

PhD Course Linköping, 2012

# Semantic CMS: Introduction and Overview of IKS

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#### Interactive Knowledge Stack

Reference Implementation

User-centered interaction with Knowledge Objects

Presentation, Modality & Discourse Patterns

Knowledge Representation for dynamic models (Rules & Reasoning)

Knowledge Representation for static models (Schemas & Ontologies) Methodology

Distribution: Transactions & Services

Data Access: High-level DDLs, Query languages & APIs

Models of Persistence (Relational, O-O, TripleStores...)

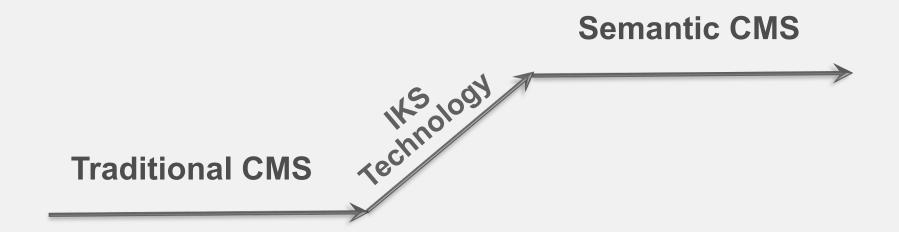
Entity Identifier Systems, Operating Systems, etc.

# IKS Goal

A Reference Architecture for Semantically Enabled Content Management Systems

# Specifications

# IKS Technology – a Path to the Semantic Level





#### What is a Semantic CMS?

#### **Traditional CMS**

VS.

#### **Semantic CMS**

- Atomic unit: Document
- Properties as meta-data
  - e.g. author
  - tags, keywords
- Keyword search for
  - strings in docs
- Document Management
  - Document types
  - Document workflow

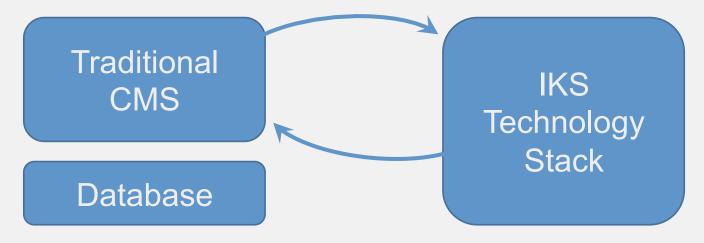
- Atomic unit: Entity
- Semantic meta-data
  - Defined entity types
  - Linked entities
- Semantic search for
  - entities and their relations
- Knowledge Management
  - Entity management
  - Ontologies



### Do Not Replace - but Extend

- No need to replace your existing technology.
- IKS components offer service oriented integration.

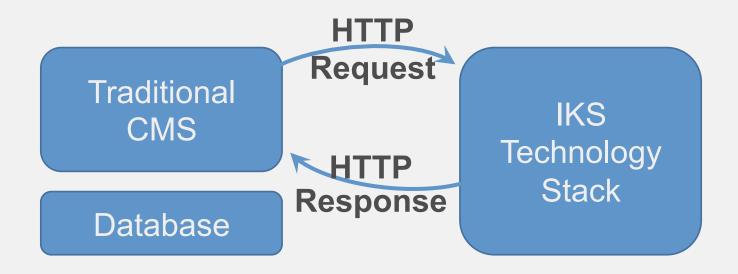
#### **Extend by Using Semantic Services**





## Rely on the Concepts of the Web

- Integration through a RESTful web service API
- Resources are identified by their URI





#### Hands on IKS 7.0

- You need: Java Runtime Environment (JRE) V1.6
- You download: <u>http://dev.iks-project.eu/downloads/iks-stack-releases/IKS-RI-7.0.zip</u>
- You execute:

