Computer Networks

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Notes derived from "*Computer Networking: A Top Down Approach*", by Jim Kurose and Keith Ross, Addison-Wesley.

The slides are adapted and modified based on slides from the book's companion Web site, as well as modified slides by Anirban Mahanti and Carey Williamson.

Scalable Content Delivery
Motivation

- Use of Internet for content delivery is massive … and becoming more so (e.g., majority of all IP traffic is video content)
- Variety of approaches: HTTP-based Adaptive Streaming (HAS), broadcast/multicast, batching, replication/caching (e.g. CDNs), P2P, peer-assisted, …
- In these slides, we only provide a few high-level examples

Why Study Multimedia Networking?

- Exciting and fun!
- Multimedia is everywhere
- Industry-relevant research topic
- Lots of open research problems

Requirements and quality of service

- Quality of Service (QoS)
  - Real-time requirements (e.g., latency, jitter)
  - Loss/stall requirements (e.g., drop rates, late packets)
  - Bandwidth requirements (e.g., throughput)
  - Service availability
- Quality of Experience (QoE)
  - Measure of the users quality of experience
  - Multimedia: Most negatively effected by stalls

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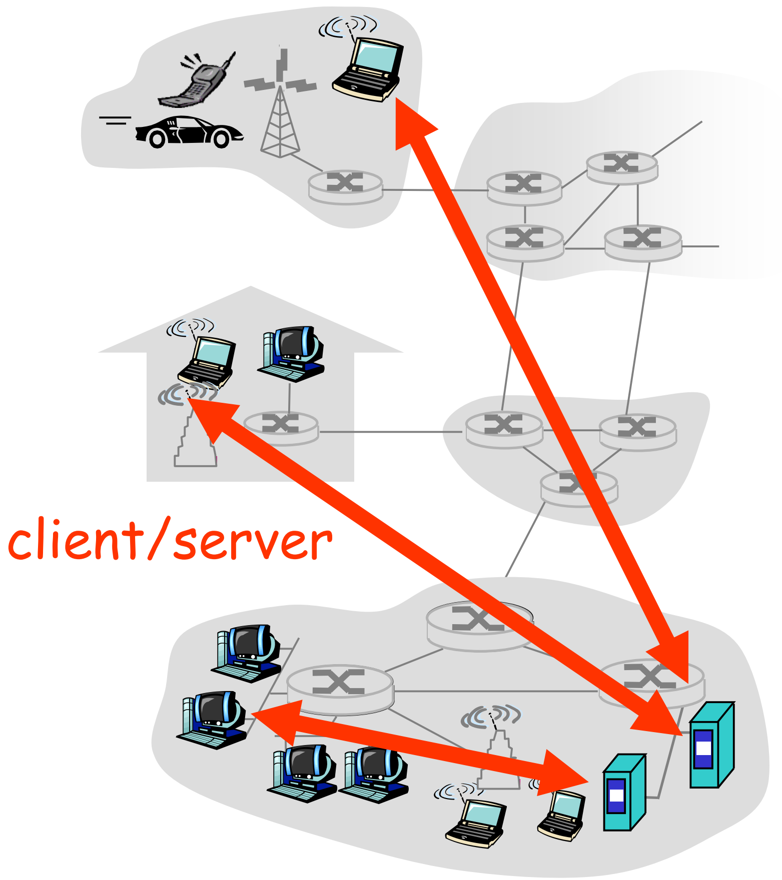
Service models

- Client-server (one-to-one)
- Peer-to-peer (machines can act as both client and server)
- Multicast/broadcast (one-to-many and many-to-many)
  - Application layer, IP-based, and down at the MAC-layer
- ISP-based caching, CDNs, cloud, and other third-party solutions

Client-server architecture

Client/server model has well-defined roles.



server:

- always-on host
- permanent IP address
- server farms for scaling

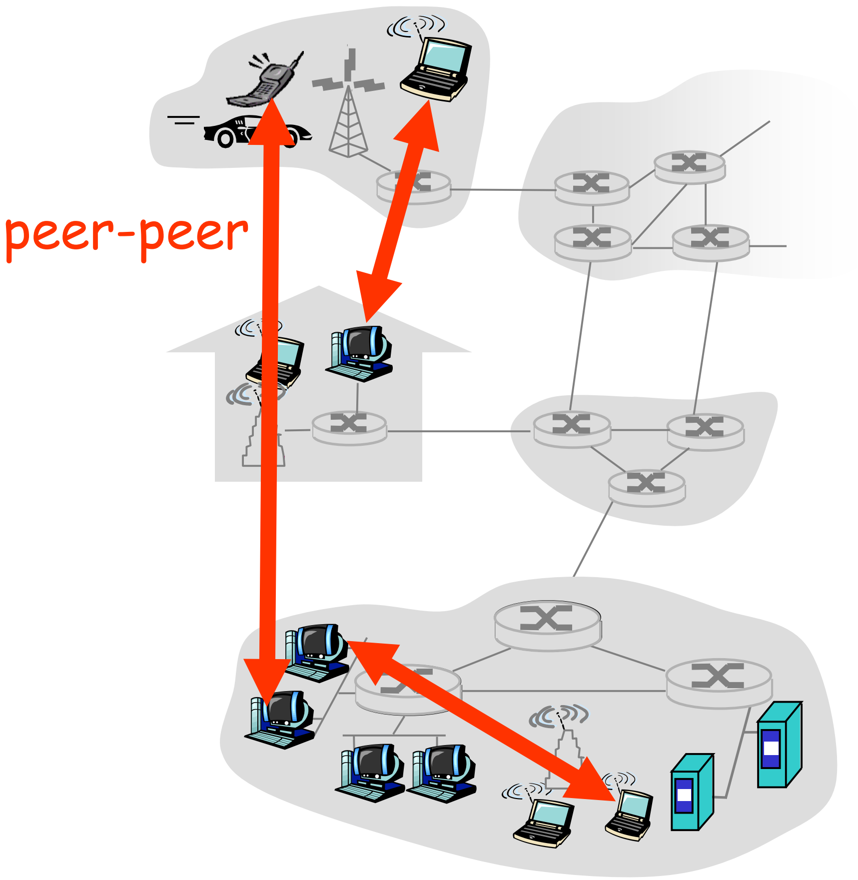
clients:

- communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other

Pure P2P architecture

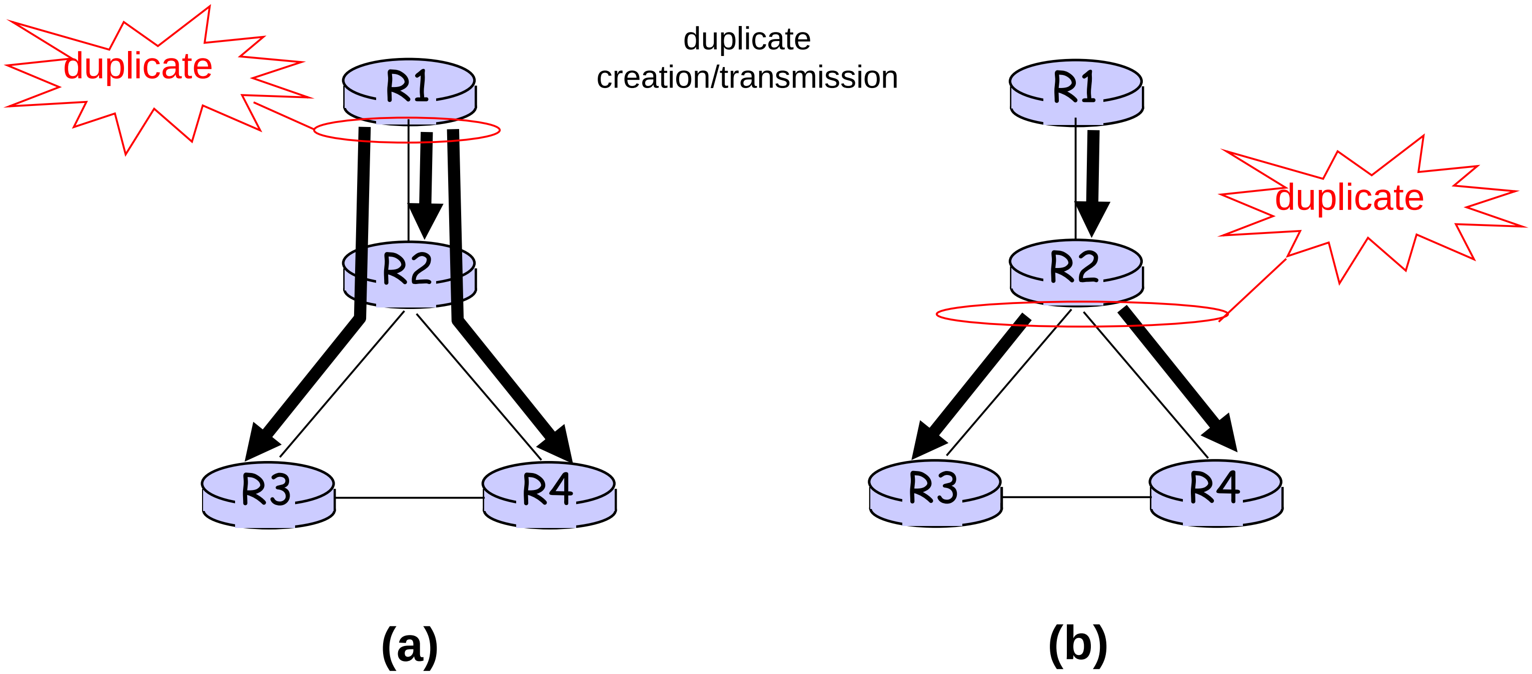
No fixed clients or servers: Each host can act as both client and server at any time

- *no* always-on server
- arbitrary end systems directly communicate
- peers are intermittently connected and change IP addresses



Additional Multimedia Support

Multicast/Broadcast



Source-duplication versus in-network duplication.
(a) source duplication, (b) in-network duplication

