_{out}(X)$ (sup)
  $P_{out}(A) = P_{in}(A) - kill(A) \cup gen(A)$ ($f_K(P_{in}(A))$)
- backward and must: $P_{out}(A) = \bigsqcup_{X \in Succ(A)} P_{in}(X)$
  $P_{in}(A) = P_{out}(A) - kill(A) \cup gen(A)$
- forward and may: $P_{in}(A) = \bigsqcup_{X \in Pre(A)} P_{out}(X)$
  $P_{out}(A) = P_{in}(A) - kill(A) \cup gen(A)$
- backward and may: $P_{out}(A) = \bigsqcup_{X \in Succ(A)} P_{in}(X)$
  $P_{in}(A) = P_{out}(A) - kill(A) \cup gen(A)$

Initialization

- Assume a Power Set Lattice $\mathcal{P}^S$
- Initialization with the smallest element
- General initializations with $\perp$ for all but start node n^1:
  - may: Initialization with $[(n^1, \iota, \emptyset), (n^2, \emptyset, \emptyset), \ldots, (n^{|N|}, \emptyset, \emptyset)]$ as empty set $\emptyset$ is the smallest element for each position
  - must: Initialization with $[(n^1, \iota, S), (n^2, S, S), \ldots, (n^{|N|}, S, S)]$ as universe of values S is the smallest element for each position in the inverse lattice
- Special (problem specific) initializations ι:
  - forward: $[(n^1, \iota, \perp), \ldots]$ as general initialization not defined before the start node
  - backward: $[\ldots, (n^e, \perp, \iota)]$ as general initialization not defined after the end node

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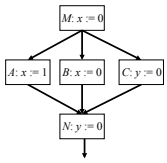
Example I

- Property P: $x = 1$ guaranteed?
- Universe Boolean, CPO Boolean Lattice
- Transfer functions: *true, false, id*
  - Statement A: $f_A = true$
  - Statement B: $f_B = false$
  - Statement C: $f_C = id$ i.e. does not change
- Let P_A, P_B, P_C, P_N be values of P after statements A,B,C, (P_{out}) and Let $\underline{P}_A$, $\underline{P}_B$, $\underline{P}_C$, $\underline{P}_N$ be values of P before statements A,B,C, (P_{in})
- Forward – must problem
  - it holds $\underline{P}_N = P_A \wedge P_B \wedge P_C$
  - Begin with $\underline{P}_{A,B,C,N} = true$ before statements (assumption $x = 1$)
  - Initialization $\underline{P}_M = false$ before statement M is $x := 1$
  - Iteration leads to fixed point $P_N = false$
- $x := 1\text{-}x$ more difficult:
  - Obviously, a naive transfer function *not* is not monotone
  - Conservative transfer function: $f = false$
  - Conservatively, $x = 1$ is not guaranteed any more by analysis in some cases where we (as humans) could see it holds



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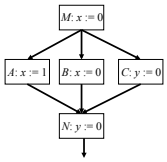
Example II

- Property P: $x = 1$ possible?
- Universe Boolean, CPO Boolean Lattice
- Transfer functions identical
- Forward – may problem
  - $\underline{P}_N = P_A \vee P_B \vee P_C$
  - Begin with $\underline{P}_{A,B,C,N} = false$ (assumption $x \neq 1$)
  - Initialization $\underline{P}_M = false$
  - Iteration leads to fixed point $P_N = true$
- Generalization:
  - Compute properties of several (all) variables in each step
  - Property: are variables equal to a specific constant or are variables actually compile time constants at a certain program point
  - Universe: Bit vector with a vector element for each variable
  - CPO induced by bit vector lattice



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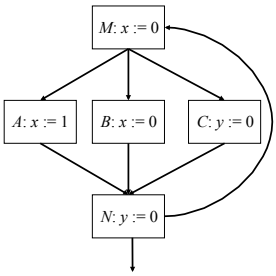
What does Data Flow Analysis?



