



Automated Planning

Landmark Heuristics for State Space Planning

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Landmarks (1)

Landmark:

"a **geographic feature** used by explorers and others to **find their way** back or through an area"



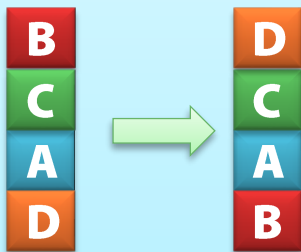
Landmarks in planning:

Something you must *achieve* or *use* in *every solution* to a problem instance

Assume we are considering a state s ...

Fact Landmark for s :

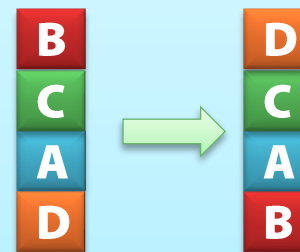
A fact that must be true
at some point
in every solution starting in s



$\text{clear}(A)$
 $\text{holding}(C)$
...

Formula Landmark for s :

A formula that must be true
at some point
in every solution starting in s

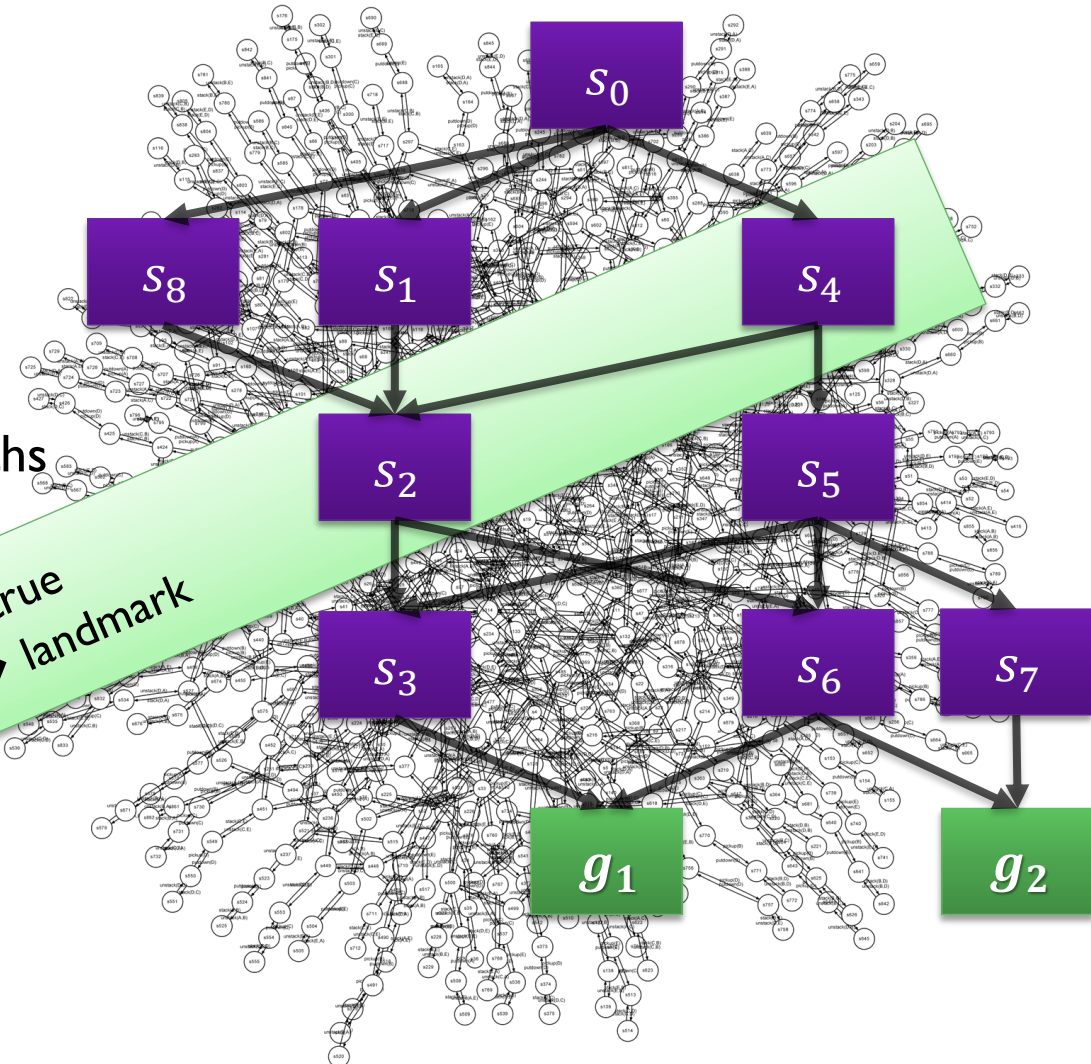


$\text{clear}(A) \wedge \text{handempty}$
...

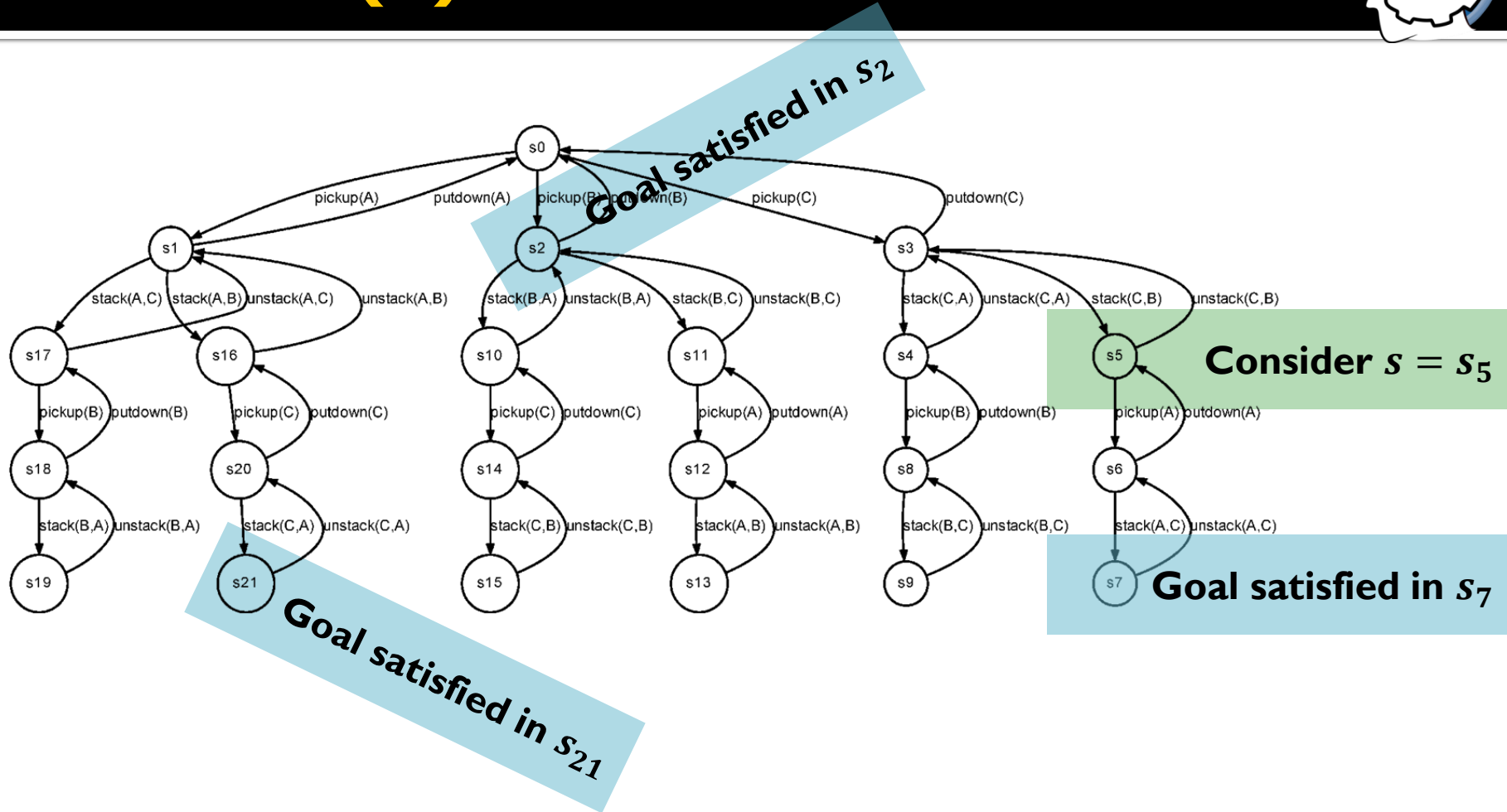
Facts and formulas, not states! Why?

