

# Algorithmic Problem Solving Practice Problem Solving Session/Seminar

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# Solving a Problem



## 1. Analyze the problem

- What data do you receive? What data should you produce?
- How large is the input? What is a feasible time complexity?

## 2. Find a solution

- Can you solve it by hand/in your head? What steps do you perform?
- Could you alter the problem somehow to make it easier?
- When you have a solution in mind, are there any corner cases?
- How can you store the data efficiently? Is your time complexity feasible?

## 3. Implementation

- Start with pseudo code (code comments) if your solution is complex.
- Use good variable names! Invest a few seconds writing descriptive names, and in return you'll save a lot of debugging time.
- Are there large numbers? Are 32-bit integers sufficient?
- Is there a lot of input/output? See Kattis' help section for advice.
- Learn tools for debugging (valgrind, gdb, jdb, pdb, ...)

# Problem A - Babelfish



- We receive a dictionary of  $n$  ( $0 \leq n \leq 100,000$ ) word pairs  $\langle a_i, b_i \rangle$  followed by  $m$  ( $0 \leq m \leq 100,000$ ) words  $c_i$ .
- **Problem:** For each  $c_i$ , if  $c_i = b_i$  for some  $i$ , print  $a_i$ . Otherwise, print “eh”.
- **Solution:** Straight-forward lookup of  $c_i$ , but make sure to use an appropriate data structure.
  - Linear search takes  $O(n)$  per lookup => too slow.
  - Use a map (TreeMap/std::map in Java/C++) with lookup in  $O(\log(n))$ .
  - Use a hashmap (HashMap/std::unordered\_map in Java/C++ or Python dictionary) with lookup in  $O(1)$ .
- **Time complexity:** Assuming  $n \geq m$ , the time is bounded by parsing the input.  $O(n \log(n))$  with map or  $O(n)$  with hashmap.

# C++ example (0.09 sec)



## babelfish.cpp

```
1 #include <iostream>
2 #include <sstream>
3 #include <string>
4 #include <unordered_map>
5
6 using namespace std;
7
8 int main() {
9     // Improves IO speed without having to use the cstdio routines,
10    // which are particularly cumbersome for string data. Note call
11    // order, this is important!
12    ios_base::sync_with_stdio(false);
13    cin.tie(NULL);
14
15    unordered_map<string, string> dict;
16    while (true) {
17        string line, a, b;
18        getline(cin, line);
19        if (line.size() == 0) break;
20        istringstream iss(line);
21        iss >> a >> b;
22        dict[b] = a;
23    }
24
25    string c;
26    while (cin >> c) {
27        if (dict.find(c) == dict.end()) {
28            cout << "eh" << "\n";
29        } else {
30            cout << dict[c] << "\n";
31        }
32    }
33    // Avoid std::endl. It flushes the buffer, which reduces performance
34    // when the output is large. Prefer "\n" for line breaks and do a
35    // single std::flush at the end.
36    cout << flush;
37
38    return 0;
39 }
```

# Java example (0.57 sec)



## Babelfish.java

```
1 import java.util.HashMap;
2 import java.io.*;
3
4 public class Babelfish {
5     public static void main(String[] args) throws IOException{
6         // Buffered IO to improve speed. For numeric input, use
7         // the Kattio class available on the Kattis help pages.
8         BufferedReader input = new BufferedReader(new InputStreamReader(System.in));
9         PrintWriter output = new PrintWriter(new BufferedOutputStream(System.out));
10
11         HashMap<String, String> dict = new HashMap<>();
12         while (true) {
13             String[] line = input.readLine().split(" ");
14             if (line.length < 2) break;
15             dict.put(line[1], line[0]);
16         }
17         while (true) {
18             String c = input.readLine();
19             if (c == null) break;
20             if (dict.containsKey(c)) {
21                 output.println(dict.get(c));
22             } else {
23                 output.println("eh");
24             }
25         }
26         output.flush(); // Crucial when using buffered output!
27     }
28 }
```

# Python example (0.62 sec)



## babelfish.py

```
1 dict = {}
2 line = input()
3 while line:
4     a, b = line.split()
5     dict[b] = a
6     line = input()
7 while True:
8     try:
9         c = input()
10        print('eh' if c not in dict else dict[c])
11    except:
12        break
13
```

# Notes on Input/Output



- Fast I/O: See Kattis help section.
- The IO specification describes the content of ONE test file, i.e. what your program must process in a single run. Your program might be executed several times with different input files.
- Make sure to read the IO specification carefully! In particular:
  - How are sections of the input data separated?
  - How do you know how many datapoints there are?
  - How do you know that there is no more input? Common variants: An integer  $n$  at the beginning, blank lines, special input tokens (such as "o o" before EOF in CD), or simply EOF.
  - Are there multiple test cases in the same input?
- What are the limits of the problem? Is it e.g. allowed to have 0 datapoints?
- Note that the sample input/output doesn't always cover all possible input/output formats!

