

Database Technology

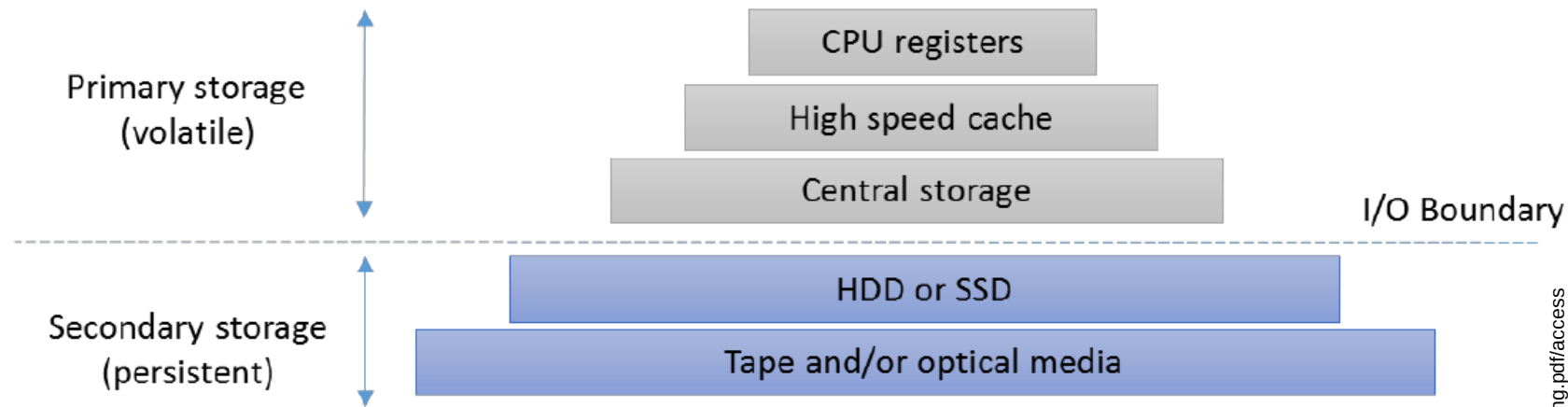
Topic 8: Data Structures for Databases

Olaf Hartig

olaf.hartig@liu.se

Storage Hierarchy

Quiz



Which of the following statements *is correct*?

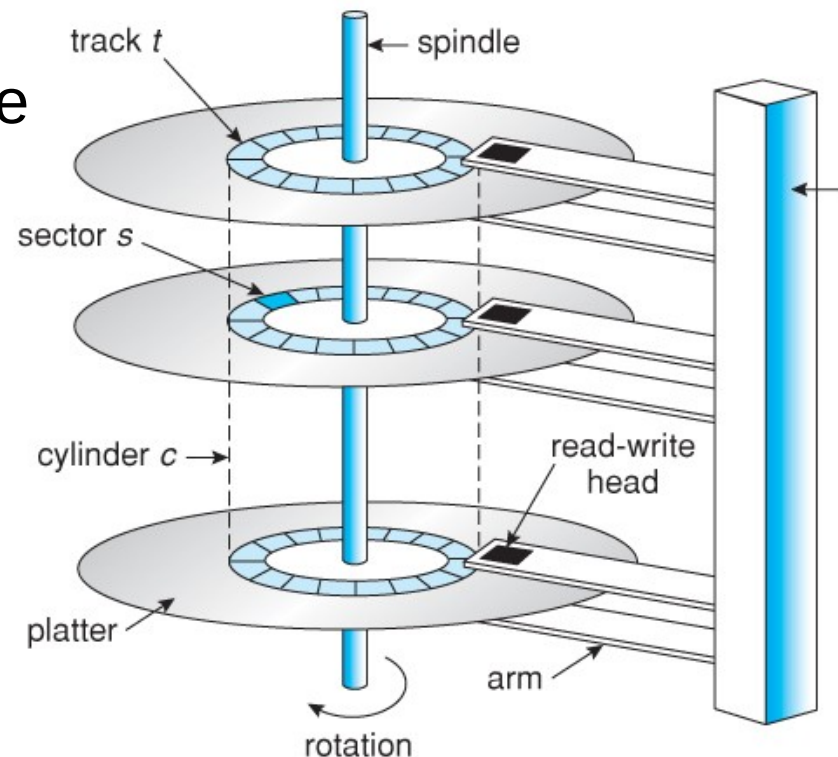
- 1) Secondary storage devices are usually faster than primary storage devices.
- 2) Data in a primary storage device may be lost when switching of the power.
- 3) The CPU may operate directly on data that is in a secondary storage device.
- 4) A piece of data (e.g., a record) may not be held both in a primary storage device and in a secondary storage device at the same time.

Storage Hardware

Quiz

Which of the following statements on an HDD is *not* correct?

- 1) The platters containing magnetic particles are secured on a spindle that rotates at a constant speed.
- 2) An HDD needs three-dimensional movements in order to access all of its data.
- 3) All data on the same cylinder can be read without moving the actuator.
- 4) The tracks represent concentric circles of magnetic particles; each track consists of individual sectors.



Record Organization

(Organizing Fields / Data Items in Records)

Quiz

T1

A1	A2	A3
alice	3	100
bob	5	23

T2

A1	A2	A3
alice	NULL	41
bob	NULL	NULL

Assume we have two tables, T1 and T2, such that the rows in T1 do not contain NULL values whereas rows in T2 may contain several NULL values.