- Usually many paths lead from s to goal states
- Few states are shared among all paths
- Many facts occur along all paths
 - Landmarks!

In S_2 and S_4 , **holding(B)** is true
 → occurs along all paths → landmark



Landmarks (4)



Is there a landmark state s_{lm} we must pass to reach some goal from s_5 ?

No! But we may have to pass different states satisfying the same facts f_{lm} !

Landmarks (5): Misunderstandings



Not "we must reach (pass through)
the landmark state"

Instead "we must reach
some state that satisfies
the fact/formula landmark"

Not "A landmark fact is a state that..."

A fact is not a state.
A state *consists of many* facts.
("A word is a sentence that...")

A landmark fact is not
"a fact that is true in every solution"

A solution is a plan.
Facts are true in *states*.

A landmark fact is
"a fact that is true in some state
along every path
from the initial state to any goal state".

Can you be "close" to a landmark?

You **can** be in a state s that is close to
another state s' satisfying the landmark.

Problem: How to know?

Distance is "number of edges" or "cost of
reaching", not $\Delta x / \Delta y$. And the graph may
not even be expanded yet.

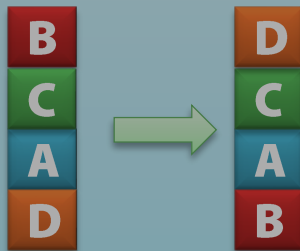
Landmarks in planning:

Something you must *pass by/through* in every solution to a specific planning problem

Assume we are currently in state s ...

Fact Landmark for s :

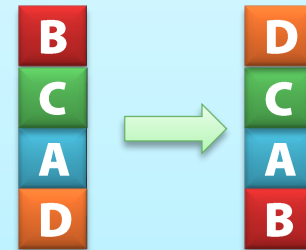
A **fact** that must be true
at some point
in every solution starting in s



clear(A)
holding(C)
...

Action Landmark for s :

An **action** that must be used
in every solution starting in s



unstack(B,C)
putdown(B)
stack(D,C)

...so the effects of
action landmarks
are *fact landmarks*,
and so are their
preconds

(except those facts
that are already true
in s)

- Generalization:
 - **Disjunctive** action landmark $\{a_1, a_2, a_3\}$ for state s
 - Every solution starting in state s and reaching a goal must use *at least one* of these actions

Finding Landmarks: A (Too) General Technique

Finding Landmarks: General Technique

- One general technique for discovering landmarks:

Current planning problem, P

Initial state does not include atom A



Modified planning problem, P'

*Removed all actions
that add atom A*



...then **every** solution to P
must use one of the removed actions

→ Action set is a disj. act. landmark

→ Atom A is a fact landmark



If this problem (P') is **unsolvable**...

Test:

Delete relaxation of P' is
unsolvable,

or $h_m(s_0) = \infty$, or ...

→ P' is unsolvable

Unsolvable when removing a set of actions

→ some action in the set must be used → **disjunctive action landmark!**

Finding Landmarks: General Technique (2)



- This technique is very general
 - Applicable to *any* planning problem, *any* atom
- General techniques tend to be widely applicable but slow...

Verifying Landmarks (1)



- How difficult is it to verify that an action is an action landmark, in the general case?
 - Suppose we can verify this
 - Then given any STRIPS problem P , we can determine if it has a solution:
 - Add a new action:
 - cheat
 - :precond true
 - :effects goal-formula
 - If cheat is an action landmark, then it is *needed* in order to solve the problem
 - ➔ the original problem was *unsolvable*
 - ➔ As difficult as solving the planning problem (PSPACE-complete)

Verifying Landmarks (2)



- How difficult is it to verify that a fact is a fact landmark, in the general case?
 - Suppose we can verify this
 - Then given any STRIPS problem P , we can determine if it has a solution:
 - Add a new fact:
 - cheated (false in the initial state)
 - Add new action:
 - cheat
 - :precond true
 - :effects (and cheated goal-formula)
 - If cheated is a fact landmark,
then cheat was necessary → the original problem was unsolvable
 - → Again , as difficult as solving the planning problem

But of course there are special cases...

Finding Landmarks: Efficiently

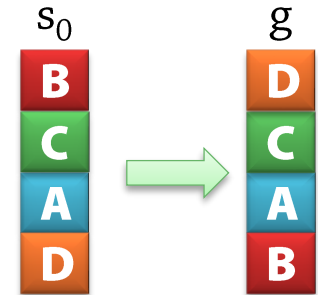
Means-Ends Analysis

- Discover landmarks using means-ends analysis

Unachieved goal facts

are (obviously) fact landmarks:

clear(D), on(D,C), on(A,B), ontable(B)



$\text{fact-landmarks} \leftarrow g - s$

Suppose $\text{on}(D,C)$ is a landmark,
 $\text{on}(D,C)$ is not true in the current state (s)
 $\rightarrow$ we must *cause* **on(D,C)** with an action
 $\rightarrow$ compute $\text{achievers} = \{ \text{stack}(D,C) \}$

do {
 for each p in fact-landmarks {
 // Create *disjunctive* action landmark
 $\text{achievers} \leftarrow \{a \in A \mid p \in \text{eff}(a)\}$

All achievers require these candidates =
 $\{ \text{holding}(D), \text{handempty}, \text{clear}(C), \dots \}$

$\text{candidates} \leftarrow \bigcap_{a \in \text{achievers}} \text{pre}(a)$

handempty is already true, but
 $\text{new} = \{ \text{holding}(D), \text{clear}(C) \}$ are *not*

$\text{new} \leftarrow \text{candidates} - s$
 $\text{fact-landmarks} \leftarrow \text{fact-landmarks} \cup \text{new}$

Maybe we can find more landmarks
related to achieving *those*!

