

# ASSIGNMENT SET 1

Multi Agent Systems (TDDE13), Linköping University  
By Fredrik Prántare, Autumn Semester 2023

**Directions:** *Individually* (not in groups or pairs) solve the assignments below and send your solutions (with a clear and precise line of reasoning!) to your TA's mail address (Daniel de Leng, [daniel.de.leng@liu.se](mailto:daniel.de.leng@liu.se)) before the deadline. It is important that you

1. use the course's *LaTeX*/Word template for the answers (only submit the compiled .pdf; bad formats and file types will be rejected);
2. use "TDDE13: Assignment Set 1" as the header in your mail; and
3. send the answers from your *LiU student account*.

**Deadline:** See the course's webpage. After the deadline, you receive only half the points for correct answers.

**Prerequisites:** Course lectures + the following chapters in the course's textbook (*Multia-  
gent Systems: Algorithmic, Game-Theoretic, and Logical Foundations*):

- Chapter 3: Introduction to Noncooperative Game Theory: Games in Normal Form;
- Chapter 4: Computing Solution Concepts of Normal-Form Games;
- Chapter 5: Games with Sequential Actions;
- Chapter 12: Teams of Selfish Agents: An Introduction to Coalitional Game Theory.

## Exercises:

- 1** (0.5p) What is a von Neumann and Morgenstern *utility function* and what does it represent? Define such a utility function for a situation/context relevant to your own daily life, and describe what it represents.
- 2** (0.5p) If any, what are the key differences between *common-payoff games*, *zero-sum games* and *constant-sum games*? Describe and define an example you know of for two of them. Your examples must be distinct (i.e., they cannot be the same example) and *not* taken from the course's book!
- 3** (1.5p) Find all existing *pure* **and** *mixed Nash equilibria* (NE) and the players' expected payouts (in the NEs) in the following normal-form games:

|   | L     | R    |
|---|-------|------|
| U | 1, 1  | 1, 0 |
| D | -1, 1 | 0, 1 |

|   | X    | Y    |
|---|------|------|
| A | 1, 0 | 0, 1 |
| B | 0, 1 | 1, 0 |

|    | C1   | C2   |
|----|------|------|
| R1 | 1, 3 | 2, 1 |
| R2 | 2, 1 | 1, 4 |

- 4 (0.5p) Find the row player's *pure maxmin solution* (both value and strategy) in the following (simultaneous move) normal-form game:

|   |       |       |
|---|-------|-------|
|   | A     | B     |
| X | 4, -3 | -2, 1 |
| Y | 0, 0  | 5, -6 |

- 5 (0.5p) In game theory, what are the key differences between *cooperative* and *non-cooperative* settings? Give an example you know of for each. Your examples must be distinct and *not* taken from the course's textbook.
- 6 Design and define your own *non-convex 3-player (characteristic function) coalitional game* (with transferable utility), for which each coalition (except the empty coalition) must have a distinct (= unique) non-zero value. The coalitional game that you define needs to represent a (potentially fictional) scenario that you can think of.
- (a) (0.5p) Describe what the coalitional game that you defined represents.
- (b) (0.5p) Compute one player's *marginal contributions* and her *Shapley value*. Also, clearly describe (with your scenario as basis) what her Shapley value represents.
- (c) (0.5p) Prove that your game is **not** *convex*.
- (d) (0.5p) Show that the game that you defined is (or that it is not, depending on the game that you defined) *additive* and/or *superadditive*.

- 7 (1.0p) Recall that a coalitional *partition function game* is defined as a tuple  $\langle N, v \rangle$  where:

- $N = \{1, \dots, n\}$  is a set of agents;
- $v((C, CS)) \mapsto \mathbb{R}$  is a function that maps a value (e.g., potential utility) to every *embedded coalition*  $(C, CS)$  over  $N$ ; see Definition 1.

**Definition 1.** An *embedded coalition* (over the agents  $N$ ) is a pair  $(C, CS)$ , where  $CS$  is a coalition structure (see Definition 2) over  $N$ , and  $C$  is a coalition with  $C \in CS$ .

**Definition 2.** A *coalition structure*  $CS = \{C_1, \dots, C_{|CS|}\}$  over the set of agents  $N$  is a set of coalitions with  $C_i \subseteq N$  for  $i = 1, \dots, |CS|$ ,  $C_i \neq \emptyset$  for  $i = 1, \dots, |CS|$ ,  $C_i \cap C_j = \emptyset$  for all  $i \neq j$ , and  $\bigcup_{i=1}^{|CS|} C_i = N$ .

Now, suppose we have a partition function game  $\langle \{1, 2, 3\}, v \rangle$  with:

$$v((C, CS)) = \begin{cases} 2 & \text{if } 1 \in C \text{ and } 2 \in C; \\ 1 & \text{else if } |CS| = 2 \text{ and } |C| = 1; \\ 3 & \text{else if } |CS| = 2, C \neq \emptyset \text{ and } C \neq \{2, 3\}; \\ 1.5 & \text{else if } C = \{3\}; \\ 0 & \text{otherwise.} \end{cases}$$

Your task is to compute  $\max_{CS \in \Pi} \left\{ \sum_{C \in CS} v((C, CS)) \right\}$  and  $\arg \max_{CS \in \Pi} \left\{ \sum_{C \in CS} v((C, CS)) \right\}$ , where  $\Pi$  denotes the set of all coalition structures over  $N$ .

**8** (0.5p) Discuss and critique using the Shapley value for dividing payoff in **partition** function games.