Also, application-layer multicast …

Evolved Multimedia Broadcast/Multicast Service (eMBMS) in LTE-advanced

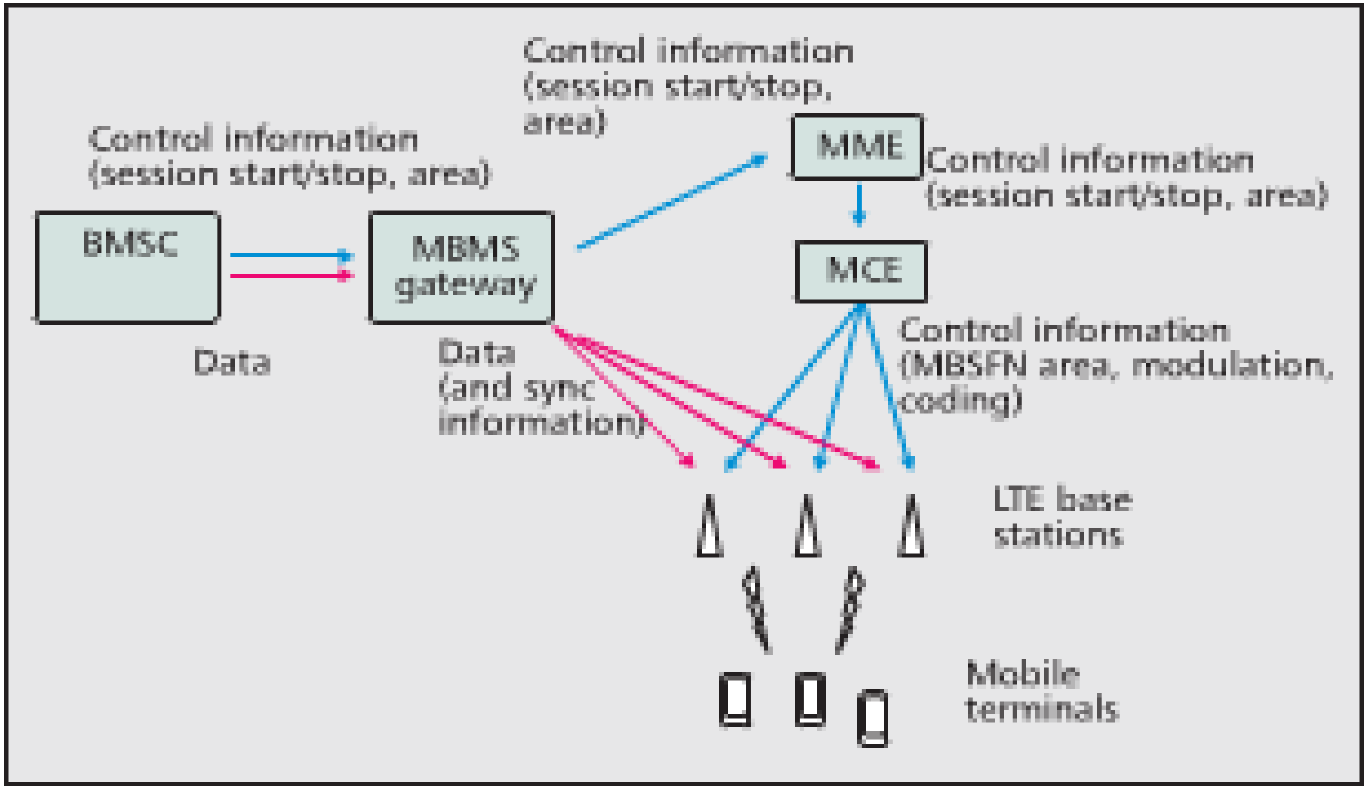


Figure 4. *RAN architecture for SFN across LTE base stations.*

- Separation of control plane and data plane

Image from: Lecompte and Gabin, Evolved Multimedia Broadcast/Multicast Service (eMBMS) in LTE-Advanced: Overview and Rel-11 Enhancements, IEEE Communications Magazine, Nov. 2012.

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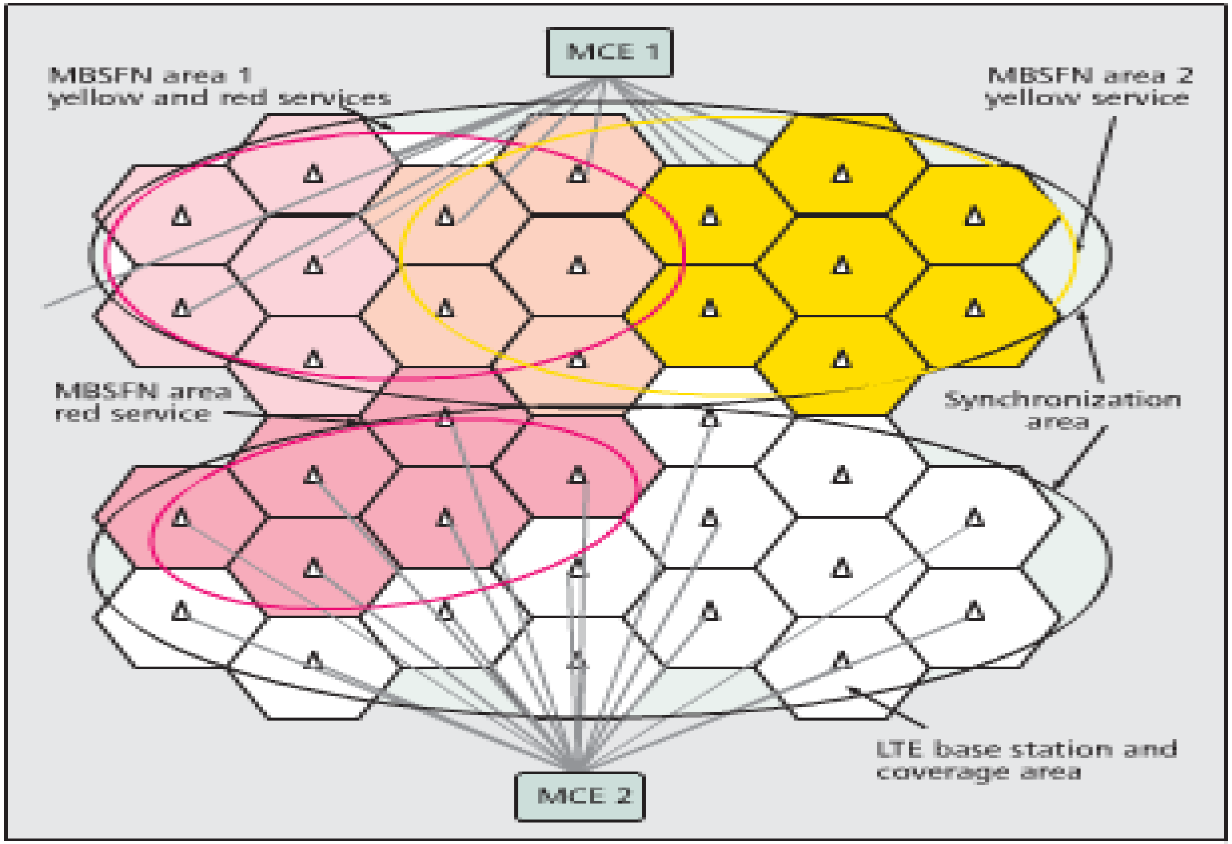