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Path Graph

- For nodes $n \in N$ of $G=(N, E)$ define path graph $G'(n)=(N', E')$
  - For every path Π ending in n: $n' \in \Pi \Leftrightarrow n' \in N'$
  - $(n', n'') \in E' \Leftrightarrow (n', n'') \in \Pi$
- The path graph acyclic by definition
- Since the set of paths to a node n in G is possibly countable (iff G contains loops) the graph $G'(n)$ is in general not finite

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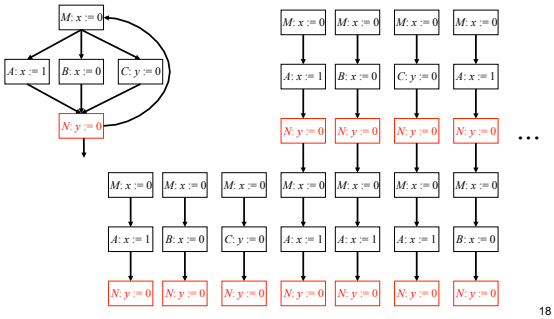
Example: Path Graph



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MFP and *MOP*

For a monotone DFA problem (set of equations) $DFE = (U, \sqsubseteq, F, \iota)$ and G

- Define: Minimum Fixed Point *MFP* is computed by iteratively applying F beginning with the smallest element in U

Let $DFE'(n) = (U, \sqsubseteq, F, \iota)$ and $G'(n)$ (same equations as DFE, applied to path graphs)

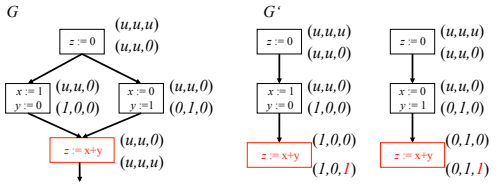
- Define: Meet Over all Paths *MOP* of DFE in (any arbitrary) node n is the minimum fix point *MFP* of $DFE'(n)$ in node n.
- MFP is equivalent with MOP, if f are distributive over $\sqcup$ in U,
- MFP is a conservative approximation of the MOP otherwise

Attention:

- It is not decidable if a path is actually executable
- Hence, MOP is already conservative approximation of the actual analysis result since some path may be not executable in any program run
- $MOP \neq MOEP$ (meet over all executable paths)

Example for $MFP(G) \neq MOP(G)$

Constant propagation: value vector: $(x,y,z) \in \{0,1,unknown\}^3$



Errors due to our DFA Method

- Call Graphs:
  - Nodes – Procedures, Edges – calls
  - Only a conservative approximation of actually possible calls, some calls represented in the call graph might never occur in any program run
  - Allows impossible paths like call → procedure → another call
- Data flow graph of a procedure:
  - Nodes – Statements (Expressions), Edges – (initial or essential) dependencies between them
  - Application of a monotone DFA framework computes *MFP* not *MOP*

Outline

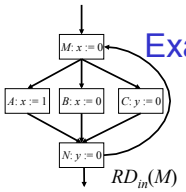
- Summary of Data Flow Analysis (yesterday's lecture)
- Problems left open
- Abstract interpretation idea

Problems left open

- How to derive the transfer functions for a DFA
- How to make sure they compute the intended result, i.e.,
  - MOP approximates the intended question, and
  - $MOP \sqsubseteq MFP$?

Example: Reaching Definitions (Must)

- Which „Definitions" (assignments) are guaranteed to reach a node A?
- Solution: Subset of all nodes $\{A_1 \ldots A_N\}$ containing an assignment to a variable reaching a access from this variable in node A
- Universe: e.g. a bit-vector representation $[\{false, true\}_1 \ldots \{false, true\}_N]$ of all possible subsets of nodes for each variable
- Forward, Must
- Schema:
$$RD_{in}(A) = \bigcap_{X \in Pre(A)} RD_{out}(X)$$
$$RD_{out}(A) = RD_{in}(A) - kill(A) \cup gen(A)$$
- Initialization:
  - Universe, i.e. all definitions ($\{A_1 \ldots A_N\}$) reach each program point
  - Start node: no definition reaches, i.e. $RD_{in}(A_1) = \emptyset$ for all variables
- A definition is generated ($gen(A)$) by assignment `x:=expr`
- A definition is removed ($kill(A)$) by an assignment `x:=expr'` to the same variable in another node



Example: MFP (for x)