java -Xmx1024m -jar iks-7.0-launcher.jar

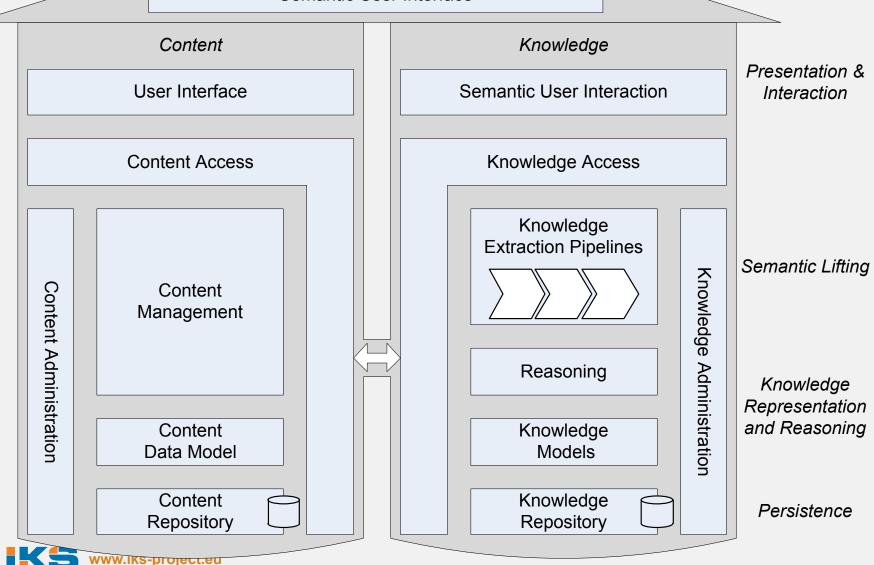
and open

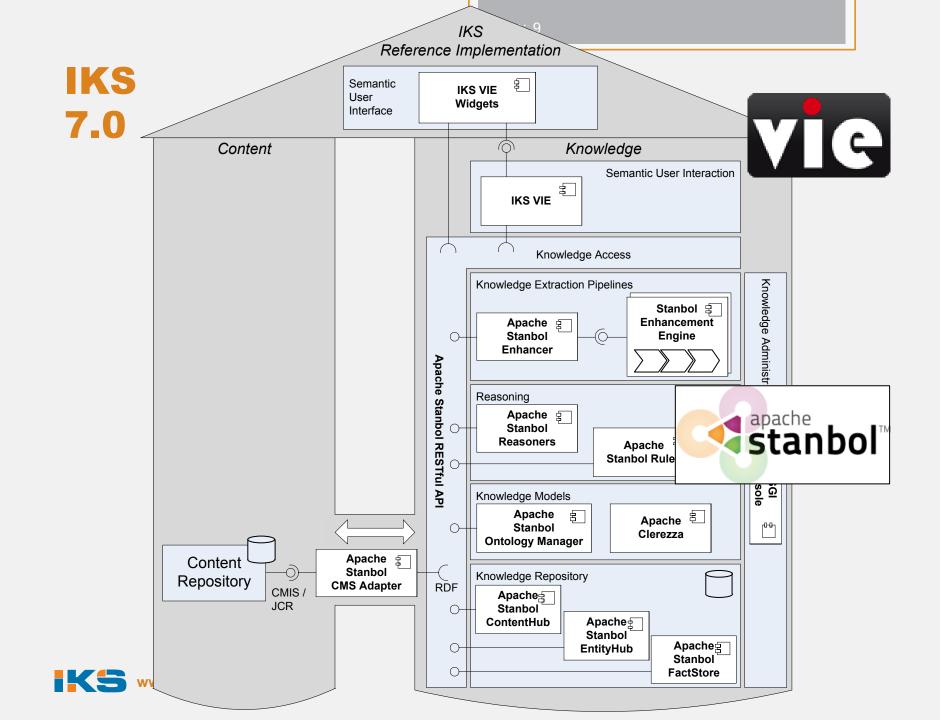
http://localhost:8080



#### SCMS -Semantic Content Management System

#### Semantic User Interface





#### **VIE Quick Facts**

- VIE is a utility library for semantic maintenance in JavaScript
- Offers semantic web developers a DSL to ease recurring tasks
  - Easy access to embedded semantic annotations in HTML (RDFa)
  - Easy loading of properties for entities from external services
  - Easy saving of knowledge about entities
  - Easy querying of semantic services
- o VIE Widgets are web user interface components based on VIE.



### **Apache Stanbol Quick Facts**

Modular (OSGi) components implemented in Java

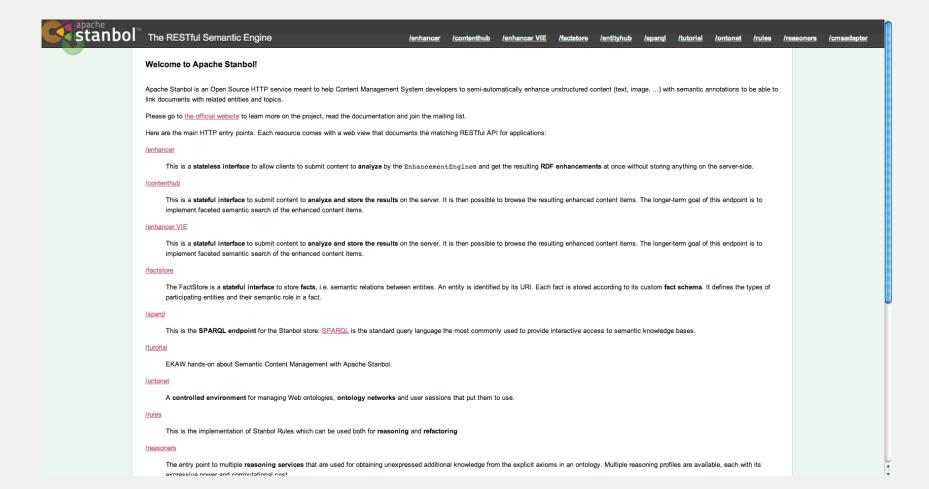
#### **Semantic Lifting**

- o Enhance content
- Link to Linked Open Data (LOD) sources
- Store and index enhanced content for search

- Manage ontologies
- Apply rules to ontologies
- Reasoning over managed ontologies