Figure 5. *Example with two MBMS services with different services areas.*

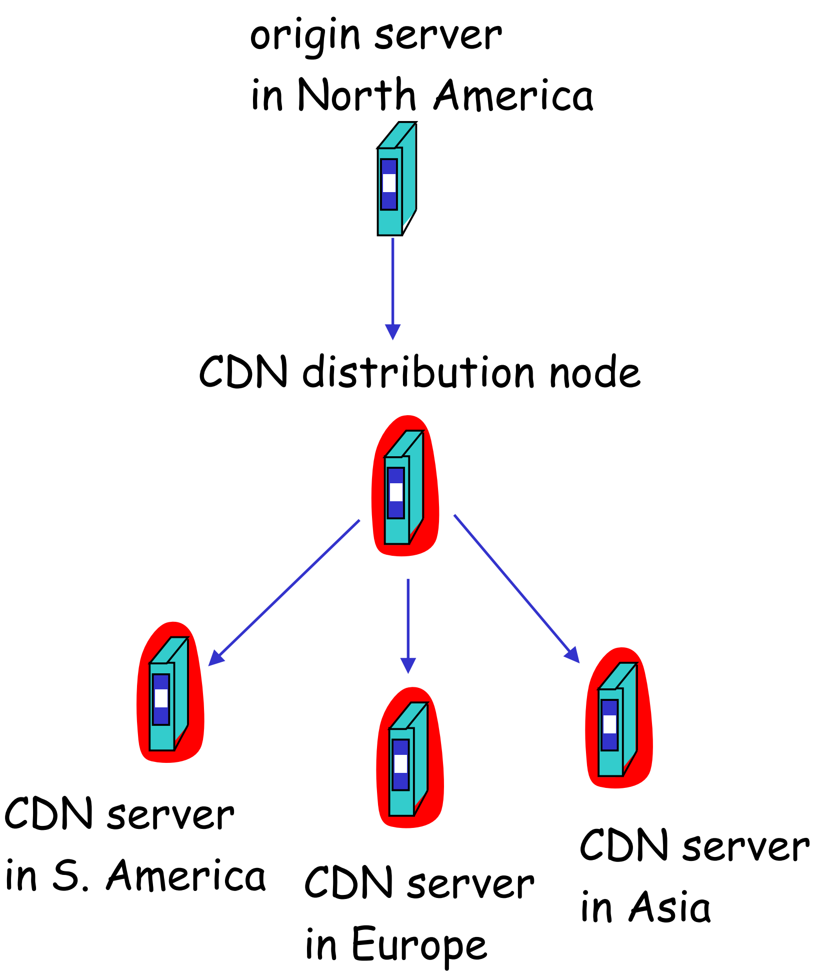
- MBMSFN and use of services areas

Image from: Lecompte and Gabin, Evolved Multimedia Broadcast/Multicast Service (eMBMS) in LTE-Advanced: Overview and Rel-11 Enhancements, IEEE Communications Magazine, Nov. 2012.

Content distribution networks (CDNs)

Content replication

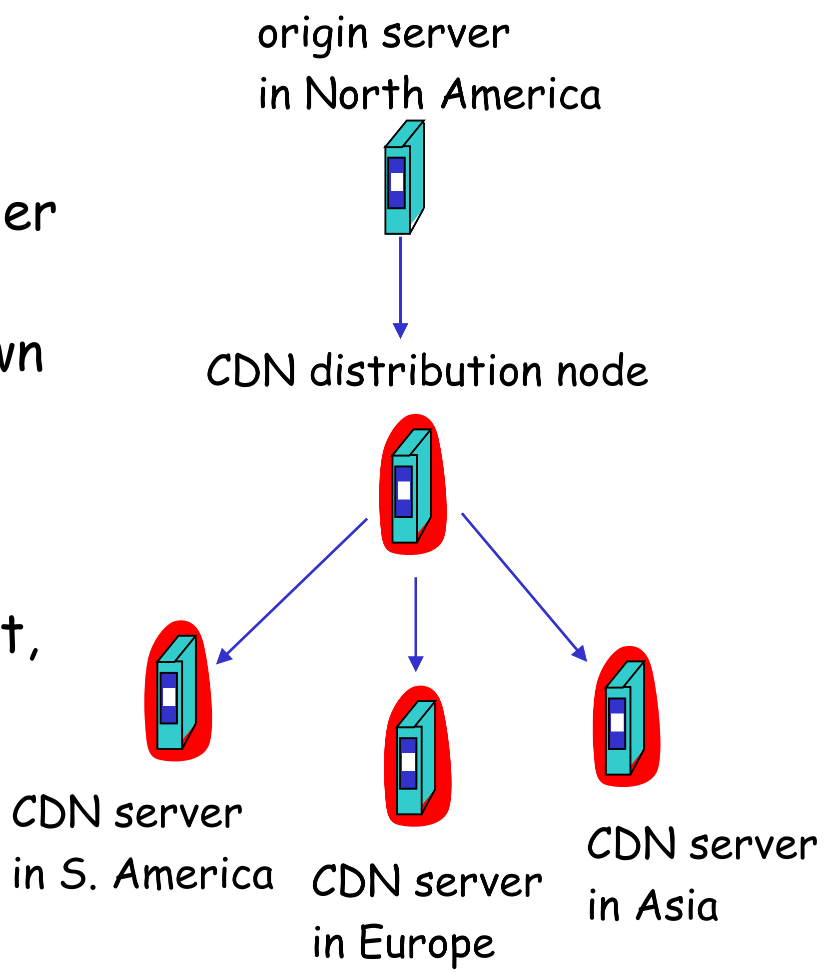
- replicate content at hundreds of servers throughout Internet (often in edge/access network)
- content "close" to user reduce impairments (loss, delay) of sending content over long paths



Content distribution networks (CDNs)

Content replication

- CDN (e.g., Akamai, Limewire) customer is the content provider (e.g., CNN)
- Other companies build their own CDN (e.g., Google)
- CDN replicates customers' content in CDN servers.
- When provider updates content, CDN updates servers



DNS redirect example

Multimedia Networking

Principles

- Classify multimedia applications
- Identify the network services the apps need
- Making the best of "best effort" service
- Mechanisms for providing QoS

Protocols and Architectures

- Specific protocols for best effort delivery
- Architectures for QoS

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Multimedia Networking Applications

