

# Performance Estimation for Embedded Systems with Data and Control Dependencies

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# Motivation and Characteristics



## Performance estimation.

- Worst case delay on the system execution time.

## Characteristics:

- Distributed hard real-time applications.
- Heterogeneous system architectures.
- Fixed priority preemptive scheduling.
- Systems with data and control dependencies.

# Schedulability Analysis

## Schedulability test:

- The *worst case response time* of each process is compared to its *deadline*.

## Process models:

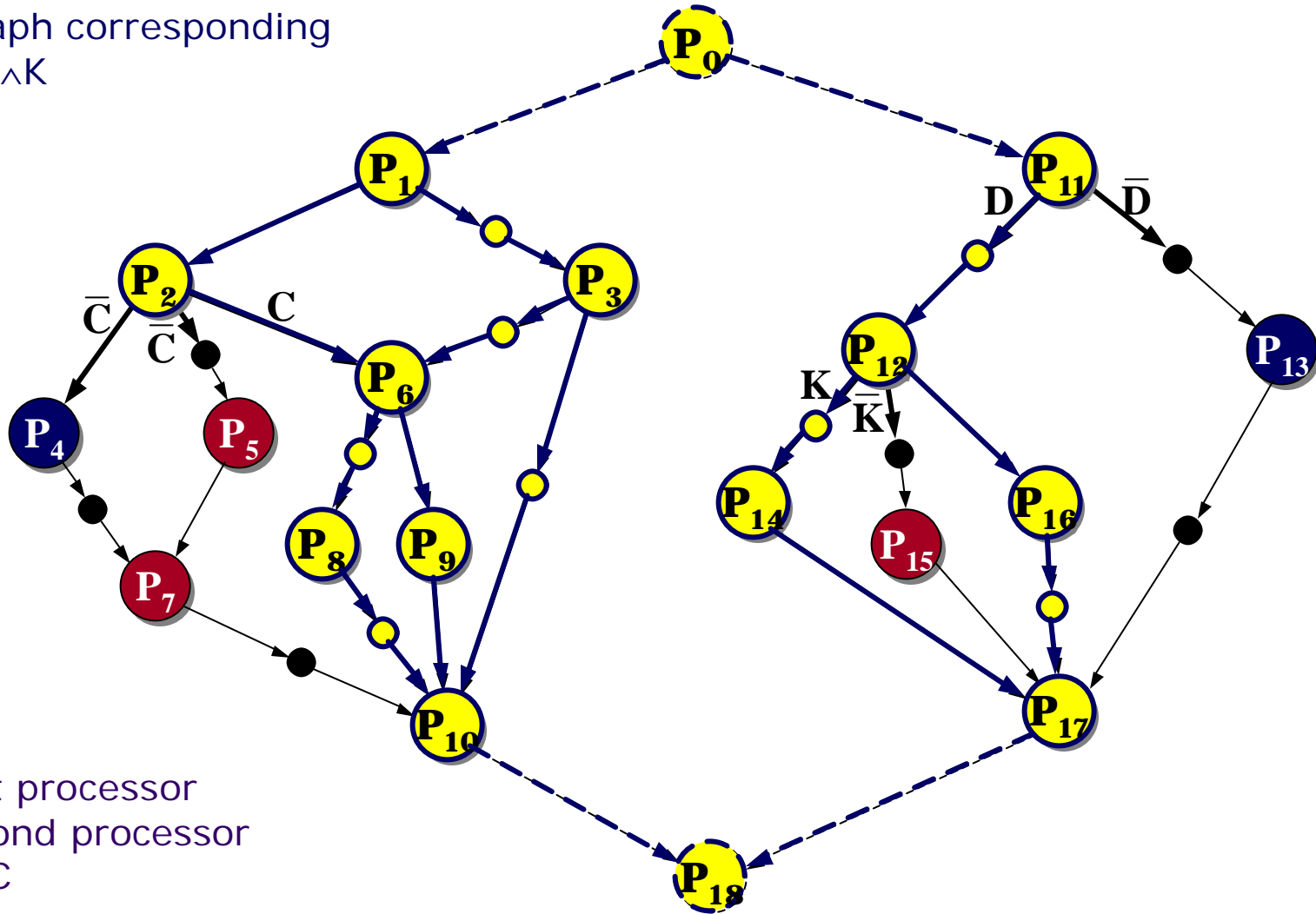
- Independent processes;
- Data dependencies: *release jitter, offsets, phases*;
- **Control dependencies:** *modes, periods, recurring tasks*.

## Message:

- The pessimism of the analysis can be significantly reduced by considering the conditions during the analysis.

