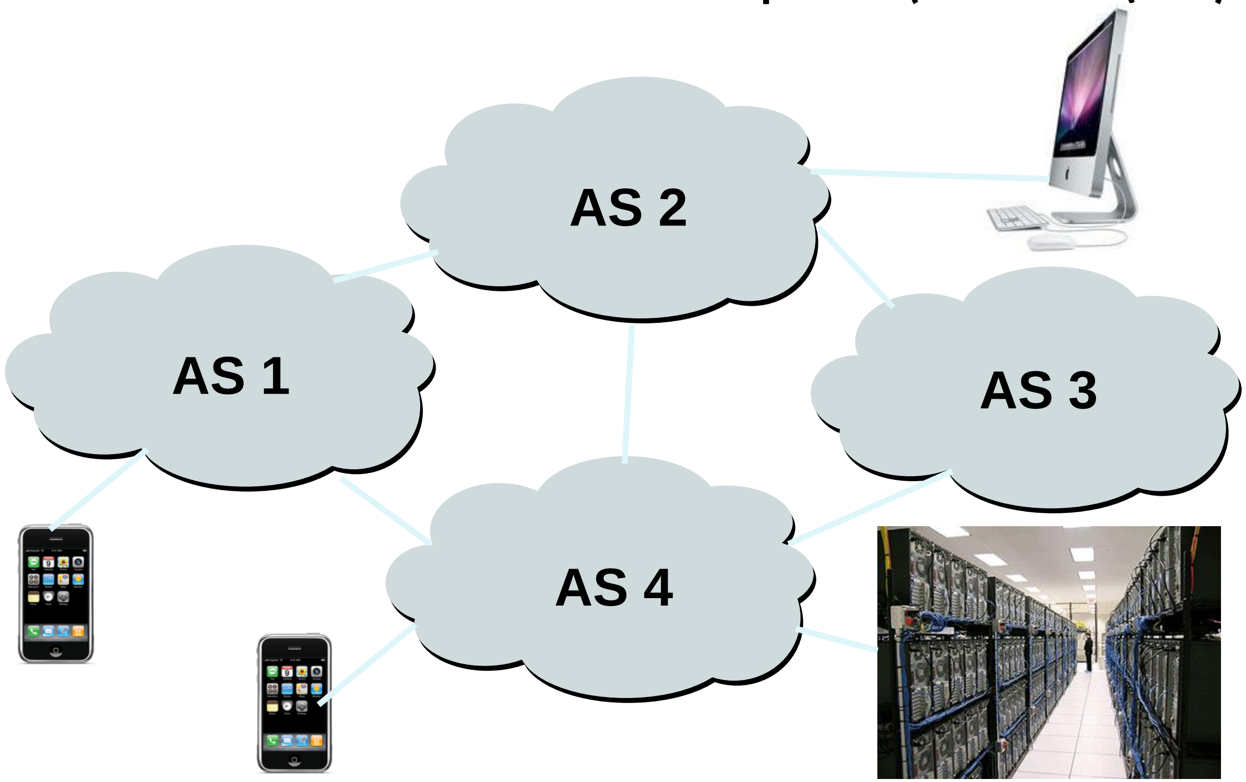
TDTS21 Advanced Networking

Lecture: Transport, including TCP and congestion control …

Based on slides from D. Choffnes, P. Gill, and S. Katti
Revised Spring 2019, 2021 by N. Carlsson

Holding the Internet Together

- Distributed cooperation for resource allocation
  - BGP: what end-to-end *paths* to take (for ~O(100K) ASes)
  - TCP: what *rate* to send over each path (for ~O(1B) hosts)



What Problem Does a Protocol Solve?

- BGP path selection
  - Select a path that each AS on the path is willing to use
  - Adapt path selection in the presence of failures
- TCP congestion control
  - Prevent congestion collapse of the Internet
  - Allocate bandwidth fairly and efficiently

What Problem Does a Protocol Solve?

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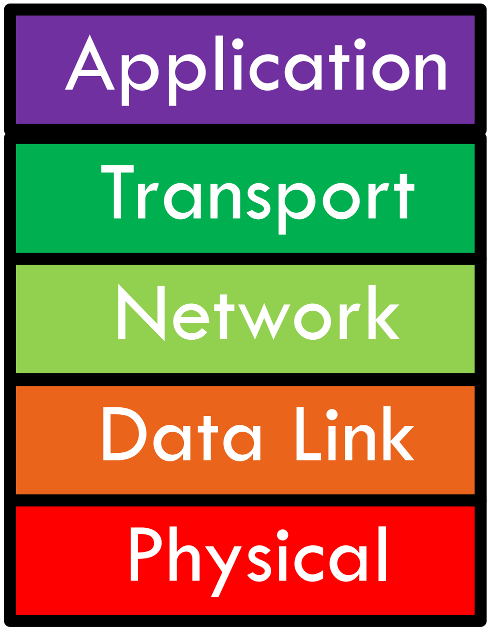
Today, we will focus on TCP (and UDP) …

Transport Layer

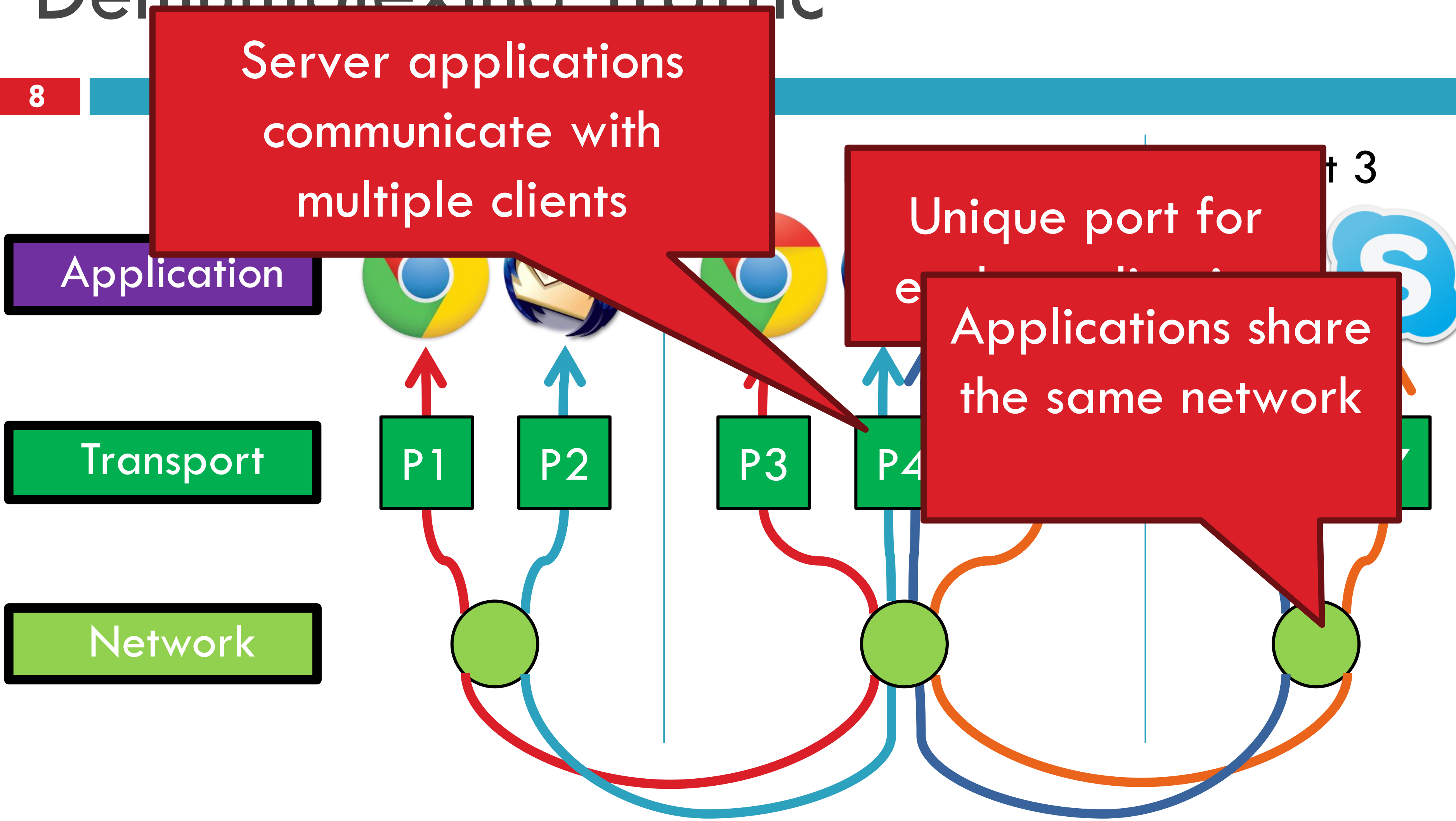


- Function:
  - Demultiplexing of data streams
- Optional functions:
  - Creating long lived connections
  - Reliable, in-order packet delivery
  - Error detection
  - Flow and congestion control
- Key challenges:
  - Detecting and responding to congestion
  - Balancing fairness against high utilization

Outline

- UDP
- TCP
- Congestion Control
- Evolution of TCP
- Problems with TCP

Demultiplexing Traffic



Endpoints identified by <*src_ip, src_port, dest_ip, dest_port*>

User Datagram Protocol (UDP)

0	16 31
Source Port	Destination Port
Payload Length	Checksum

- Simple, connectionless datagram
- Port numbers enable demultiplexing
  - 16 bits = 65535 possible ports
  - Port 0 is invalid
- Checksum for error detection
  - Detects (some) corrupt packets
  - Does not detect dropped, duplicated, or reordered packets

Uses for UDP

- Invented after TCP
  - Why?
- Not all applications can tolerate TCP
- Custom protocols can be built on top of UDP
  - Reliability? Strict ordering?
  - Flow control? Congestion control?
- Examples
  - DNS, …
  - RTMP, real-time media streaming (e.g. voice, video)
  - Facebook datacenter protocol
  - QUIC

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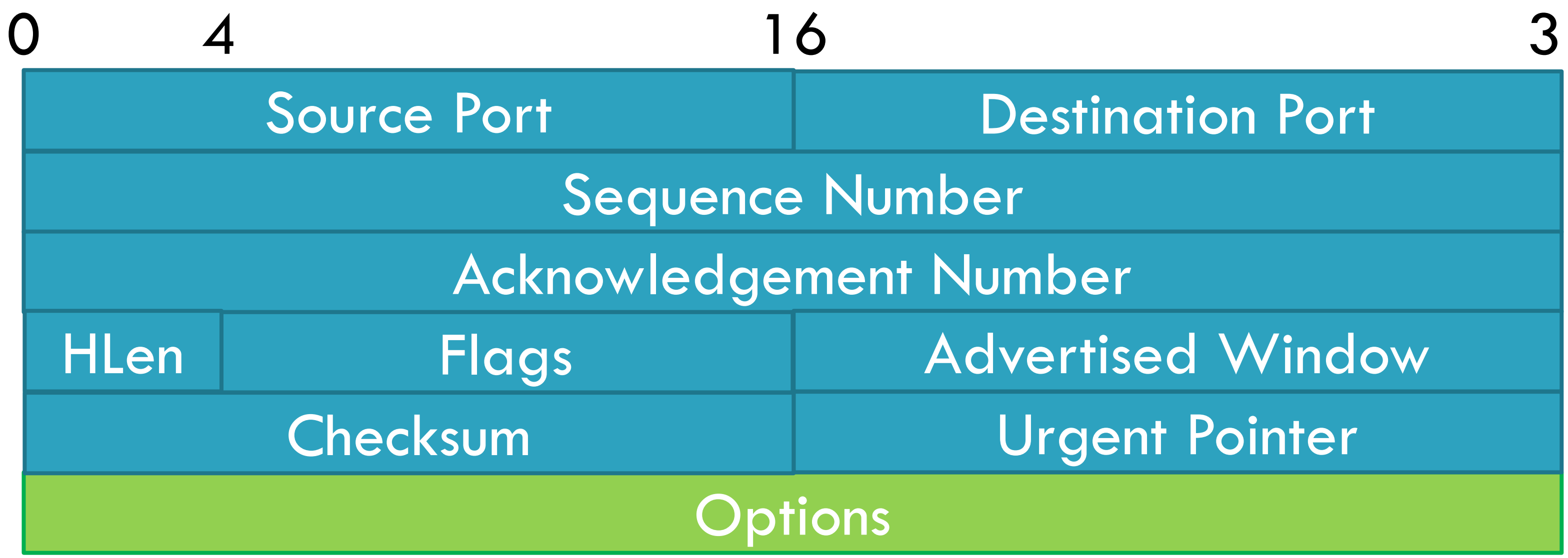
Transmission Control Protocol

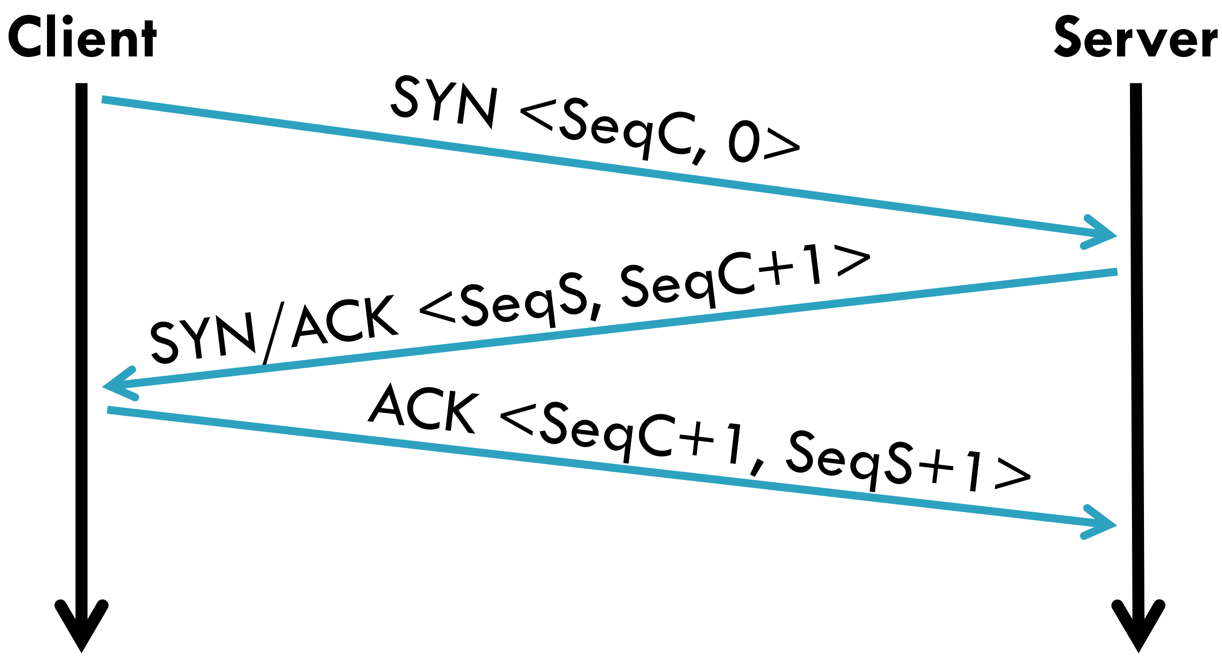
- Reliable, in-order, bi-directional byte streams
  - Port numbers for demultiplexing
  - Virtual circuits (connections)
  - Flow control
  - Congestion control, approximate fairness



Connection Setup

- Why do we need connection setup?
  - To establish state on both hosts
  - Most important state: sequence numbers
    - Count the number of bytes that have been sent
    - Initial value chosen at random
- Important TCP flags (1 bit each)
  - SYN – synchronization, used for connection setup
  - ACK – acknowledge received data
  - FIN – finish, used to tear down connection

Three Way Handshake



- Each side:
  - Notifies the other of starting sequence number
  - ACKs the other side's starting sequence number

Connection Setup Issues

- Connection confusion
  - How to disambiguate connections from the same host?
  - Random sequence numbers
- Source spoofing
  - Kevin Mitnick
  - Need good random number generators!
- Connection state management
  - Each SYN allocates state on the server
  - SYN flood = denial of service attack
  - Solution: SYN cookies

U.S. Department of Justice
United States Marshals Service

WANTED
BY U.S. MARSHALS

NOTICE TO ARRESTING AGENCY: Before arrest, validate warrant through National Crime Information Center (NCIC).
United States Marshals Service NCIC entry number: (NIC/ W721460021).