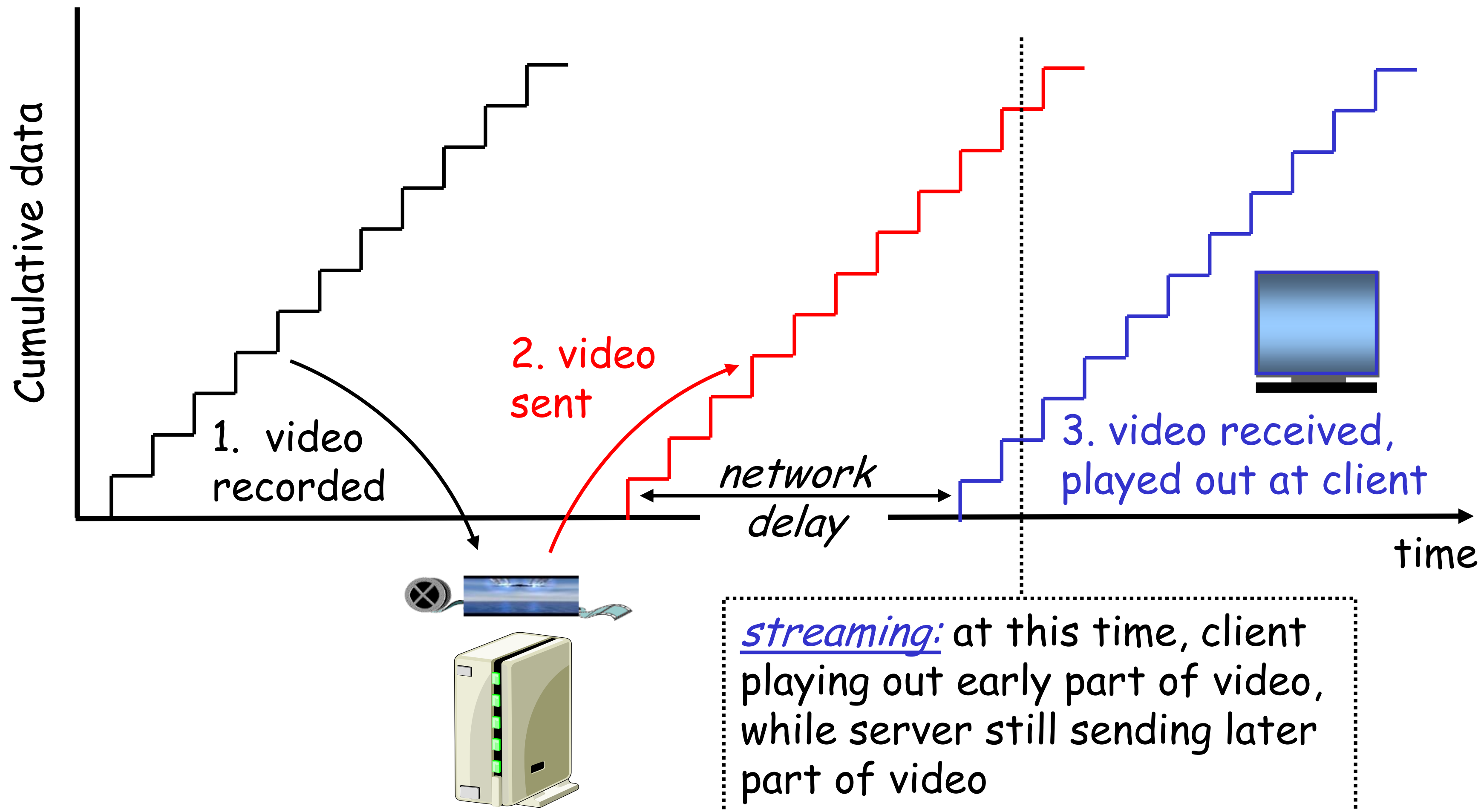
Classes of MM applications:

Multimedia Networking Applications

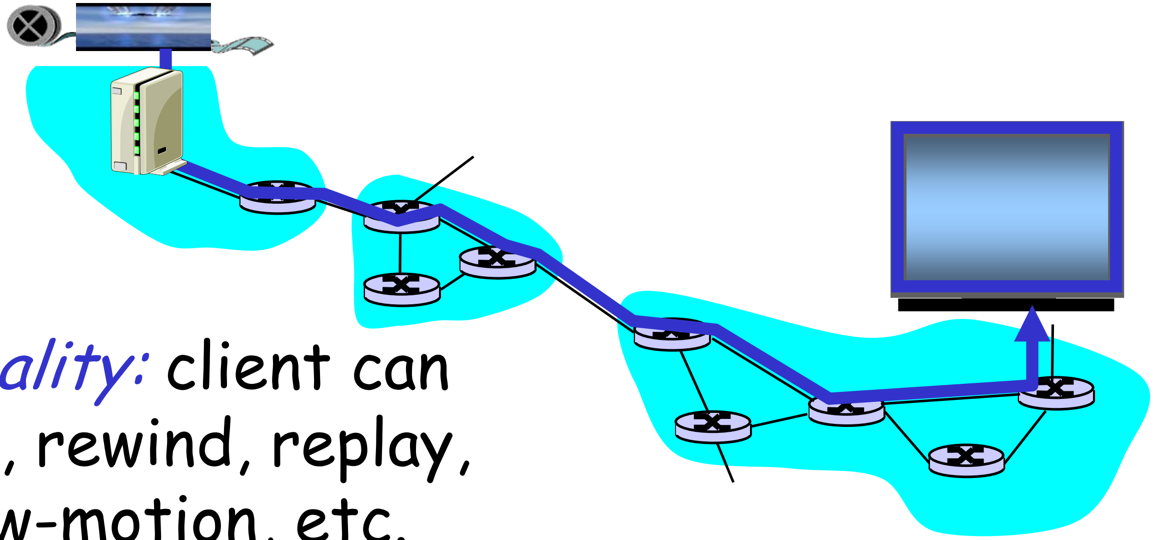
Classes of MM applications:

1) Streaming stored audio and video
2) Streaming live audio and video
3) Real-time interactive audio and video