Each table should be stored in a separate physical file.

Which *record organization technique* should we choose for these files if we want to minimize storage space efficiently?

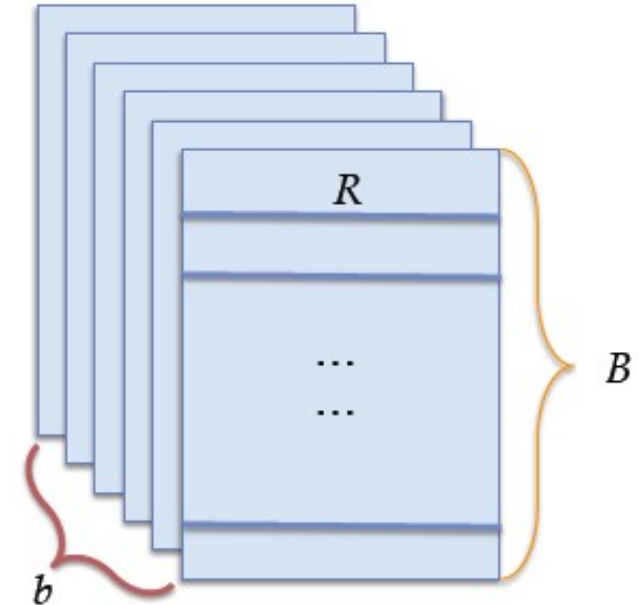
- 1) *Embedded identification* for the file of T1 and *relative location* for the file of T2
- 2) *Embedded identification* for the file of T2 and *relative location* for the file of T1
- 3) *Embedded identification* for both files.
- 4) *Relative location* for both files.

Record Allocation

(Allocating Record to File Blocks)

Quiz

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,

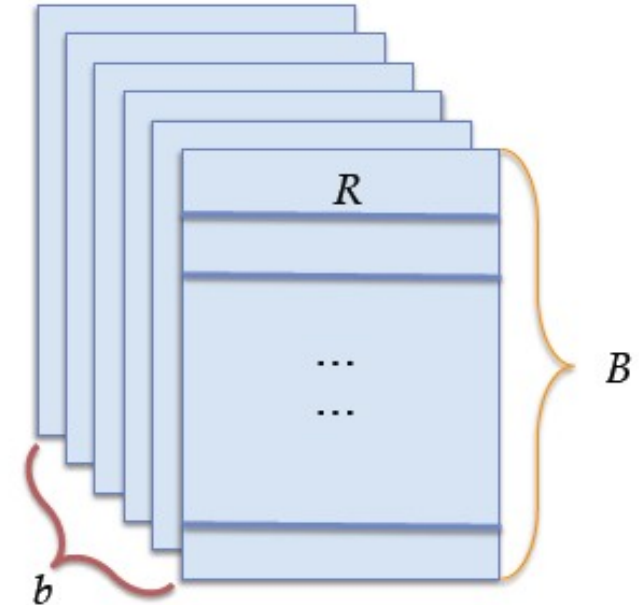


- How many blocks are needed to store the file?

1) $b = 10$ 2) $b = 20$ 3) $b = 100$ 4) $b = 200$

Quiz

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,



- How many blocks are needed to store the file?

1) $b = 10$ 2) $b = 20$ 3) $b = 100$ 4) $b = 200$

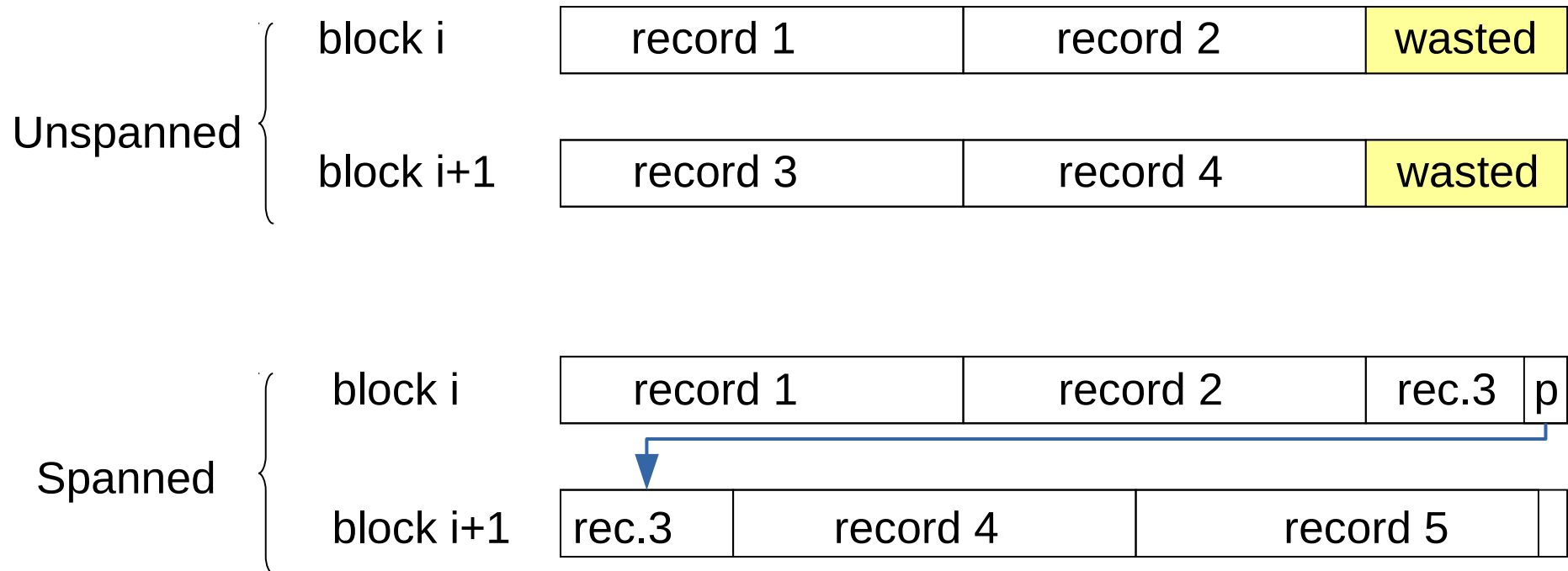
$$bfr = \left\lceil \frac{B}{R} \right\rceil \quad b = \left\lceil \frac{r}{bfr} \right\rceil$$

