## TDDD55

## 1 EXercises in Formal Languages and Automata Theory

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This material presents problems with partial solutions related to formal languages and automata theory for TDDB44 and TDDD55

## 2 Problems

### 2.1 Transforming a language to a DFA

Construct a DFA that accepts the language for each of the following languages:
A. $L_{1}=\left\{x \in\{0,1\}^{*} \mid x\right.$ should end with 00$\}$
B. $L_{2}=\left\{x \in\{0,1\}^{*} \mid x=(01)^{n}, n \geq 0\right\}$
C. $L_{3}=\left\{x \in\{0,1\}^{*} \mid\right.$ Every 0 is immediately followed by 1$\}$

### 2.2 From NFA to DFA

Given the NFA in Figure 1, see below, construct an equivalent DFA.


Figure 1: $M_{8}$

### 2.3 NFA Construction

Construct an NFA that accepts the language defined by the following regular expression:
A. $0+1$
B. $[0]^{+}+1$
C. $[0-1]^{+}+1$
D. $10+(0+11) 0^{*} 1$

## 2.4 (Bonus) Describing Regular Expressions using Grammars

Provide Context-Free Grammar for the following regular expressions:
A. $00(1+0) * 1$
B. $101(101) * 010(010)^{*}$
C. $(11+010) * 11(00+11) *$

## Partial Solutions

### 2.1 A)

An example of a DFA M15 such that $\mathbf{L}\left(\mathbf{M}_{\mathbf{1 5}}\right)=\mathbf{L} \mathbf{1}$ is given below


## 2.1 в)

An example of a DFA $\mathrm{M}_{16}$ such that $\mathrm{L}\left(\mathrm{M}_{16}\right)=\mathrm{L}_{2}$ is below


### 2.1 C)

An example of a DFA $\mathrm{M}_{16}$ such that $\mathrm{L}\left(\mathrm{M}_{16}\right)=\mathrm{L}_{2}$ is provided below


## 2.2

Hint: the subset construction results in a DFA with reachable states:

- $\{q 0, q 1, q 3, q 5, q 6, q 7, q 10\}$
- $\{q 2, q 9, q 10$,
- $\quad\{q 4, q 10\}$
- $\}$
- $\{q 9\}$

The initial state is $\{q 0, \mathrm{q} 1, \mathrm{q} 3, \mathrm{q} 5, \mathrm{q} 6, \mathrm{q} 7, \mathrm{q} 10\}$. The final states are those containing q 10 .

## 2.3 (C)

By decomposing the regular expression syntactically according to the recursive definition of regular expressions, an NFA can be constructed systematically in a bottom-up fashion. By successively joining NFA: s corresponding to subexpressions according to the standard operators:

$$
(*,+,+)
$$

The resulting NFA is depicted in Figure 5: $\mathrm{M}_{23}$

## 2.4

A)


Figure 5: $M_{23}$

$$
\begin{aligned}
& S \rightarrow 00 A 1 \\
& A \rightarrow \epsilon|0 A| 1 A
\end{aligned}
$$

B)
$S \rightarrow 101 A 010 B$
$A \rightarrow \epsilon \mid 101 A$
$B \rightarrow \epsilon \mid 010 B$
C)
$S \rightarrow A 11 B$
$A \rightarrow \epsilon|11 A| 010 A$
$B \rightarrow \epsilon|00 B| 11 B$