Streaming Stored Multimedia (1/2)

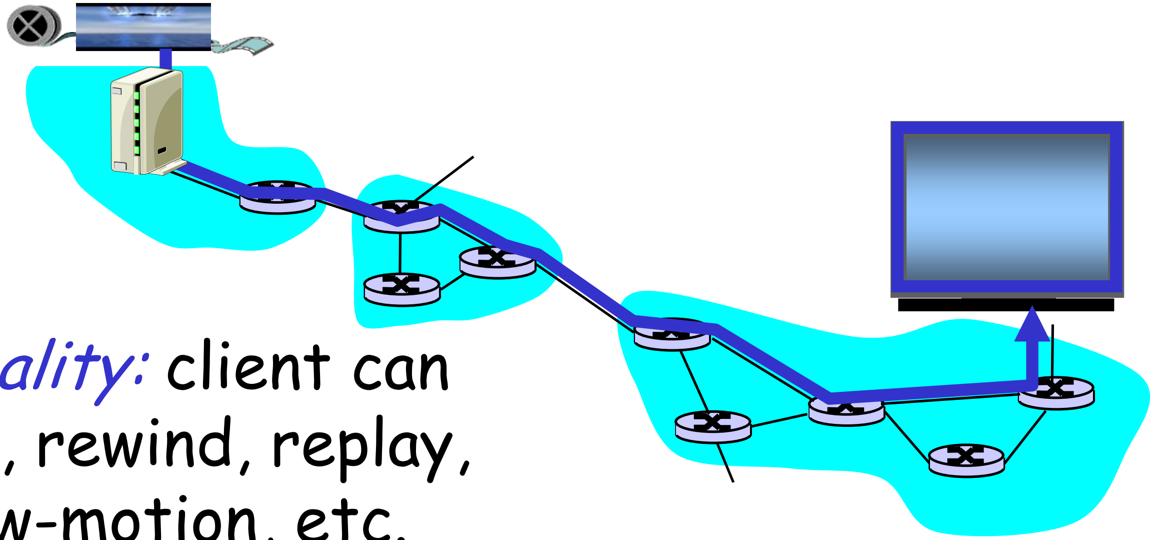


Streaming Stored Multimedia (2/2)



- *VCR-like functionality:* client can start, stop, pause, rewind, replay, fast-forward, slow-motion, etc.

Streaming Stored Multimedia (2/2)



- *VCR-like functionality:* client can start, stop, pause, rewind, replay, fast-forward, slow-motion, etc.
  - 10 sec initial delay OK
  - 1-2 sec until command effect OK
  - need a separate control protocol?
- timing constraint for data that is yet to be transmitted: must arrive in time for playback

Streaming Live Multimedia

Examples:

- Internet radio talk show
- Live sporting event

Streaming Live Multimedia

Examples:

- Internet radio talk show
- Live sporting event

Streaming

- playback buffer
- playback can lag tens of seconds after transmission
- still have timing constraint

Interactivity

- fast-forward is not possible
- rewind and pause possible!

Interactive, Real-time Multimedia

- applications: IP telephony, video conference, distributed interactive worlds

Interactive, Real-time Multimedia

- applications: IP telephony, video conference, distributed interactive worlds
- end-end delay requirements:
  - audio: < 150 msec good, < 400 msec OK
    - includes application-layer (packetization) and network delays
    - higher delays noticeable, impair interactivity
- session initialization
  - callee must advertise its IP address, port number, frame rate, encoding algorithms

And the future has more …

Multimedia Networking Applications

Fundamental characteristics:

Multimedia Networking Applications

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- But **loss-tolerant**: infrequent losses cause minor transient glitches

- Unlike data apps, which are often delay-tolerant but loss-sensitive.

