Lectures

TDDA69 Data and Program Structure Imperative Programming Techniques *Cyrille Berger*



¹Introduction

2Concepts and models of programming languages **3**Declarative Computation Model 4Declarative Programming Techniques 5 Declarative Computation Implementation 6Declarative Concurrency 7 Message Passing Concurrency 8Explicit State and Imperative Model 9Imperative Programming Techniques 10Imperative Programming Implementation 11Shared-State Concurrency 12Relational Programming 13Constraint Programming 14Macro 15Running natively and IIT **16**Garbage Collection 17Summary

2/53

Lecture content

- Programming with states: imperatively
- Data abstraction
- Modular Programming
- Object Oriented Programming
 Object-Oriented Concepts
- Memory and threading

Programming with states: imperatively

Imperative vs declarative

Declarative

- ° No state
- Expresses logic of computation without control flow

Imperative

- ^o States
- Describes computation in term of statements that change the internal state

Imperative syntax construction

- A state is a sequence of values in time that contains the intermediate results of a desired computation.
- Would not it make sense to express computation also as a sequence?
- Syntax abstractions to manipulate state: loop, if

5/53

Loops



Side-effects

• Generally no-pure functions in imperative programming,

but: function factorial(n) { var r = 1; for(var i = 2; i <= n; ++i) { r *= i; } return r; } </pre>

 No side-effect, always return the same result for the same arguments



7/53

Random function

<pre>are_rand(seed) -> rec s = stream(); s.output(); A = 333667; B = 213453321; RAND_MAX = 10000000; on p(X) and(X = A*X+B) % M;);</pre>

9/53

Reasoning with state

- In the general case, we cannot ignore side-effects, how can we reason?
- Simple with simple programs, impossible with larger ones
- What if a state is visible through the entire program? cell a = 0; function () -> y { y = a; } a = 2;

}

10 / 53

Reasoning with state: invariant

- An invariant is a condition that is assumed to hold at a specific point in the program
- Invariants are relationships among the variables of a program that always hold. cell a = 0;

```
function () -> y { assert(a < 2); y
= a; }</pre>
```

```
a = 2;
```

```
• Avoid unwanted side effects!
```

Loop invariants

• A loop invariant is a relationship among the variables in a loop that holds at the begining and at the end of every iteration. for(var i = 0; i < 10; ++i) { assert(i >= 0 && i < 10); }



²roving properties in imperative programming

Also used for proving:

```
• function power(b, n) {
    int result = 1;
    for(int i = 0; i < n; ++i)
    {
        result *= b;
    }
    return result;
}
• Devise a loop invariant:
    ^ (n ≥ i) ∧ (result = b')
    Prove that it is true for the first loop iteration
    Prove that (n ≥ i) ∧ (result = b')
    Prove that (n ≥ j) ∧ (result = b')
</pre>
```

13 / 53

Data structure invariant

 Data structure invariants are the invariant relations among the fields of any object of a particular class

• List:

- ^o operations for creating, empty, append, prepend: no invariant
- ^o operations for head of list, pop: invariant is that the list is non empty

14 / 53

Encapsulation

- How to get protection from sideeffects
- In the rand function, only rand can update the state
- It is called encapsulation!

Data abstraction

The goals of data abstraction

- Encapsulation
 ^o Hide internals from the interface
- Compositionality
 - ^o Combine parts to make new parts
- Instantiation/invocation
 - ° Create new instances of parts

Program Organizing Techniques



18/53

17 / 53

Non-structured Programming

 The main program directly operates on global data

> Program **main** program / data

- Impractical when the program gets sufficiently large
- The same statement sequence must be copied if it is needed several times

Procedural Programming

- Combine a sequence of statements into a procedure with calls and returns
- The main program coordinates calls to procedures and hands over appropriate data as parameters



• Tasks are a collection of states, structures and procedures



Limitations of Structured Programming



- Attributes and behaviors are seperated
- Not extensible

21 / 53

Modular Programming

- Procedures of a common functionality are grouped together into separate modules.
- The main program coordinates calls to procedures in separate modules and hands over data as parameters.