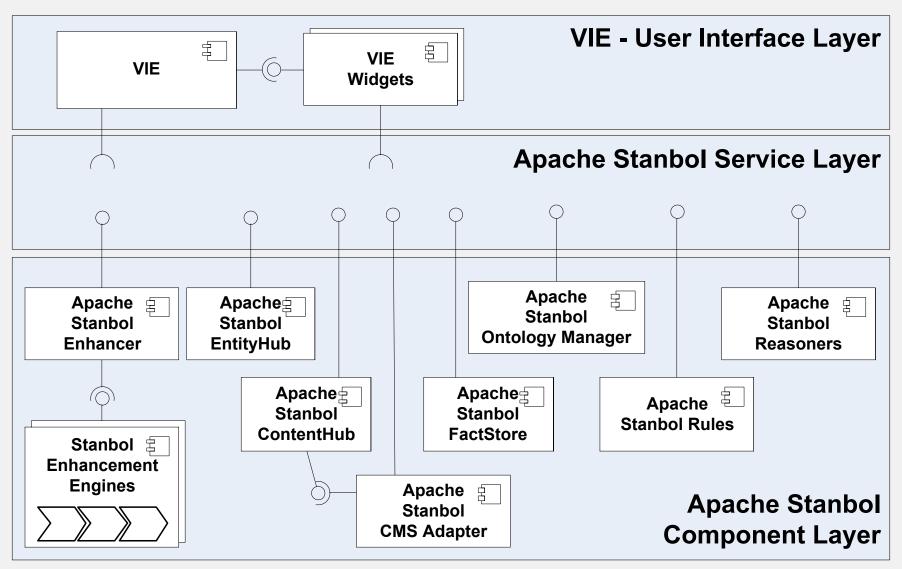


#### **Stanbol GUI**



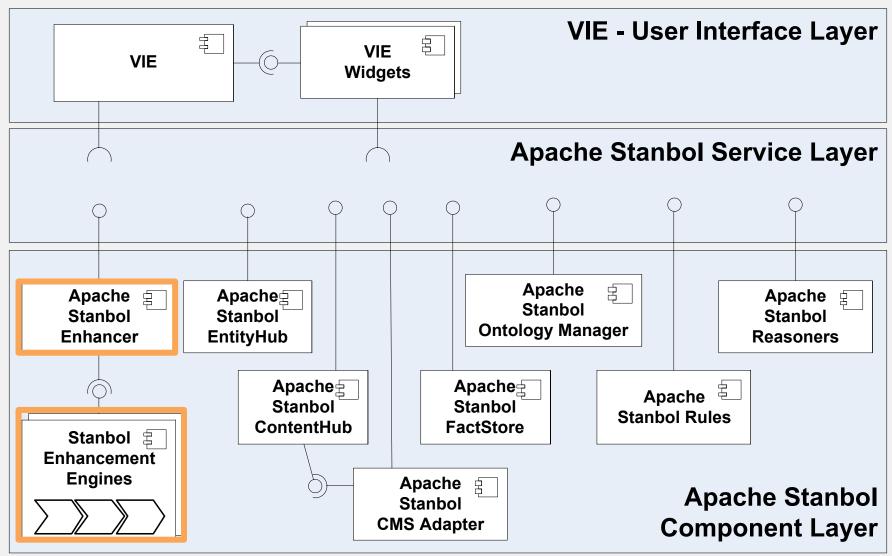


#### **Service-Oriented View**



**Semantic Lifting** 

## **Enhancer & Enhancement Engines**



**Semantic Lifting** 

# **Enhancer & Engines Features**

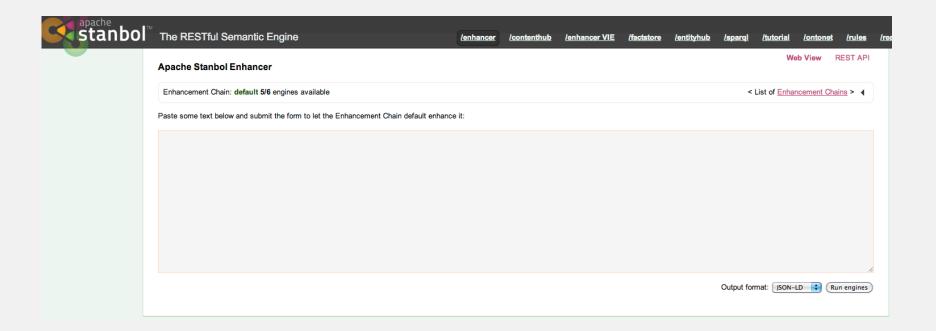


- Semantic lifting by automatically extracting entities from textual content
- Different enhancement engines for specific tasks
- Engines are arranged in customizable enhancement chains where one engine may rely on the output of another engine
- Examples
  - Language Identification Engine
  - Named Entity Extraction Engine
  - Geonames Engine to annotate places with additional information from geonames.org



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#### **Stanbol Enhancer GUI**





#### **Stanbol Enhancer**

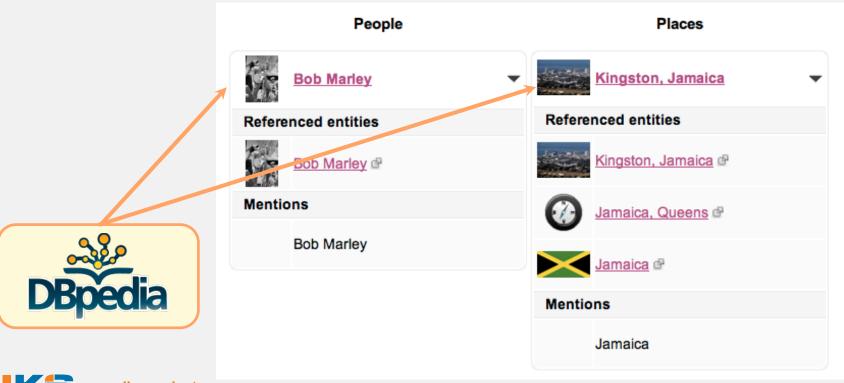
- o Input
  - "Bob Marley was a famous musician from Jamaica."
- Output





#### **Stanbol Enhancer**

- o Input
  - o "Bob Marley was a famous musician from Jamaica."
- Output



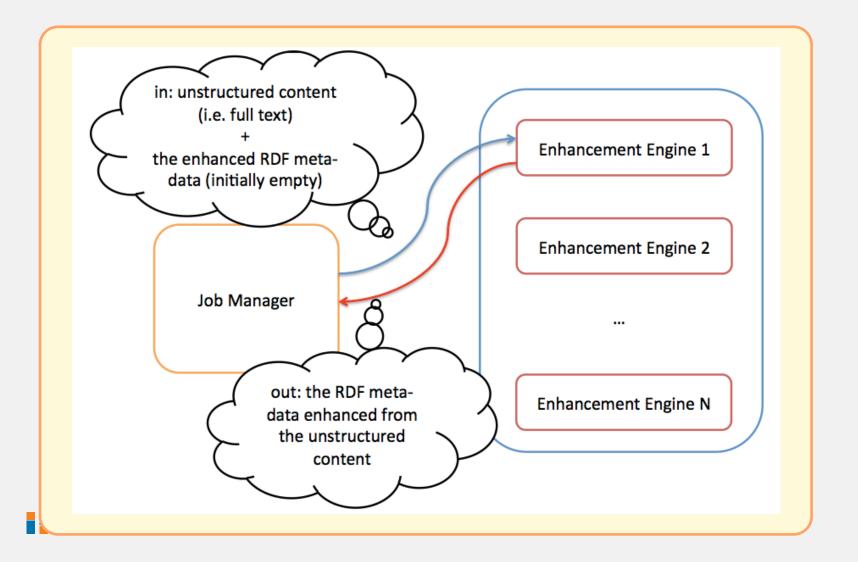


#### **Stanbol Enhancer**

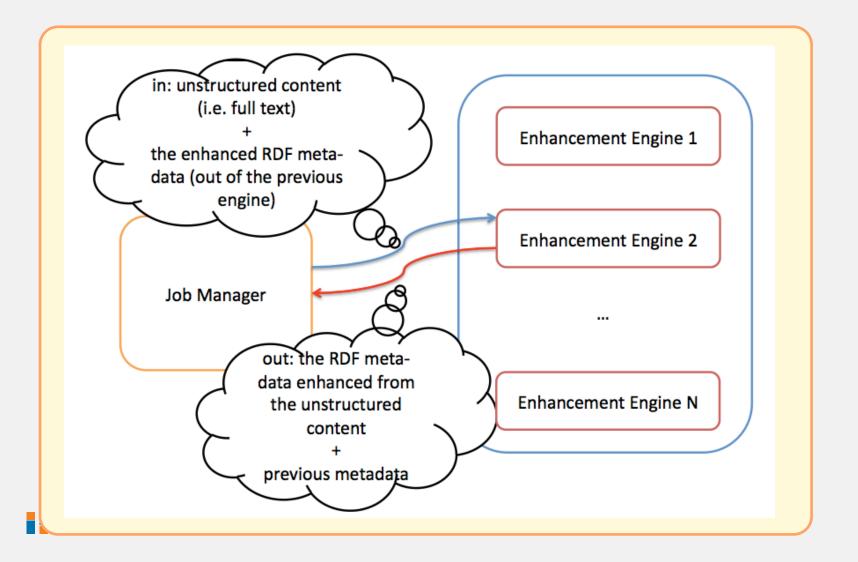
- o Input
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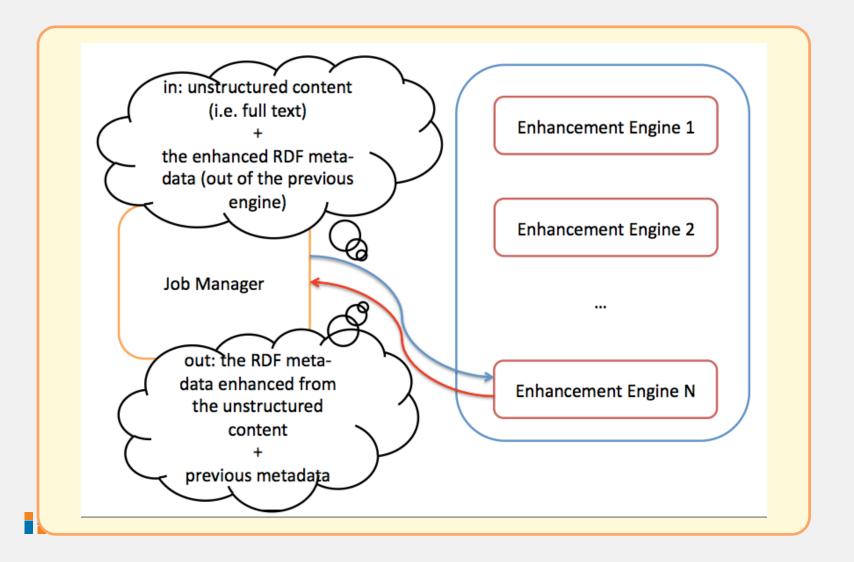
#### **How the Enhancer works**