# Conditional Process Graph

Subgraph corresponding to  $D \wedge C \wedge K$



- First processor
- Second processor
- ASIC

# Problem Formulation

## Input

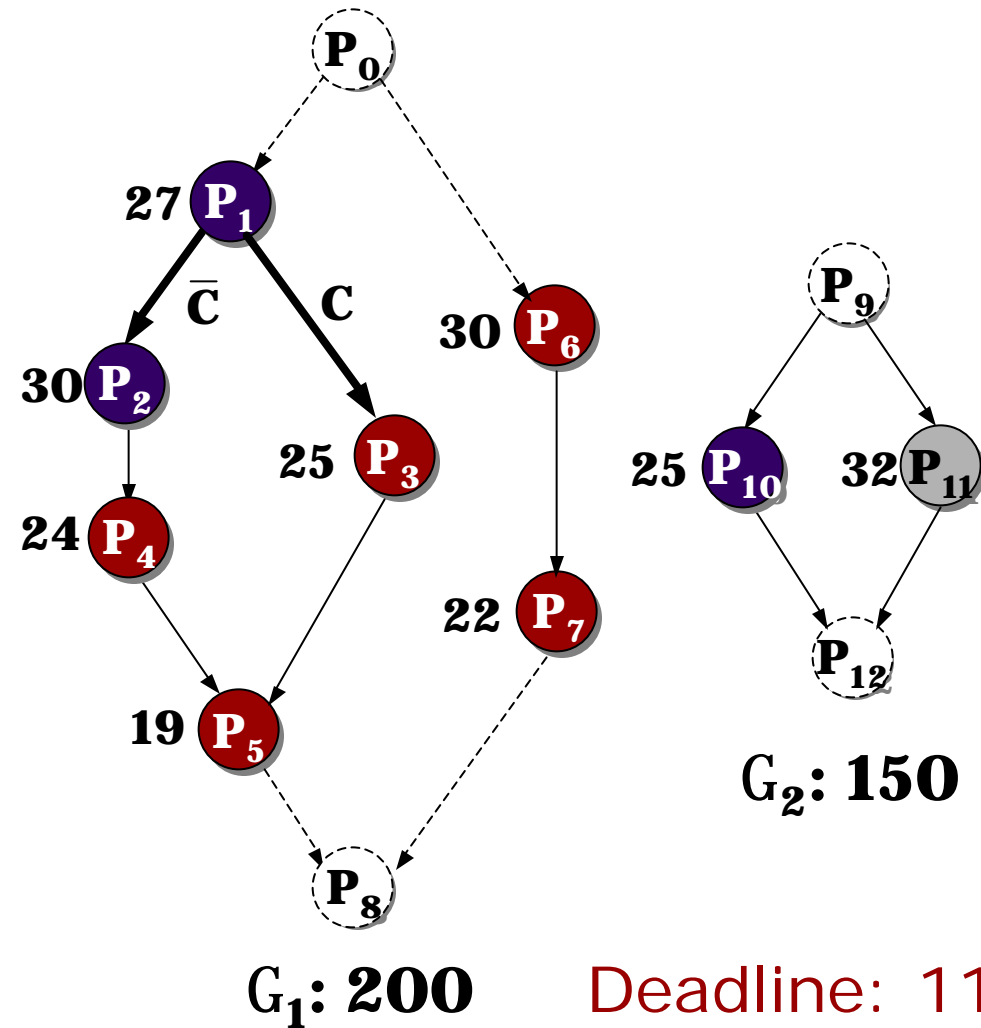
- An application modelled using **conditional process graphs** (CPG).
- Each CPG in the application has its own independent *period*.
- Each process has a *worst case execution time*, a *deadline*, and a *priority*.
- The system architecture and mapping of processes are given.

## Output

- Worst case response times for each process.

**Performance estimation for systems modelled using conditional process graphs.**

# Example



CPG	Worst Case Delays	
	No conditions	Conditions
$G_1$	120	100
$G_2$	82	82

# Task Graphs with Data Dependencies



- **K. Tindell: Adding Time-Offsets to Schedulability Analysis, Research Report**  
Offset: fixed interval in time between the arrival of sets of tasks.  
Can reduce the pessimism of the schedulability analysis.  
Drawback: how to derive the offsets?
  
- **T. Yen, W. Wolf: Performance Estimation for Real-Time Distributed Embedded Systems, IEEE Transactions On Parallel and Distributed Systems**  
Phase (similar concept to offsets).  
Advantage: gives a framework to derive the phases.

# Schedulability Analysis for Task Graphs

```
DelayEstimate(task graph G, system S)
```

```
  for each pair  $(P_i, P_j)$  in G
```

```
    maxsep $[P_i, P_j] = \infty$ 
```

```
  end for
```

```
  step = 0
```

```
  repeat
```

```
    LatestTimes(G)
```

```
    EarliestTimes(G)
```

```
    for each  $P_i \in G$ 
```

```
      MaxSeparations( $P_i$ )
```

```
    end for
```

```
  until maxsep is not changed or step < limit
```

```
  return the worst case delay  $d_G$  of the graph G
```

```
end DelayEstimate
```

worst case response times and upper bounds for the offsets

lower bounds for the offsets

maximum separation:  
 $\text{maxsep}[P_i, P_j] = 0$  if the execution of the two processes never overlaps



# Schedulability Analysis for CPGs, 1



## Two extreme solutions:

- Ignoring Conditions (IC)

Ignore control dependencies and apply the schedulability analysis for the (unconditional) task graphs.