}
until no more fact-landmarks found

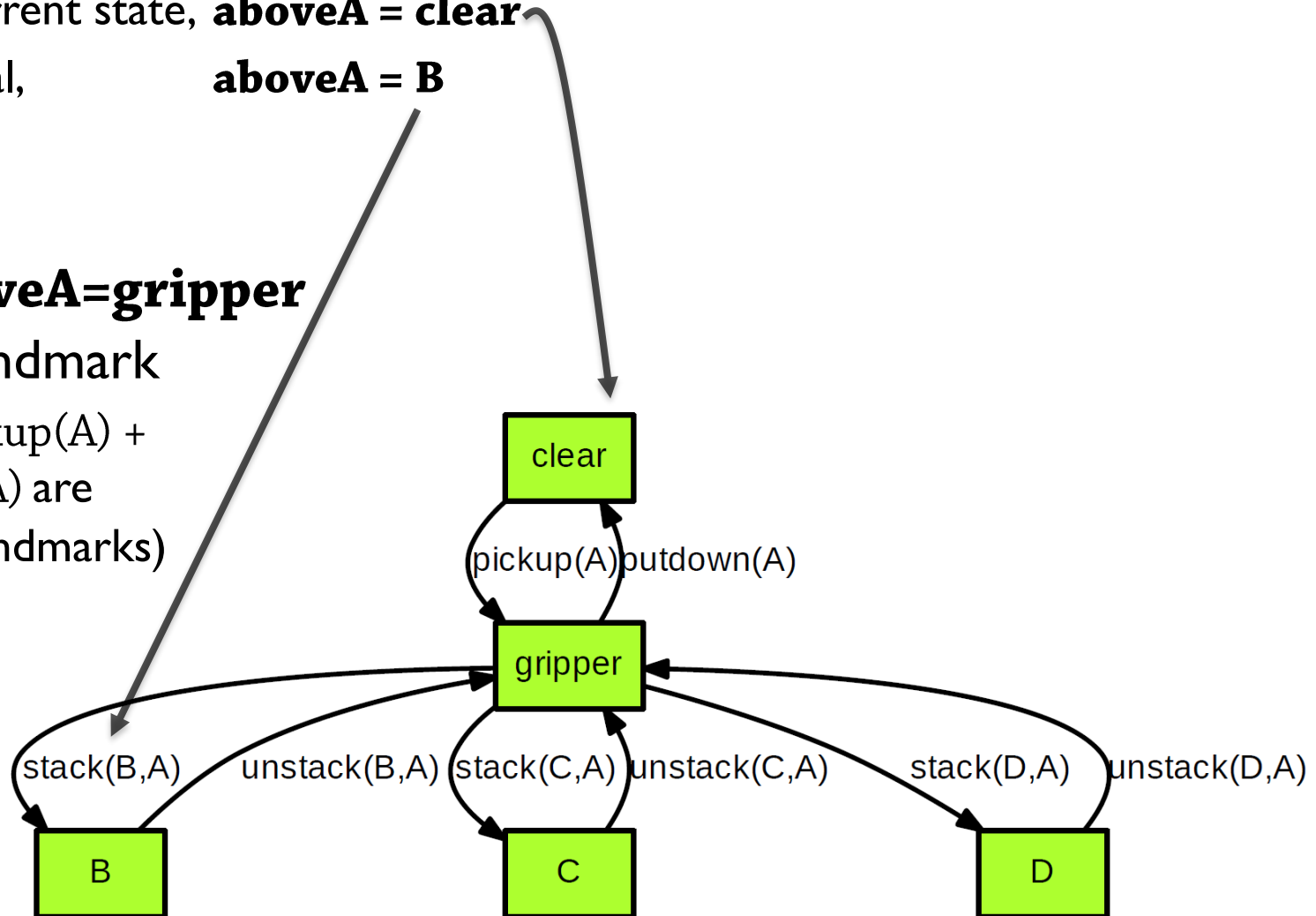
- Extensions to backwards means-ends analysis:
 - **Effects of disjunctive action landmarks:**
 - All shared **effects** must also take place regardless of the "chosen" action, similarly to shared *preconditions* on the previous page
 - Given a disjunctive action landmark,
every fact in $(\bigcap \{\text{eff}(a) \mid a \in \text{landmark}\} - s)$ is a fact landmark for s

- Another method: Use domain transition graphs:

- In the current state, **aboveA = clear**
- In the goal, **aboveA = B**

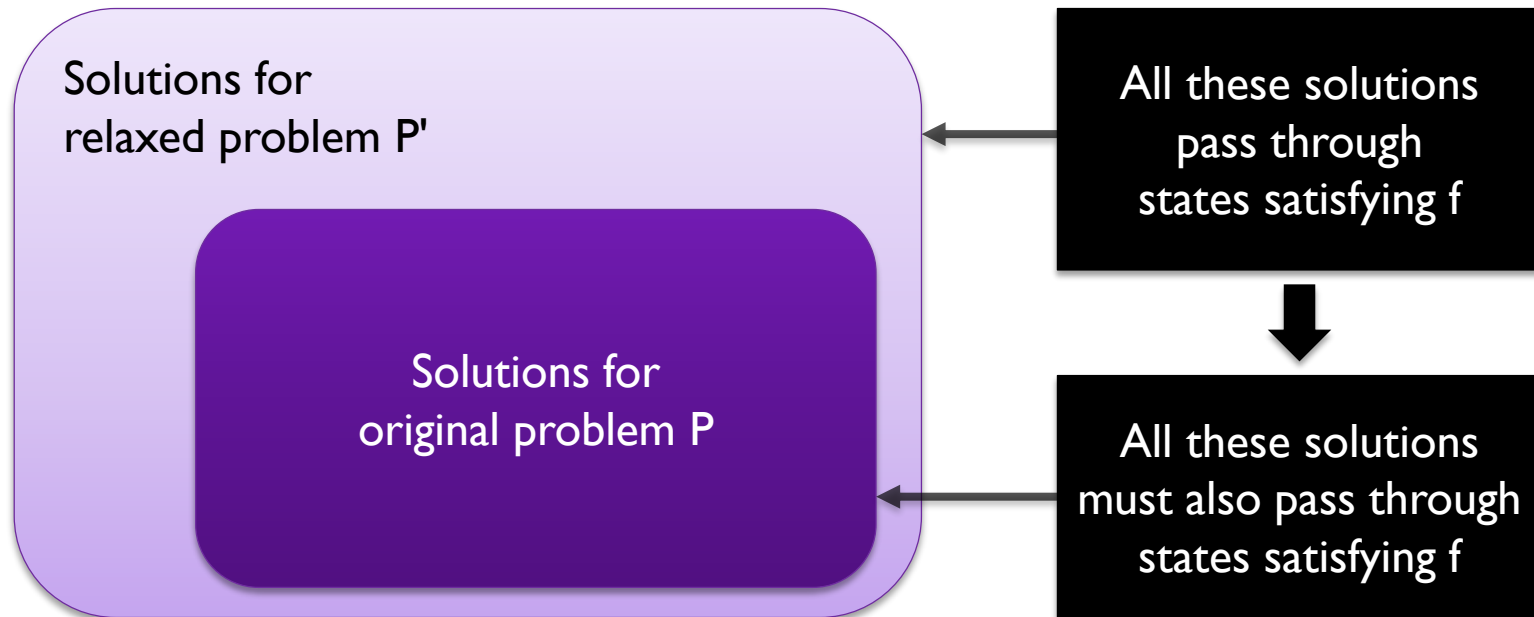
- Then **aboveA=gripper** is a fact landmark

- (And pickup(A) + stack(B,A) are action landmarks)



Landmarks and Relaxation

- Assume a problem P , and a relaxed problem P'
 - Suppose f is a fact landmark for P'



- Then f is a fact landmark for the original problem as well!
- Similarly for action landmarks, etc.