- Space wasted per block = $B - bfr * R$

blocking factor

Spanned Records

... avoid wasting space

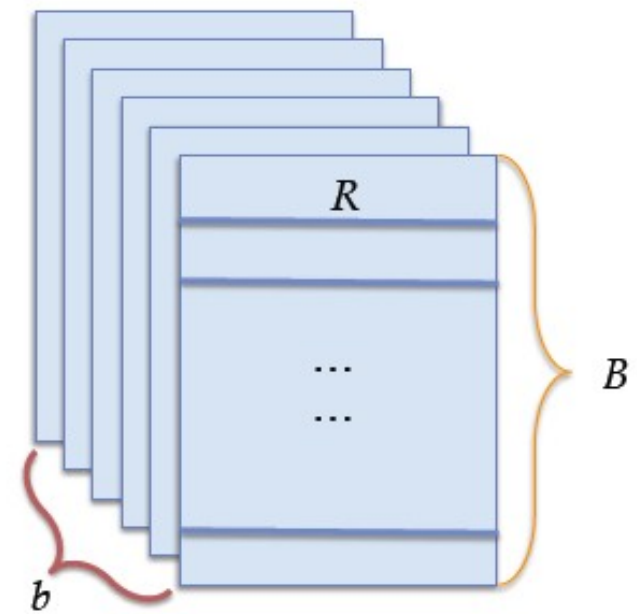


File Organization

(Organizing Records in Files)

Exercise: Heap File

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **heap file**
 - i.e., new records are always appended to the end of the file
- How many blocks do we need to read?



Name	ID	Salary
Andersson	12	2000
Svensson	13	4000
...		
...		

Block 1

..		
...		
...		
...		

Block 2

..		
...		
...		
...		

Block 3

⋮

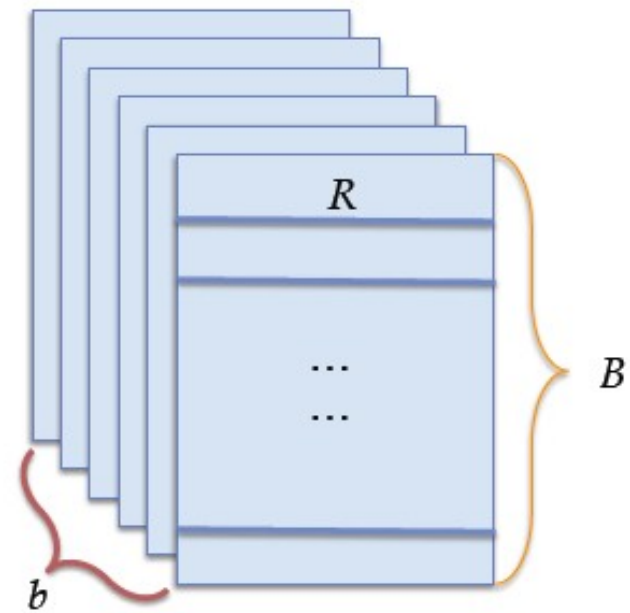
⋮

⋮

	search field = ID value = 43 (unique)	search field = Name value = Smith (non-unique)
worst case		
best case		
average case		

Exercise: Heap File

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **heap file**
 - i.e., new records are always appended to the end of the file
- How many blocks do we need to read?



Name	ID	Salary
Andersson	12	2000
Svensson	13	4000
...		
...		

Block 1

..		
...		
...		
...		

Block 2

..		
...		
...		
...		

Block 3

⋮

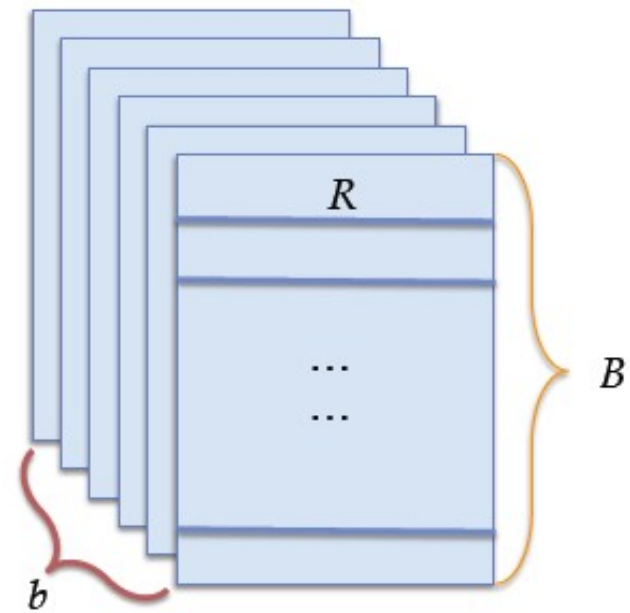
⋮

	search field = ID value = 43 (unique)	search field = Name value = Smith (non-unique)
worst case	200	
best case	1	
average case	100	

$$\left\lceil \frac{b}{2} \right\rceil$$

Exercise: Heap File

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **heap file**
 - i.e., new records are always appended to the end of the file
- How many blocks do we need to read?