$RD_{in}(M)$	$= \emptyset \cap RD_{out}(N)$	$= \emptyset$
$RD_{out}(M)$	$= RD_{in}(M) - \{M,A,B\} \cup \{M\}$	$= \{M\}$
$RD_{in}(A)$	$= RD_{out}(M)$	$= \{M\}$
$RD_{out}(A)$	$= RD_{in}(A) - \{M,A,B\} \cup \{A\}$	$= \{A\}$
$RD_{in}(B)$	$= RD_{out}(M)$	$= \{M\}$
$RD_{out}(B)$	$= RD_{in}(B) - \{M,A,B\} \cup \{B\}$	$= \{B\}$
$RD_{in}(C)$	$= RD_{out}(M)$	$= \{M\}$
$RD_{out}(C)$	$= RD_{in}(C)$	$= \{M\}$
$RD_{in}(N)$	$= RD_{out}(A) \cap RD_{out}(B) \cap RD_{out}(C)$	$= \emptyset$
$RD_{out}(N)$	$= RD_{in}(N)$	$= \emptyset$

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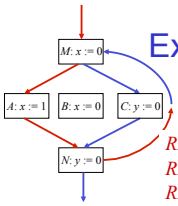

Example: Run (for x)

$RD_{in}(M)$	$= \emptyset$	$= \emptyset$
$RD_{out}(M)$	$= RD_{in}(M) - \{M,A,B\} \cup \{M\}$	$= \{M\}$
$RD_{in}(A)$	$= RD_{out}(M)$	$= \{M\}$
$RD_{out}(A)$	$= RD_{in}(A) - \{M,A,B\} \cup \{A\}$	$= \{A\}$
$RD_{in}(N)$	$= RD_{out}(A)$	$= \{A\}$
$RD_{out}(N)$	$= RD_{in}(N)$	$= \{A\}$
$RD_{in}(M)$	$= RD_{out}(N)$	$= \{A\}$
$RD_{out}(M)$	$= RD_{in}(M) - \{M,A,B\} \cup \{M\}$	$= \{M\}$
$RD_{in}(C)$	$= RD_{out}(M)$	$= \{M\}$
$RD_{out}(C)$	$= RD_{in}(C)$	$= \{M\}$
$RD_{in}(N)$	$= RD_{out}(C)$	$= \{M\}$
$RD_{out}(N)$	$= RD_{in}(N)$	$= \{M\}$

Problem left open

- How to make sure RD computes the correct result?
  - As intended by the problem
  - Exact result or a conservative approximation
- Actually, in the example program and the specific run RD behaves correctly:
  - Static analysis: $RD_{out}(N) = \emptyset$
  - Example run: $RD_{out}(N) = \{A\}$, $RD_{out}(N) = \{M\}$
  - $\{A\} \sqsubseteq \emptyset$ and $\{M\} \sqsubseteq \emptyset$
    - Recall that RD was a must problem, ascending on the downwards CPO induced by the lattice power set lattice
    - Hence $\sqsubseteq$ relation is the inverse set inclusion $\supseteq$ on the label sets
- How does this generalize?
  - For all runs, all programs, and for all dataflow problems
  - We cannot test all (countable) paths of all (countable) programs and all (infinitely many) possible dataflow problems

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Outline

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Abstract Interpretation

- Relates semantics of a programming language to an abstract analysis semantics
- Allows to compute or prove correct data flow equations (transfer functions)
  - Define abstraction of the execution semantics wrt. analysis problem
  - Define abstraction of execution traces to program points (in general, finite many contexts for each program point)
  - Prove that they are abstractions indeed.
- Idea even generalizes to other than dataflow analyses, as well (e.g., control flow analysis)

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Program Traces

- Each program run is defined by a trace $tr \in Labels^*$
- Traces are defined by the programming language semantics, e.g.,
  - tr[`stats;stat`] = tr[`stats`] $\oplus$ tr[`stat`]
  - tr[`assign`] := *label*(`assign`)
  - tr[`if expr then stats`$_1$ `else stats`$_2$] := *eval*[`expr`] = *true* ? tr[`stats`$_1$] : tr[`stats`$_2$]
  - tr[`while expr do stats od`] := *eval*[`expr`] = *true* ? tr[`stats`] $\oplus$ tr[`while … od`] : ε
- For instance, actual reaching definitions RD_{act} is a mapping RD_{act}: $Tr \rightarrow \mathcal{P}^{Labels}$

  i.e., for each trace (tr) a subset of definitions ($\subseteq \mathcal{P}^{Labels}$) reaches the end point of that trace

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Analysis Execution Semantics

- Non-standard semantics: actually expected analysis results are defined for traces as an abstraction of the program's execution semantics wrt. the analysis problem
- Standard semantics: a program's execution semantics is defined by the semantics of each programming language construct and their composition in the program
- There are only finitely many such constructs

