

Lectures

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TDDA69 Data and Program Structure

Concepts and models of programming languages

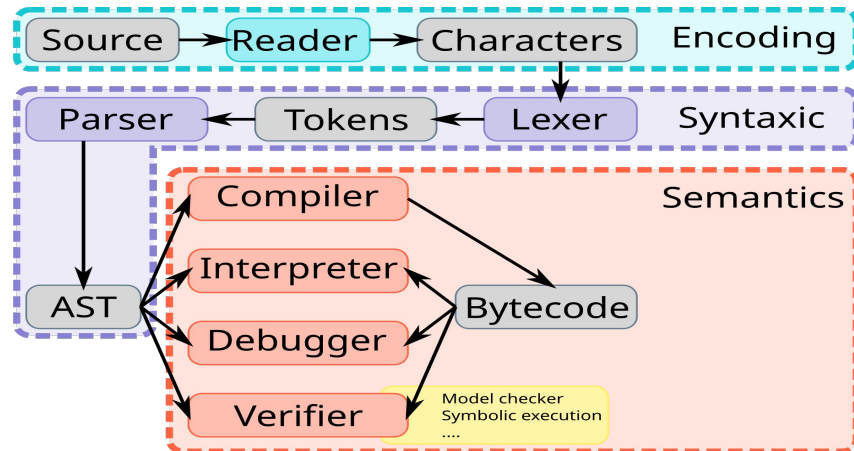
Cyrille Berger

Lecture content

- Defining a programming language
- Classification of programming languages
- The different programming paradigms
- Why different paradigms?
- Introduction to KL

Defining a programming language

Defining a programming language (1/2)



Defining a programming language (2/2)

- Syntax


```
SELECT * FROM table WHERE id=1
SELECT * FROM table WHERE
```
- Semantics
 - What operation does a SELECT do?

Formal Grammar

Defined by $G=(N, \Sigma, P, S)$:

- A finite set N of nonterminal symbols
- A finite set Σ of terminal symbols, disjoint from N
- A finite set P of production rules, each of the form

$$(\Sigma \cup N)^* N (\Sigma \cup N)^* \rightarrow (\Sigma \cup N)^*$$
- A start symbol $S \in N$
 - $\langle \text{digit} \rangle ::= '0' \mid '1' \mid '2' \mid '3' \mid '4' \mid '5' \mid '6' \mid '7' \mid '8' \mid '9'$
 - $\langle \text{integer} \rangle ::= ['-'] \langle \text{digit} \rangle \{ \langle \text{digit} \rangle \}$

Context-Free Grammar

Defined by $G=(N, \Sigma, P, S)$:

- A finite set N of nonterminal symbols
- A finite set Σ of terminal symbols, disjoint from N
- A finite set P of production rules, each of the form

$$N \rightarrow (\Sigma \cup N)^*$$
- A start symbol $S \in N$
 - $\langle \text{digit} \rangle ::= '0' \mid '1' \mid '2' \mid '3' \mid '4' \mid '5' \mid '6' \mid '7' \mid '8' \mid '9'$
 - $\langle \text{integer} \rangle ::= ['-'] \langle \text{digit} \rangle \{ \langle \text{digit} \rangle \}$

LR(k) Parser

- LR parser has a stack and input
- It uses two operations: *shift* and *reduce*
- It builds a *parse tree*

LR(1) and Chomsky Normal Form Grammar

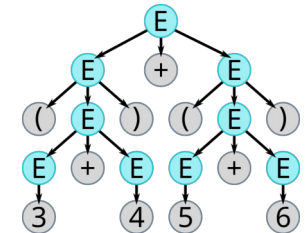
- A Chomsky Normal Form Grammar is a CFG such as, the production rules have the form

$$\begin{array}{ll} A \rightarrow BC & A, B, C \in N^3 \\ A \rightarrow a & A \in N, a \in \Sigma \end{array}$$

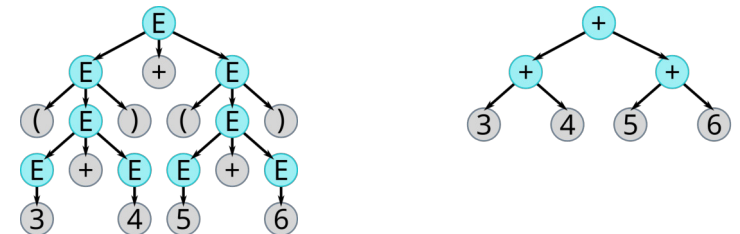
- Any CFG can be transformed into a CFN
- Most generated parser use a variant of LR(1) which is defined from a CNF