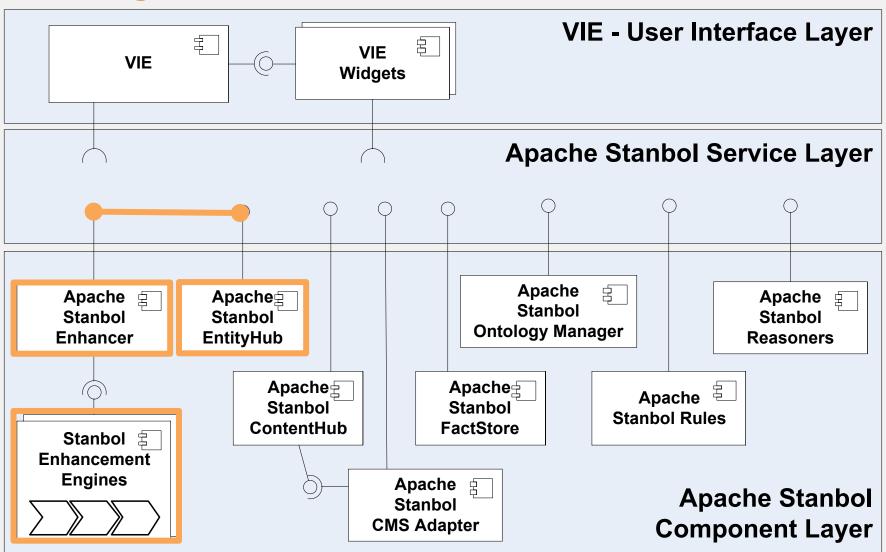
#### **How the Enhancer works**



#### **How the Enhancer works**



### **Entity Hub**



**Semantic Lifting** 

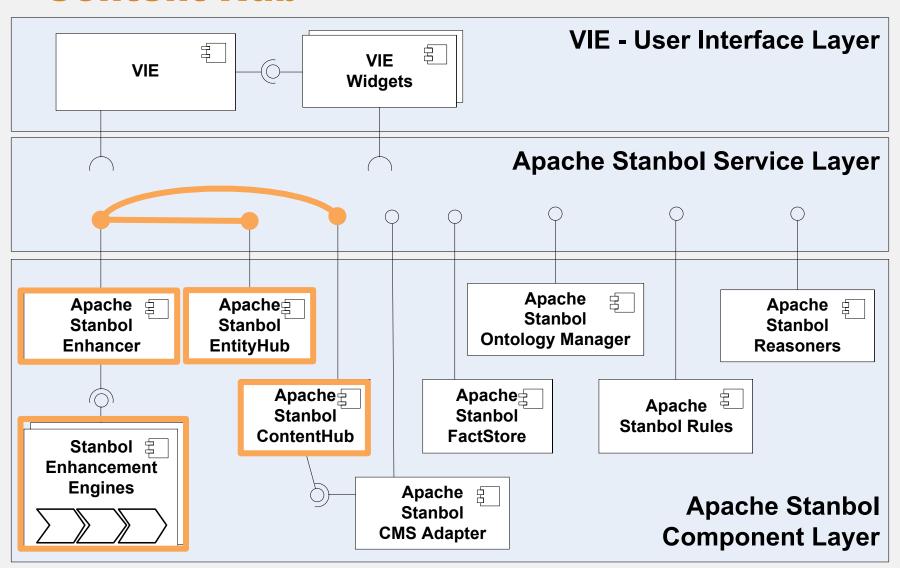
# **Entityhub Features**



- Manage a network of remote sites for fast entity lookup
- Caching of externally retrieved entity information
- CRUD management of local entities
- Examples
  - Use DBPedia linked open data source to retrieve additional information for entities
  - Use a customized vocabulary for local entities



#### **Content Hub**



**Semantic Lifting** 

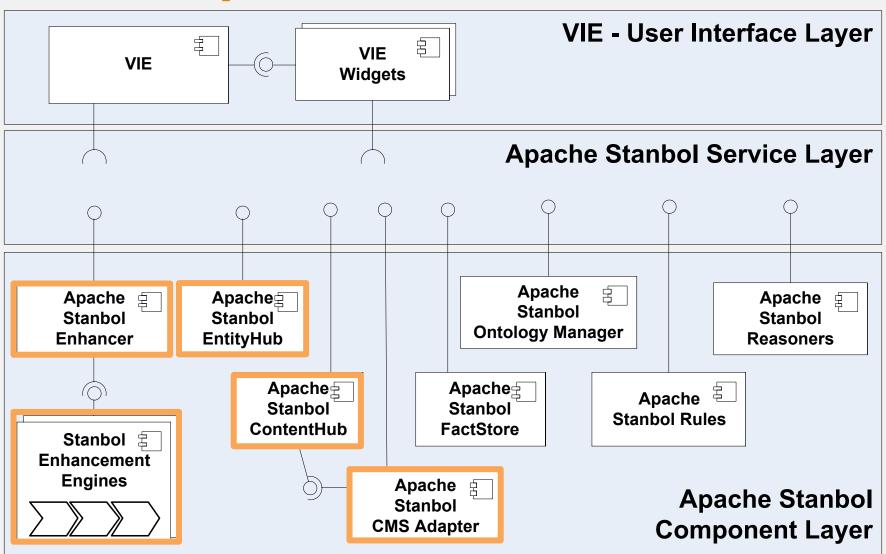
# **Contenthub Features**



- Document repository by indexing retrieved documents
- Supports indexing of additional semantic metadata provided along the content
- Search facilities
  - Keyword Search
  - Faceted Search based on available semantic metadata