Encapsulation

Hide internals from the interface
 Can be implemented by bundling of data with methods



- Private: only accessible to the object
 Public: accessible to the world
- How to implement abstraction in KL?
- +state+closure

Modular Programming



Linked-list

function new_list()



No encapsulation

Rundling data with method

function new list() return { cell head = null; cell tail = null; function push(_value) cell node = { cell value = _value; cell next = null }; if(list.tail) tail.next = node; tail = node; } else { head = node: tail = node; 3 } 3; 3

No encapsulation

}

25 / 53

26/53

Private data attemnt

function new_list()



No

encapsulation: cell list = new_list(); list.priv.head = 'breaking the data structure'

Fncansulation in C

• Using int_list_t*:

struct int_list_t; int_list_t* new_int_list(); void int_list_push(int_list_t*, int _value);

Internaly: typedef struct { int value;

. . .

int_list_node* next;
int_list_node;
typedef struct
{
 int_list_node* head;
 int_list_node* tail;

int_list_node* t } int_list;

```
int_list_t* new_int_list()
{
    int_list* 1 =
    malloc(sizeof(int_list));
    l->head = NULL;
    l->tail = NULL;
    return 1;
}
void int_list_push(int_list_t*
    l, int_value)
{
    ...
}
```






Private data attemnt

function new list()
{
cell priv = { cell head = null; cell tail =
null; };
cell self = {
function push(_value)
{
cell node = { cell value = _value, cell
<pre>next = null };</pre>
<pre>if(list.tail)</pre>
{
<pre>priv.tail.next = node;</pre>
<pre>priv.tail = node;</pre>
} else {
priv.head = node;
<pre>priv.tail = node;</pre>
}
}
};
clear(priv);
return self;
}

Does not work: cell list = new_list(); list.push(1); // triggers an error that priv is not defined

Private data, attempt #3

cell priv = { cell head = null; cell tail = null; function push(value) { cell node = { cell value; cell next = null }; if(list.tail) tail.next = node; tail = node; } else { head = node; tail = node; } ;{ cell self = { push = priv.push }; clear(priv); return self; 3

29 / 53

Private data, alternative

function new_list()

{
cell self = {};
(function() {
cell head = null;
cell tail = null;
define self.push = function(value)
{
cell node = { value, next = null };
if(list.tail)
{
<pre>tail.next = node;</pre>
tail = node;
} else {
head = node;
tail = node;
}
}
})();
return self;
}

Object Oriented Programming

30 / 53





function new_list()

Object-Oriented Concepts

What is an Object?

- Real-world objects are composed of attributes and behaviors
 - ^o Attributes: color, size,
 - ^o Behaviors: accelerate, brake,
- Program objects
 - ^o Attributes are states/
 - ^o Behaviors are

Object-Oriented Concepts

- Class
 - ° Instantiation/invocation
- Encapsulation
 - ° Compositionality
- Polymorphism
- Inheritance

34 / 53

Class

- An object is a specific *instance* (ie this car, this computer...)
- A class is a type of *object* (ie a car, a computer...) • It is a blueprint of object
- The class defines:
- ° The default set of variables
- ° The default set of functions
- ^o An initialisation function (called *constructor*
- The object is an instance of a class
- $^{\circ}$ It points to a specific memory location with the actual values for variables



Class Car



• object:

red_audi := { color:'red', max_speed: 200, max_acceleration: 10 }
 blue_volvo := { color:'blue', max_speed: 180, max_acceleration: 7 }

37 / 53

Inheritance from the class Car

- class Car: color, speed, max_speed and max_acceleration
- class Audi: class Car with max_speed := 200 and max_acceleration := 10
- class Volvo: class Car with max_speed
 := 180 and max_acceleration := 7