Name	ID	Salary
Andersson	12	2000
Svensson	13	4000
...		
...		
..		
...		
...		
...		
..		
...		
...		
...		

Block 1

Block 2

Block 3

⋮

	search field = ID value = 43 (unique)	search field = Name value = Smith (non-unique)
worst case	200	200
best case	1	200
average case	100	200

linear search
until last block

$$\left\lceil \frac{b}{2} \right\rceil$$

Exercise: Sorted File (a.k.a. Sequential File)

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **sorted file** by using the ID field as the *sorting field*
 - i.e., records inserted based on their ID value
- How many blocks do we need to read?

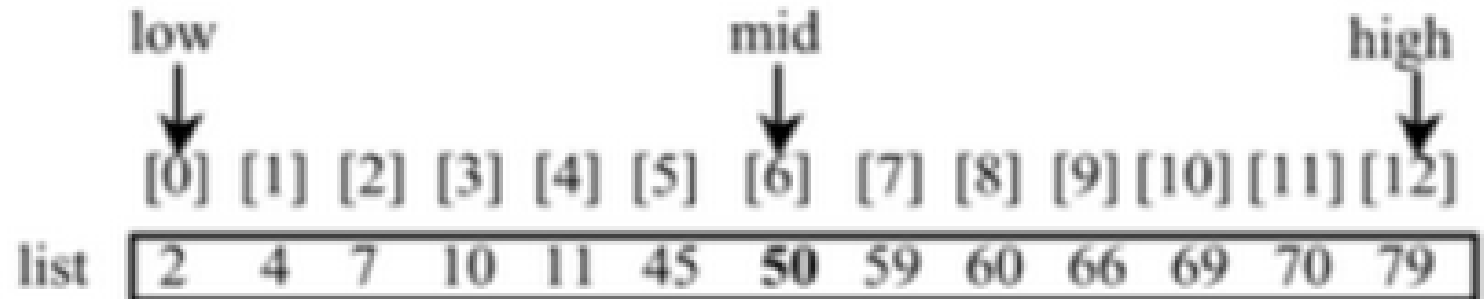
Name	ID	Salary	
Andersson	12	2000	} Block 1
Svensson	13	4000	
...			
...			
..			} Block 2
...			
...			
...			
..			} Block 3
...			
...			
...			
⋮			⋮
⋮			⋮
⋮			⋮

	search field = ID value = 43 (unique)	search field = Name value = Smith (non-unique)
worst case		
best case		
average case		