- Many other techniques exist...
 - Beyond the scope of the course

Landmark Ordering

Landmark Ordering (1)

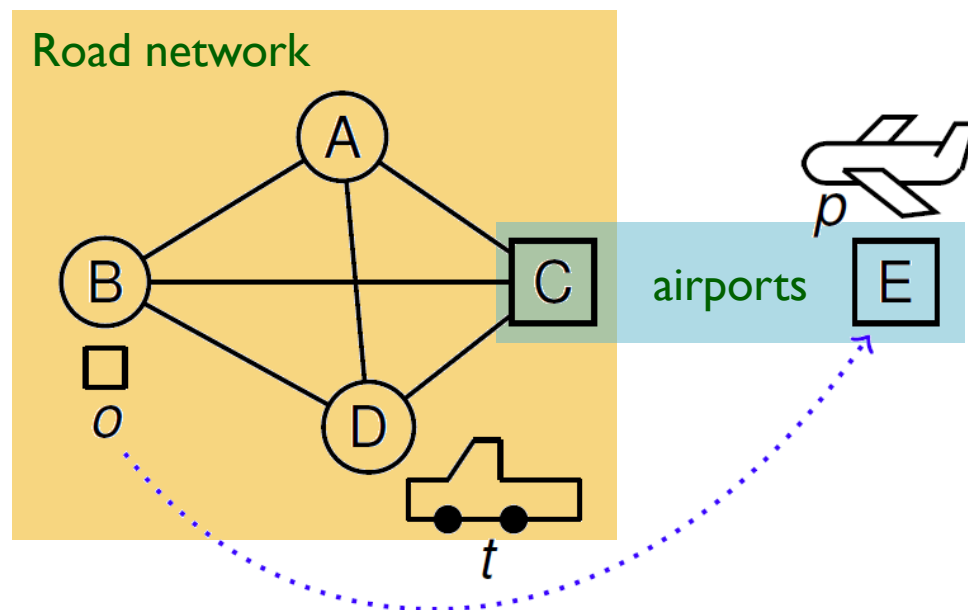


- Sometimes we can find or approximate necessary orderings
 - We must achieve $\text{holding}(A)$, *then* $\text{holding}(B)$

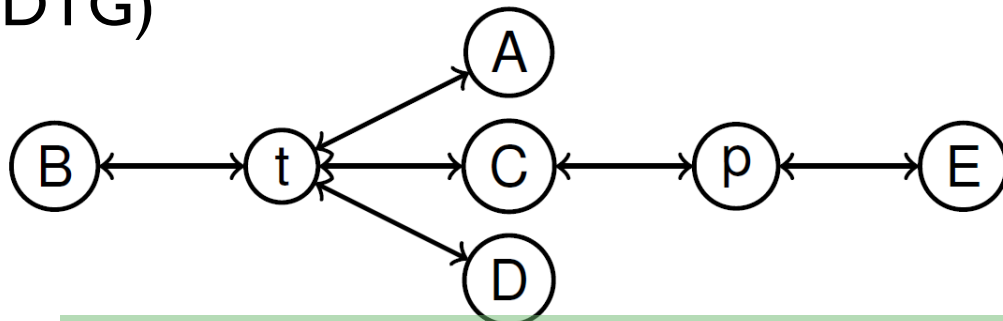
Landmark Ordering (2): Example Problem

■ Example Problem:

- Truck t transports object o within **road network** A/B/C/D
- Airplane p transports object between **airports** C/E
- Goal: Object at E

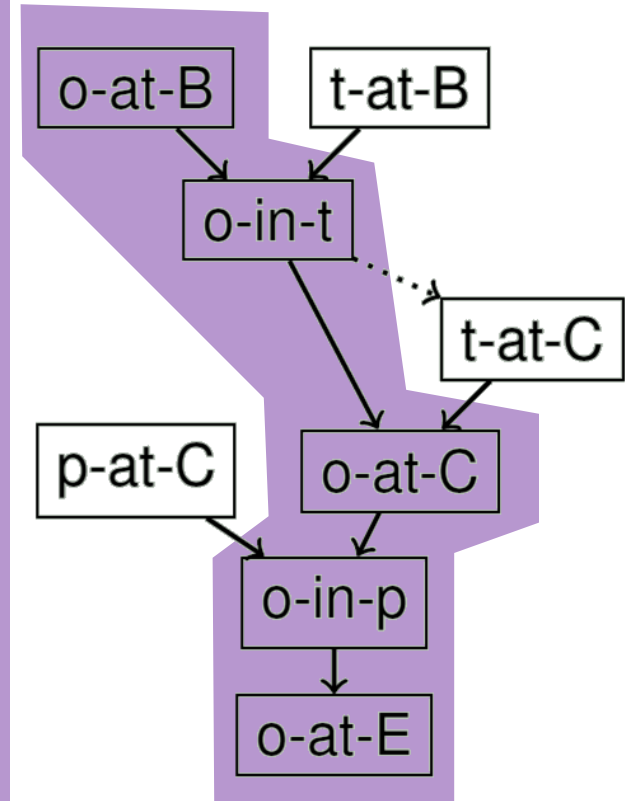
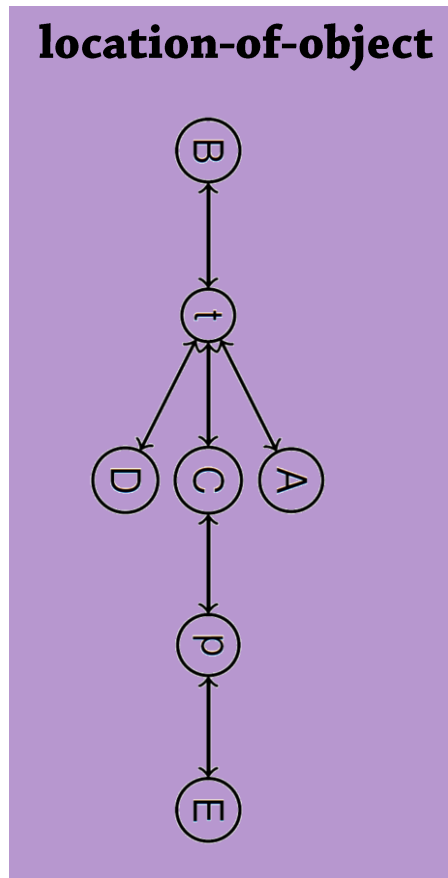
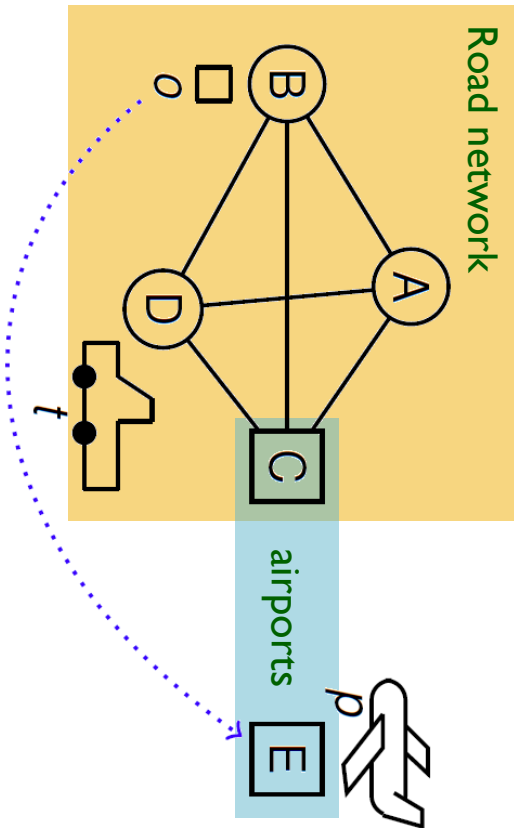