### **CMS Adapter**



Semantic Lifting

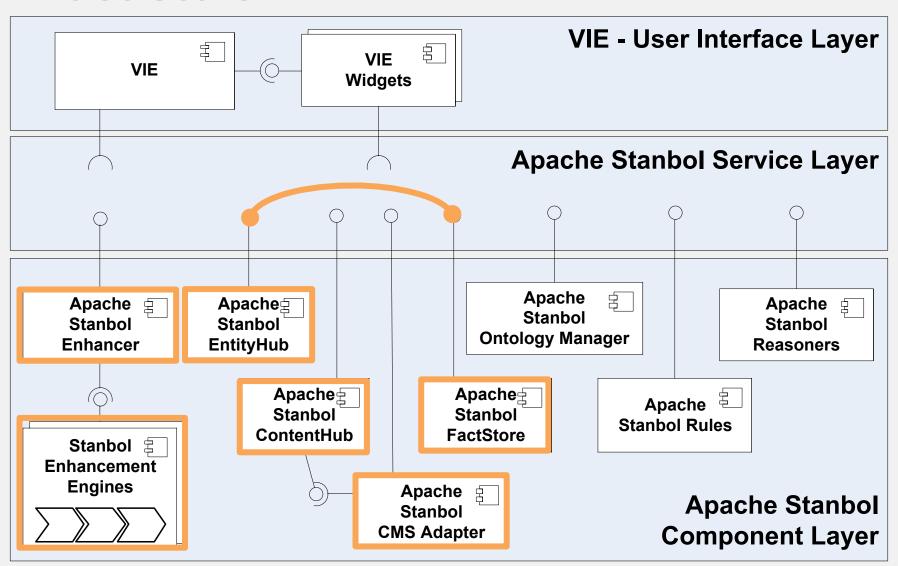
# **CMS Adapter Features**



- Bootstrapping component to import content from a CMS into Apache Stanbol
- Import content from a CMIS/JCR compliant CMS into the Apache Stanbol Contenthub



#### **Fact Store**



**Semantic Lifting** 

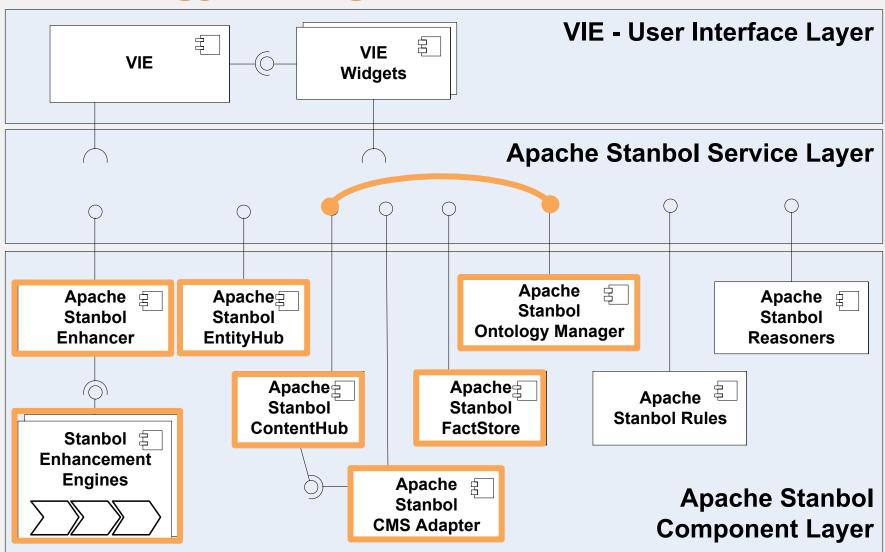
# Fact Store Features



- Simple storage for relations between entities, i.e. facts
- Definition of custom semantic relations, i.e. fact schemata
- Not limited to triples support for N-ary relations
- Simple query language for facts, no SPARQL



### **Ontology Manager**



**Semantic Lifting** 

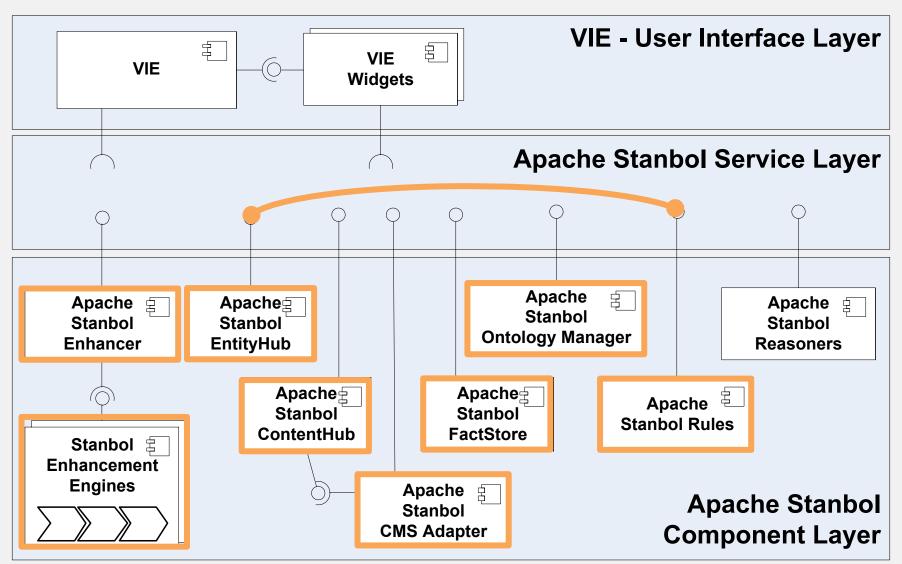
# **Ontology Manager Features**



- Controlled environment for managing ontologies
- Manage ontology networks to activate/deactivate parts of complex ontologies
- Manage user sessions for ontologies allowing local user changes



#### Rules



**Semantic Lifting** 

# Rules Features



- Construction and execution of inference rules
- Inference rules, also called transformation rules, take premises and return conclusions
- Rules can be organized in recipes which allow to execute a set of rules as a whole
- Example
  - Define rules for doing integrity checks on data fetched from heterogeneous external data sources



#### What is a rule?

- In logic, a transformation rule or rule of inference is a syntactic rule or function which takes premises and returns a conclusion
- The rule is **sound** with respect to the semantics of classical logic in the sense that if the premises are interpreted to be true then so is the conclusion.