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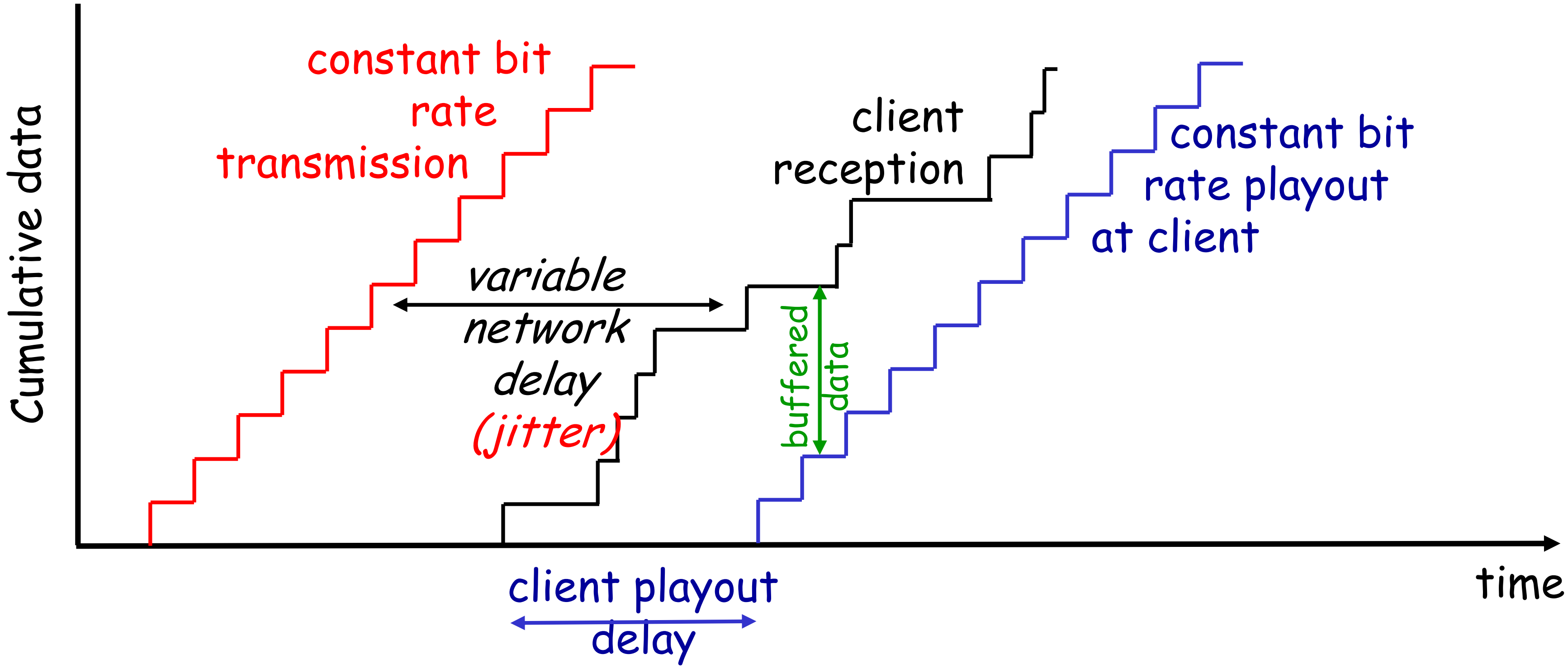
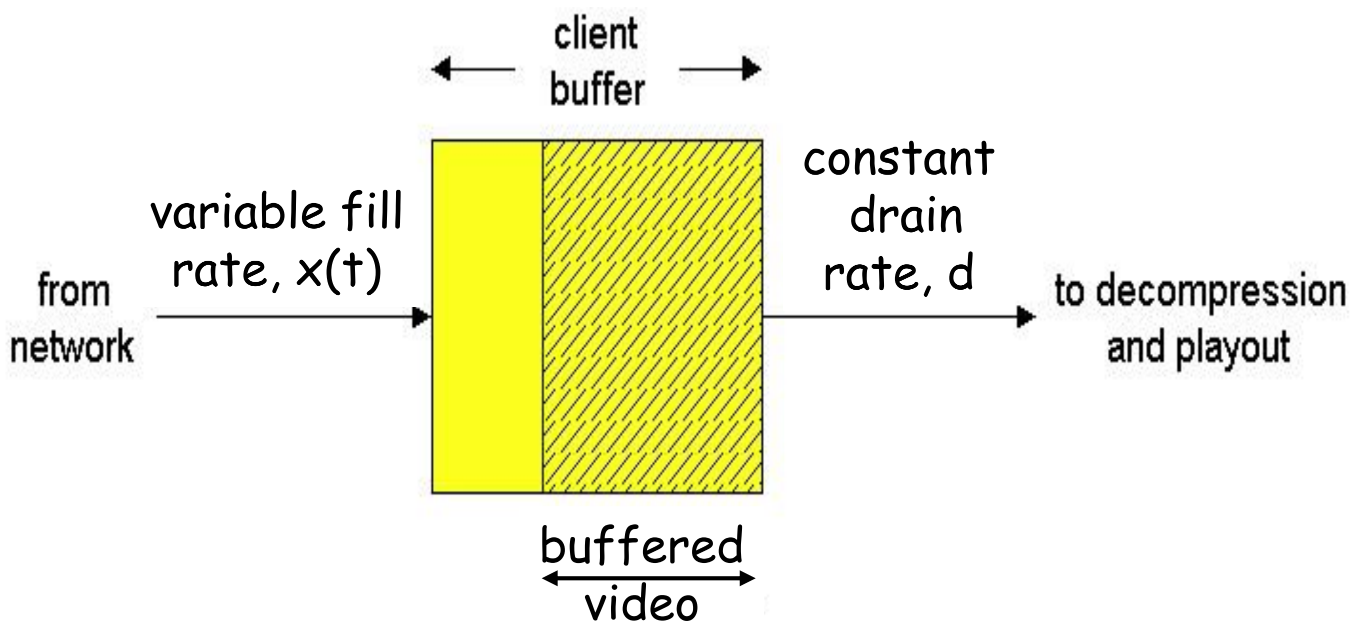
Jitter is the variability of packet delays within the same packet stream

Delay Jitter



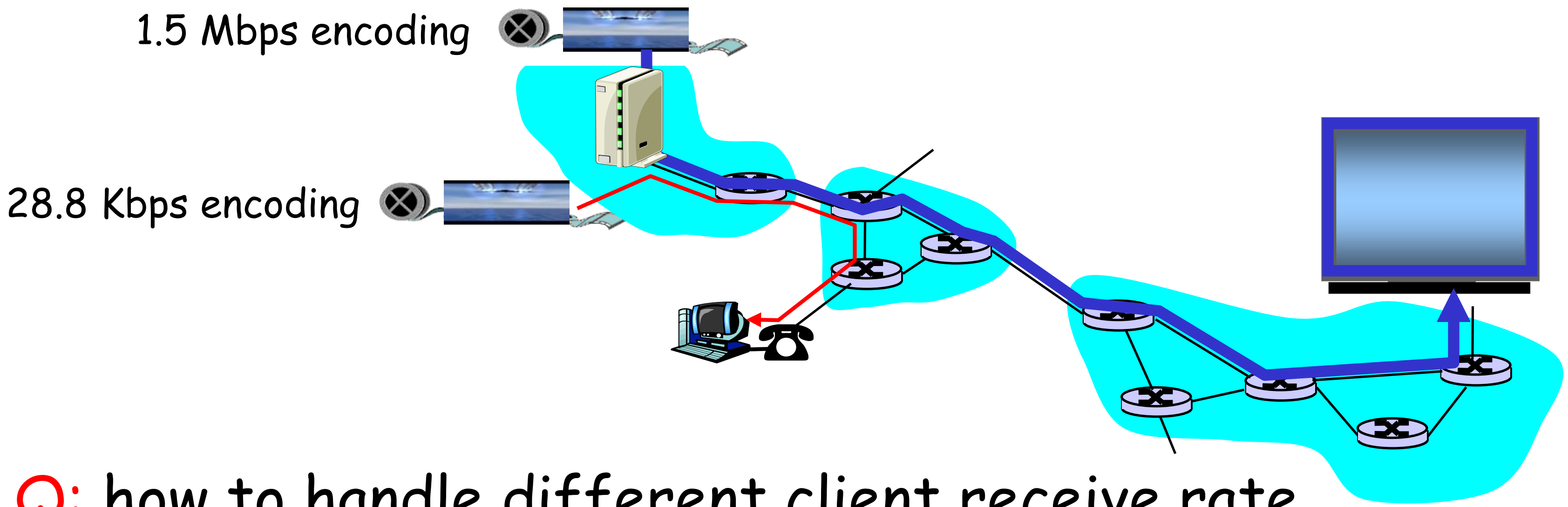
- consider end-to-end delays of two consecutive packets: difference can be more or less than 20 msec (transmission time difference)