RD_{act} Execution Semantics

- Given a program $G = (N, E, n^1)$
- $RD_{act} : Tr \to \mathcal{P}^{Labels}$
- Basis for recursive definitions: empty trace
  - no definition reaches the end of the empty trace
  - $RD_{act}(\varepsilon) := \emptyset$
- Analysis execution semantics of $tr \oplus label$ (trace tr expanded by the next dynamic step $label$)
  - recursively defined on analysis execution semantics of trace tr and analysis execution semantics of the static programming language construct of step $label$
  - $RD_{act}(tr \oplus label : S)$:=
    if (S = "`x:=expr`") then
    $RD_{act}(tr) - \{l \mid (l : \texttt{x:=expr}') \in N\} \cup \{label\}$
    else
    $RD_{act}(tr)$

Observation

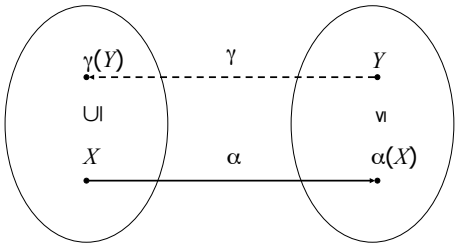
- Traces and semantics analysis values define a CPO $(U, \sqsubseteq)$
  - Universe U defined by $Tr \to \mathcal{P}^{Labels}$
  - Partial order $\sqsubseteq$: same program G, hence $Labels$, and same traces and subset of $\mathcal{P}^{Labels}$
  - Smallest element $\varepsilon \to \emptyset$
- Universe U is countable not finite
- Even if the semantic analysis function (e.g., RD_{act}) is monotone it is in general not continuous as universe not finite since $Tr(G)$ is not
- Then a solution to the analysis problem may exist, but cannot computed iteratively by applying the analysis function on the smallest element to fix point
  - Non-terminating program runs
  - Infinitely many different inputs

Solution

- Define an abstraction α of traces to make universe finite
- Perform an abstract analysis on the abstraction of traces
- Define a inverse concretization function γ to map results back to the semantic domain of the programming language
- α and γ should form a so called adjunction, or Galois connection, i.e.
  $\alpha(X) \leq Y \Leftrightarrow X \subseteq \gamma(Y)$
  - Mind the different domains of α and γ
  - Consequently, there are different partial order relations $\sqsubseteq$
  - $\subseteq$ on analysis execution semantics domain, countable and
  - $\leq$ on abstract analysis domain, finite
- Showing that abstraction and concretization form a Galois connection is one of our proof obligations to prove the analysis correct

Galois Connections

$\alpha(X) \leq Y \Leftrightarrow X \subseteq \gamma(Y)$



Countable Execution $CPO\ (U, \subseteq)$ Finite Abstraction $CPO\ (U', \leq)$

How to define the analysis?

- Take (any) abstract analysis F function abstracting (i.e., gives larger results that) $\alpha \bullet Act \bullet \gamma : U' \to U'$ where Act is the actual analysis execution semantics function
- Analysis terminates if $(U', \leq)$ a CPO and F monotone
- Analysis is conservative if Act is monotone and (α, γ) a Galois connection
- Then conservative approximation computable by fixed point iteration
- It holds for the minimum fix points MFP:
  $\alpha(MFP(Act)) \leq MFP(\alpha \bullet Act \bullet \gamma) \leq MFP(F)$

Reaching Definitions (α)

- Let Tr_{label} be the set of all traces ending with program point *label*
  $Tr_{label} = \{ tr \oplus label \mid label \in Label \wedge tr \oplus label$ is an admissible trace of $G \}$
- We abstract a set $Tr_{label} \in \mathcal{P}^{Tr}$ with that program point $label \in Label$

  $\alpha: \mathcal{P}^{Tr} \rightarrow Label$
  $\alpha(Tr_{label}) = label$
- Concrete and abstract analysis value domains $\mathcal{P}^{Label}$ are the same:
  - Let $RD_{act}(tr \oplus label) \in \mathcal{P}^{Label}$ be the set of definitions reaching *label*
  - Let $RD(label) \in \mathcal{P}^{Label}$ be the set of reaching definitions analyzed for *label*
- We abstract the analysis execution semantics of the trace $tr \in Tr_{label}$: $RD(tr)$ with the abstract analysis results $RD(label)$ of the program point *label*