NAME: .... MITNICK, KEVIN DAVID
AKS (S): .... MITNIK, KEVIN DAVID
MERRILL, BRIAN ALLEN

DESCRIPTION:
Sex: .... MALE
Race: .... WHITE
Place of Birth: .... VAN NUYS, CALIFORNIA
Date(s) of Birth: 08/06/63; 10/18/70
Height: .... 5'11"
Weight: 190
Eyes: .... BLUE
Hair: .... BROWN
Skintone: .... LIGHT
Scars, Marks, Tattoos: .... NONE KNOWN
Social Security Number (s): 550-39-5695
NCIC Fingerprint Classification: DOPM20PM13DIPM19PM09

ADDRESS AND LOCALE: KNOWN TO RESIDE IN THE SAN FERNANDO VALLEY AREA OF CALIFORNIA AND LAS VEGAS, NEVADA

WANTED FOR: VIOLATION OF SUPERVISED RELEASE
ORIGINAL CHARGES: POSSESSION UNAUTHORIZED ACCESS DEVICE; COMPUTER FRAUD
Warrant Issued: CENTRAL DISTRICT OF CALIFORNIA
Warrant Number: 9312-1112-0154-C

DATE WARRANT ISSUED: NOVEMBER 10, 1992

MISCELLANEOUS INFORMATION: SUBJECT SUFFERS FROM A WEIGHT PROBLEM AND MAY HAVE EXPERIENCED WEIGHT GAIN OR WEIGHT LOSS
VEHICLE/TAG INFORMATION: NONE KNOWN OFTEN USES PUBLIC TRANSPORTATION

If arrested or whereabouts known, notify the local United States Marshals Office, (Telephone: 213-894-2485).
If no answer, call United States Marshals Service Communications Center in McLean Virginia.
Telephone (800)336-0102: (24 hour telephone contact) NLETS access code is VAUSMOOOO.

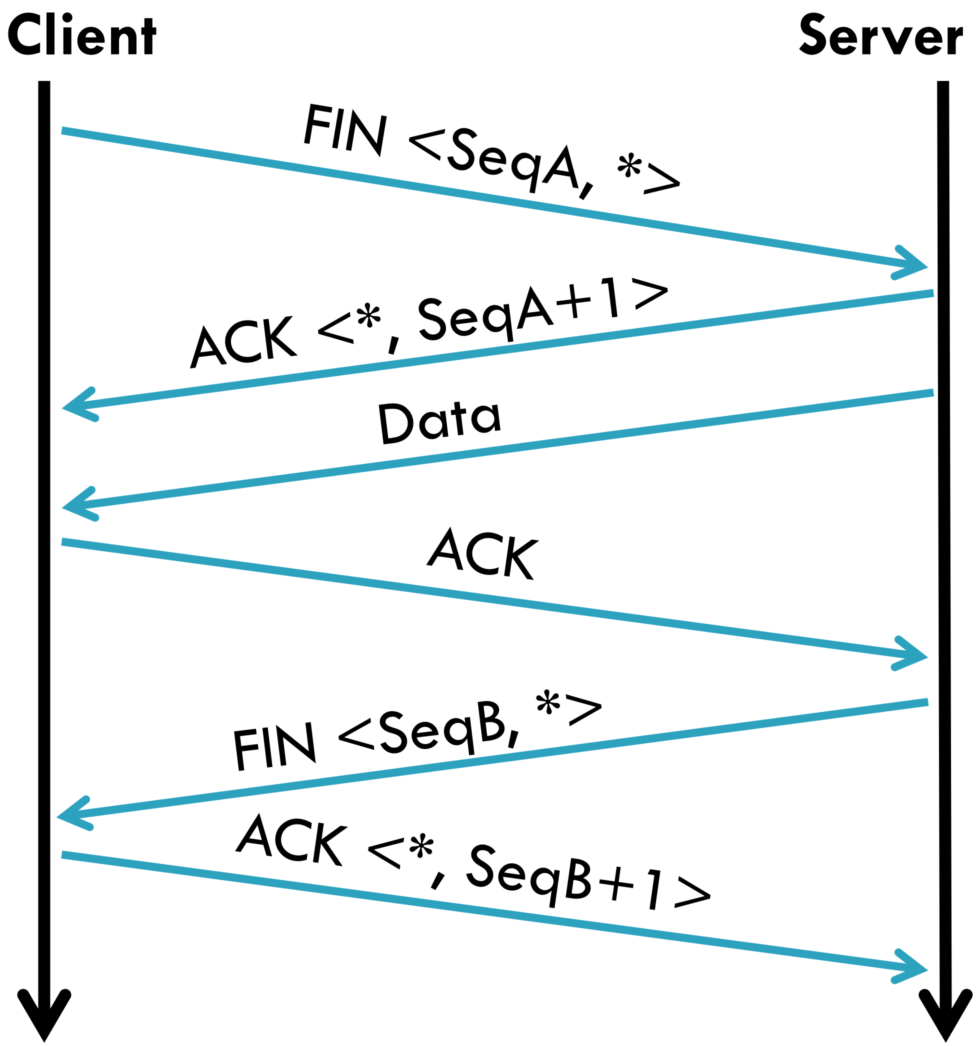
Form USM-132
(Rev. 3/2/82)

PRIOR EDITIONS ARE OBSOLETE AND NOT TO BE USED

November 1992

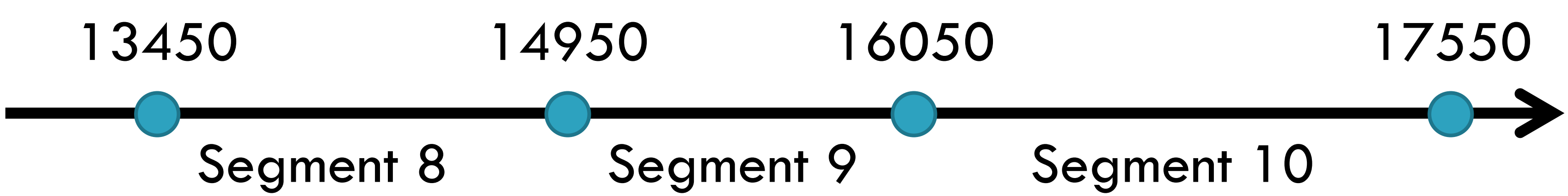
Connection Tear Down

- Either side can initiate tear down
- Other side may continue sending data
  - Half open connection
  - *shutdown()*
- Acknowledge the last FIN
  - Sequence number + 1

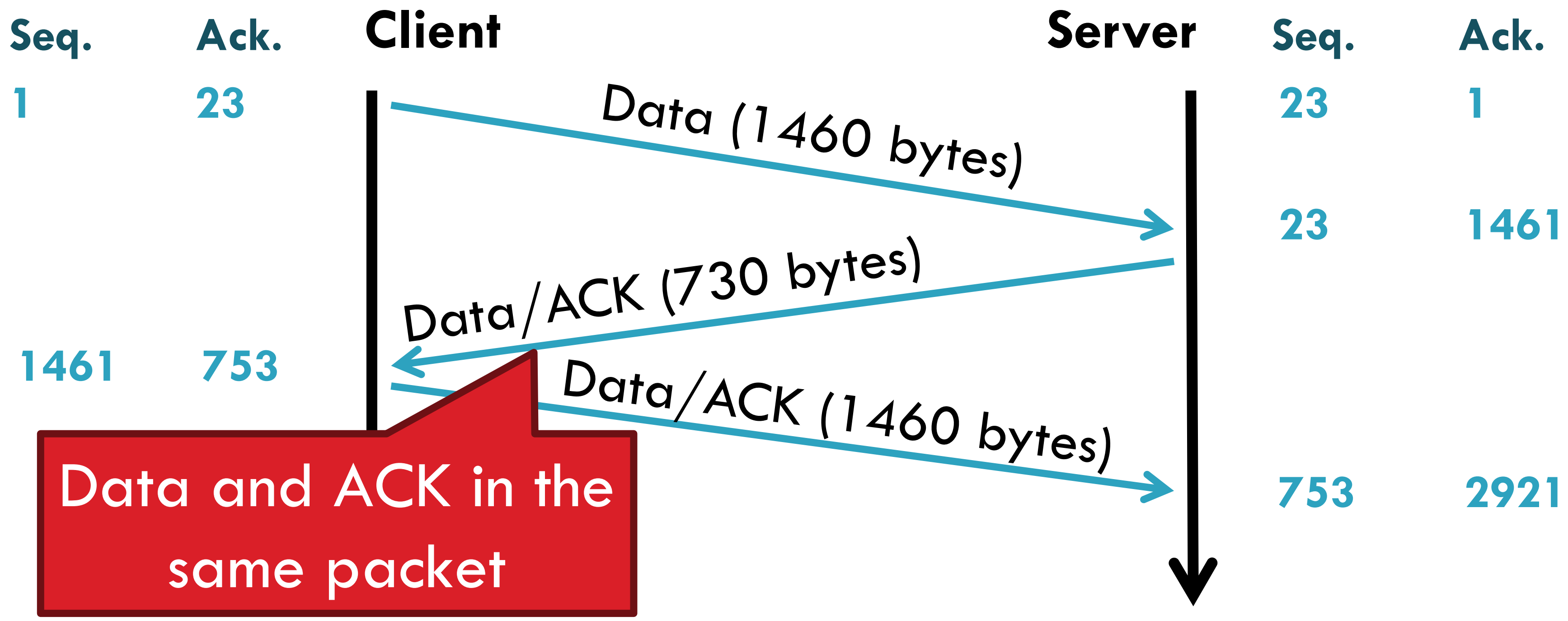


Sequence Number Space

- TCP uses a byte stream abstraction
  - Each byte in each stream is numbered
  - 32-bit value, wraps around
  - Initial, random values selected during setup
- Byte stream broken down into segments (packets)
  - Size limited by the Maximum Segment Size (MSS)
  - Set to limit fragmentation
- Each segment has a sequence number



Bidirectional Communication

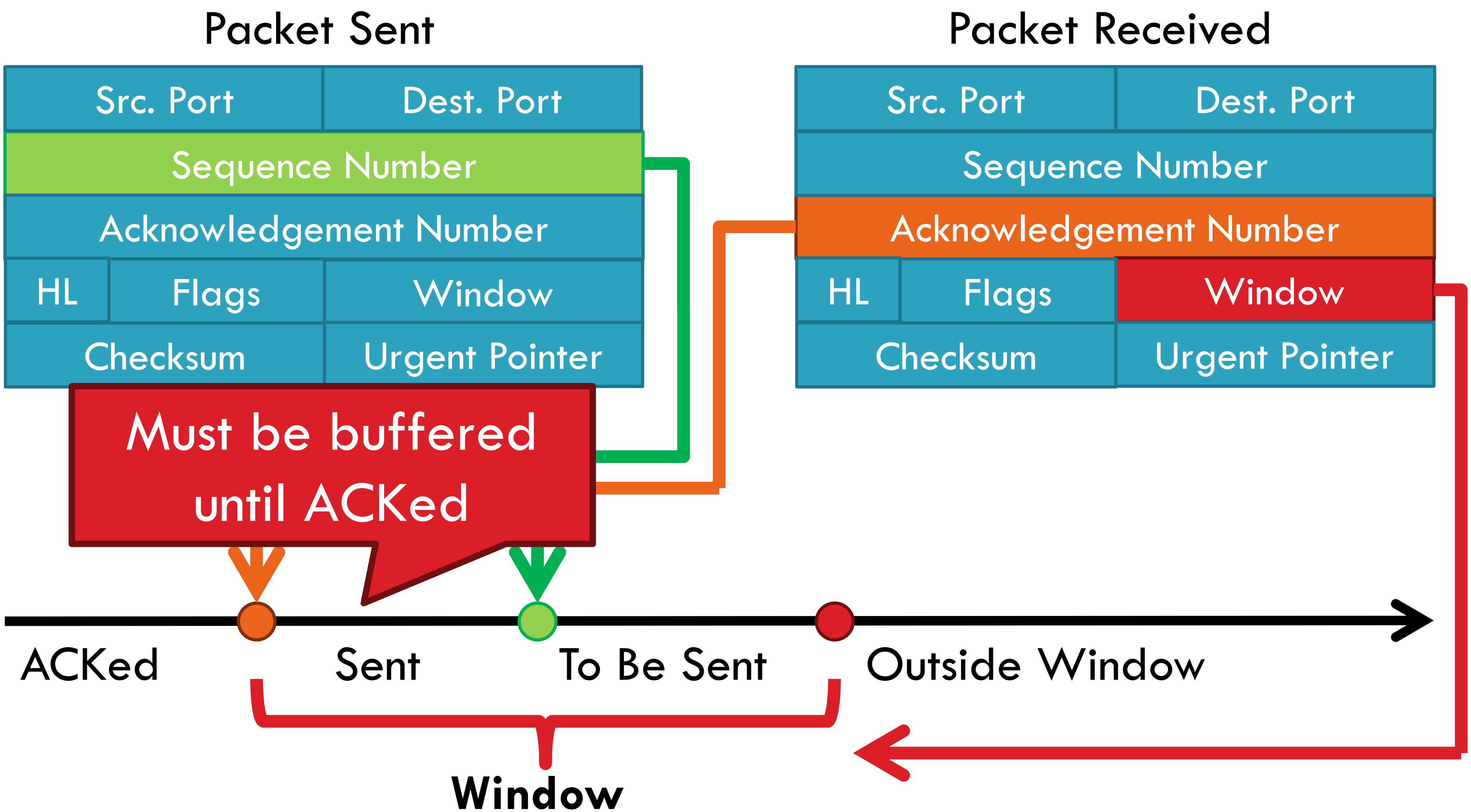


- Each side of the connection can send and receive
  - Different sequence numbers for each direction

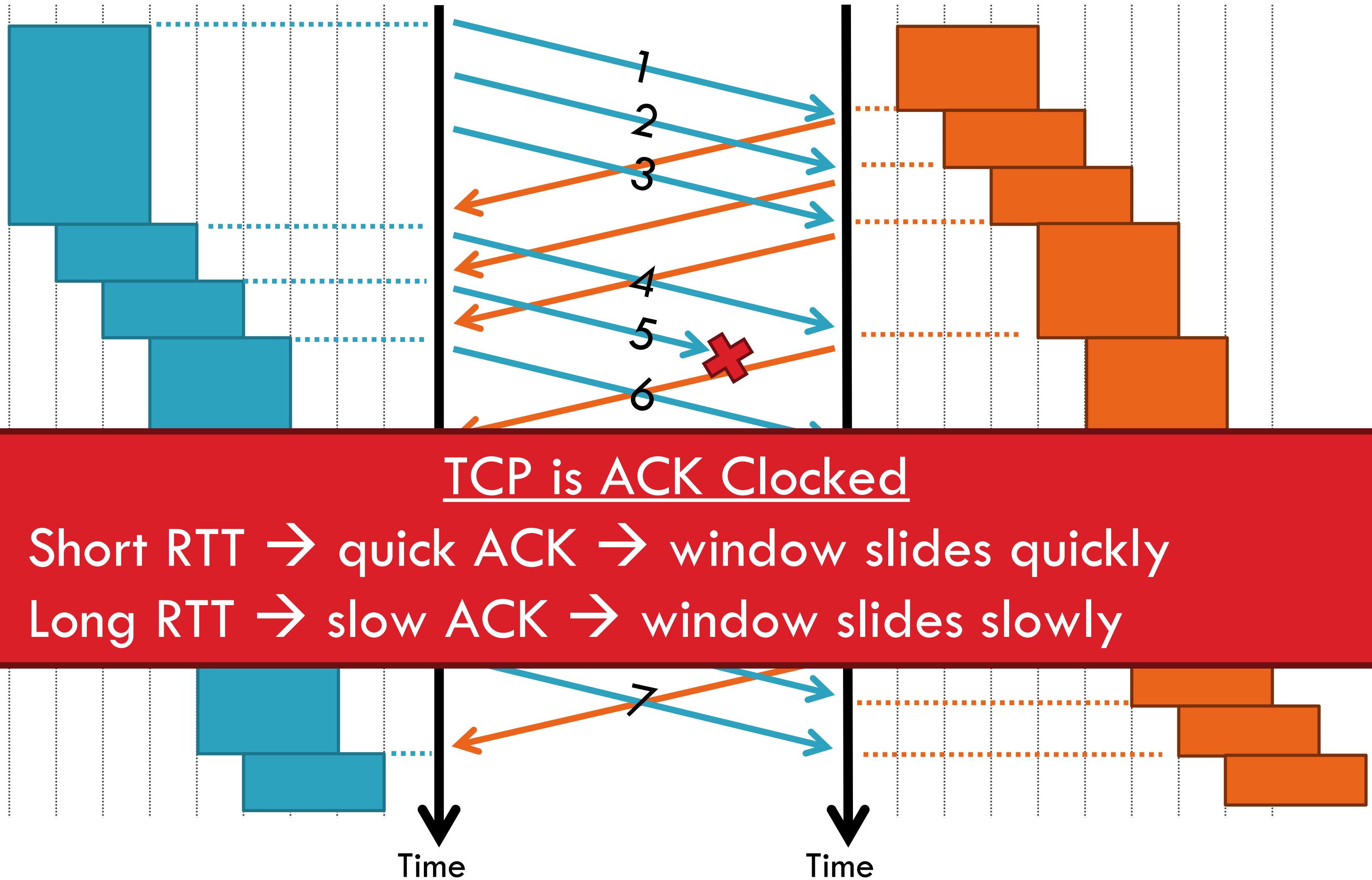
Flow Control

- Problem: how many packets should a sender transmit?
  - Too many packets may overwhelm the receiver
  - Size of the receiver’s buffers may change over time
- Solution: sliding window
  - Receiver tells the sender how big their buffer is
  - Called the advertised window
  - For window size *n*, sender may transmit *n* bytes without receiving an ACK
  - After each ACK, the window slides forward