■ Domain transition graph (DTG) for **location-of-object**:



Note: Every **edge** in the road network corresponds to a **path** through t in the DTG!

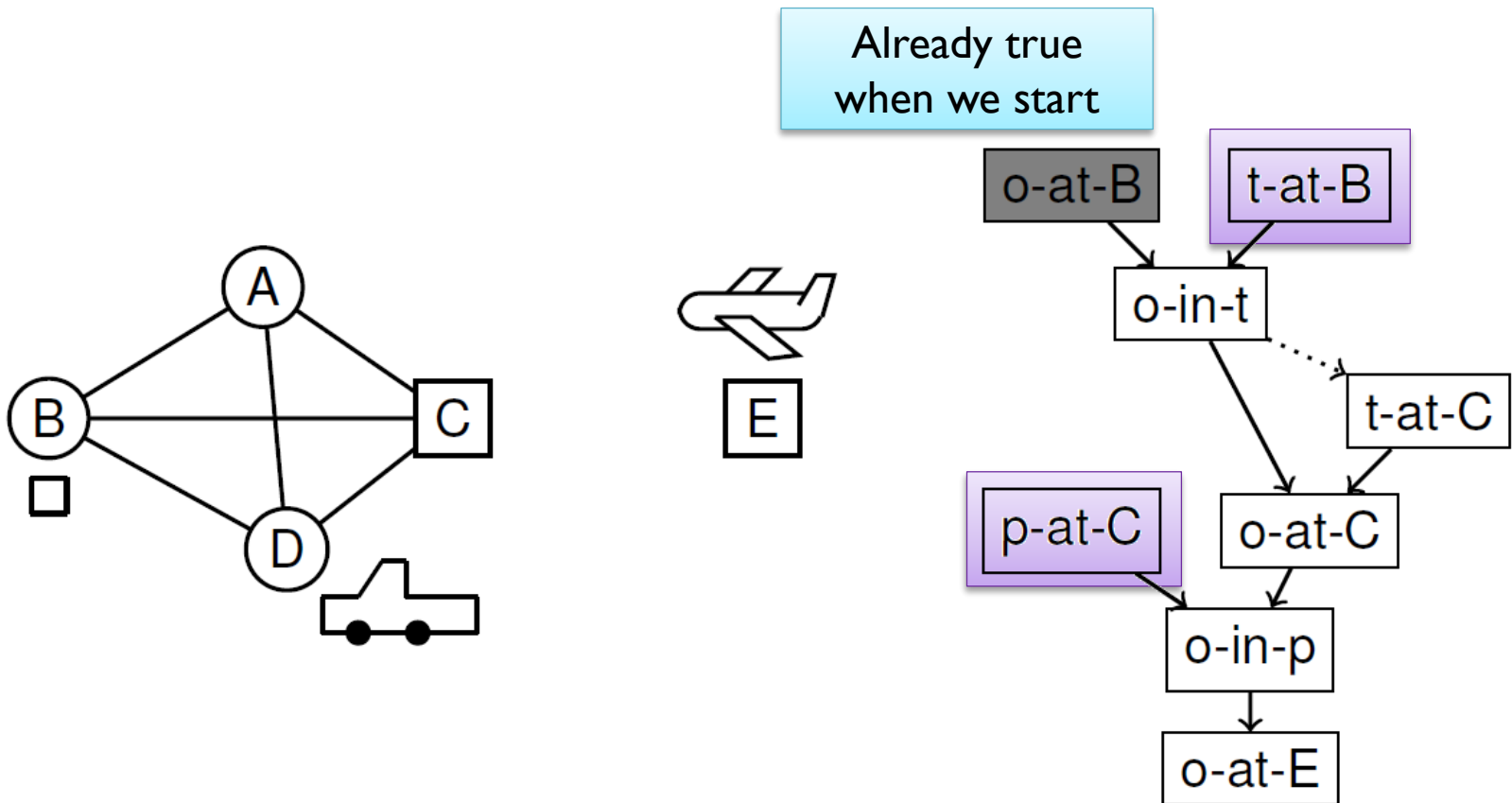
- **Directly from the DTG!**



Using Landmarks: As Subgoals

Landmarks as Subgoals (1)

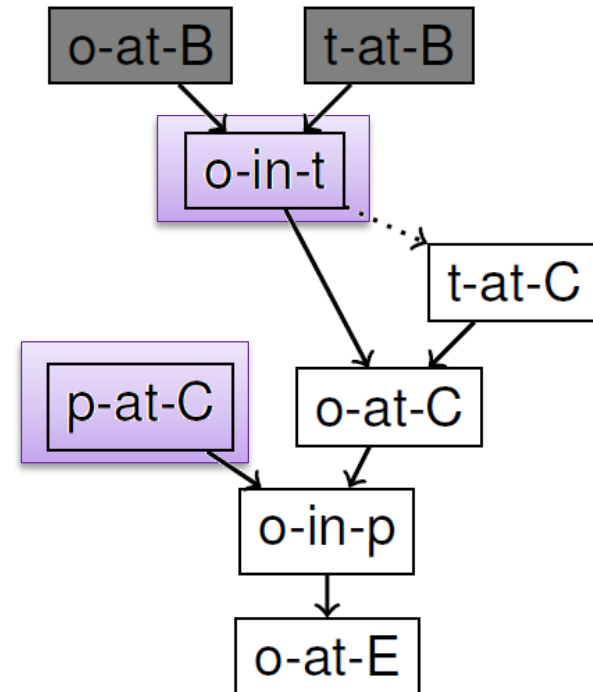
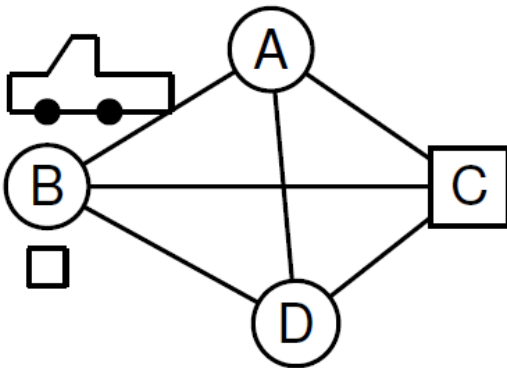
- One use of **ordered** landmarks:
 - As **subgoals**: Try to plan for each landmark **separately** in the inferred **order**



Two landmarks could be "first" (all predecessors achieved)
Current goal: $t\text{-at-B} \vee p\text{-at-C}$ (disjunctive!)