  $\alpha: \mathcal{P}^{Label} \rightarrow \mathcal{P}^{Label}$
  $\alpha(RD_{act}(tr)) = RD(label)$ *iff* $tr \in Tr_{label}$

Reaching Definitions (γ)

- Conversely, we concretize each program point *label* with the set of all traces ending in *label*
- The concretization function on labels is

  $\gamma: Label \rightarrow \mathcal{P}^{Tr}$
  $\gamma(label) = Tr_{label}$
- Consequently, we concretize the abstract analysis results $RD(label)$ of a program point *label* by assuming it holds for any of the traces $tr \in Tr_{label}$:

  $\gamma: \mathcal{P}^{Label} \rightarrow \mathcal{P}^{Label}$
  $\gamma(RD(label)) = \forall\ tr \in Tr_{label} : RD(tr) = RD(label)$

RD Analysis Semantics

- Given a program $G = (N, E, n^1)$
- $RD : Lable \rightarrow \mathcal{P}^{Labels}$
- Basis for recursive definitions:
  - Empty trace abstraction: starting point of the program n^1
  - no definition reaches n^1
  - $RD_{in}(n^1) := \emptyset$
- Analysis semantics at *label* which is actually $\alpha \bullet RD_{act} \bullet \gamma$
  - recursively defined on analysis semantics of abstractions of predecessor traces *tr* (predecessor labels)
    $RD_{in}(label : S) := \bigcap_{p \in Pre(label)} RD_{out}(p)$
  - analysis abstraction of the execution semantics of the static programming language construct of step *label* (transfer function)
    $RD_{out}(label : S)$:=
    if ($S =$ "`x:=expr`") then
    $RD_{in}(lable) - \{l \mid (l : \texttt{x:=expr}') \in N\} \cup \{label\}$
    else
    $RD_{in}(label)$

Correctness of Analysis Abstraction

- Use structural induction over all programs
- Compare execution semantics and analysis semantics (transfer functions) of program constructs
- Basis:
  - Claim holds for the empty trace: each program's starting point is abstracted correctly: $RD_{in}(n^1) = \emptyset$, $Rd_{act}(\varepsilon) = \emptyset$
- Step:
  - Given a trace $tr \oplus label$ and its abstraction *label*
  - Provided $RD_{in}(label : S)$ is a correct abstraction of $RD_{act}(tr)$
  - Then $RD_{out}(label : S)$ is a correct abstraction of $RD_{act}(tr \oplus label)$:
    $\forall\ tr \in \gamma(label)$: $\alpha(RD_{act}(\gamma(RD_{in}(label)))) \subseteq RD_{out}(label)$
  - Distinguish cases of each program construct and transfer function
  - Here trivial as RD_{act} and RD are identical (and monotone)

RD Proof Obligations

- To show (i): (α, γ) is a Galois connection (obvious)
- To show (ii): $\alpha \bullet RD_{act} \bullet \gamma$ is abstracted with RD i.e., $\alpha \bullet RD_{act} \bullet \gamma \leq RD$
- Proof (sketch): for each node n of G
  - By our definition of γ, $\gamma(label) = Tr_{label}$ of corresponds to path graph of G in $n = (label{:}S)$
  - By our definition of RD_{act}, $\alpha \bullet RD_{act} \bullet \gamma$ in a node n is MFP of RD of path graph of G in n
  - MFP of RD of path graph of G in n is MOP of G in n
  - $MOP \leq MFP$ of RD

General Proof Obligations

- To show (i): (α, γ) is a Galois connection
- To show (ii): $\alpha \bullet Act \bullet \gamma$ is abstracted with F i.e., $\alpha \bullet Act \bullet \gamma \leq F$
- Proof (sketch): for each node n of G
  - By our definition of γ, $\gamma(label) = Tr_{label}$ of corresponds to path graph of G in $n = (label{:}S)$
  - By our definition of Act and F, $\alpha \bullet Act \bullet \gamma\,(n) \subseteq F(n)$ in every node n (sufficient to show this for every $f_K(n)$)
  - Then $\alpha \bullet Act \bullet \gamma$ in a node n is MFP of F of path graph of G in n
  - MFP of F of path graph of G in n is MOP of G in n
  - $MOP \leq MFP$ of F