Flow Control: Sender Side



Sliding Window Example



<u>TCP is ACK Clocked</u>

- Short RTT → quick ACK → window slides quickly
- Long RTT → slow ACK → window slides slowly

Observations

- Throughput is ~ w/RTT
- Sender has to buffer all unacknowledges packets, because they may require retransmission
- Receiver may be able to accept out-of-order packets, but only up to buffer limits

Outline

- UDP
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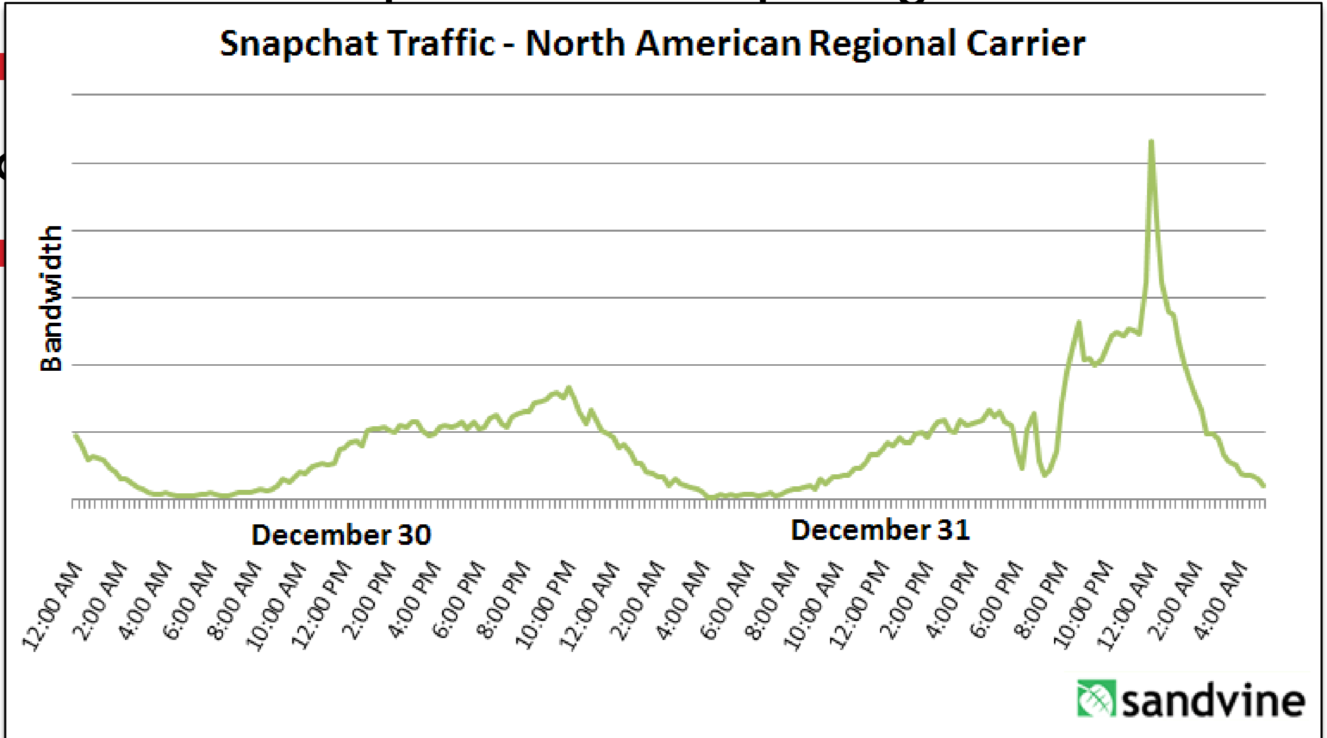
What is Congestion?

- Load on the network is higher than capacity
  - Capacity is not uniform across networks
    - Modem vs. Cellular vs. Cable vs. Fiber Optics
  - There are multiple flows competing for bandwidth
  - L



Why is Congestion Bad?

- Results in packet loss
  - Routers have finite buffers
  - Internet traffic is bursty, no buffer can prevent all drops
  - When routers get overloaded, packets will be dropped
- Practical consequences
  - Router queues build up, delay increases
  - Wasted bandwidth from retransmissions
  - Low network "goodput"

CONGESTION AVOIDANCE AND CONTROL

VAN JACOBSON '88

Main contributions

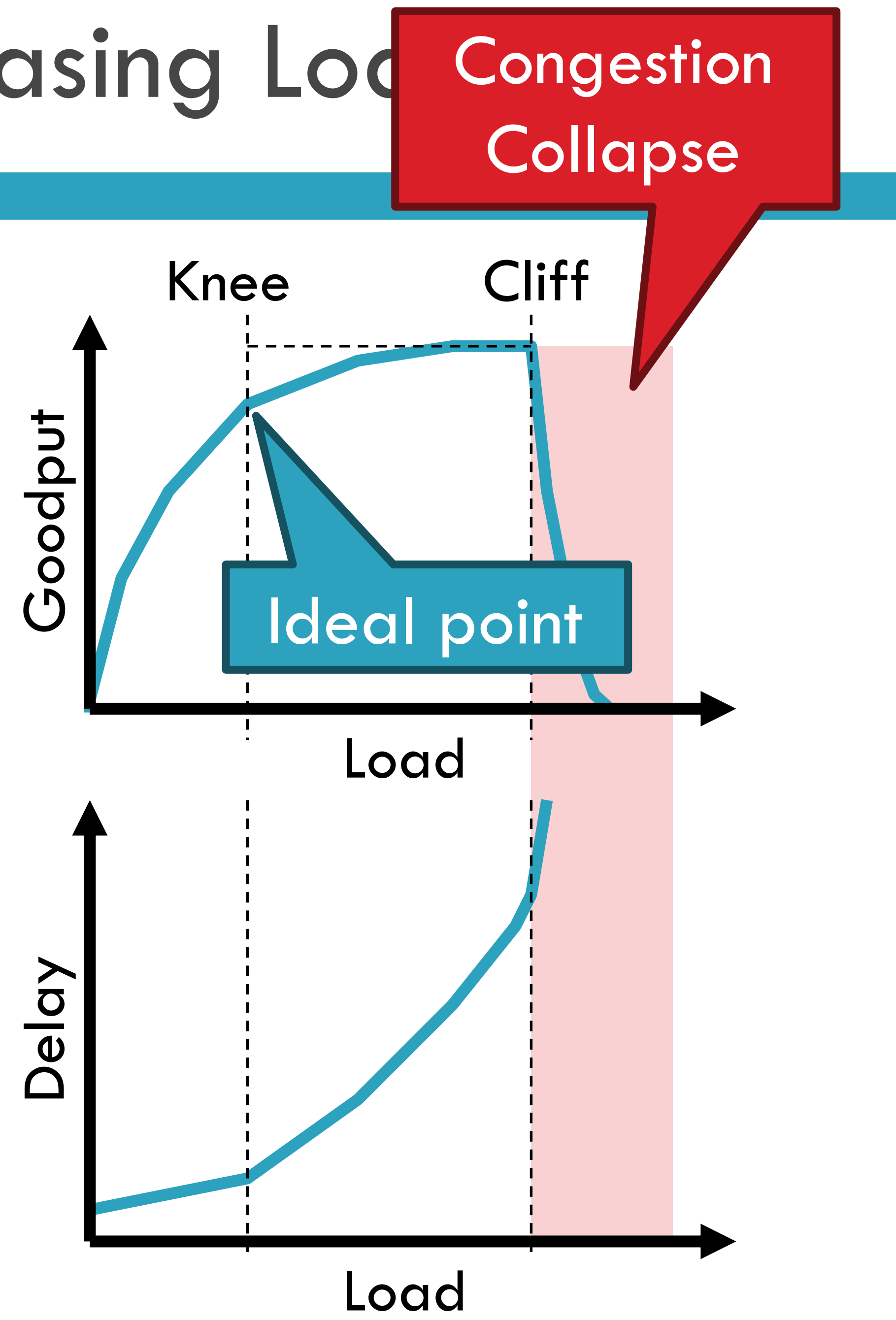
Seven new algorithms:

1. RTT Variance estimation
2. Exponential retransmit timer backoff
3. Slow-start
4. More aggressive receiver ack policy
5. Dynamic window sizing on congestion
6. Karn's algorithm
7. Fast retransmit

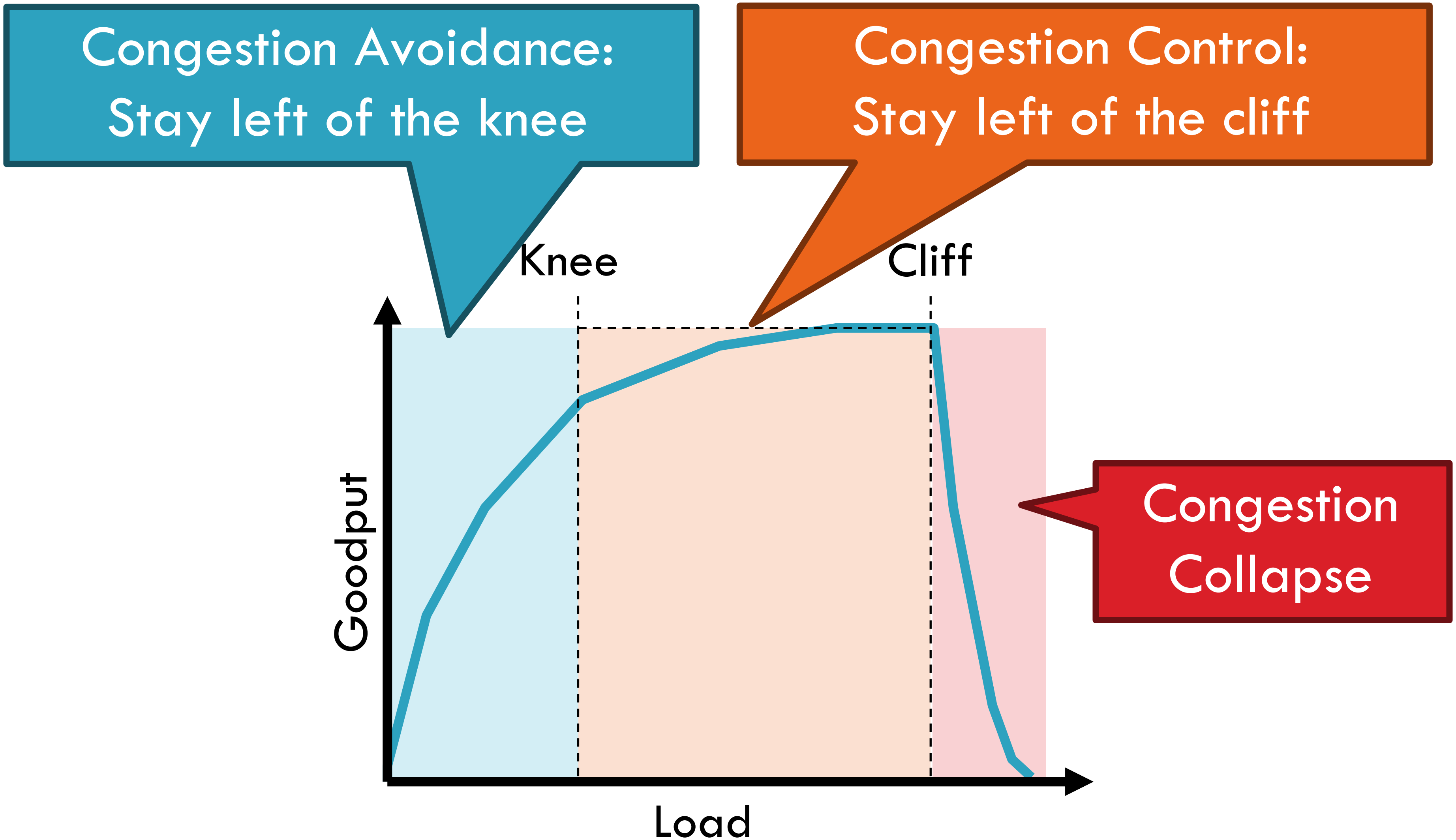
Paper explores the first 5.

The Danger of Increasing Load

- Knee – point after which
  - Throughput increases very slow
  - Delay increases fast
- In an M/M/1 queue
  - Delay = 1/(1 – utilization)
- Cliff – point after which
  - Throughput → 0
  - Delay → ∞



Cong. Control vs. Cong. Avoidance



Advertised Window, Revisited

- Does TCP's advertised window solve congestion?

NO

- The advertised window only protects the receiver
- A sufficiently fast receiver can max the window
  - What if the network is slower than the receiver?
  - What if there are other concurrent flows?
- Key points
  - Window size determines send rate
  - Window must be adjusted to prevent congestion collapse

Goals of Congestion Control

1. Adjusting to the bottleneck bandwidth
2. Adjusting to variations in bandwidth
3. Sharing bandwidth between flows
4. Maximizing throughput

General Approaches

- Do nothing, send packets indiscriminately
  - Many packets will drop, totally unpredictable performance
  - May lead to congestion collapse
- Reservations
  - Pre-arrange bandwidth allocations for flows
  - Requires negotiation before sending packets
  - Must be supported by the network
- Dynamic adjustment
  - Use probes to estimate level of congestion
  - Speed up when congestion is low
  - Slow down when congestion increases
  - Messy dynamics, requires distributed coordination

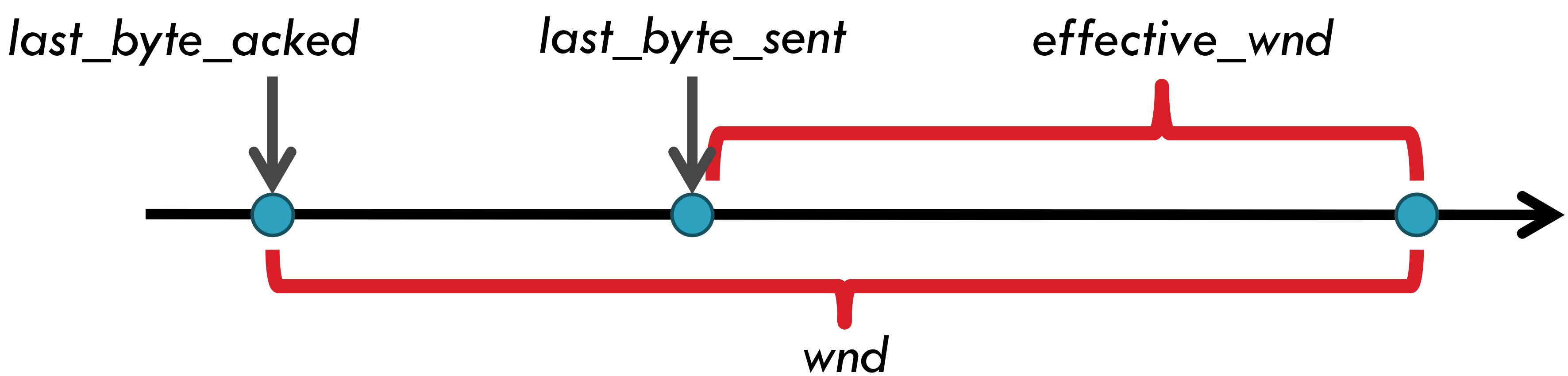
TCP Congestion Control

- Each TCP connection has a window
  - Controls the number of unACKed packets
- Sending rate is ~ window/RTT
- Idea: vary the window size to control the send rate
- Introduce a congestion window at the sender
  - Congestion control is sender-side problem

