

# **Notation for Processes**

- C = Worst-case execution time
- B = Worst-case blocking time
- D =Relative deadline
- n = Number of processes
- T = Period
- R = Worst-case response time
- J =Release jitter

### **Schedulability test for Rate Monotonic:**

$$\sum_{i=1}^n \left(\frac{C_i}{T_i}\right) \le n(2^{1/n}-1)$$

### **Schedulability test Earliest Deadline First:**

$$\sum_{i=1}^{n} \left( \frac{C_i}{T_i} \right) \le 1$$

### **RMS Response time analysis**

$$w_{i} = C_{i} + B_{i} + \sum_{\forall P_{j} \in lip(P_{i})} \left[ \frac{w_{i} + J_{j}}{T_{j}} \right] C_{j}$$
$$R_{i} = w_{i} + J_{i}$$

 $hp(P_i)$  is the set of processes with a higher priority than process  $P_i$ .





## Timing Analysis of CSMA/CR

B = blocking time C = transmission time of entire frame T = period  $\tau_{bit} = transmission time of one bit$  w = response time for the first bit of a frame to be sent R = total response time J = Jitter t = Longest busy interval lp(m) = set of messages with lower priority than m. hp(m) = set of messages with higher priority than m. hep(m) = set of messages with higher or equal priority than m.

$$\begin{split} R_m &= \max_{q=0..Q_m-1}(R_m(q))\\ R_m(q) &= J_m + w_m(q) - q \cdot T_m + C_m\\ w_m(q) &= B_m + q \cdot C_m + \sum_{\forall j \in hp(m)} \left[ \frac{w_m(q) + J_j + \tau_{bit}}{T_j} \right] \cdot C_j \end{split}$$

(with  $w_m^0(q) = B_m + C_m q$ )

$$\begin{aligned} \mathcal{Q}_m &= \left[ \frac{t_m + J_m}{T_m} \right] \\ t_m &= B_m + \sum_{j \in hep(m)} \left[ \frac{t_m + J_j}{T_j} \right] \cdot C_j \quad (\text{with } t_m^{0} = C_m^{-1}) \\ C_m &= \left( 8n + 47 + \left\lfloor \frac{34 + 8n - 1}{4} \right\rfloor \right) \tau_{bit} \\ B_m &= \max_{j \in lp(m)} (C_j^{-1}) \end{aligned}$$