Inheritance

- Some classes have common part
- ^o Audi is a car, Volvo is a car, but Volvo and Audi are two different type of cars
- Inheritance allows to define a class in term of an other super-class
- Sub-classes
 - ^o can add new variables and functions
 - $^{\circ}$ override variables and functions from the super-class

38/53

Encapsulation

- Encapsulation with inheritance
- Three common levels of hiding:
 - ^o Private: only accessible to the object
 - Protected: only accessible to the object and through inheritance
 - ^o Public: accessible to the world



Car example

•function Car(self, protected, _color, max_speed, max_acceleration)
{
 cell protected.color = _color;
 cell speed = 0.0;
 function accelerate()
 {
 speed = Math.min(max_speed, speed + max_acceleration)
 }
 function deccelerate()
 {
 speed = Math.max(0, speed - max_acceleration)
 }
 };
 define self.accelerate = priv.accelerate;
 define self.deccelerate = priv.deccelerate;
 clear(priv);
 ;
 var obj = {);
 car(obj, {, 'red', 200, 10)
 }
}

Convenient object creation

```
• A new operator:
    function new(constructor, arguments...)
    {
        cell self = {}
        constructor(self, {}, arguments...);
        return self;
    }
    new(Car, 'red', 200, 10)
```


41 / 53

Inheritence

• For inheritance, all we need is to call the super class constructor

```
function Audi(self, protected, color)
{
    Car(self, protected, color, 200, 10)
})
new(Audi, 'blue');
```

Polymorphism

- *polymorphism* is the ability to appear in many forms
- It is the ability to call a function on an *data structure* that will execute different code for different data structure
- Define generic interface
 - ^o Example:
 - List push, pop
 - Stack push, pop
- It is one way to define generic and composable code
- With KL, it is achieved through a dynamic type system

43 / 53



Polymorphism in practise

```
• Just overwrite function in the child class
function Shape(self, protected)
{
    self.area = function() { throw 'Error'; }
}
function Square(self, protected, width,
height)
{
    define self.area =
        function() { return width * height };
}
```

45 / 53

Class syntax

 A nicer syntax for class definition would be: *class* Audi inherits Car
 {
 Audi(color) : Car(color, 200, 10)
 {
 };
 };
 Can be achieved with macros (future lecture)

46 / 53

Encapsulation in Python (1/2)

```
class Car:
    def __init__(self, max_speed):
        self.__max_speed = max_speed
    def __update(self):
        ...
c = Car(200)
c._Car__max_speed = 10
c._Car__update()
```

• Protection against 'accident', but still possible to break encapsulation

Encapsulation in Python (2/2)

```
class Car:
    def __init__(self, max_speed):
    _current_speed = 0
    def accelerate():
        nonlocal _current_speed
        _current_speed =
            min(_current_speed + 1,
            max_speed)
    def current_speed():
        return _current_speed
        self.accelerate = accelerate
        self.current_speed = current_speed
```



Closure in C++ pre 11

- If Object-Orientation is defined as +state+closure, does that means that all object-oriented programming language support closure?
- Closure is a reference to a function that has an embedded, persistent, hidden and unseparable context

```
• The following define a closure (with a state):
  class sum {
  private:
      int m sum;
```

```
public:
    sum() : m_sum(0) {}
   void operator () (int x) { m_sum += x; }
    int get sum() const { return m sum; }
```

}

49/53

Memory and threading

Memory Management

- Value store is an abstract concept, during execution the value of a binding can be replaced immediately
- With state, comes the need to track when a value is needed or not

Concurrency and state

- cell x: thread() { x = 1;}(); thread() { x = 3;}();
- No failure, but what is the value of x?
- In some programming languages, can lead to random crashes



Conclusion

- Conscequence of adding state
- Data abstraction
 - ° Encapsulation
 - ° Compositionality with polymorphism
 - ° Instantiation/invocation