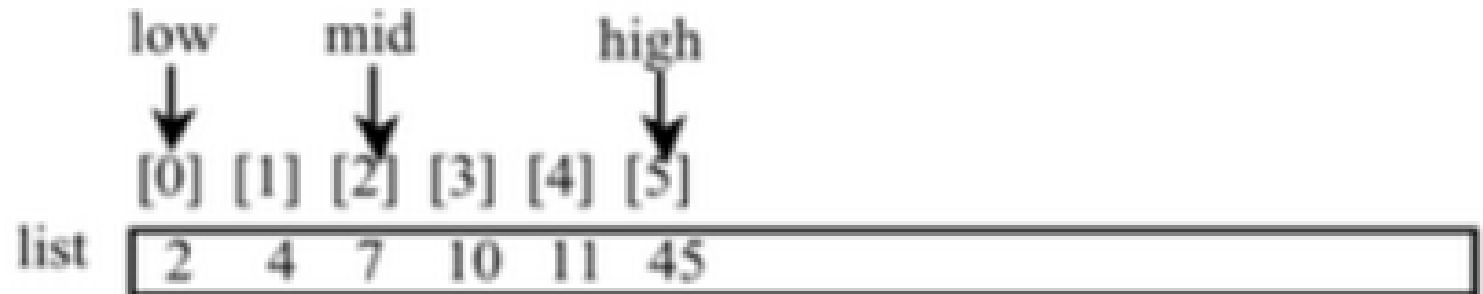
Binary Search

key is 11

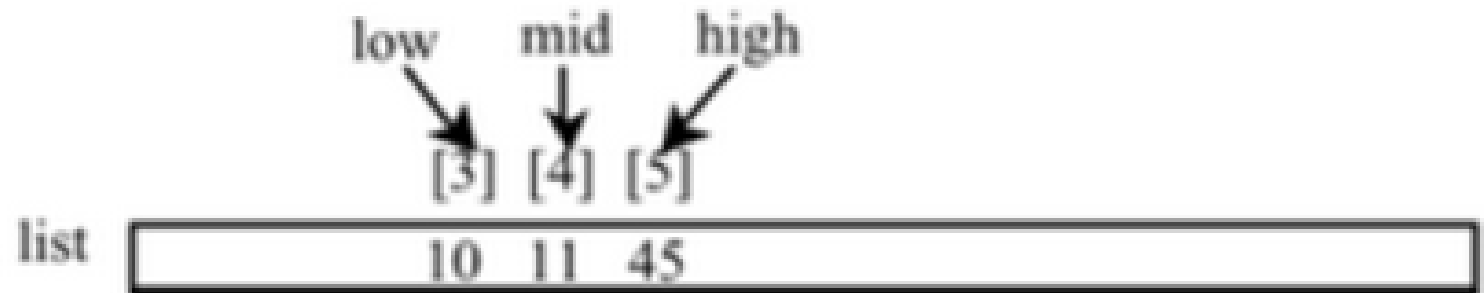
key < 50



key > 7

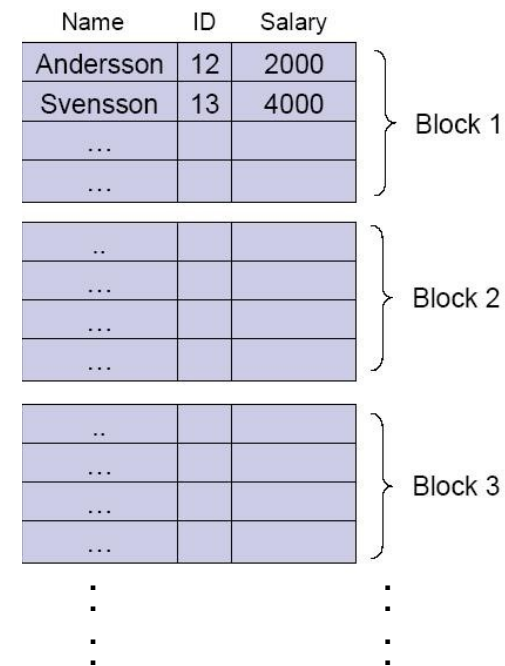


key == 11



Exercise: Sorted File (a.k.a. Sequential File)

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **sorted file** by using the ID field as the *sorting field*
 - i.e., records inserted based on their ID value
- How many blocks do we need to read?



	search field = ID value = 43 (unique)	search field = Name value = Smith (non-unique)
worst case	8	200
best case	1	200
average case	8	200

$$\log_2(128) = 7$$

$$\log_2(256) = 8$$

$$\lceil \log_2 b \rceil$$

Exercise: Hash File (a.k.a. Random File Orga.)

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **hash file** by using the ID field as the *hash field* and 120 buckets with 2 blocks per bucket
- How many blocks do we need to read?*

Name	ID	Salary
Andersson	12	2000
Svensson	13	4000
...		
...		

} Block 1

..		
...		
...		
...		

} Block 2

..		
...		
...		
...		

} Block 3

*assuming there are no collisions

	search field = ID value = 43 (unique)	search field = Name value = Smith (non-unique)
worst case		
best case		
average case		

Exercise: Hash File (a.k.a. Random File Orga.)

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **hash file** by using the ID field as the *hash field* and 120 buckets with 2 blocks per bucket
- How many blocks do we need to read?*

Name	ID	Salary
Andersson	12	2000
Svensson	13	4000
...		
...		

Block 1

..		
...		
...		
...		

Block 2

..		
...		
...		
...		

Block 3

*assuming there are no collisions

	search field = ID value = 43 (unique)	search field = Name value = Smith (non-unique)
worst case	2	≥ 200
best case	1	≥ 200
average case	1.5	≥ 200

Exercise: Hash File (a.k.a. Random File Orga.)

- Assume a file with
 - $r = 2000$ records,
 - $R = 100$ bytes per record, and
 - $B = 1000$ bytes per block,
- Hence, $b = 200$ blocks needed to store the file
- Assume we organize the file as a **hash file** by using the ID field as the *hash field* and 120 buckets with 2 blocks per bucket
- What if we want to retrieve all records with an ID value and smaller than 10?*

Name	ID	Salary
Andersson	12	2000
Svensson	13	4000
...		
...		

Block 1

..		
...		
...		
...		

Block 2

..		
...		
...		
...		