LR(k) Parser - Example

- Grammar:
 $E \rightarrow \text{int}$
 $E \rightarrow (E)$
 $E \rightarrow E + E$
- Stack:
E
- Input:
- Next action:



Abstract Syntax Tree



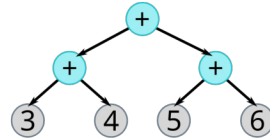
Semantic

- The grammar defines what is valid or not

$(3+4)+(5+6)$
 $3+)(+6)$

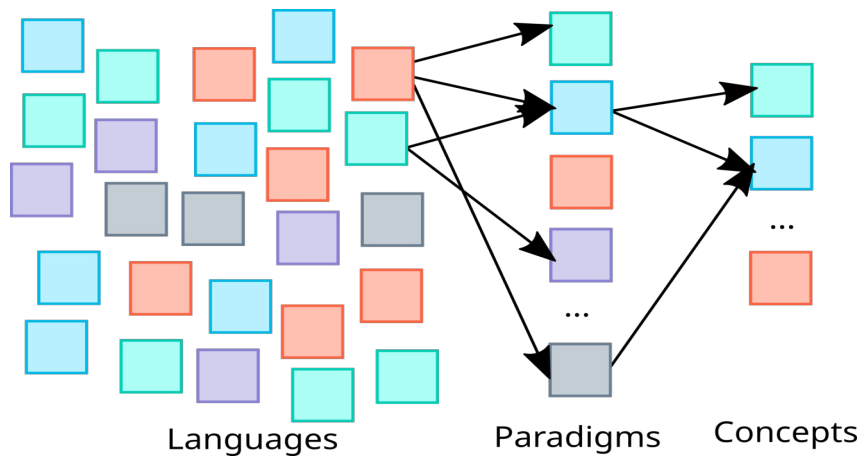
- The semantic defines what is the expected result

