# **Database Technology**

# Topic 12: Query Processing and Optimization

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### **General Concepts**



The purpose of query processing is to ...

- (1) ... prepare the database such that we can retrieve data from it
- (2) ... produce a logical plan for a given query
- (3) ... produce a physical plan for a given query
- (4) ... produce the result of a given query



Only one of the following statements is correct. Which one?

- (1) Query validation ensures that the cost of the given query is not too high according to the cost model.
- (2) Parsing converts an SQL query string into a program that produces the result of the query.
- (3) There are some physical operators without any counterpart among the logical operators.
- (4) For every query there exists exactly one logical plan.



#### **Logical Plans and Logical Optimization**



- Consider the following SQL query
  - **SELECT** Student.PN, Grade.CourseCode, Grade.Grade **FROM** Grade **JOIN** Student **ON** Grade.StPN = Student.PN **WHERE** Student.IsInternational = TRUE;
- Here are two logical plans for this query, which are semantically





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Which of these two logical plans is likely more efficient?
(1) the plan on the left
(2) the plan on the right



# Let's Calculate

- Consider the following SQL query
  - **SELECT** Student.PN, Grade.CourseCode, Grade.Grade **FROM** Grade **JOIN** Student **ON** Grade.StPN = Student.PN **WHERE** Student.IsInternational = TRUE;
- Here are two logical plans for this query, which are semantically



• Assume 10,000 students, each of them has a grade in about 10 courses (i.e., ca 100,000 rows in Grade); 1,000 are international



### Physical Operators and Physical Query Optimization



- Suppose we want to join two relations, *Student* and *Grade* 
  - the file for *Student* consists of 100 pages
  - the file for *Grade* consists of 1,000 pages
- Assume we use the nested loops join (basic version) where the outer loop iterates over the *Student* relation
- What is the I/O cost in terms of page reads?
  - (1) 1,100
  - (2) 100,100
  - (3) 101,000
  - (4) 110,000



- Suppose we want to join two relations, *Student* and *Grade* 
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- Assume we use the nested loops join (basic version) where the outer loop iterates over the *Student* relation
- What is the I/O cost in terms of page reads?
  - (1) 1,100
  - (2) **100,100**
  - (3) 101,000
  - (4) 110,000
- Formula: pages(outer) + pages(outer) · pages(inner)



# **Block Nested Loops Join**

- Suppose we want to join two relations, *Student* and *Grade* 
  - the file for *Student* consists of 100 pages
  - the file for *Grade* consists of 1,000 pages
- Assume we use the nested loops join (basic version) where the outer loop iterates over the *Student* relation
- What is the I/O cost in terms of page reads?
  - (1) 1,100
  - (2) **100,100**
  - (3) 101,000
  - (4) 110,000
- Formula: pages(outer) + pages(outer) · pages(inner)
- Block-NLJ with 10 buffers for Student: 10,100 page reads



#### Exercise

- Suppose we want to join two relations, *Student* and *Grade* 
  - the file for *Student* consists of 100 pages
  - the file for *Grade* consists of 1,000 pages
- Assume we use the sort-merge join
  - external merge sort for the sorting, i.e., the I/O cost of sorting a relation is  $2 \times p \times \lceil \log_m(p) \rceil$  page reads & writes
  - assume we have only 2+1 buffers for it, i.e., m=2
- Calculate the I/O cost needed for the sort-merge join
  - ignore the cost for writing the result

 $[\log_2(100)] = 7$  $[\log_2(1,000)] = 10$ 



# Solution

- Suppose we want to join two relations, *Student* and *Grade* 
  - the file for *Student* consists of 100 pages
  - the file for *Grade* consists of 1,000 pages
- Assume we use the sort-merge join
  - external merge sort for the sorting, i.e., the I/O cost of sorting a relation is  $2 \times p \times \lceil \log_m(p) \rceil$  page reads & writes
  - assume we have only 2+1 buffers for it, i.e., m=2
- Calculate the I/O cost needed for the sort-merge join
  - sorting of *Student* relation: 1,400 page reads and writes
  - sorting of *Grade* relation: 20,000 page reads and writes
  - merge phase: 100 + 1,000 = 1,100 page reads
  - total: 22,500 page reads and writes

 $[\log_2(100)] = 7$  $[\log_2(1,000)] = 10$ 



# Let's use more buffers!

- Suppose we want to join two relations, *Student* and *Grade* 
  - the file for *Student* consists of 100 pages
  - the file for *Grade* consists of 1,000 pages
- Assume we use the sort-merge join
  - external merge sort for the sorting, i.e., the I/O cost of sorting a relation is  $2 \times p \times \lceil \log_m(p) \rceil$  page reads + writes
  - assume we have 10+1 buffers for it, i.e., m=10
- Calculate the I/O cost needed for the sort-merge join
  - sorting of *Student* relation: 400 page reads and writes
  - sorting of *Grade* relation: 6,000 page reads and writes
  - merge phase: 100 + 1,000 = 1,100 page reads
  - total: 7,500 page reads and writes

 $[\log_{10}(100)] = 2$  $[\log_{10}(1,000)] = 3$ 



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