Congestion Window (*cwnd*)

- Limits how much data is in transit
- Denominated in bytes

1. *wnd* = min(*cwnd*, *adv_wnd*);
2. *effective_wnd* = *wnd* –
   (*last_byte_sent* – *last_byte_acked*);



Two Basic Components

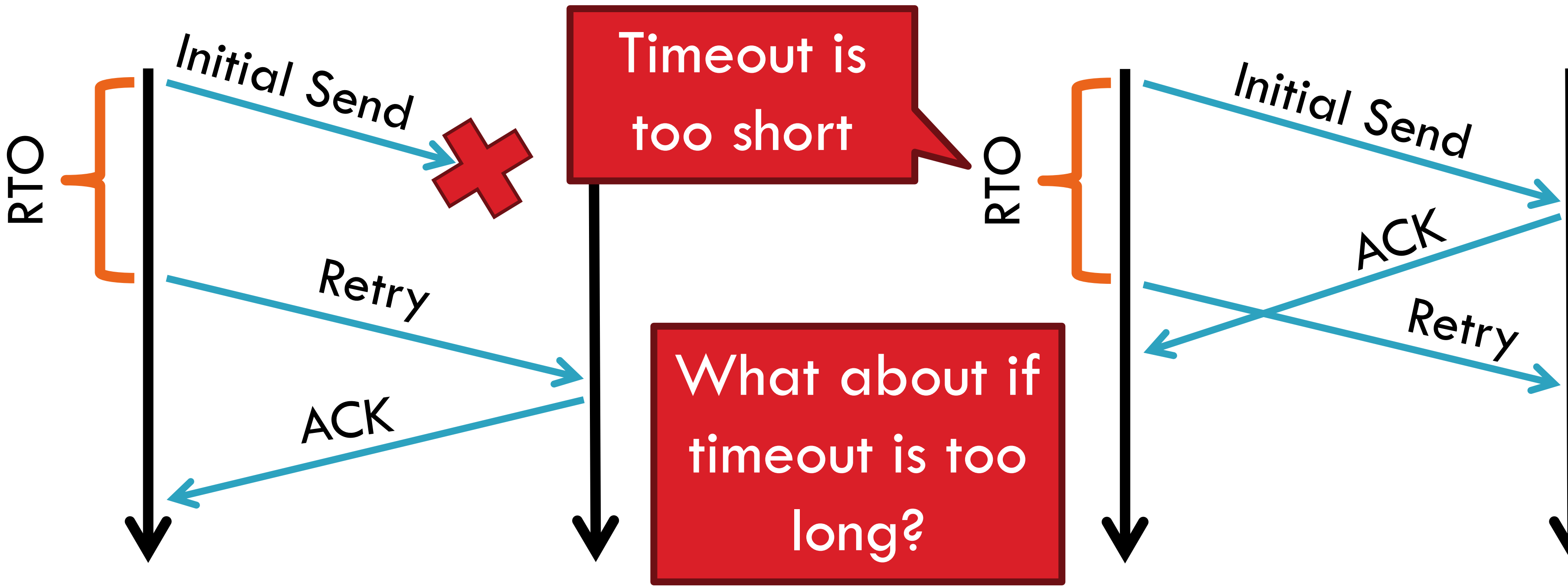
1. Detect congestion
   - Packet dropping is most reliably signal
     - Delay-based methods are hard and risky
   - How do you detect packet drops? ACKs
     - Timeout after not receiving an ACK
     - Several duplicate ACKs in a row (ignore for now)
2. Rate adjustment algorithm
   - Modify *cwnd*
   - Probe for bandwidth
   - Responding to congestion

Error Detection

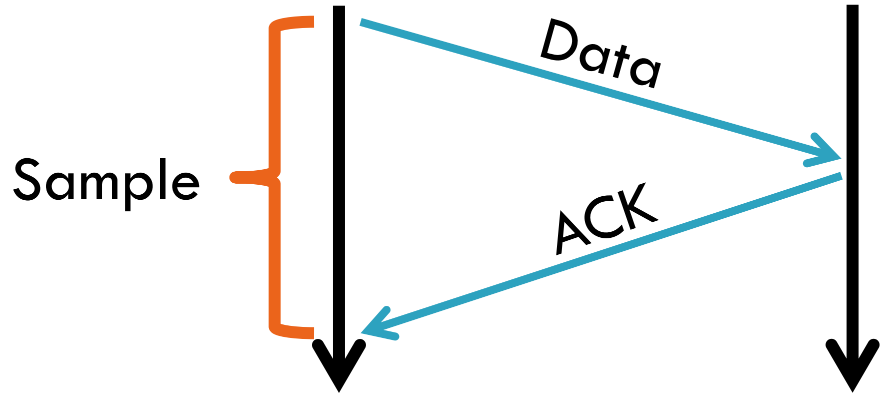
- Checksum detects (some) packet corruption
  - Computed over IP header, TCP header, and data
- Sequence numbers catch sequence problems
  - Duplicates are ignored
  - Out-of-order packets are reordered or dropped
  - Missing sequence numbers indicate lost packets
- Lost segments detected by sender
  - Use timeout to detect missing ACKs
  - Need to estimate RTT to calibrate the timeout
  - Sender must keep copies of all data until ACK

Retransmission Time Outs (RTO)

- Problem: time-out is linked to round trip time

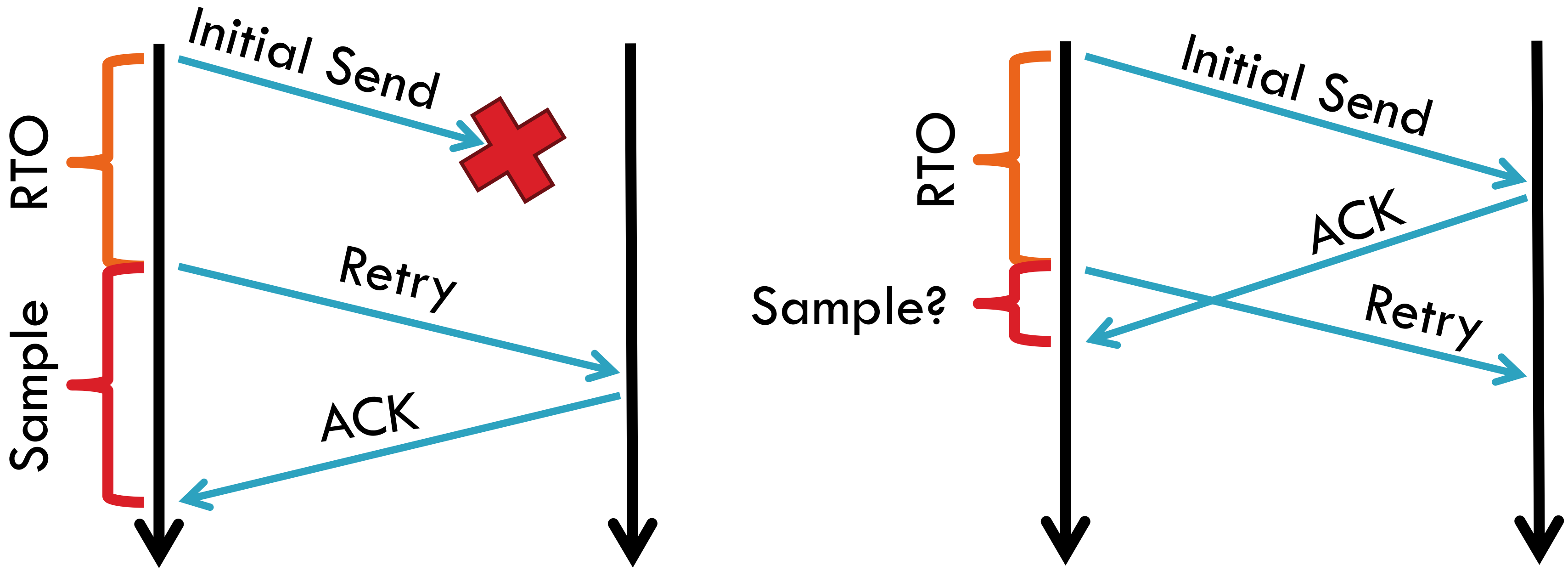


Round Trip Time Estimation



- Original TCP round-trip estimator
  - RTT estimated as a moving average
  - new_rtt = α (old_rtt) + (1 – α)(new_sample)
  - Recommended α: 0.8-0.9 (0.875 for most TCPs)
- RTO = function of new_rtt and new_dev_rtt

RTT Sample Ambiguity

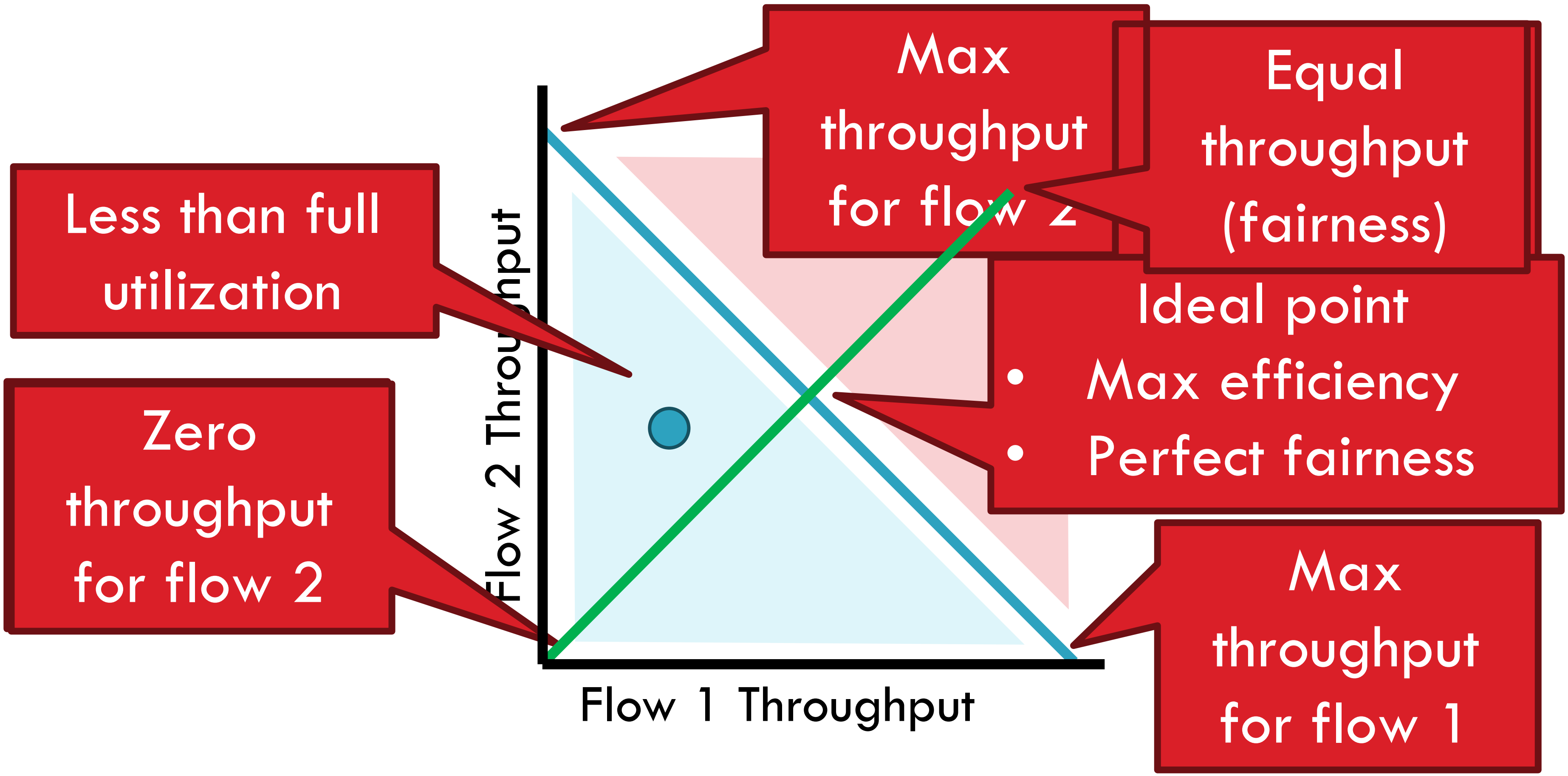


- Karn's algorithm: ignore samples for retransmitted segments

Rate Adjustment

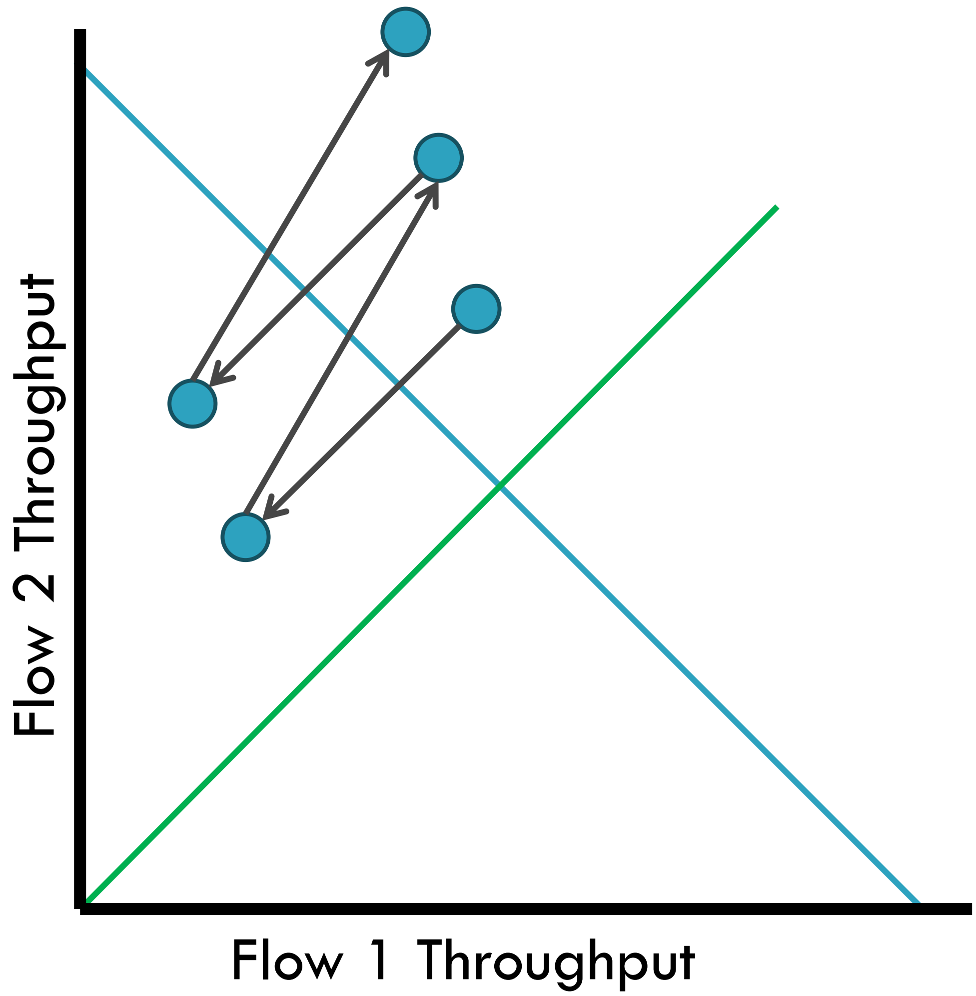
- Recall: TCP is ACK clocked
  - Congestion = delay = long wait between ACKs
  - No congestion = low delay = ACKs arrive quickly
- Basic algorithm
  - Upon receipt of ACK: increase cwnd
    - Data was delivered, perhaps we can send faster
    - *cwnd* growth is proportional to RTT
  - On loss: decrease cwnd
    - Data is being lost, there must be congestion
- Question: increase/decrease functions to use?