18

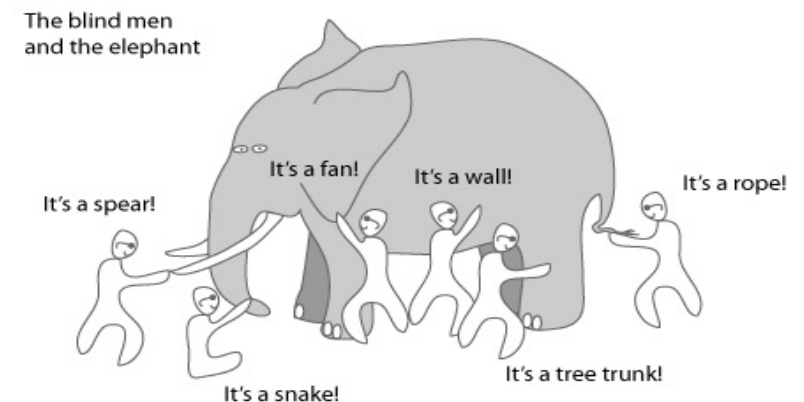


Classification of programming languages

Classification of programming languages



Limitation of a paradigm classification



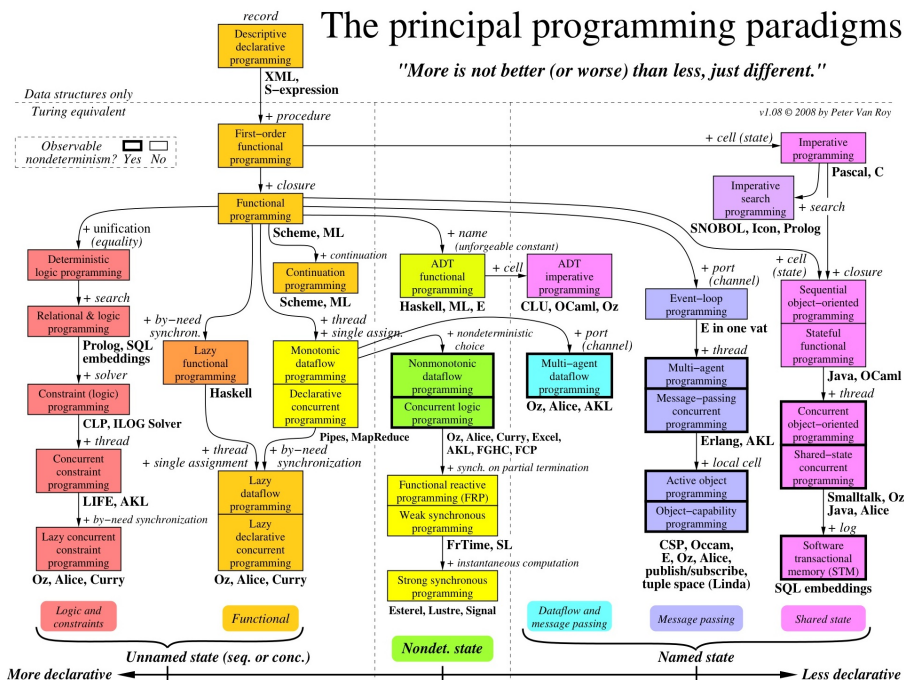
The different programming paradigms

Creative extension principle

- No exceptions


```
function f1 -> x1()
{
  set x2 = f2()
  if(!x2) { ... }
  else x1 = ...
}
function f2 -> x2()
{
  set x3 = f3()
  if(!x3) { x2 = null }
  else x2 = ...
}
function f3 -> x3()
{
  set x4 = f4()
  if(!x4) { x3 = null }
  else x3 = ...
}
function f4 -> x4()
{
  if(valid) { f4 = ... }
  else { f4 = null }
}
```
- With exceptions


```
function f1 -> x1()
{
  set x2 = f2()
  x1 = ...
}
function f2 -> x2()
{
  set x3 = f3()
  x2 = ...
}
function f3 -> x3()
{
  set x4 = f4()
  x3 = ...
}
function f4 -> x4()
{
  if(valid) { f4 = ... }
  else { throw Exception() }
}
```



Declarative

- Expresses logic of computation without control flow:
 - What should be computed and not how it should be computed.
- Examples: XML/HTML, antlr4/yacc/regular expressions, make/ants, SQL, ...

Declarative - Examples

- `Hello world!`
- `SELECT name FROM student WHERE course eq 'TDDA69'`
- grammar Hello;
 `r : 'hello' ID;`
 `ID : [a-z]+ ;`
 `WS : ['\t\r\n']+ -> skip ;`

Functional

- Computation are treated as mathematical function
 - without changing any internal state
- Examples: Lisp, Scheme, Haskell...

Functional - Examples

- `(print "Hello World")`
- `(take 25 (squares-of (integers)))`
 `-> (1 4 9 16 25 36 49 64 ... 576 625)`

Imperative

- Express how computation are executed
 - Describes computation in term of statements that change the internal state
- Examples: C/C++, Pascal, Java, Python, JavaScript...

Imperative - Examples

```
• for(var i = 1; i < 26; ++i)
{
    var sq = i*i;
    console.log(sq)
}
• #include <stdio.h>
int main()
{
    char ch;
    printf("Enter a character\n");
    scanf("%c", &ch);
    if (ch == 'a' || ch == 'A' || ch == 'e' || ch == 'E' || ch == 'i' || ch == 'I'
    || ch == 'o' || ch == 'O' || ch == 'u' || ch == 'U')
        printf("%c is a vowel.\n", ch);
    else
        printf("%c is not a vowel.\n", ch);
    return 0;
}
```

Object-Oriented - Programming

```
• #include <iostream>
class Character : public Symbol
{
public:
    Character(char _c) : m_c(_c) {}
    bool isVowel() const
    {
        return ch == 'a' || ch == 'A' || ch == 'e' || ch == 'E' || ch == 'i'
        || ch == 'I' || ch == 'o' || ch == 'O' || ch == 'u' || ch == 'U';
    }
private:
    char m_c;
};
int main()
{
    char c;
    std::cout << "Enter a character:\n";
    std::cin >> c;
    Character ch(c);
    if(ch.isVowel())
    {
        std::cout << c << " is a vowel.\n";
    }
    else {
        std::cout << c << " is not a vowel.\n";
    }
}
```

Object-Oriented

- Based on the concept of *objects*, which are *data structures* containing *fields* and *methods*
 - Programs are designed by making objects interact with each others
- Examples: C++, Java, C#, Python, Ruby, JavaScript...

Logic and symbolic

- Logic
 - Based on Formal logic: expressing facts and rules
- Symbolic
 - A program can manipulate its own formulas and components as if they are data
- Example: prolog

Logic programming

- likes(mary,food).
likes(mary,wine).
likes(john,wine).
likes(john,mary).
- | ?- likes(mary,food).
yes.
| ?- likes(john,wine).
yes.
| ?- likes(john,food).
no.