Landmarks as Subgoals (2)

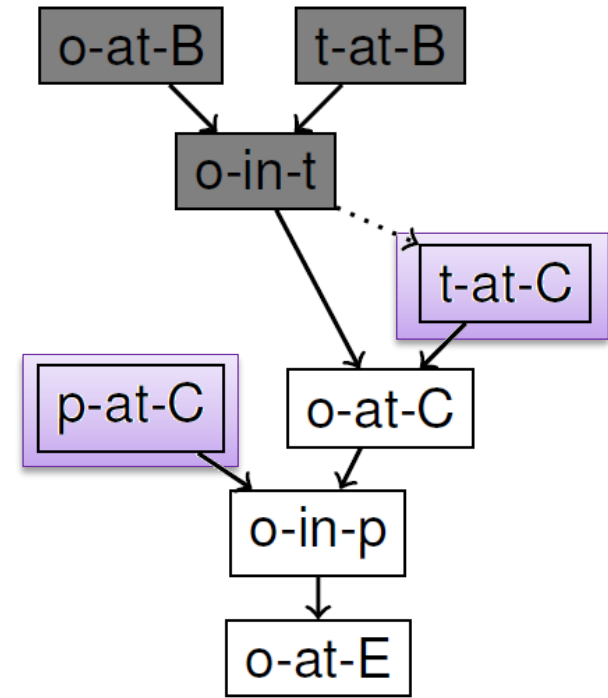
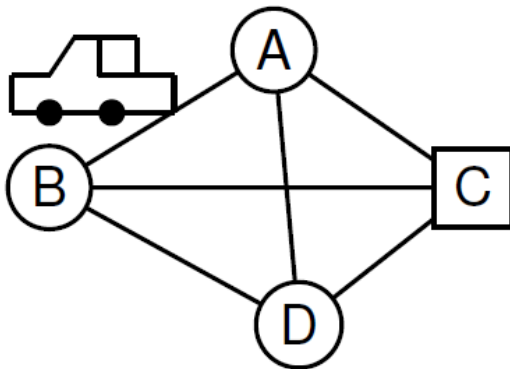
Suppose we begin by achieving t-at-B:
Simple planning problem,
results in a single action -- drive(t, B)



Current goal: o-in-T or p-at-C

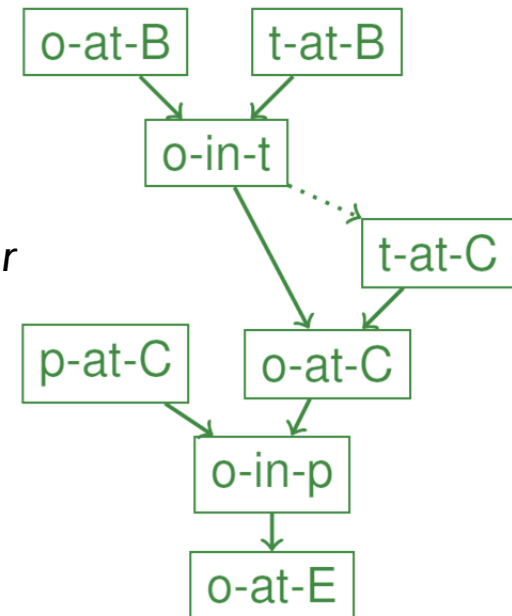
Landmarks as Subgoals (3)

Suppose we continue by achieving o-in-T:
Simple planning problem,
results in a single action -- load-truck(o,t,B)



Landmarks as Subgoals (4)

- Sometimes very helpful, but:
 - There are still **choices** to be made – backtrack points!
 - Optimizing for one **part** of the overall goal at a time:
 - Can't see the whole picture
 - Can miss opportunities:
Cheapest solution *here* → **more expensive** solution *later*
 - Can be incomplete:
Cheapest solution *here* → **impossible** to solve *later*

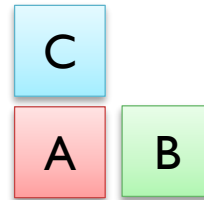


The Problem with Separating Subgoals

- The Sussman Anomaly (Gerald Sussman)

- Goal is **on(A,B)**, **on(B,C)**

- Now:

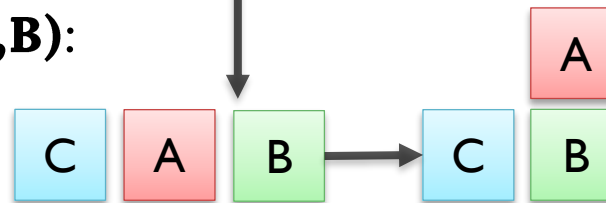


- Idea: Achieve one at a time

- First, plan only for **on(A,B)**
- Then, plan only for **on(B,C)**

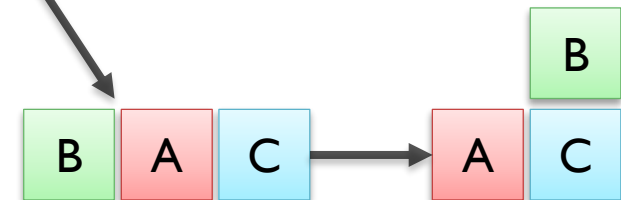
- Achieve first subgoal, **on(A,B)**:

- unstack(C,A); putdown(C);
pickup(A); stack(A,B)



- Achieve second subgoal, **on(B,C)**:

- unstack(A,B); putdown(A);
pickup(B); stack(B,C) → original goal destroyed!

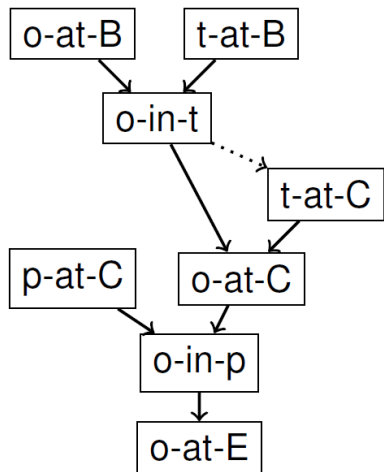


Using Landmarks: To Define Heuristic Functions

Landmark Heuristics (1)

- The LAMA state space planner **counts** landmarks:
 - Landmarks that need to be achieved after reaching state s through path (action sequence) π

■ $L(s, \pi) =$



$(L \setminus \text{Accepted}(s, \pi))$

$\cup$

$\text{ReqAgain}(s, \pi)$

All discovered landmarks,
minus those that are
accepted as achieved
(have become true *after*
predecessors are achieved!)