Utilization and Fairness



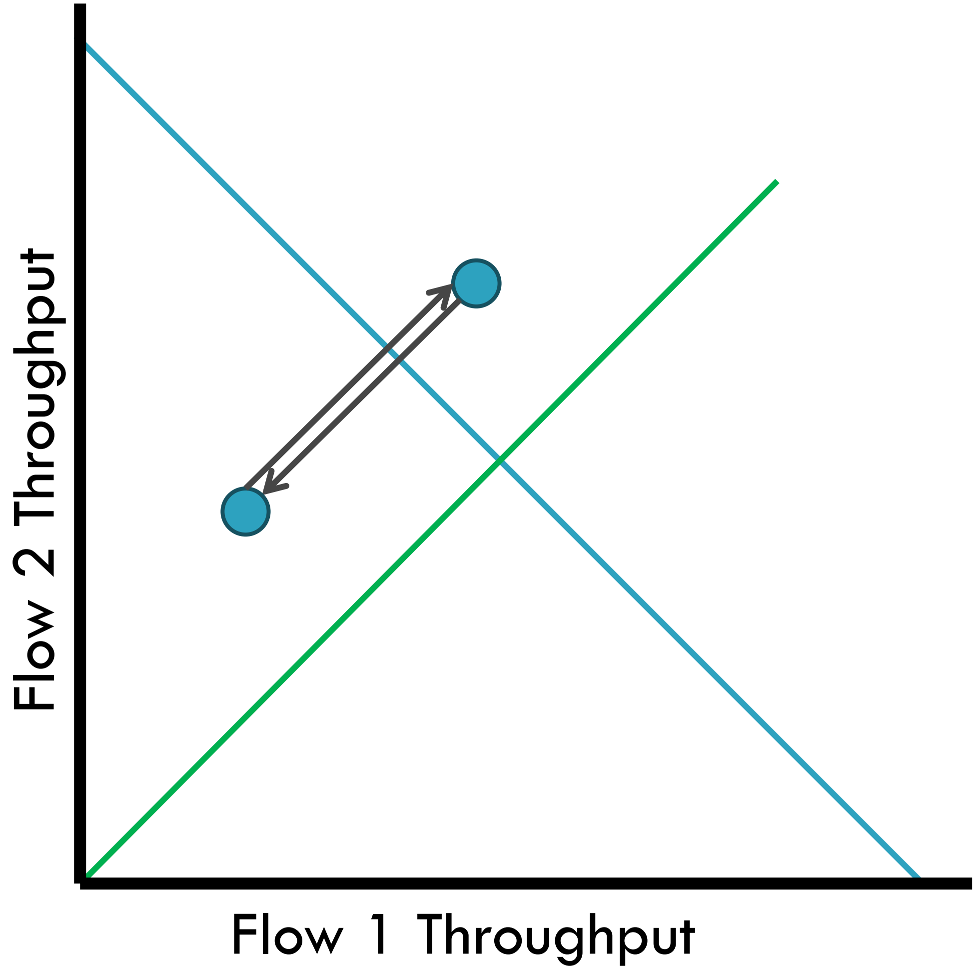
Multiplicative Increase, Additive Decrease

- Not stable!
- Veers away from fairness



Additive Increase, Additive Decrease

- Stable
- But does not converge to fairness



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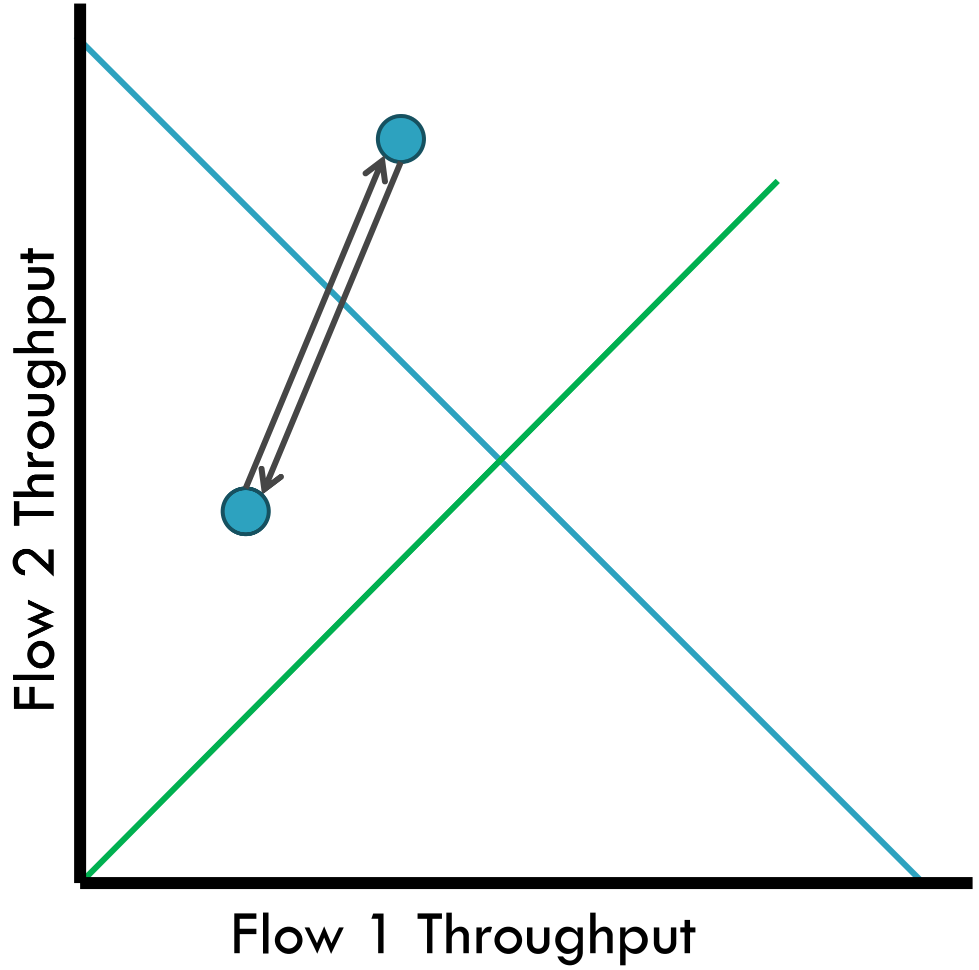
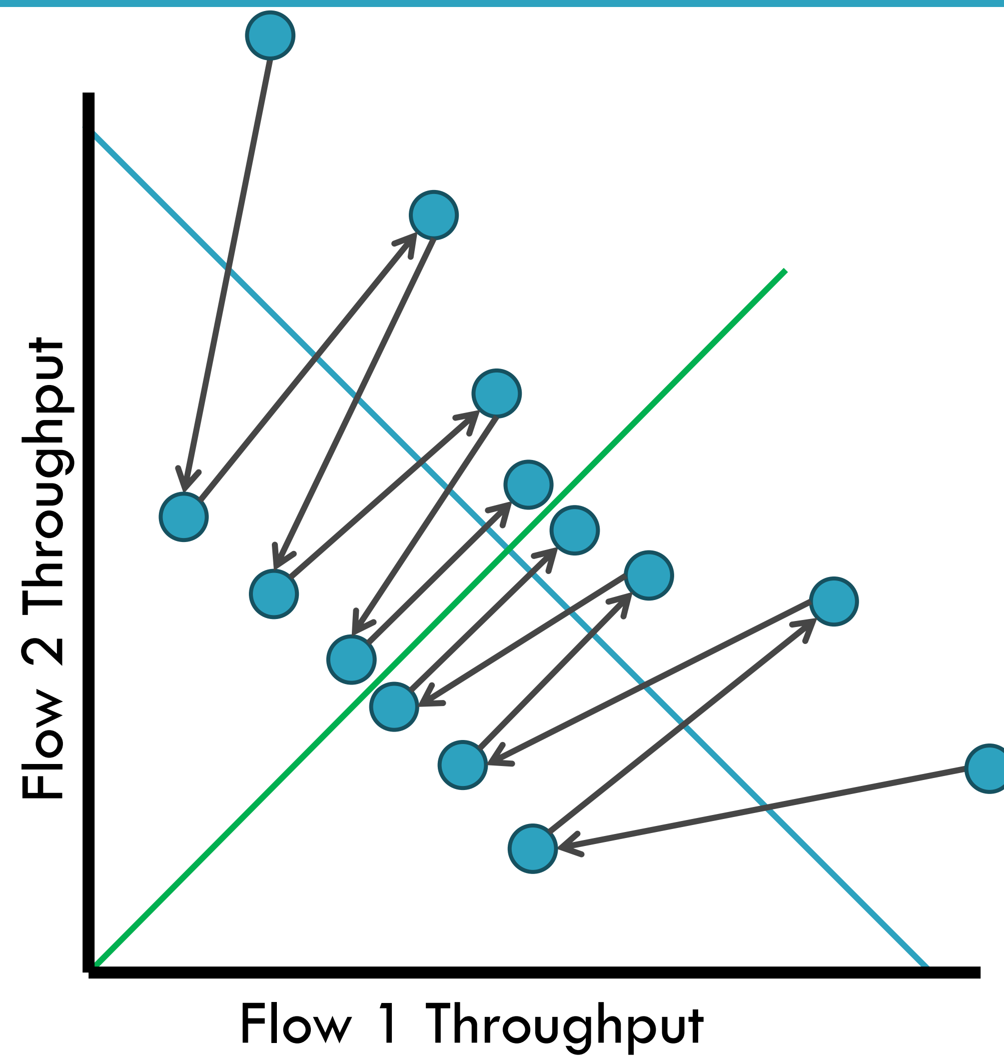


Additive Increase, Multiplicative Decrease

- Converges to stable and fair cycle
- Symmetric around $y=x$

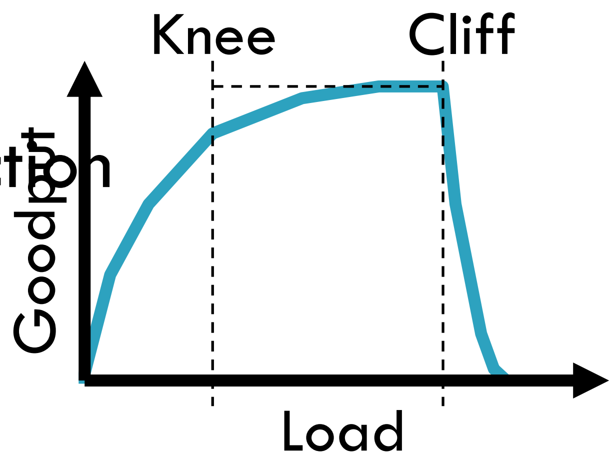


Implementing Congestion Control

- Maintains three variables:
  - *cwnd*: congestion window
  - *adv_wnd*: receiver advertised window
  - *ssthresh*: threshold size (used to update *cwnd*)
- For sending, use: *wnd = min(cwnd, adv_wnd)*
- Two phases of congestion control
  1. Slow start (*cwnd* < *ssthresh*)
     - Probe for bottleneck bandwidth
  2. Congestion avoidance (*cwnd* >= *ssthresh*)
     - AIMD

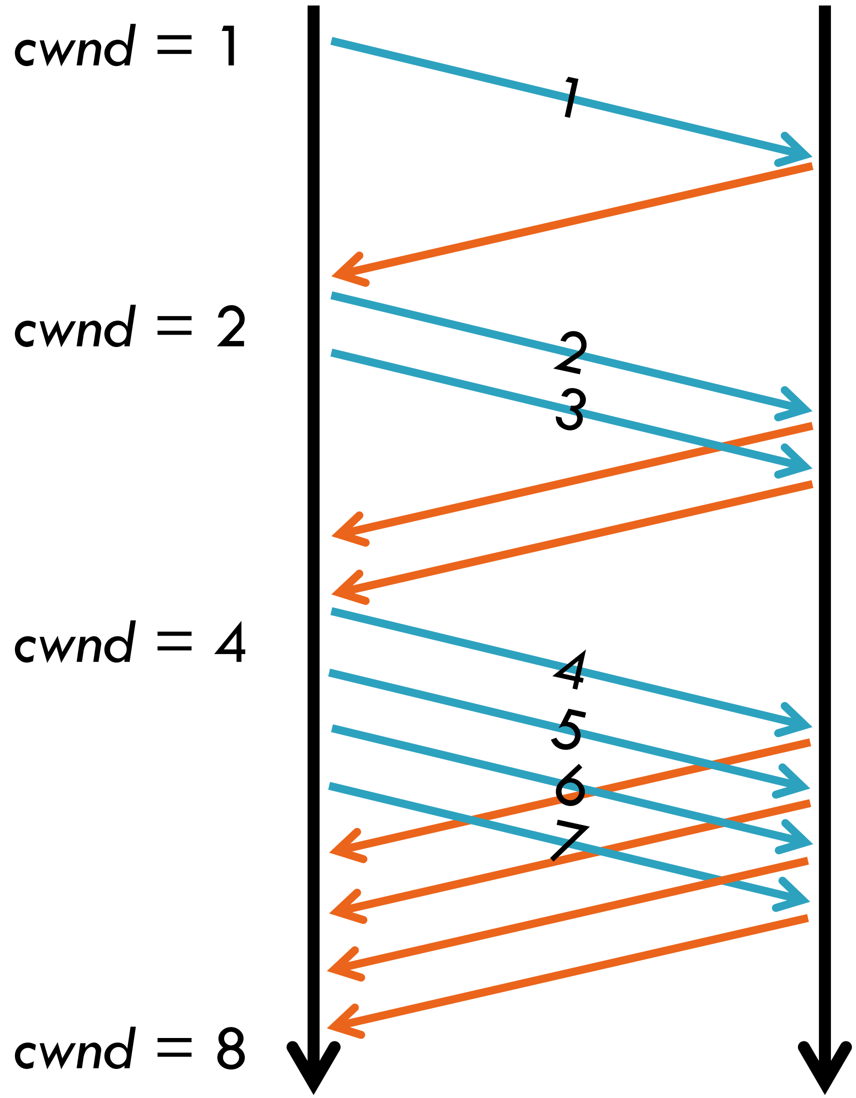
Slow Start

- Goal: reach knee quickly
- Upon starting (or restarting) a connection
  - *cwnd* =1
  - *ssthresh* = *adv_wnd*
  - Each time a segment is ACKed, *cwnd*++
- Continues until…
  - *ssthresh* is reached
  - Or a packet is lost
- Slow Start is not actually slow
  - *cwnd* increases exponentially



Slow Start Example

- *cwnd* grows rapidly
- Slows down when…
  - *cwnd* >= *ssthresh*
  - Or a packet drops



