Database Technology

Topic 12: Query Processing and Optimization

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General Concepts



The purpose of query processing is to ...

- (1) ... prepare the database such that we can retrieve data from it
- (2) ... produce a logical plan for a given query
- (3) ... produce a physical plan for a given query
- (4) ... produce the result of a given query



Only one of the following statements is correct. Which one?

- (1) Query validation ensures that the cost of the given query is not too high according to the cost model.
- (2) Parsing converts an SQL query string into program that produces the result of the query.
- (3) There are some physical operators without any counterpart among the logical operators.
- (4) For every query there exists exactly one logical plan.

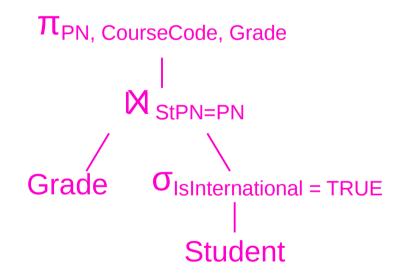


Logical Plans and Logical Optimization TODO



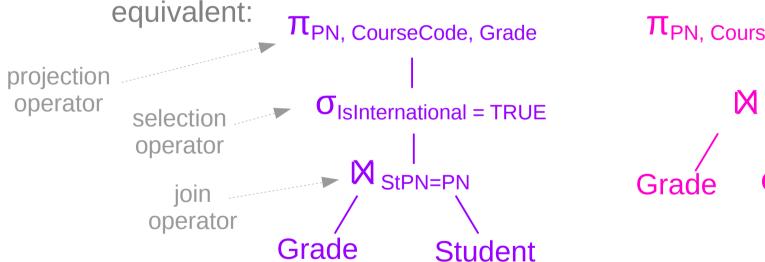
Consider the following SQL query
 SELECT Student.PN, Grade.CourseCode, Grade.Grade
 FROM Grade JOIN Student ON Grade.StPN = Student.PN
 WHERE Student.IsInternational = TRUE;

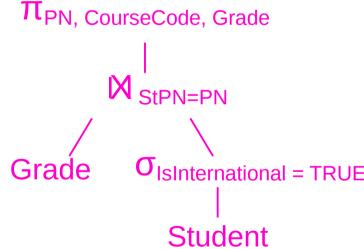
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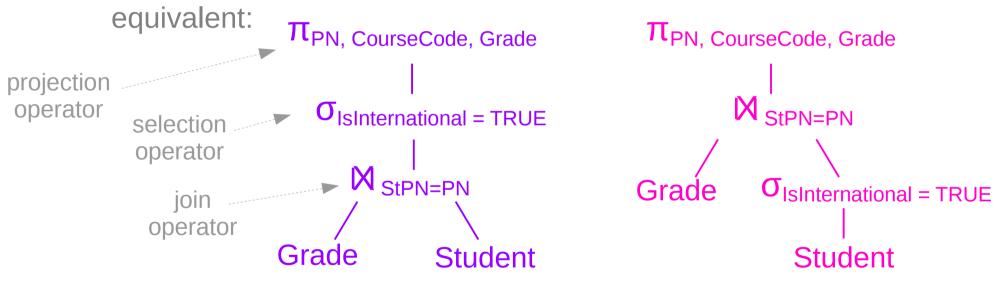
Which of these two logical plans is likely more efficient?
 (1) the plan on the left
 (2) the plan on the right



Let's Calculate

Consider the following SQL query
 SELECT Student.PN, Grade.CourseCode, Grade.Grade
 FROM Grade JOIN Student ON Grade.StPN = Student.PN
 WHERE Student.IsInternational = TRUE;

Here are two logical plans for this query, which are semantically



Assume 10,000 students, each of them has a grade in about 10 courses (i.e., ca 100,000 rows in Grade); 1,000 are international



Physical Operators and Physical Query Optimization



- Suppose we want to join two relations, Student and Grade
 - the file for Student consists of 100 pages
 - the file for Grade consists of 1,000 pages
- Assume we use the nested loops join (basic version) where Student is considered as the outer-loop relation
- What is the I/O cost in terms of page reads?
 - (1) 1,100
 - (2) 100,100
 - (3) 101,000
 - (4) 110,000



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- What is the I/O cost in terms of page reads?
 - (1) 1,100
 - (2) **100,100**
 - (3) 101,000
 - (4) 110,000
- Formula: pages(outer) + pages(outer) pages(inner)



Block Nested Loops Join

- Suppose we want to join two relations, Student and Grade
 - the file for Student consists of 100 pages
 - the file for Grade consists of 1,000 pages
- Assume we use the nested loops join (basic version)
 where Student is considered as the outer-loop relation
- What is the I/O cost in terms of page reads?
 - (1) 1,100
 - (2) **100,100**
 - (3) 101,000
 - (4) 110,000
- Formula: pages(outer) + pages(outer) pages(inner)
- Block-NLJ with 10 buffers for Student: 10,100 page reads



Exercise

- Suppose we want to join two relations, Student and Grade
 - the file for Student consists of 100 pages
 - the file for Grade consists of 1,000 pages
- Assume we use the sort-merge join
 - external merge sort for the sorting, i.e., the I/O cost of sorting a relation is $2 \times p \times \lceil \log_m(p) \rceil$ page reads + writes
 - assume we have only 2+1 buffers for it, i.e., m=2
- Calculate the I/O cost needed for the sort-merge join
 - ignore the cost for writing the result

$$\lceil \log_2(100) \rceil = 7$$

 $\lceil \log_2(1,000) \rceil = 10$



Solution

- Suppose we want to join two relations, Student and Grade
 - the file for Student consists of 100 pages
 - the file for Grade consists of 1,000 pages
- Assume we use the sort-merge join
 - external merge sort for the sorting, i.e., the I/O cost of sorting a relation is $2 \times p \times \lceil \log_m(p) \rceil$ page reads + writes
 - assume we have only 2+1 buffers for it, i.e., m=2
- Calculate the I/O cost needed for the sort-merge join
 - sorting of Student relation: 1,400 page reads and writes
 - sorting of Grade relation: 20,000 page reads and writes
 - merge phase: 100 + 1,000 = 1,100 page reads
 - total: 22,500 page reads and writes

$$\lceil \log_2(100) \rceil = 7$$

 $\lceil \log_2(1,000) \rceil = 10$



Let's use more buffers!

- Suppose we want to join two relations, Student and Grade
 - the file for Student consists of 100 pages
 - the file for Grade consists of 1,000 pages
- Assume we use the sort-merge join
 - external merge sort for the sorting, i.e., the I/O cost of sorting a relation is $2 \times p \times \lceil \log_m(p) \rceil$ page reads + writes
 - assume we have 10+1 buffers for it, i.e., m=10
- Calculate the I/O cost needed for the sort-merge join
 - sorting of Student relation: 400 page reads and writes
 - sorting of Grade relation: 6,000 page reads and writes
 - merge phase: 100 + 1,000 = 1,100 page reads
 - total: 7,500 page reads and writes

$$\lceil \log_{10}(100) \rceil = 2$$

 $\lceil \log_{10}(1,000) \rceil = 3$



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