Plus those we can show will
have to be *re-achieved*

(Example:
Landmarks that were reached,
are no longer true,
but are required by the goal)

- $h(s) = |L(s, \pi)|$
- Not admissible: One action may achieve multiple landmarks!

Landmark Heuristics (2)

- To achieve **admissible** heuristic estimates:

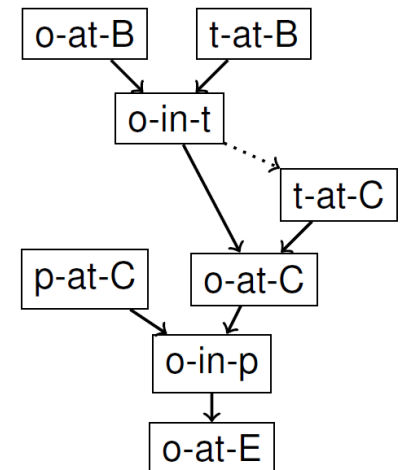
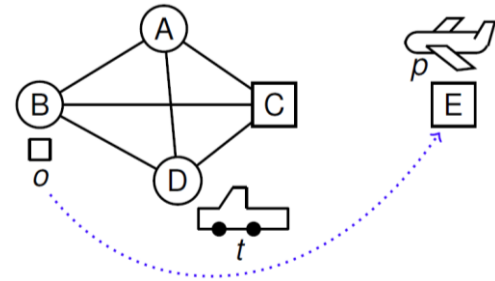
- Idea: The cost of each action is *divided* across the landmarks it achieves

- **Simplified example:**

- Suppose there is a goto-and-pickup action of cost 10, that achieves both t-at-B and o-in-t
- Suppose *no other action* can achieve these landmarks
- One can then let (for example)
 $\text{cost}(\text{t-at-B})=3$ and $\text{cost}(\text{o-in-t})=7$

- The sum of the cost of remaining landmarks is then an **admissible heuristic**

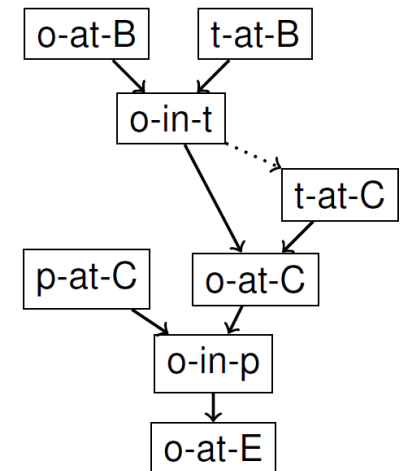
- Must decide how to split costs across landmarks
- Optimal split *can* be computed polynomially, but is still expensive



Using Landmarks: For Problem Modification

■ Landmarks as a basis for a modified planning problem

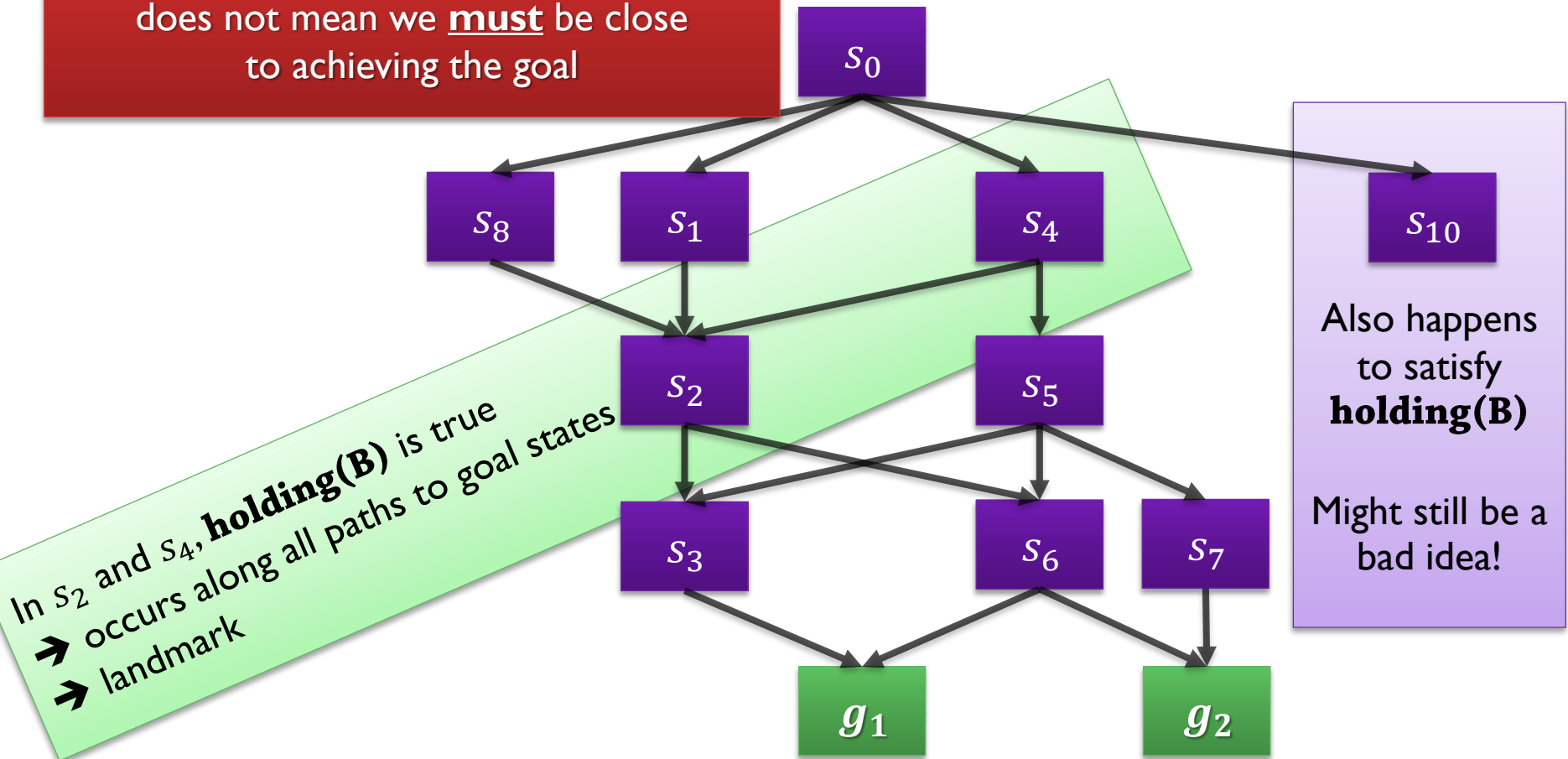
- Add new facts "achieved-landmark-*n*"
 - Concretely: *object-has-been-in-plane*
- An action achieving a landmark makes the corresponding facts true
 - (load object plane) → *object-has-been-in-plane* := true
- The goal requires all such facts to be true
 - (:goal *object-has-been-in-plane* ...)
- → Any *other* heuristic can be applied to the modified problem!
 - $h_{PDB}(s)$ – pattern databases: What is the cost of achieving *object-has-been-in-plane*?