- Brute Force Algorithm (BF)

Apply the schedulability analysis after each of the CPGs in the application have been decomposed in their constituent unconditional subgraphs.

# Schedulability Analysis for CPGs, 2

## In between solutions:

### ■ Conditions Separation (CS)

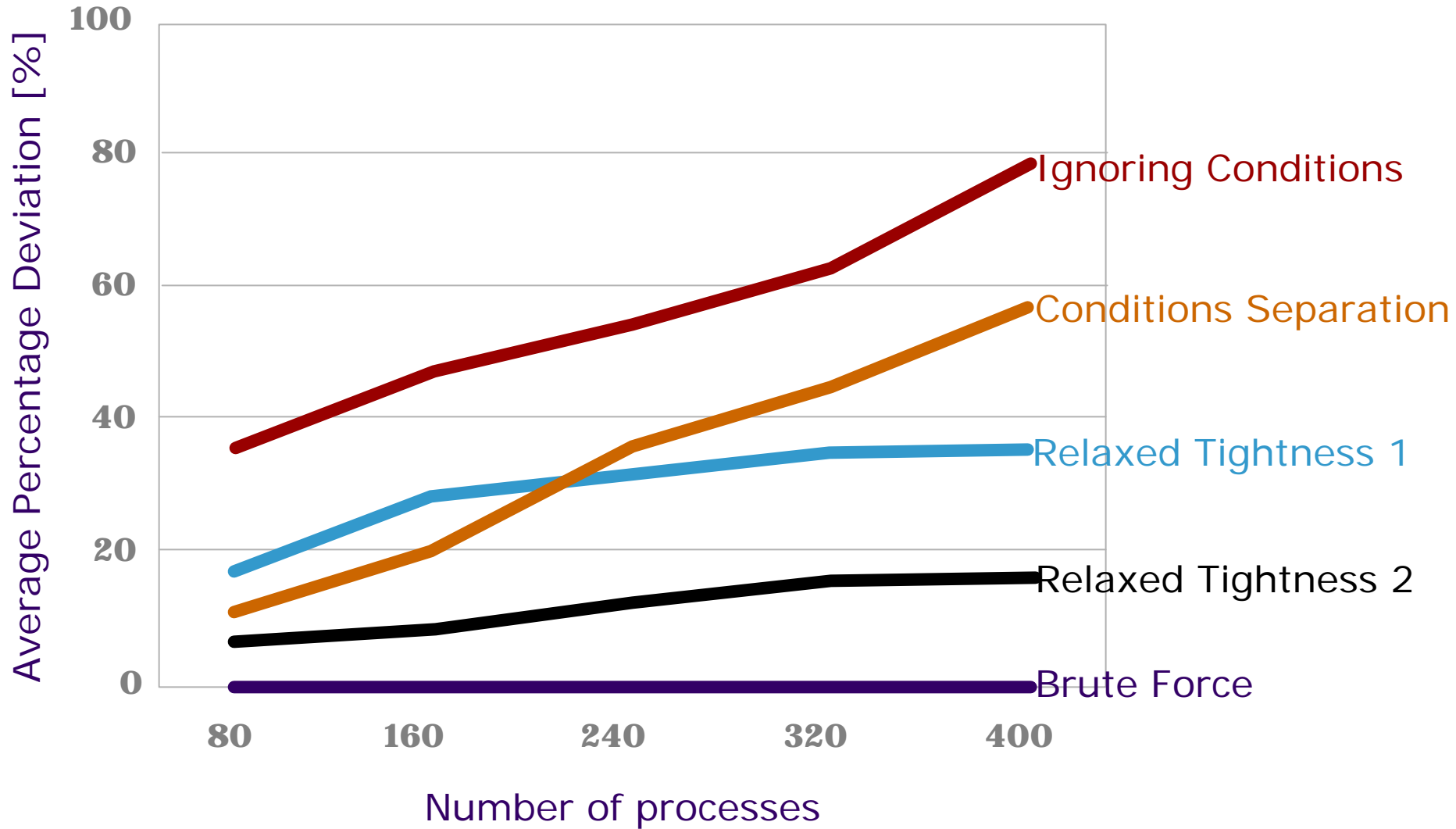
Similar to *Ignoring Conditions* but uses the knowledge about the conditions in order to update the **maxsep** table:

$\text{maxsep}[P_i, P_j] = 0$  if  $P_i$  and  $P_j$  are on different conditional paths.

### ■ Relaxed Tightness Analysis (two variants: RT1, RT2)

Similar to the *Brute Force Algorithm*, but tries to reduce the execution time by removing the iterative tightening loop (relaxed tightness) in the **DelayEstimation** function.

# Experimental Results



- **Performance estimation** for hard real-time systems with control and data dependencies.
- Modelling using **conditional process graphs** that capture both the flow of data and that of control.
- Heterogeneous architectures, **fixed priority scheduling**.
- Five approaches to the schedulability analysis of such systems.
- Extensive experiments and a real-life example show that:  
The pessimism of the analysis can be significantly reduced by considering the conditions during the analysis.