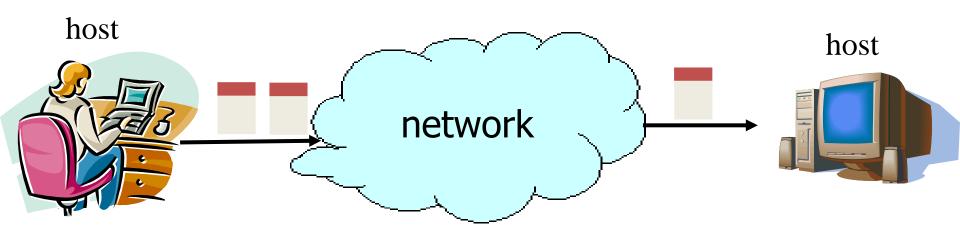
Best-effort Packet-delivery Service

1

This subset of slides is based on slides by Jennifer Rexford

Host-network Division of Labor

- Packet switching
 - Divide messages into a sequence of packets
 - Headers with source and destination address
- Best-effort delivery
 - Packets may be lost
 - Packets may be corrupted
 - Packets may be delivered out of order



Host-Network Interface: Why Packets?

- Data traffic is bursty
 - Logging in to remote machines
 - Exchanging e-mail messages
 - Request webpage
- Don't want to waste bandwidth

 No traffic exchanged during idle periods
- Better to allow multiplexing

 Different transfers share access to same links
- Packets can be delivered by most anything

Host-Network Interface: Why Best-Effort?

- Never having to say you're sorry...
 - Don't reserve bandwidth and memory
 - Don't do error detection & correction
 - Don't remember from one packet to next
- Easier to survive failures
 - Transient disruptions are okay during failover
- Can run on nearly any link technology

- Greater interoperability and evolution

Intermediate Transport Layer

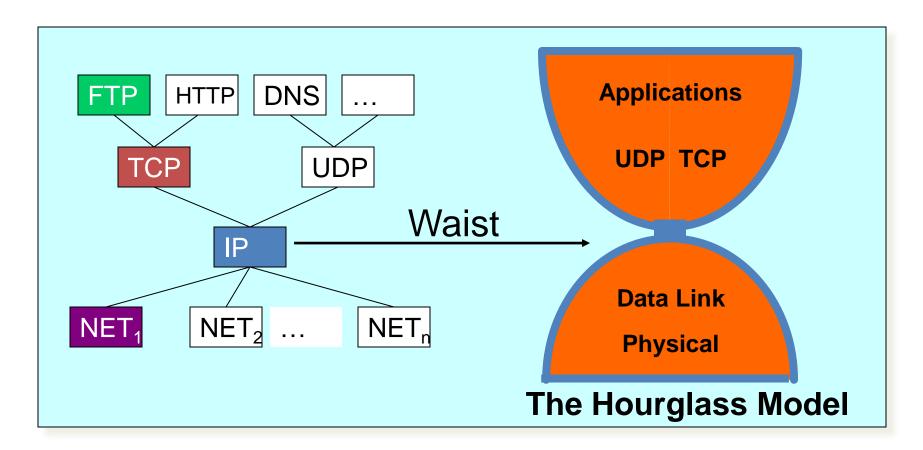
- But, *applications* want efficient, accurate transfer of data in order, in a timely fashion
 - Let the end hosts handle all of that
 - (An example of the "end-to-end argument")
- Transport layer can optionally...
 - Detect and retransmit lost packets
 - Put out-of-order packets back in order
 - Detect and handle corrupted packets
 - Avoid overloading the receiver
 - <insert your requirement here>

Modularity Through Layering

IP Protocol Stack

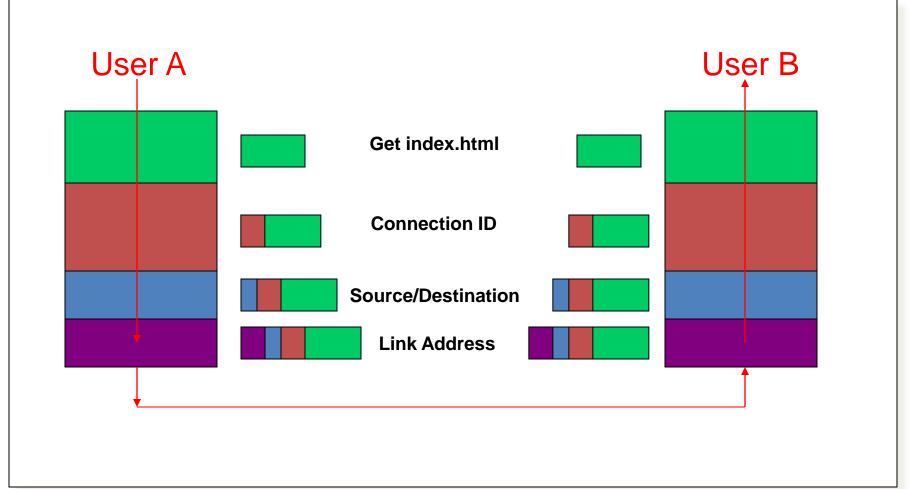
Application	Applications	
Transport	Reliable streams	Messages
Network	Best-effort global packet delivery	
Link	Best-effort local packet delivery	

The "Narrow Waist" of IP

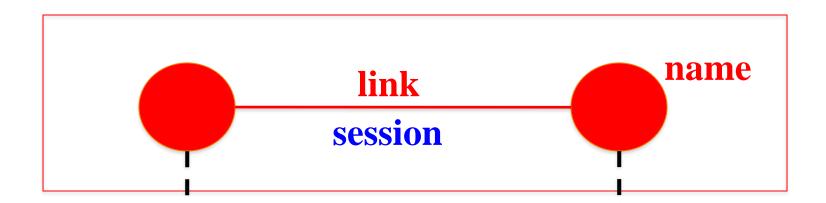


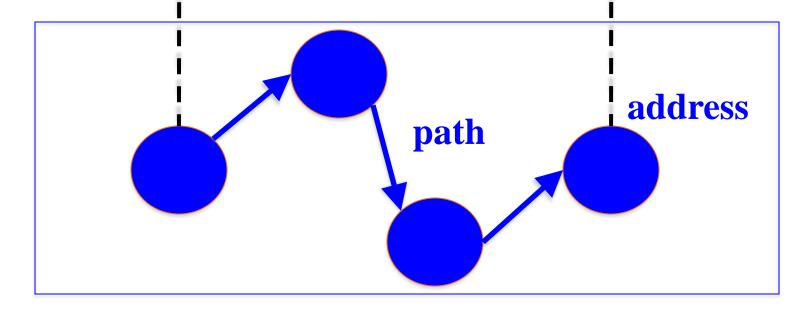
The waist facilitates interoperability





Relationship Between Layers





IP Suite: End Hosts vs. Routers



host