Landmarks: Ideas that Won't Work

Landmarks: More Misunderstandings

Satisfying a landmark
does not mean we must be close
to achieving the goal

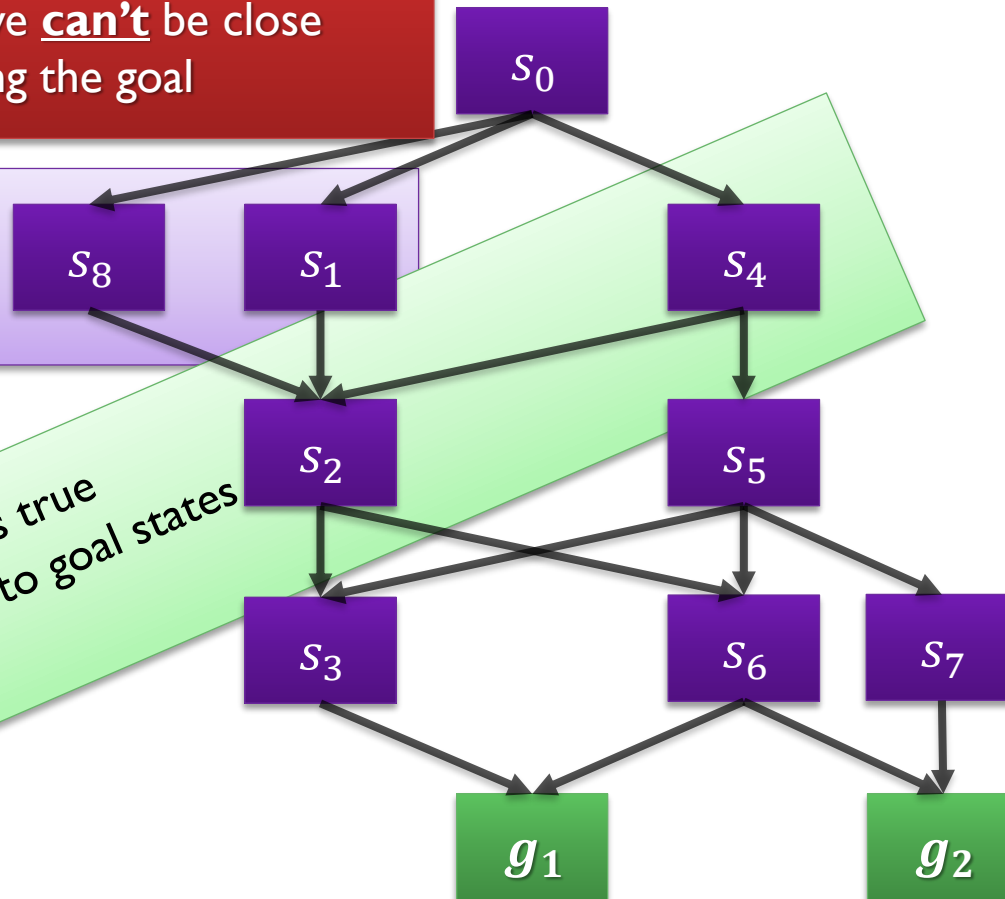


Landmarks: More Misunderstandings

Not satisfying a landmark
does not mean we can't be close
to achieving the goal

Satisfied no
landmarks yet?
Still a good idea

In s_2 and s_4 , **holding(B)** is true
→ occurs along all paths to goal states
→ landmark



Landmarks: More Misunderstandings

We rarely talk about "the" landmark

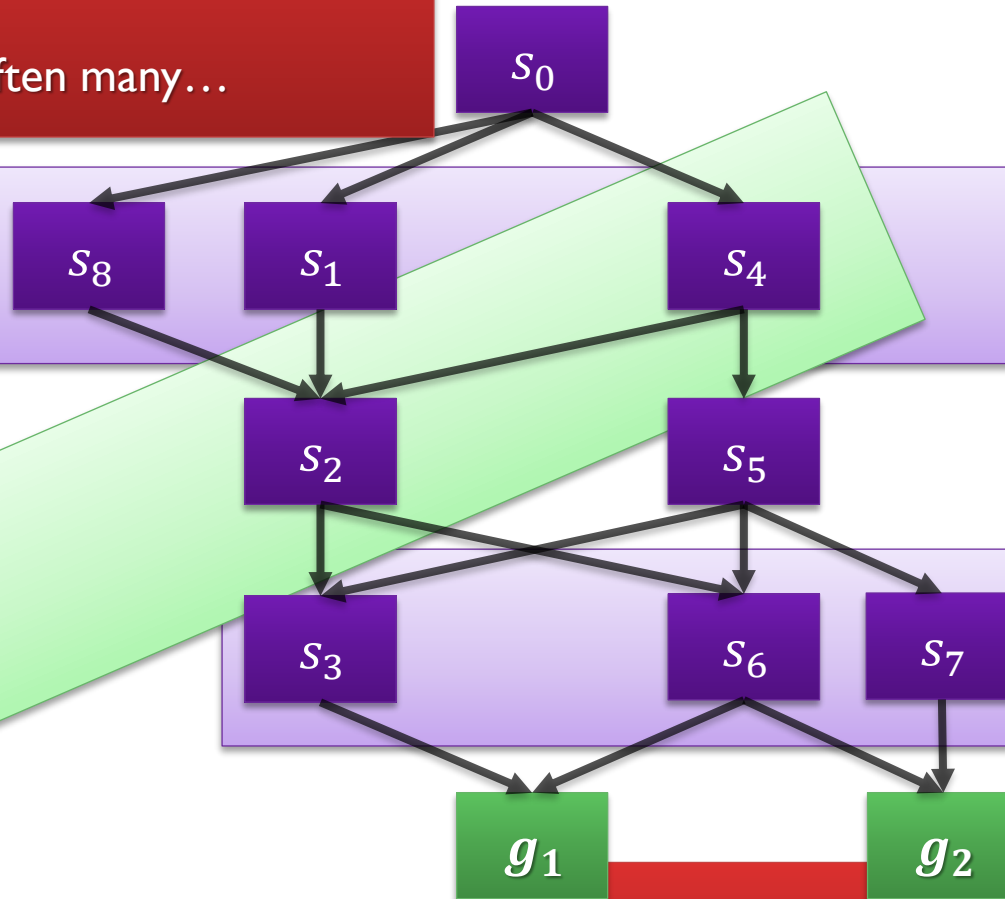
There are often many...

Satisfying **on(D,E)**

Satisfying **holding(B)**

Satisfying ...

A landmark is **not** a unique location
(unlike real-world landmarks):
The intuition breaks down...



Landmarks: Ideas that Don't Work

"If I reach a state satisfying a landmark, I won't have to backtrack"