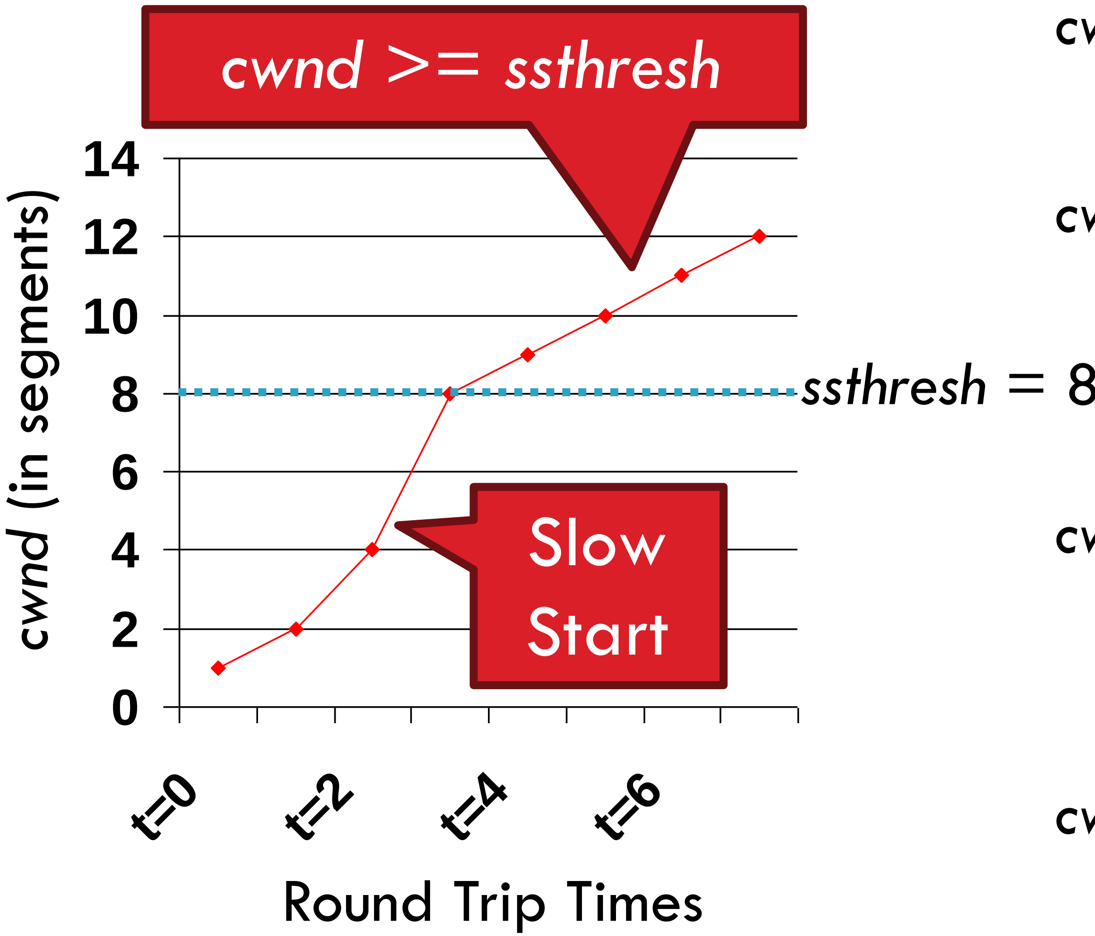
Congestion Avoidance

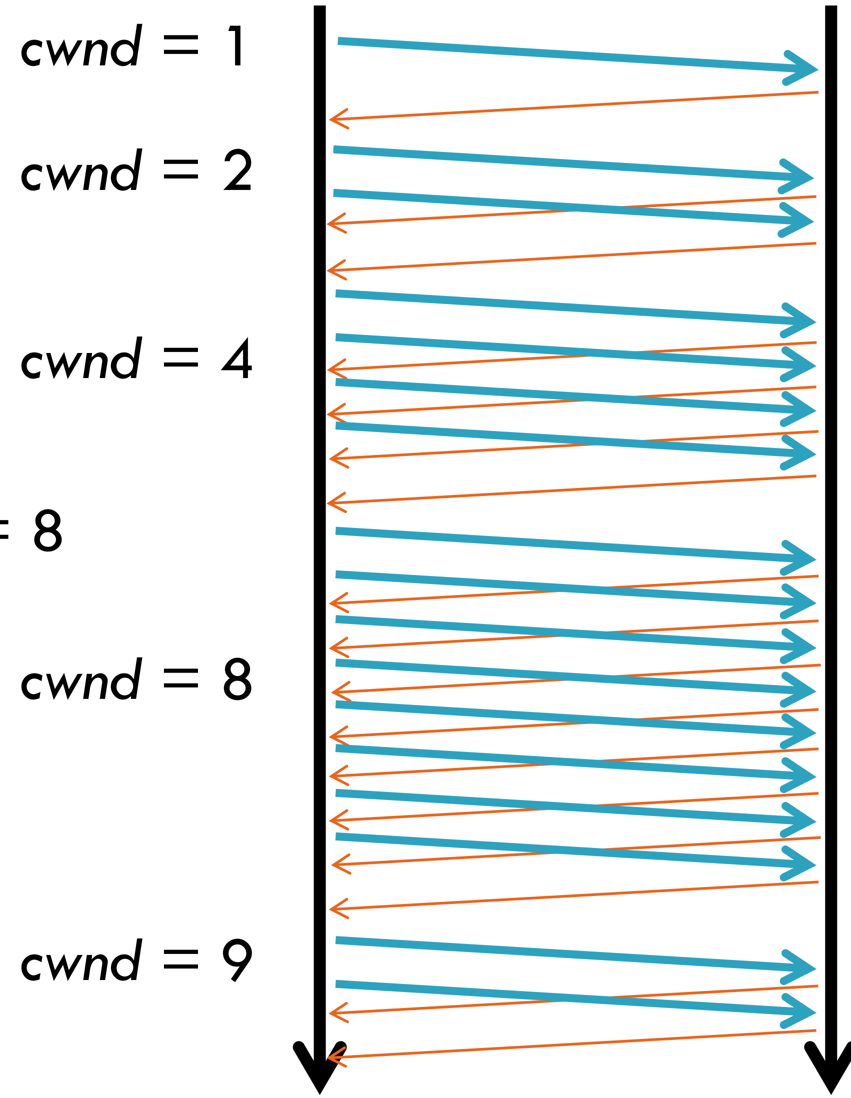
- AIMD mode
- *ssthresh* is lower-bound guess about location of the knee
- **If** *cwnd* >= *ssthresh* **then**

  each time a segment is ACKed

  increment *cwnd by 1/cwnd* (*cwnd* += *1/cwnd*).
- So *cwnd* is increased by one only if all segments have been acknowledged

Congestion Avoidance Example

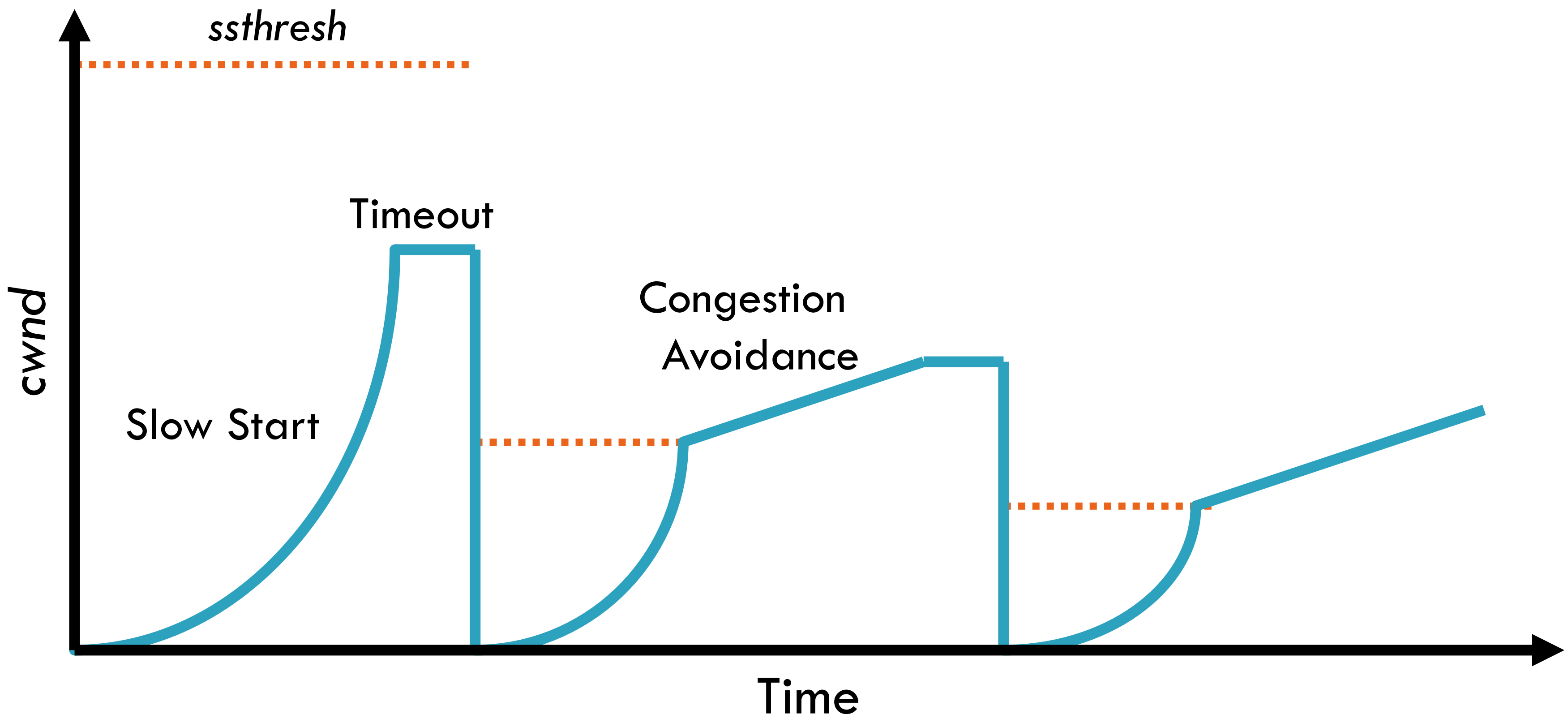




TCP Pseudocode


```
Initially:
        cwnd = 1;
        ssthresh = adv_wnd;
New ack received:
        if (cwnd < ssthresh)
            /* Slow Start*/
            cwnd = cwnd + 1;
        else
            /* Congestion Avoidance */
            cwnd = cwnd + 1/cwnd;
Timeout:
        /* Multiplicative decrease */
        ssthresh = cwnd/2;
        cwnd = 1;
```

The Big Picture



Outline

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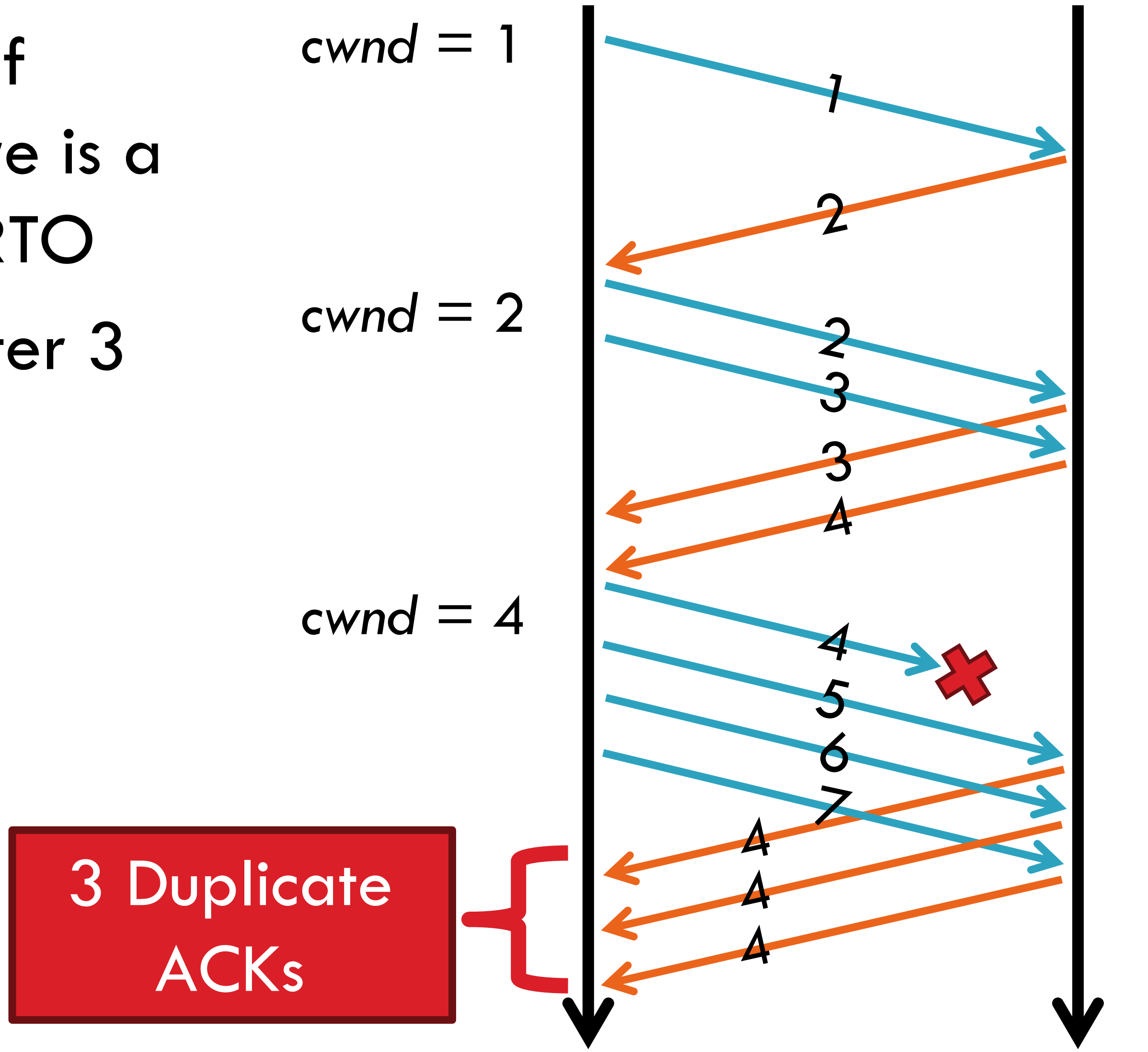
- Congestion Control
- Evolution of TCP
- Problems with TCP

The Evolution of TCP

- Thus far, we have discussed TCP Tahoe
  - Original version of TCP
- However, TCP was invented in 1974!
  - Today, there are many variants of TCP
- Early, popular variant: TCP Reno
  - Tahoe features, plus…
  - Fast retransmit
  - Fast recovery

TCP Reno: Fast Retransmit

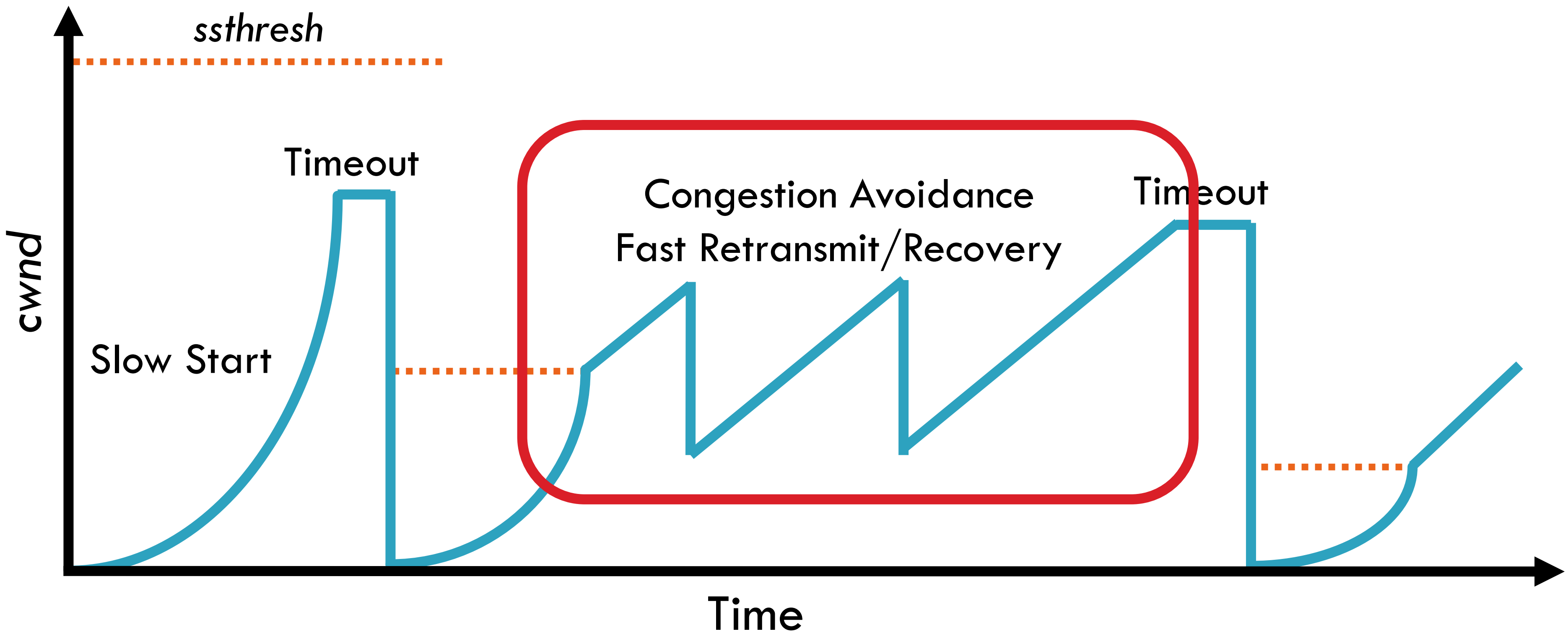
- Problem: in Tahoe, if segment is lost, there is a long wait until the RTO
- Reno: retransmit after 3 duplicate ACKs



TCP Reno: Fast Recovery

- After a fast-retransmit set *cwnd* to *ssthresh*/2
  - i.e. don't reset *cwnd* to 1
  - Avoid unnecessary return to slow start
  - Prevents expensive timeouts
- But when RTO expires still do *cwnd* = 1
  - Return to slow start, same as Tahoe
  - Indicates packets aren't being delivered at all
  - i.e. congestion must be really bad

Fast Retransmit and Fast Recovery



- At steady state, *cwnd* oscillates around the optimal window size
- TCP always forces packet drops

Many TCP Variants…

- Tahoe: the original
  - Slow start with AIMD
  - Dynamic RTO based on RTT estimate
- Reno: fast retransmit and fast recovery
- NewReno: improved fast retransmit
  - Reduce number of retransmissions
  - Window inflation
- Vegas: delay-based congestion avoidance
- And many, many, many more…