Block 3

*assuming there are no collisions

...and IDs cannot be smaller than 1

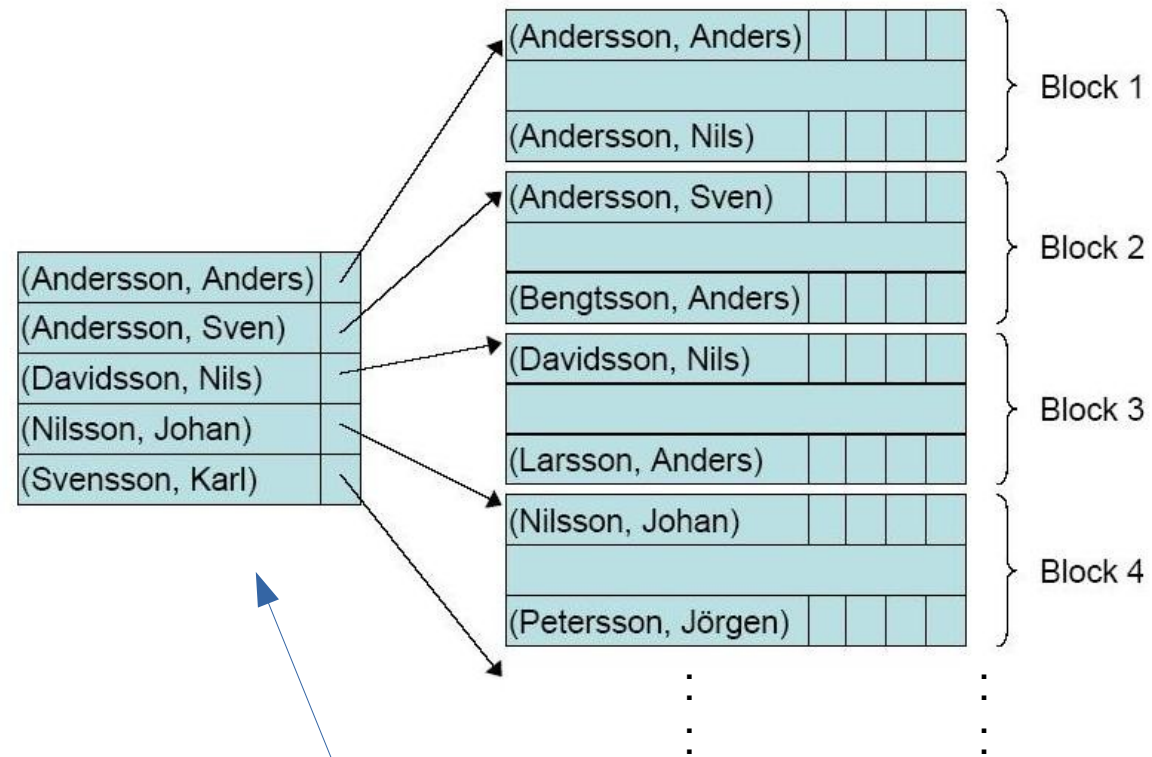
	search field = ID value = 43 (unique)
worst case	$9 \cdot 2 = 18$
best case	1
average case	depends

an ID value and smaller than 10?*

Index Sequential File Organization

Primary Index

- Why is it faster to find a random record via a binary search in the index rather than in the data file?