### **Rule Examples**

- Rule pattern (modus ponens)
  - if condition then consequent
- A rule example
  - if X is a person then X has a father
     (i.e. every person has a father)
  - if Y is the father of X and Z the brother of Y then Z is the uncle of X
    - (i.e. the brother of the father is the uncle)



## **Stanbol Rule Syntax**

In Stanbol a rule is defined as

ruleName[body -> head]

#### where:

- The ruleName identifies the rule
- The body is a set of atoms that must be satisfied by evaluating the rule
- The head or consequent is a set of atoms that must be true if the condition is evaluated to be true



## **Stanbol Rule Syntax**

## ruleName[body -> head]

#### Where

- Both body and head consist of a list of conjunctive atoms
  - body = atom1 . atom2 . . . . atomN
  - head = atom1 . atom2 . ... atomM
- The conjunction ∧ in Stanbol Rules is expressed with the symbol "."



## **Example of rule**

Considering Stanbol Rules, the formula  $hasFather(x,y) \land hasBrother(y,z) \Rightarrow hasUncle(x,z)$ 

expresses in predicate calculus becomes

```
myRule[ has(<http://myont.org/hasFather>, ?x, ?y) . has(<http://myont.org/hasBrother>, ?y, ?z)
```

->

has(<http://myont.org/hasUncle>, ?x, ?z) ]



## **Rule Atoms**

- An atom is the smallest unit of the interpretation of a rule
  - e.g.: in predicate calculus

```
Person(x) \Rightarrow hasFather(x, y)
```

Person(•) and hasFather(•,•) are two atoms

- In Stanbol basic atoms are
  - Class assertion atom
  - Individual assertion atom
  - Data value assertion atom
  - Range assertion atom
- There are also comparison atoms, string and integer manipulation atoms



## **Atom's notation**

- The atoms may contain
  - Constants: they consist of URI (we are in Web context) or Literal (values)
    - e.g. http//dbpedia.org/resource/Bob\_Marley is a constant, but "Bob Marley"^^xsd:string is a constant too
  - Variables: any identifier preceded by ?
    - o e.g. ?x is a variable, but also ?y is a variable



#### Class assertion atom

A class assertion atom is identified by the operator

is(classPredicate, argument)

#### where

- o classPredicate is a URI that identifies a class
- argument is the resource that has to be proven as typed with the classPredicate. It can be both a constant (a URI) or a variable

e.g. is(<http://xmlns.com/foaf/0.1/Person>, ?x) returns true if the concrete value associated to ?x is typed as http://xmlns.com/foaf/0.1/Person



## Individual assertion atom

has(properyPredicate, arg1, arg2)

#### where

- propertyPredicate is the object property that has to be evaluated. It can be a constant (URI) or a variable (?x)
- arg1 and arg2 are the two arguments of the property. They can be either constants (URI) or variables (?x)



#### Data value assertion atom

values(properyPredicate, arg1, arg2)

#### where

- propertyPredicate is the data property that has to be evaluated. It can be a constant (URI) or a variable (?x)
- arg1 can be either a constant (i.e. URI) or a variable (i.e. ?x)
- arg2 can be either a constant (i.e. a literal) or a variable (i.e. ?x)



## **Namespace Prefixes**

- URIs are useful, but sometime too long for humans
- We could use namespace prefixes instead of full URIs in rule atoms



## **Comparison atoms**

- same(arg1, arg2): returns true if arg1 is equal to arg2
- different(arg1, arg2): returns true if arg1 is different from arg2
- greaterThan(arg1, arg2): returns true if arg1 > arg2
- lessThan(arg1, arg2): returns true if arg1 < arg2</p>
- startsWith(arg1, arg2): returns true if the string associated to arg1 starts with the string associated to arg2
- endsWith(arg1, arg2): returns true if the string associated to arg1 ends with the string associated to arg2



## **String manipulation**

- concat(arg1, arg2): returns a string that is the concatenation of arg1+arg2
- substring(arg, start, length): returns the sub-string of arg from position start for length chars
- lowercase(arg): returns the lower case representation of arg
- uppercase(arg): returns the upper case representation of arg
- str(arg): returns the literal value of any RDF object
- namespace(arg): returns the namespace as a string of any URI
   e.g. namespace(<http://www.foo.org#obj>) ->
   "http://www.foo.org#"
- localname(arg): returns the local as a string of any URI
   e.g. localname(<http://www.foo.org#obj>) -> "obj"



## **Production atoms**

#### newIRI(arg1, arg2)

#### where

- o arg1 is a variable
- arg2 is an expression that returns a literal
- e.g: newIRI(?x, "http://stlab.istc.cnr/Aldo Gangemi") binds the variable ?x to the URI obtained from the literal http://stlab.istc.cnr/Aldo\_Gangemi, namely
  - <http://stlab.istc.cnr/Aldo\_Gangemi>



## **Production atoms (contd)**

#### newLiteral(arg1, arg2)

#### where

- o arg1 is a variable
- arg2 is an expression that returns a literal
- e.g: newLiteral(?x, concat("Aldo ", "Gangemi")) binds the variable ?x to the string literal obtained from the literal "Aldo Gangemi", namely "Aldo Gangemi"



## **Arithmetical atoms**

- sum(arg1, arg2): returns a new integer the is equal to arg1+arg2
- sub(arg1, arg2): returns a new integer the is equal to arg1-arg2
- mult(arg1, arg2): returns a new integer the is equal to arg1\*arg2
- div(arg1, arg2): returns a new integer the is equal to arg1/arg2
- arg1 and arg2 can be numerical expression or numbers



## **Example of rule**

Considering Stanbol Rules, the formula  $hasFather(x,y) \land hasBrother(y,z) \Rightarrow hasUncle(x,z)$ 

expresses in predicate calculus becomes

```
myRule[ has(<http://myont.org/hasFather>, ?x, ?y) . has(<http://myont.org/hasBrother>, ?y, ?z)
```

->

has(<http://myont.org/hasUncle>, ?x, ?z) ]