TCP in the Real World

- What are the most popular variants today?
  - Key problem: TCP performs poorly on high bandwidth-delay product networks (like the modern Internet)
  - Compound TCP (Windows)
    - Based on Reno
    - Uses two congestion windows: delay based and loss based
    - Thus, it uses a *compound* congestion controller
  - TCP CUBIC (Linux)
    - Enhancement of BIC (Binary Increase Congestion Control)
    - Window size controlled by cubic function
    - Parameterized by the time *T* since the last dropped packet
  - Other: BBR TCP, multi-path TCP, and QUIC (UDP-based), and various data center solutions, for example, …

High Bandwidth-Delay Product

- Key Problem: TCP performs poorly when
  - The capacity of the network (bandwidth) is large
  - The delay (RTT) of the network is large
  - Or, when bandwidth * delay is large
    - b * d = maximum amount of in-flight data in the network
    - a.k.a. the bandwidth-delay product
- Why does TCP perform poorly?
  - Slow start and additive increase are slow to converge
  - TCP is ACK clocked
    - i.e. TCP can only react as quickly as ACKs are received
    - Large RTT → ACKs are delayed → TCP is slow to react

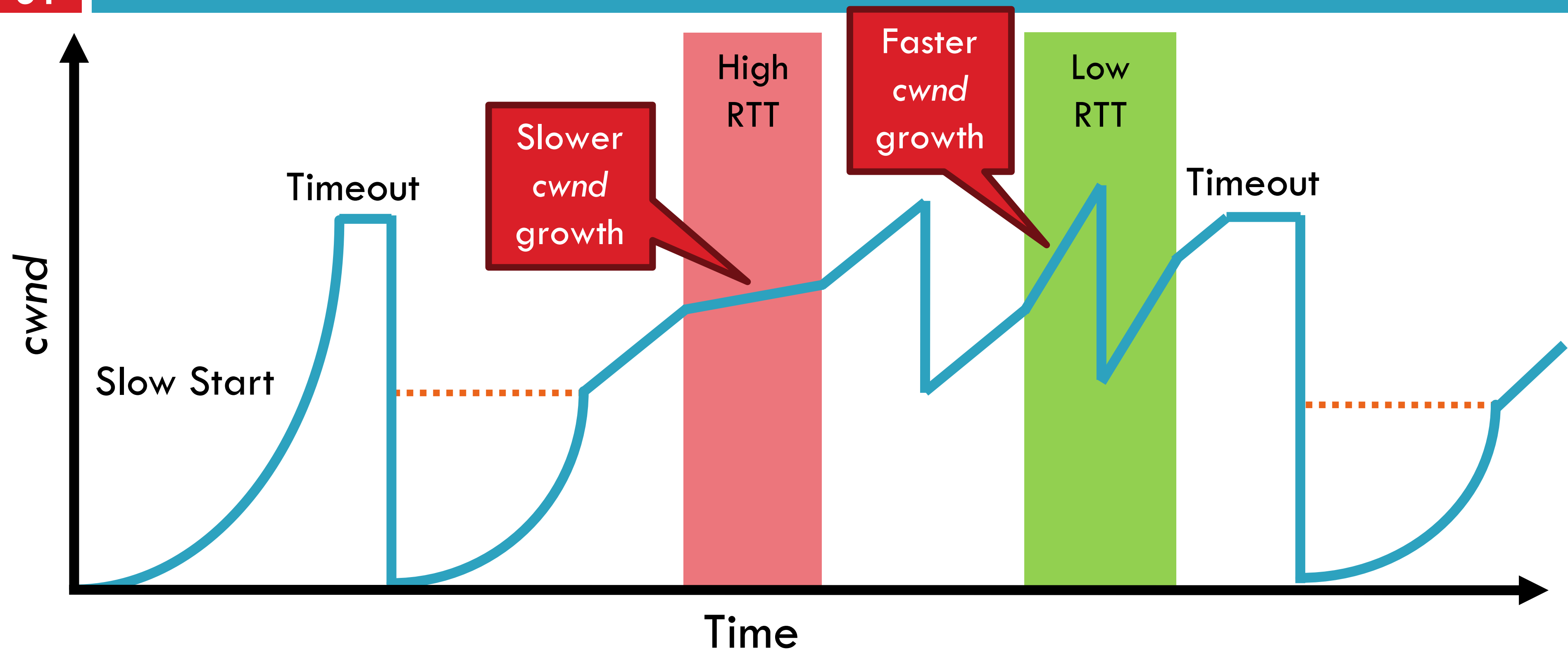
Goals

- Fast window growth
  - Slow start and additive increase are too slow when bandwidth is large
  - Want to converge more quickly
- Maintain fairness with other TCP variants
  - Window growth cannot be too aggressive
- Improve RTT fairness
  - TCP Tahoe/Reno flows are not fair when RTTs vary widely
- Simple implementation

Compound TCP Implementation

- Default TCP implementation in Windows
- Key idea: split *cwnd* into two separate windows
  - Traditional, loss-based window
  - New, delay-based window
- $wnd = \min(cwnd + dwnd, adv\_wnd)$
  - *cwnd* is controlled by AIMD
  - *dwnd* is the delay window
- Rules for adjusting *dwnd:*
  - If RTT is increasing, decrease *dwnd* (*dwnd* >= 0)
  - If RTT is decreasing, increase *dwnd*
  - Increase/decrease are proportional to the rate of change

Compound TCP Example



- Aggressiveness corresponds to changes in RTT
- Advantages: fast ramp up, more fair to flows with different RTTs
- Disadvantage: must estimate RTT, which is very challenging

TCP CUBIC Implementation

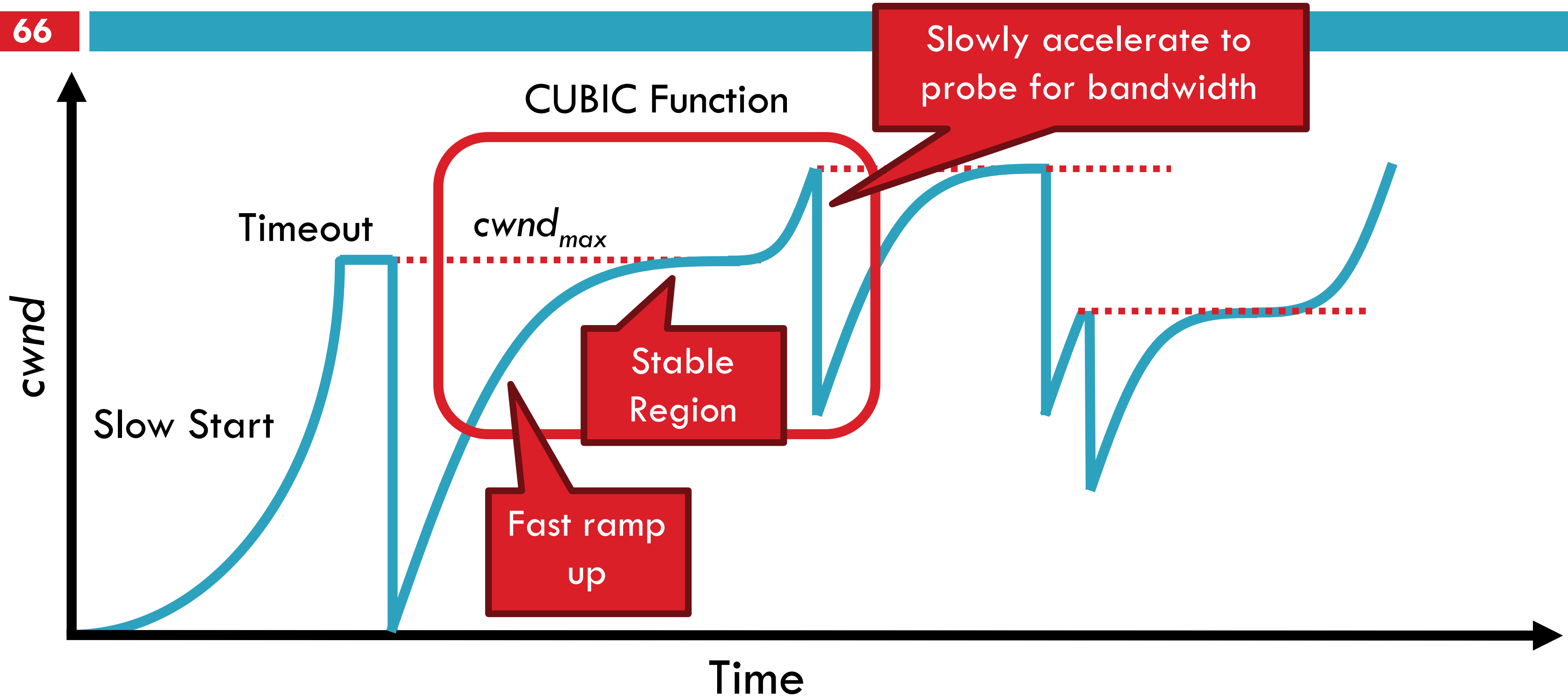
- Default TCP implementation in Linux
- Replace AIMD with cubic function

$$W_{cubic} = C(T-K)^3 + W_{max} \qquad (1)$$

C is a scaling constant, and $K = \sqrt[3]{\frac{W_{max}\beta}{C}}$

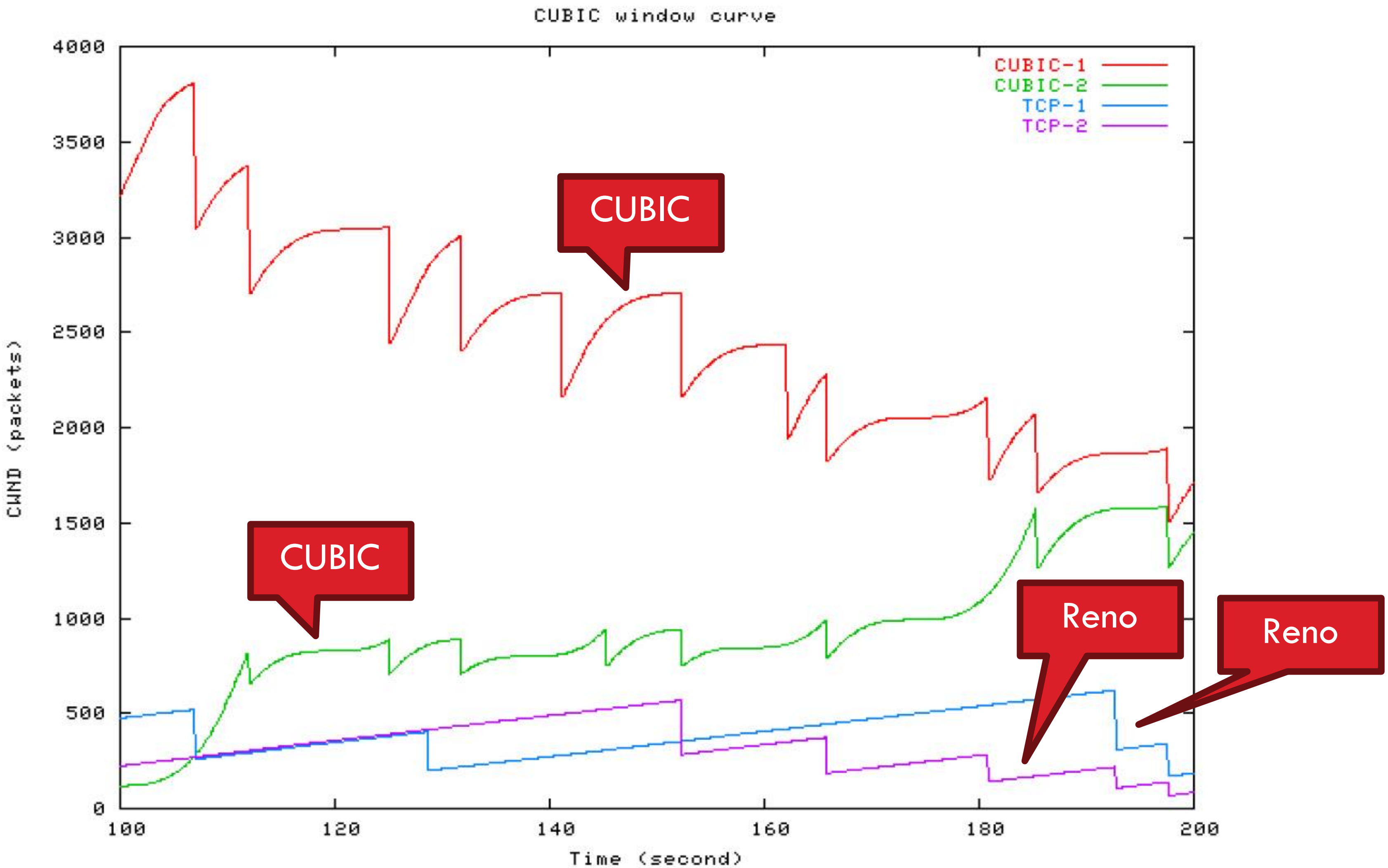
  - B → a constant fraction for multiplicative increase
  - T → time since last packet drop
  - W_max → cwnd when last packet dropped

TCP CUBIC Example



- Less wasted bandwidth due to fast ramp up
- Stable region and slow acceleration help maintain fairness
  - Fast ramp up is more aggressive than additive increase
  - To be fair to Tahoe/Reno, CUBIC needs to be less aggressive

Simulations of CUBIC Flows



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Deploying TCP Variants

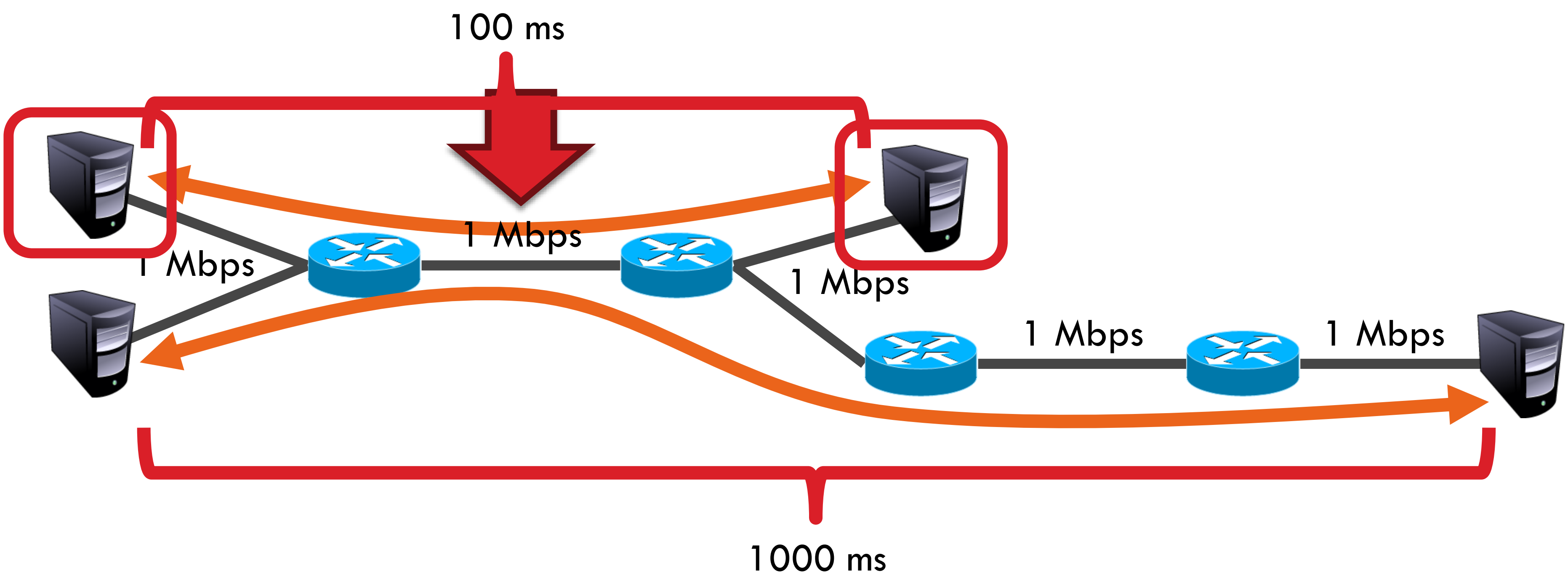
- TCP assumes all flows employ TCP-like congestion control
  - TCP-friendly or TCP-compatible
  - Violated by UDP :(
- If new congestion control algorithms are developed, they must be TCP-friendly
- Be wary of unforeseen interactions
  - Variants work well with others like themselves
  - Different variants competing for resources may trigger unfair, pathological behavior

Issues with TCP

- The vast majority of Internet traffic is TCP
- However, many issues with the protocol
  - Lack of fairness
  - Synchronization of flows
  - Poor performance with small flows
  - Really poor performance on wireless networks
  - Susceptibility to denial of service