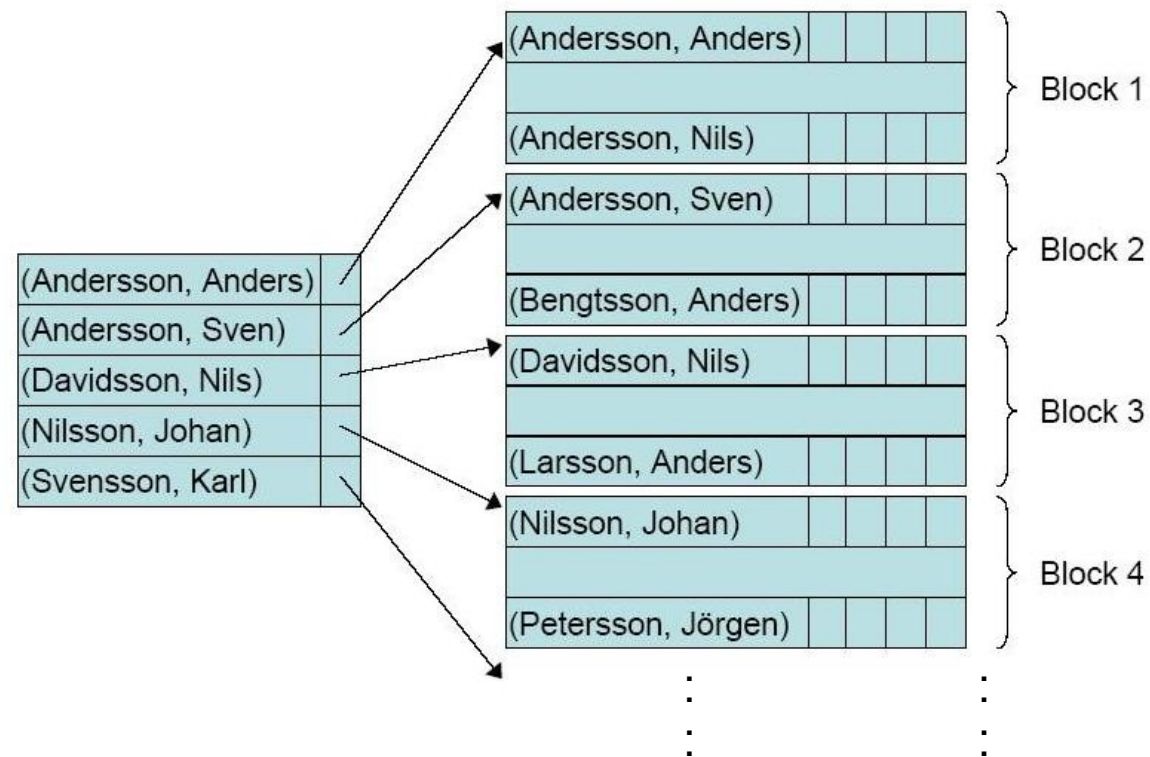
sorted file with all the records

- sorted by name, and
- name is a key

index file (sorted)
with one record per
block in the data file

Primary Index

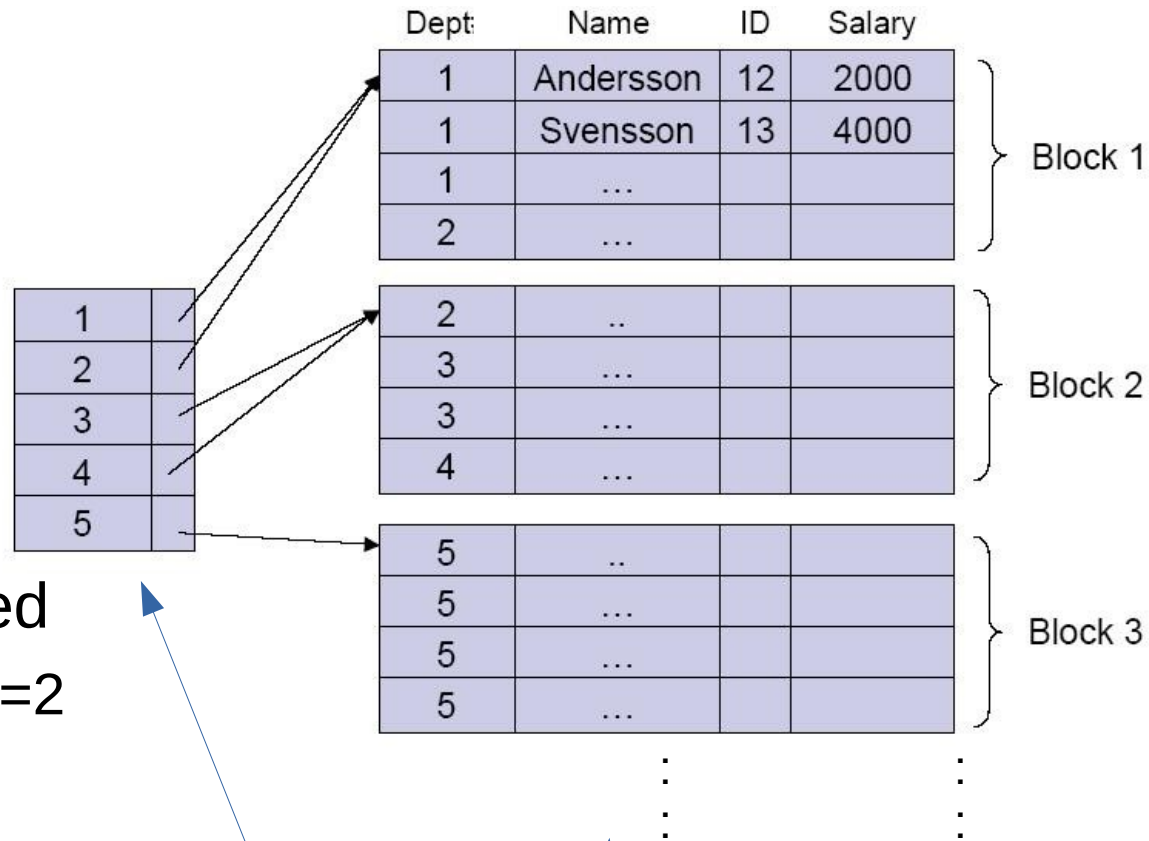
Why is it faster to find a random record via a binary search in the index rather than in the (sorted) data file?



- Index file has significantly fewer blocks because:
 - number of index records \ll number of data records
 - Index records smaller than data records (i.e., blocking factor for the index file higher than for the data file)

Clustering Index

- Attention: after binary search in the index file, multiple data file blocks may need to be accessed
 - see, for instance, Dept=2