- **Problem:** Given two sets  $A$  and  $B$  of integers ( $|A| = n, |B| = m \leq 1,000,000$ ), find the size of the intersection  $A \cap B$ .
- **Solution:** Again, the problem can be solved using lookup if we choose an appropriate data structure. There are several options:
  - Use vectors/arrays, sort them (already given in increasing order) and lookup with binary search.
  - Use set/hashset similar to map/hashmap in the previous problem.
- **Time complexity:** Assuming  $n$  in  $O(m)$ .
  - Using vectors/arrays: input in  $O(n)$ , sort in  $O(n \log n)$ , lookup in  $O(n \log n) \Rightarrow O(n \log n)$
  - Using set/hashset: both input and lookup in  $O(n \log n)/O(n)$ .
- **Warning:** The input size is quite large (up to a few MB), so make sure to use techniques for fast IO.



# Closest Sums



- Given is a set  $A = \{a_1, \dots, a_n\}$  of  $n$  ( $2 \leq n \leq 1,000$ ) integers and  $m$  ( $1 \leq m \leq 24$ ) query integers  $b_i$ .
- **Problem:** For each  $b_i$ , output  $c_i$  such that  $c_i = a_j + a_k$  where  $j \neq k$  and  $|b_i - c_i|$  is minimized.
- **Subproblem:** Find all possible sums of pairs of integers in  $A$ .
  - Iterate over  $A$  with two nested loops.  $O(n^2)$  is fine since  $n$  is small.
- **Solution:** Calculate all possible sums of pairs, store them in an efficient data structure and use lookup for all queries.
  - Here, a hash map doesn't work, since there might not be an exact match.
  - Use a sorted vector/array or a map, which allows lookup of closest match in  $O(\log(n))$ .
- **Time complexity:**  $O(n^2 \log(n^2)) = O(n^2 \log(n))$  to calculate all sums to an appropriate data structure.  $O(m \log(n^2)) = O(m \log(n))$  for lookups.

# Mathemagicians



- $n$  magicians ( $3 \leq n \leq 100,000$ ) magicians stand in a circle, each wearing either a red or a blue hat.
- The magicians can (one at a time) change the color of their hats to match one of their neighbors'.
- **Problem:** Is it possible to reach a given target configuration from a given original configuration?
- **Note:** The state space is huge ( $2^{100,000} \approx 10^{30,000}$ ), so conventional search techniques (to be covered later in the course) are not feasible.

# Mathemagicians



- Think in terms of contiguous “fields” of magicians wearing the same color hat.
- Operations on the borders extend/contract the size of adjacent fields.
- **Key conclusion**: We can only move or remove fields, not create new ones!
- **Solution**: Count the number of fields  $a$  in the original and  $b$  in the target configuration. Bar special cases, the target is reachable if  $b \leq a$ .
  - Special case 1:  $a = 1$  and all magicians wear the wrong color.
  - Special case 2: The magicians wear alternating colors and the target is the opposite configuration of alternating colors.
- **Time complexity**: Count the number of fields in  $O(n)$ .

# Problem set



- On the live sessions later during the semester, there will be six problem instead of today's four.
- Today's problems mainly involved basic data structures, but many problems on live sessions are based on more advanced algorithms and data structures implemented in the labs.
- According to the ranking at [open.kattis.com](https://open.kattis.com), today's problems ranged from 2.1 to 4.8 in difficulty on a scale from 1-10. The difficulty of the live problems are usually in the range 2.5-7.5.
- Today the problems were arranged approx. in increasing order of difficulty, but this will not be the case in general.

- Most of today's problems were about basic data structures. If you found them difficult, consider recapitulating theory from your data structures and algorithms course.
  - OpenDSA is a good free online resource:  
<https://www.ida.liu.se/opensa/Books/TDDD86F2o/html/>
  - Make sure you are familiar with how they are implemented in your chosen programming language.
- Be familiar with how Kattis handles input and output.
  - Consult the documentation for your language.
  - Know how to ensure fast IO for large datasets.