Fairness

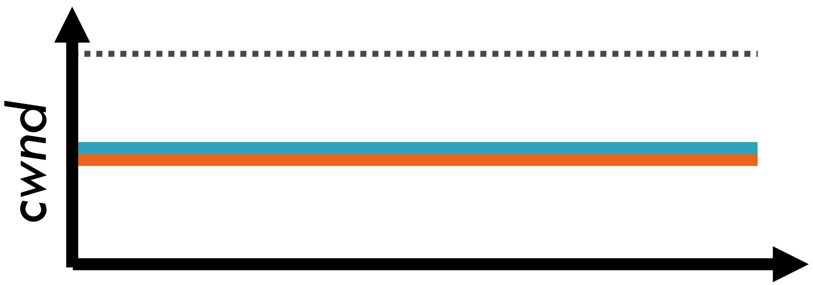
- Problem: TCP throughput depends on RTT



- ACK clocking makes TCP inherently unfair
- Possible solution: maintain a separate delay window
  - Implemented by Microsoft's Compound TCP

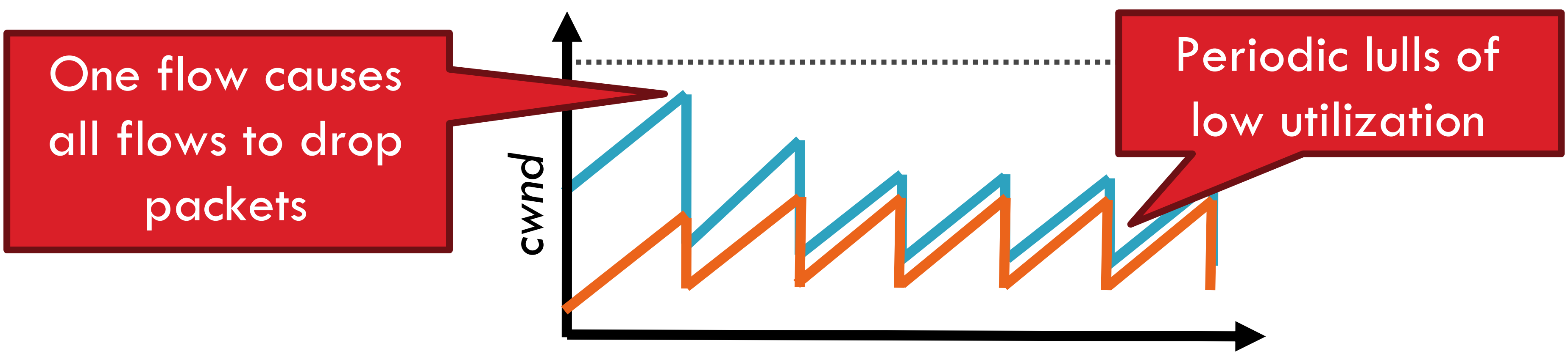
Synchronization of Flows

- Ideal bandwidth sharing



- Oscillating, but high overall utilization

- In reality, flows synchronize



Small Flows

- Problem: TCP is biased against short flows
  - 1 RTT wasted for connection setup (SYN, SYN/ACK)
  - *cwnd* always starts at 1
- Vast majority of Internet traffic is short flows
  - Mostly HTTP transfers, <100KB
  - Most TCP flows never leave slow start!
- Proposed solutions (driven by Google):
  - Increase initial *cwnd* to 10
  - TCP Fast Open: use cryptographic hashes to identify receivers, eliminate the need for three-way handshake

Wireless Networks

- Problem: Tahoe and Reno assume loss = congestion
  - True on the WAN, bit errors are very rare
  - False on wireless, interference is very common
- TCP throughput ~ 1/sqrt(drop rate)
  - Even a few interference drops can kill performance
- Possible solutions:
  - Break layering, push data link info up to TCP
  - Use delay-based congestion detection (TCP Vegas)
  - Explicit congestion notification (ECN)

Denial of Service

- Problem: TCP connections require state
  - Initial SYN allocates resources on the server
  - State must persist for several minutes (RTO)
- SYN flood: send enough SYNs to a server to allocate all memory/meltdown the kernel
- Solution: SYN cookies
  - Idea: don't store initial state on the server
  - Securely insert state into the SYN/ACK packet
  - Client will reflect the state back to the server

SYN Cookies

- Did the client really send me a SYN recently?
  - Timestamp: freshness check
  - Cryptographic hash: prevents spoofed packets
- Maximum segment size (MSS)
  - Usually stated by the client during initial SYN
  - Server should store this value…
  - Reflect the clients value back through them

SYN Cookies in Practice

- Advantages
  - Effective at mitigating SYN floods
  - Compatible with all TCP versions
  - Only need to modify the server
  - No need for client support
- Disadvantages
  - MSS limited to 3 bits, may be smaller than clients actual MSS
  - Server forgets all other TCP options included with the client's SYN
    - SACK support, window scaling, etc.

More slides …

What Should the Receiver ACK?

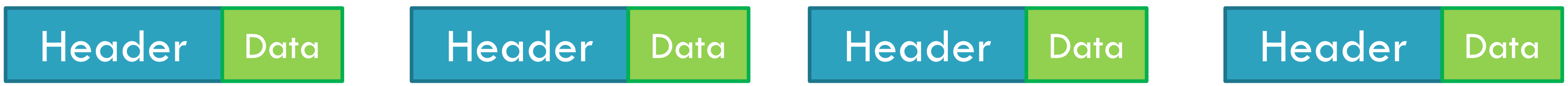
1. ACK every packet
2. Use *cumulative ACK*, where an ACK for sequence n implies ACKS for all $k < n$
3. Use *negative ACKs* (NACKs), indicating which packet did not arrive
4. Use *selective ACKs* (SACKs), indicating those that did arrive, even if not in order
   - SACK is an actual TCP extension

Sequence Numbers, Revisited

- 32 bits, unsigned
  - Why so big?
- For the sliding window you need…
  - |Sequence # Space| > 2 * |Sending Window Size|
  - $2^{32} > 2 * 2^{16}$
- Guard against stray packets
  - IP packets have a maximum segment lifetime (MSL) of 120 seconds
    - i.e. a packet can linger in the network for 2 minutes
  - Sequence number would wrap around at 286Mbps
    - What about GigE? PAWS algorithm + TCP options

Silly Window Syndrome

- Problem: what if the window size is very small?
  - Multiple, small packets, headers dominate data



- Equivalent problem: sender transmits packets one byte at a time

```
for (int x = 0; x < strlen(data); ++x)
  write(socket, data + x, 1);
```

Nagle's Algorithm

1. If the window >= MSS and available data >= MSS:
   Send the data
2. Elif there is unACKed data:
   Enqueue data in a buffer until an ACK is received
3. Else: send the data

Send a full packet

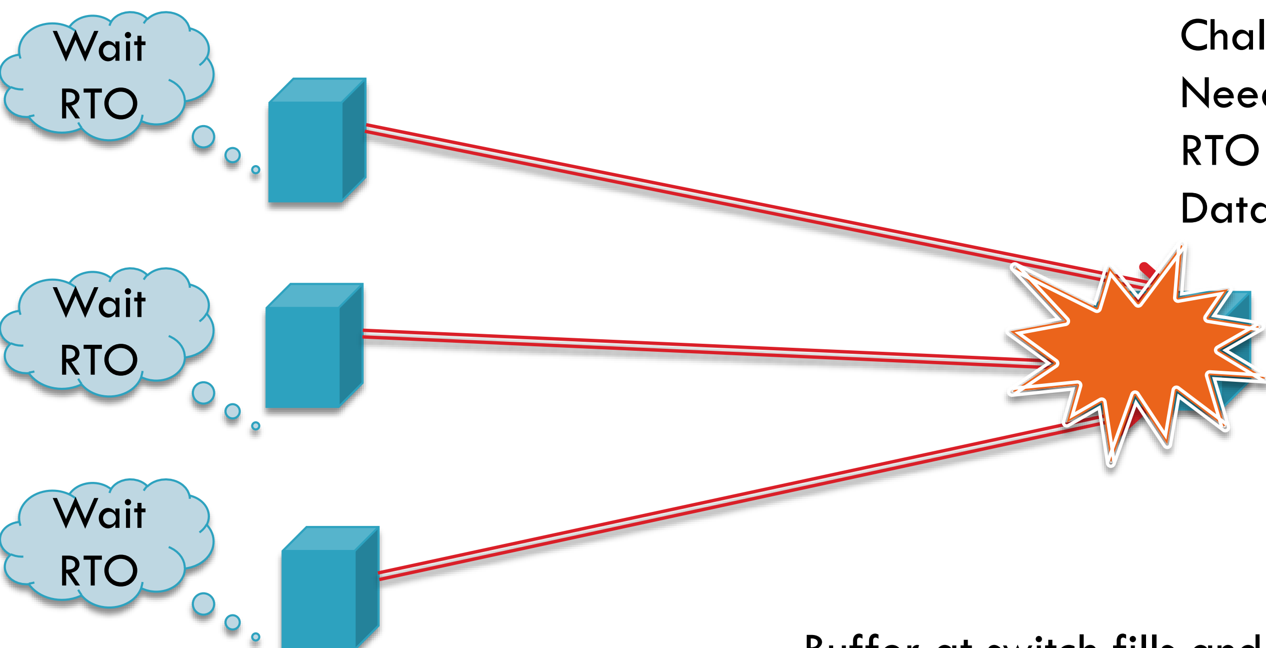
Send a non-full packet if nothing else is happening

- Problem: Nagle's Algorithm delays transmissions
  - What if you need to send a packet immediately?
  1. int flag = 1;
  2. setsockopt(sock, IPPROTO_TCP, TCP_NODELAY, (char *) &flag, sizeof(int));

Challenge of RTO in data centers

- TCP Incast problem – E.g. Hadoop, Map Reduce, HDFS, GFS

Many senders sending simultaneously to receiver



Challenges:
Need to break synchronization
RTO estimation designed for wide area
Data centers have much smaller RTT

Buffer at switch fills and packets are lost!
No ACKs will come back ☹

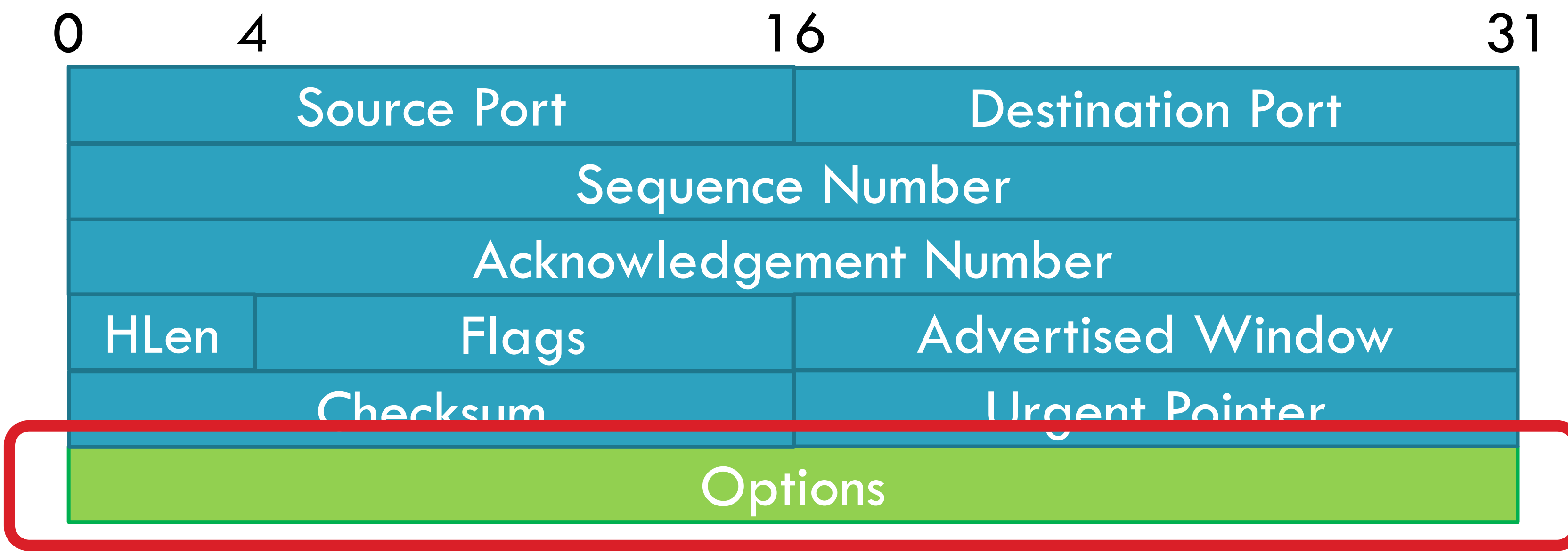
TCP Perspectives

- Cerf/Kahn
  - Provide flow control
  - Congestion handled by retransmission
- Jacobson / Karels
  - Need to avoid congestion
  - RTT estimates critical
  - Queuing theory can help
- Winstein/Balakrishnan
  - TCP is maximizing an objective function
    - Fairness/efficiency
    - Throughput/delay
  - Let a learning program pick the best fit for your environment

Outline

- UDP
- TCP
- Congestion Control
- Evolution of TCP
- Common TCP options
- Problems with TCP

Common TCP Options



- Window scaling
- SACK: selective acknowledgement
- Maximum segment size (MSS)
- Timestamp

Window Scaling

- Problem: the advertised window is only 16-bits
  - Effectively caps the window at 65536B, 64KB
  - Example: 1.5Mbps link, 513ms RTT

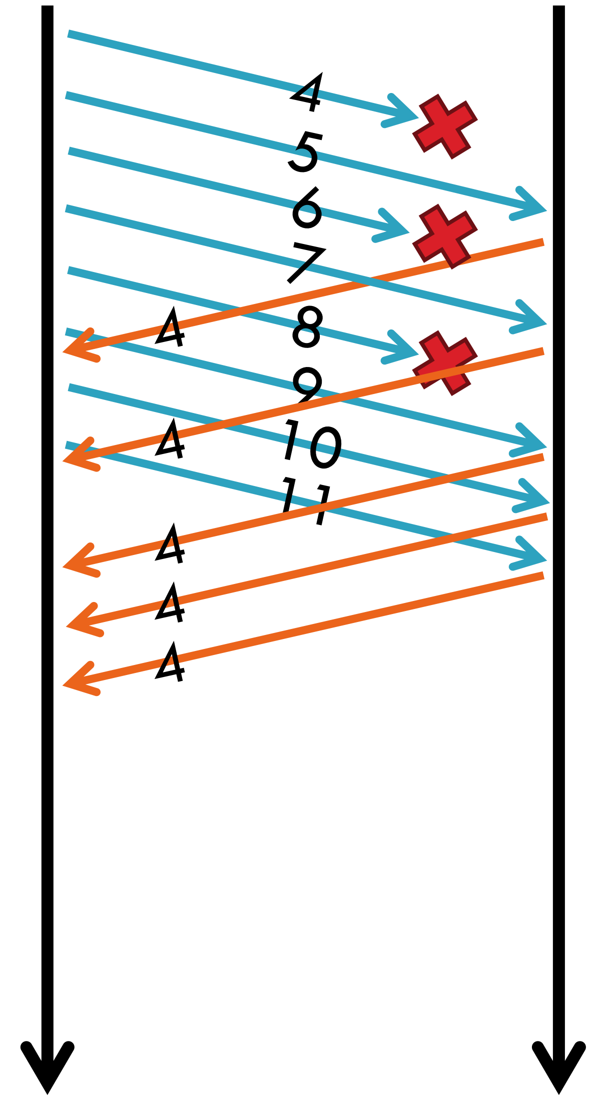
$$(1.5\text{Mbps} * 0.513\text{s}) = 94\text{KB}$$

$$64\text{KB} / 94\text{KB} = 68\% \text{ of maximum possible speed}$$

- Solution: introduce a window scaling value
  - *wnd = adv_wnd << wnd_scale;*
  - Maximum shift is 14 bits, 1GB maximum window

SACK: Selective Acknowledgment

- Problem: duplicate ACKs only tell us about 1 missing packet
  - Multiple rounds of dup ACKs needed to fill all holes
- Solution: selective ACK
  - Include received, out-of-order sequence numbers in TCP header
  - Explicitly tells the sender about holes in the sequence



Other Common Options

- Maximum segment size (MSS)
  - Essentially, what is the hosts MTU
  - Saves on path discovery overhead
- Timestamp
  - When was the packet sent (approximately)?
  - Used to prevent sequence number wraparound
  - PAWS algorithm

QUIC (Add slide(s) ...)