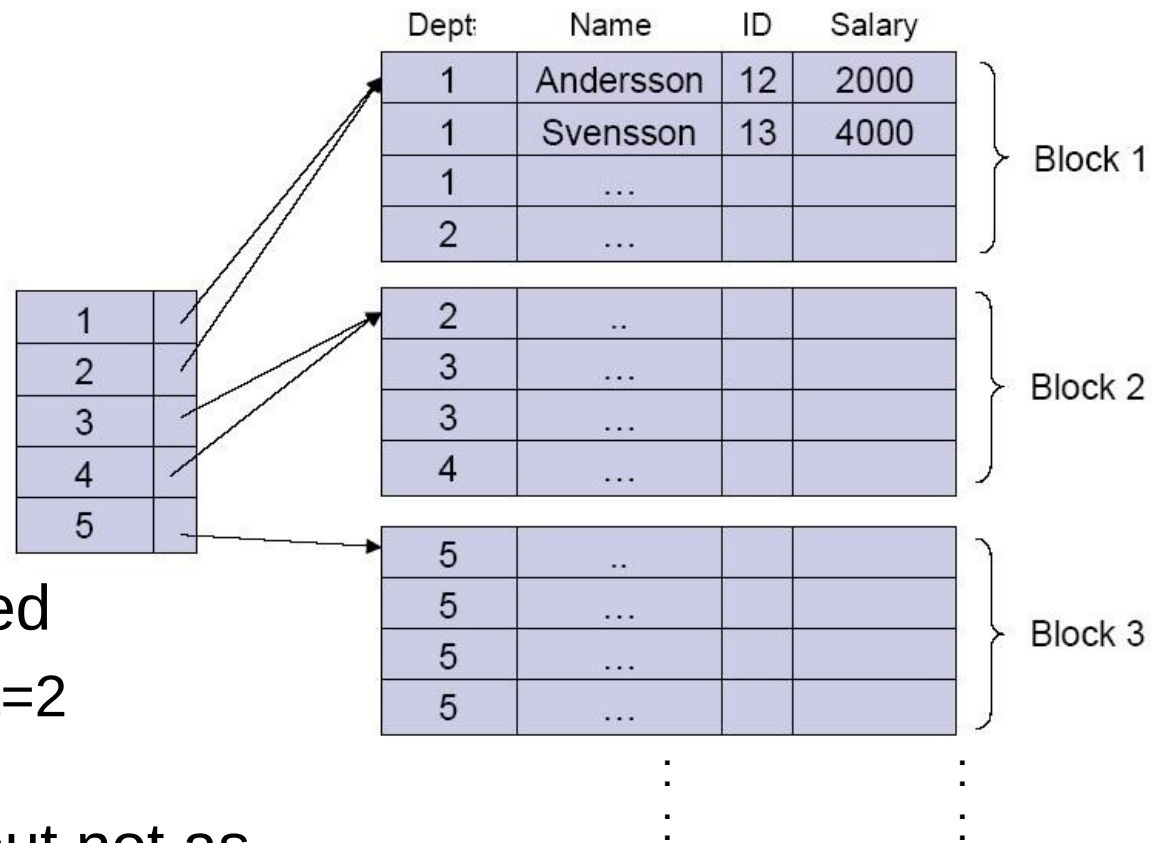
sorted file with all the records

- sorted by Dept
- Dept is *not* a key

index file (sorted)
with one record per
possible Dept value

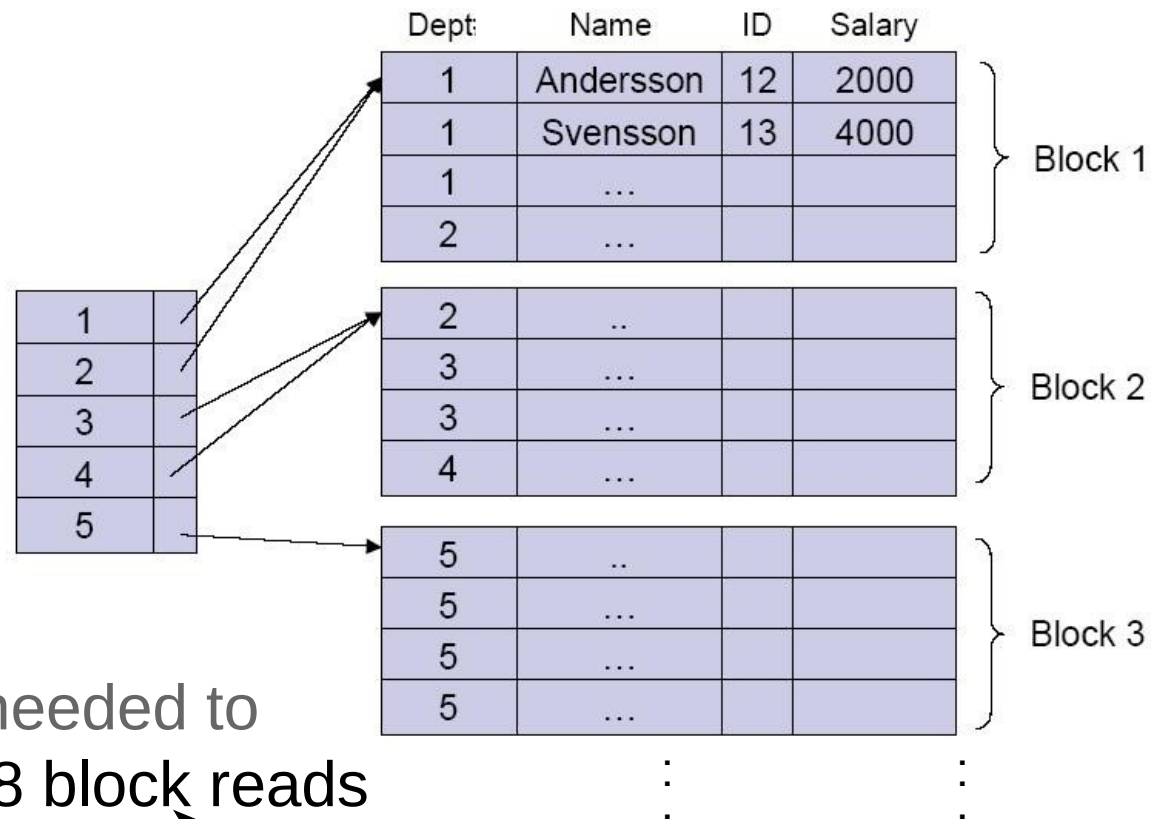
Clustering Index

- Attention: after binary search in the index file, multiple data file blocks may need to be accessed
 - see, for instance, Dept=2



- Index file also smaller, but not as much as for a primary index
 - number of index records \leq number of data records
 - at least, index records smaller than data records (like in a primary index)

Quiz



- Assume sorted file with $r = 2000$ records, $R = 100$ bytes per record, $B = 1000$ bytes per block
- Hence, $b = 200$ blocks needed to store the file and, thus, 8 block reads for a binary search on the file
- Assume $r' = 300$ different *Dept* values, $R' = 10$ bytes per index record, $B = 1000$ bytes per index block

$$\lceil \log_2 b \rceil$$

$$bfr = \left\lceil \frac{B}{R} \right\rceil \quad b = \left\lceil \frac{r}{bfr} \right\rceil$$

$$\log_2(2) = 1, \log_2(4) = 2, \log_2(8) = 3, \log_2(16) = 4$$

- How many block reads for a binary search on the index?

www.liu.se