## **Stanbol Refactor**

- Add to Stanbol a framework for RDF graph refactoring
- Rule-based refactoring



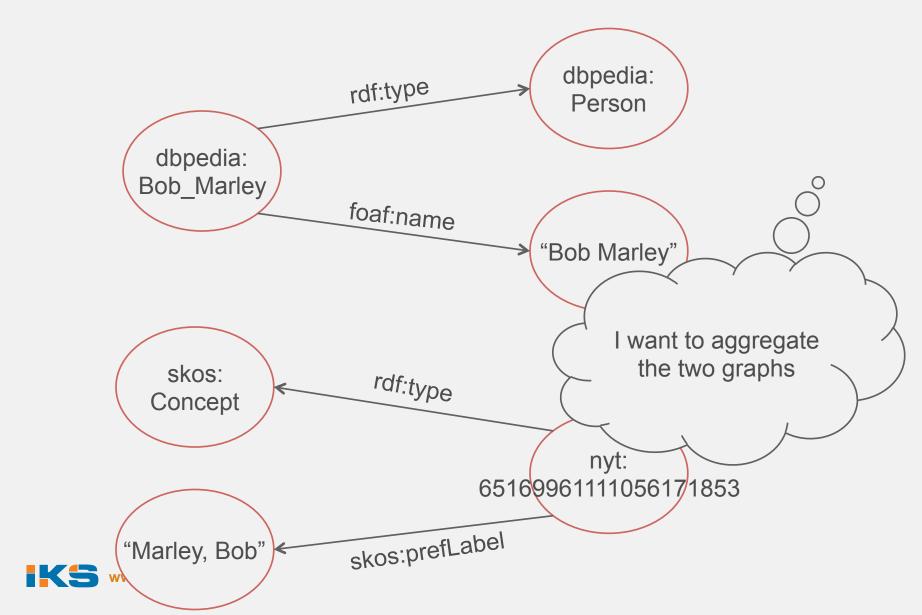
## Why do I need refactoring?

- A possible scenario
  - My system fetches knowledge from different sources in LOD
  - Each of these sources uses its own ontology/vocabulary

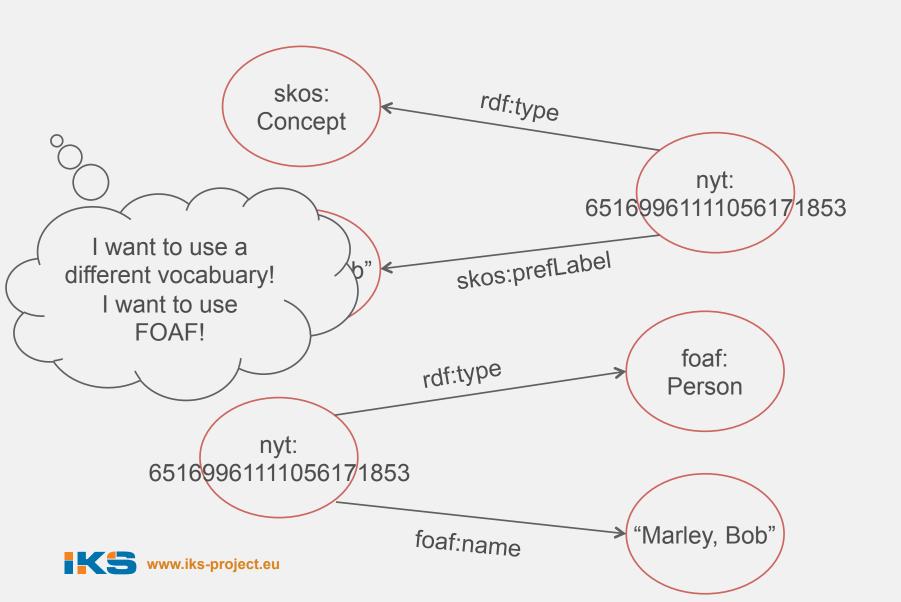
How to add a homogeneous representation of knowledge expressed with heterogeneous vocabularies?



## An example



## **Another example**

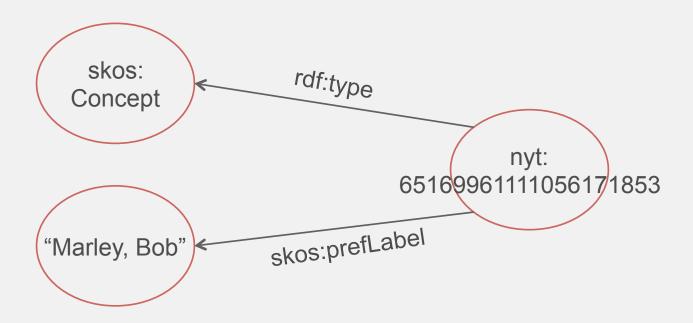


## Rule-based refactoring

- The Stanbol Refactor applies RDF graph transformation by means of transformation rules
- Rules drive the transformation
- A set of rules for a tranformation task characterize a recipe
- A recipe identifies the kind of transformation and the rules need by the transformation task



## Define a refactoring recipe



We want to use the FOAF vocabulary instead of SKOS



## Define a refactoring recipe

```
skos = <http://www.w3.org/2004/02/skos/core#> .
foaf = <http://xmlns.com/foaf/0.1/> .
```

```
conceptToPerson[ is(skos:Concept, ?x) -> is
  (foaf:Person, ?x) ] .
```

```
labelRule[ values(skos:prefLabel, ?x, ?y) ->
  values(foaf:name, ?x, ?y) ]
```



## How the refactor works

- Each rule is executed individually starting from the first in the recipe
- Each rule is interpreted and executed as a SPARQL CONSTRUCT