Symbolic programming

```
• d( X, X, 1 ):- !.           /* d(X) w.r.t. X is 1 */
d( C, X, 0 ):- atomic(C).    /* If C is a constant */
                               /* then d(C)/dX is 0 */
d( U+V, X, R ):-             /* d(U+V)/dX = A+B where */
d( U, X, A ),                 /* A = d(U)/dX and */
d( V, X, B ),
R = A + B.

...

d( sin(W), X, Z*cos(W) ):-    /* d(sin(W))/dX = Z*cos(W) */
d( W, X, Z).                  /* where Z = d(W)/dX */
d( exp(W), X, Z*exp(W) ):-    /* d(exp(W))/dX = Z*exp(W) */
d( W, X, Z).                  /* where Z = d(W)/dX */

...
• ?- d(cos(2*X+1), X, what)
what = 2*sin(2*X+1)
```

Constraint Programming

- Constraint
 - A relation between two variables are stated in the form of a constraint (can be logic or numerical)
- Example: Oz (functional), Kaleidoscope (imperative), Prolog (logic)

Constraint Programming - Examples

```
local
proc {MyScript Solution}
  X = {FD.int 1#10}
  Y = {FD.int 1#10}
  Z = {FD.int 1#10}
in
  Solution = unit(x:Y:Z)
  X + Y =: Z
  X <: Y
  %% search strategy
  {FD.distribute naive Solution}
end
in
  {Browse {SearchAll MyScript}}
end
```


Why different paradigms?

Is there a paradigm to rule them all?

- In theory you can program everything in C/C++ and imperative programming, or functional programming...
- But is that convenient?
- And is that safe?

Functional vs Imperative

- Double all the numbers in an array
`var numbers = [1,2,3,4,5]`
- Imperative:

```
var doubled = []
for(var i = 0; i < numbers.length; i++) {
  var newNumber = numbers[i] * 2
  doubled.push(newNumber)
}
```
- Functional:

```
var doubled = numbers.map(function(n) {
  return n * 2
})
```

Declarative vs Imperative

- Select all the dogs that belongs to a specific owner
- Declarative:

```
SELECT * from dogs INNER JOIN owners
WHERE dogs.owner_id = owners.id
```
- Imperative:

```
var dogsWithOwners = []
var dog, owner

for(var dog in dogs) {
  for(var owner in owners) {
    if (owner && dog.owner_id == owner.id) {
      dogsWithOwners.push({ dog: dog, owner: owner })
    }
  }
}
```

Introduction to KL

KI

- Practical language

```
function sqrt(x) {  
  return(x*x); }  
define b =  
  sqrt(sqrt(a));
```
- Kernel language

```
define sqrt = function  
(x) -> y { y = mul(x,  
x); }  
define __t1__ =  
  sqrt(a);  
define b =  
  sqrt(__t1__);
```
- Practical language
 - Provides usefull abstractions for the programmer
 - Can be extended with linguistic abstractions
- Kernel language
 - Easy to understand and reason
 - Has a precise (formal) semantic

Kernel Language

- Kernel language: is the minimal language that you need for a given paradigm
- Define a mapping between a full programming language into the kernel language

KL

- KL is a Kernel Language
- You will be developing its interpreter during the labs
- Written in RPython
- It is defined as:
 - syntax: inspired C/JavaScript
 - semantic: +procedure+closure+cell
 - In the future, based on the idea of adding concepts
- Mapping through a macro system (lecture 14)

Syntax

```
define constant_name = 10;
cell variable_name = 10;
variable_name = divide(constant_name, 2);
define f = function y -> (a, b)
{
  define _t1_ = equals(b, 0);
  define _t2_ = function () -> r { r = divide(a, b); };
  define _t3_ = function () { raise("Division by 0"); };
  y = cond(_t1_, _t2_, _t3_);
}
define _t4_ = function() { f(1, 0, 0); }
define _t5_ = function(except) { print(except); };
define _t6_ = function() { print("No problem!"); };
catch(_t4_, _t5_);
catch(_t4_, _t5_, _t6_);
```

Modules

```
import std 1.0;
std.print("Hello World");

import test 1.0;
test.case("Functions Call")
  .check(true);
```

Conclusion

- Definition of a programming language
- Defining the semantic using programming paradigms classification
- KL