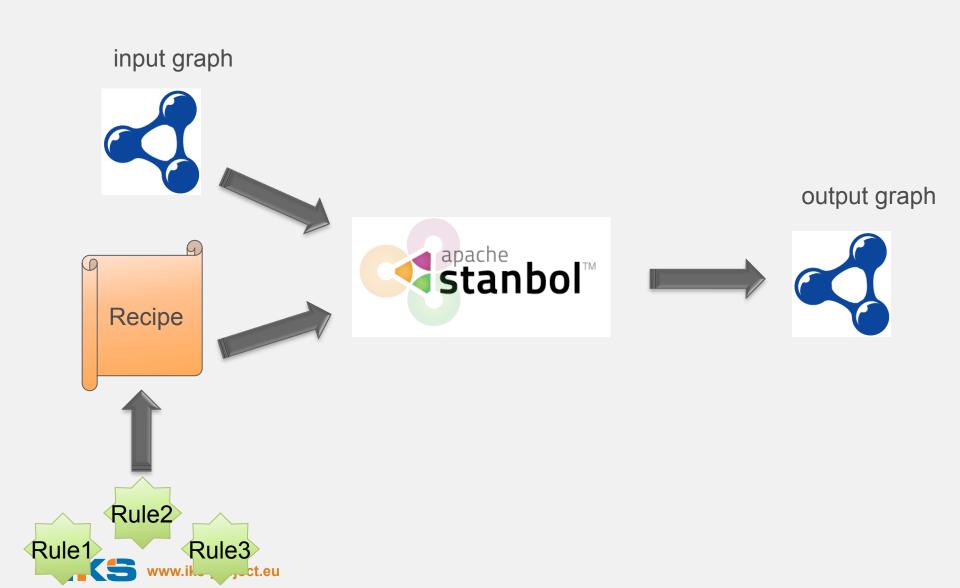
## From rules to CONSTRUCT

#### The rule

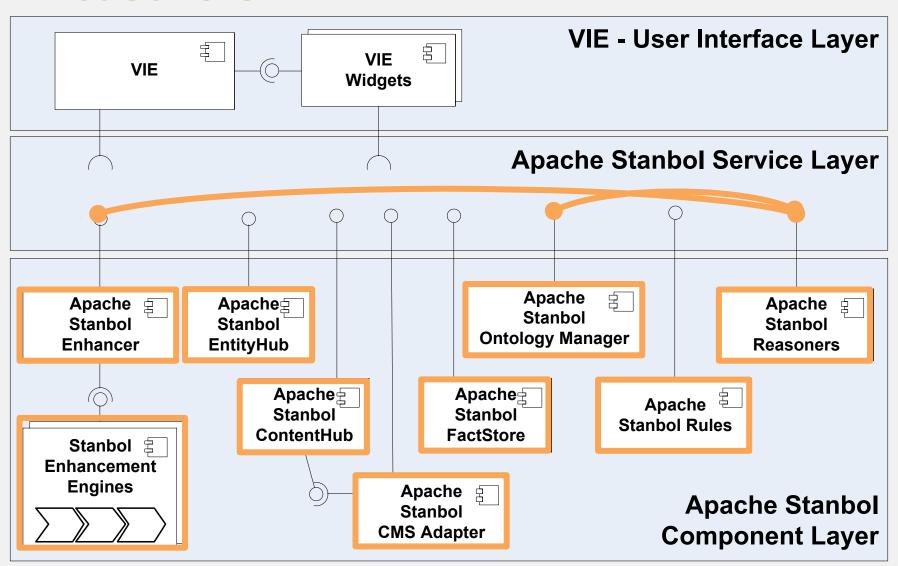
```
skos = <http://www.w3.org/2004/02/skos/core#> .
   foaf = <http://xmlns.com/foaf/0.1/> .
   conceptToPerson[is(skos:Concept, ?x) -> is(foaf:Person, ?x)]
is interpreted as
   PREFIX skos: <a href="http://www.w3.org/2004/02/skos/core#">http://www.w3.org/2004/02/skos/core#></a>
   PREFIX foaf: <a href="http://xmlns.com/foaf/0.1/">
   PREFIX rdf: <a href="http://www.w3.org/1999/02/22-rdf-syntax-ns#">http://www.w3.org/1999/02/22-rdf-syntax-ns#</a>>
   CONSTRUCT { ?x rdf:type foaf:Person }
   WHERE { ?x rdf:type skos:Concept }
```



## Recipes, rules and Stanbol



#### Reasoners



**Semantic Lifting** 

Knowledge Representation & Reasoning

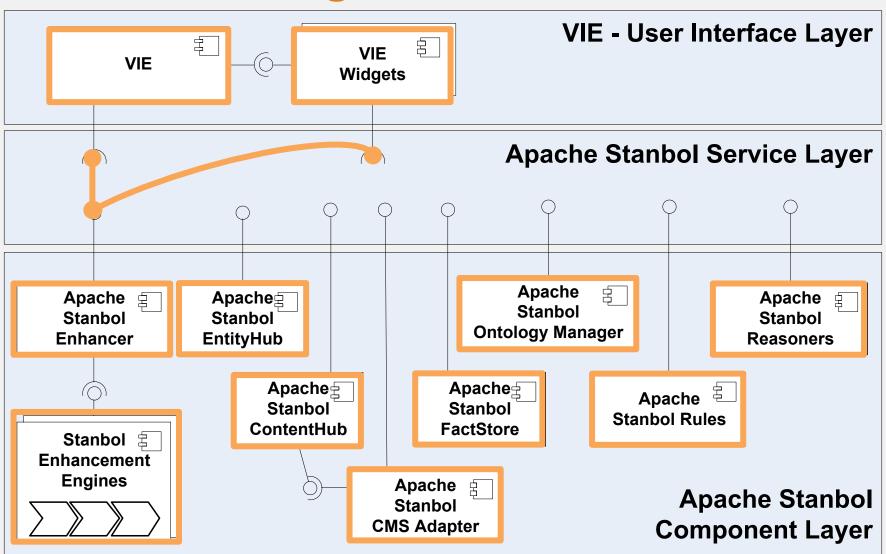
## Reasoners Features



- Common API for existing (open source) reasoning services
- Supports different reasoners and configuration in parallel
- Supported third-party reasoners
  - Jena RDFS
  - OWL
  - OWLMini
  - HermiT



## **VIE & VIE Widgets**



**Semantic Lifting** 

Knowledge Representation & Reasoning

# VIE & VIE Widgets Features



- VIE is a JavaScript library for implementing decoupled CMS and semantic interaction in web applications
- VIE provides easy access to the semantic metadata (RDFa) within a web page
- VIE Widgets are user interface components that implement semantic user interactions
- Examples
  - Semantic image search
  - Automatic tagging of entities
  - Semi-automatic content annotation



## **VIE Demo**





## Interaction with Knowledge



## License

- IKS software is licensed under business-friendly open source software licenses.
- IKS software can be freely used / changed / distributed in your products.
- For the rare cases where artifacts use a less permissive license, you will find a notice.
  - e.g. we use models for natural language processing from the Apache OpenNLP project whose licenses are not clarified, yet.



#### **Get in Contact**

- o VIE
  - Homepage http://viejs.org
  - Google User Group https://groups.google.com/forum/#!forum/viejs
- Apache Stanbol
  - Homepage http://incubator.apache.org/stanbol
  - Mailinglist subscription
     stanbol-dev-subscribe@incubator.apache.org



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