Streaming Multimedia: Client Buffering



- Client-side buffering, playout delay compensate for network-added delay, delay jitter

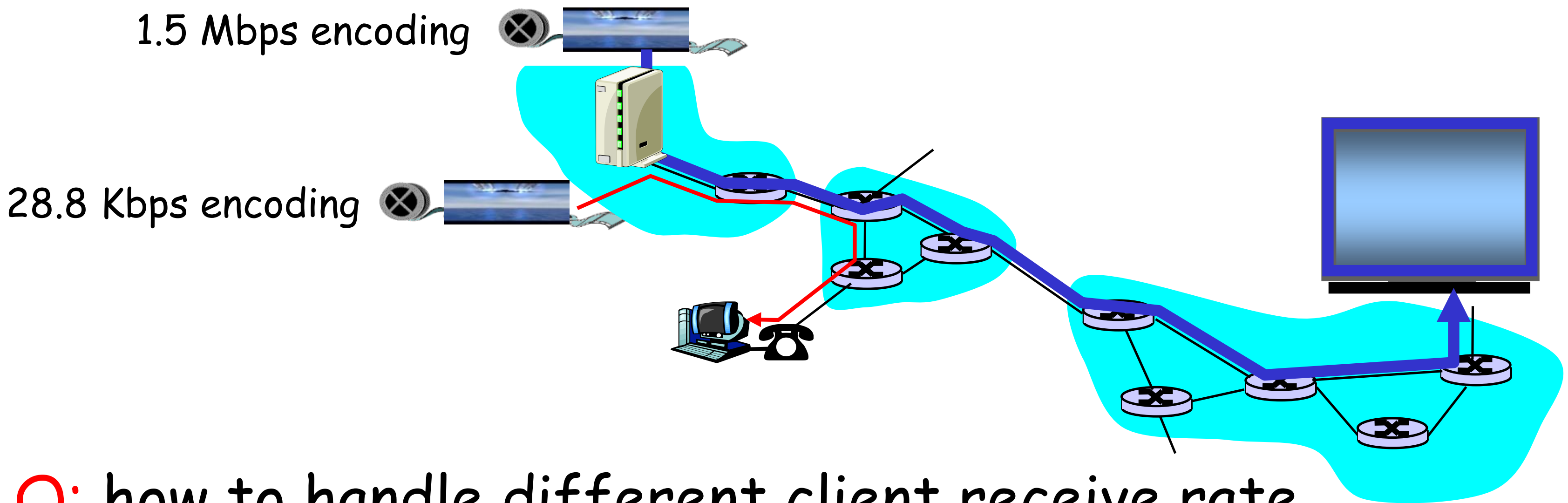
Streaming Multimedia: client rate(s)



Q: how to handle different client receive rate capabilities?

- 28.8 Kbps dialup
- 100 Mbps Ethernet

Streaming Multimedia: client rate(s)



Q: how to handle different client receive rate capabilities?

- 28.8 Kbps dialup
- 100 Mbps Ethernet

A1: server stores, transmits multiple copies of video, encoded at different rates

A2: layered and/or dynamically rate-based encoding

Consider first ...
Streaming Stored Multimedia

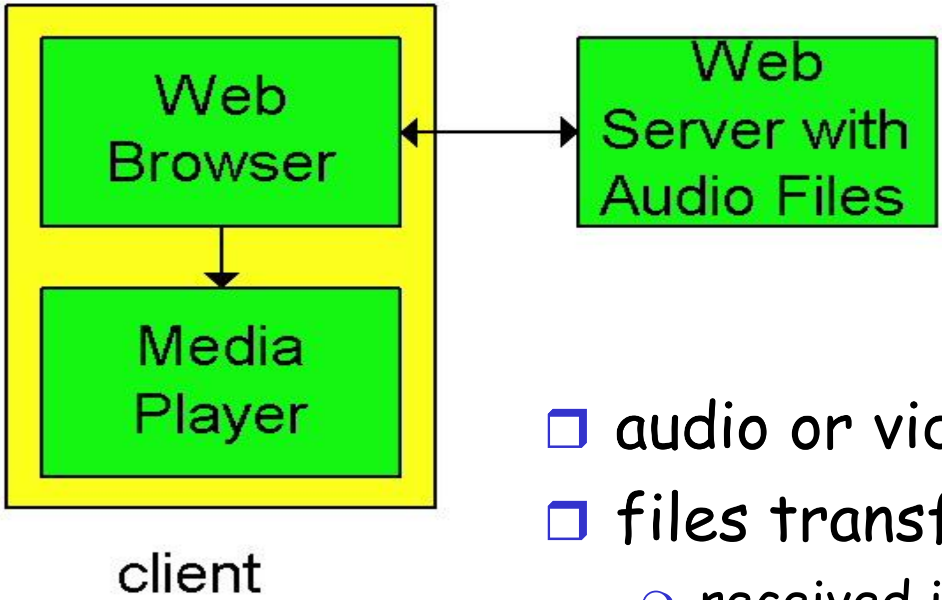
application-level streaming techniques for making the best out of best effort service:

- client-side buffering
- use of UDP versus TCP
- multiple encodings of multimedia

Media Player

- jitter removal
- decompression
- error concealment
- graphical user interface w/ controls for interactivity

Internet multimedia: simplest approach

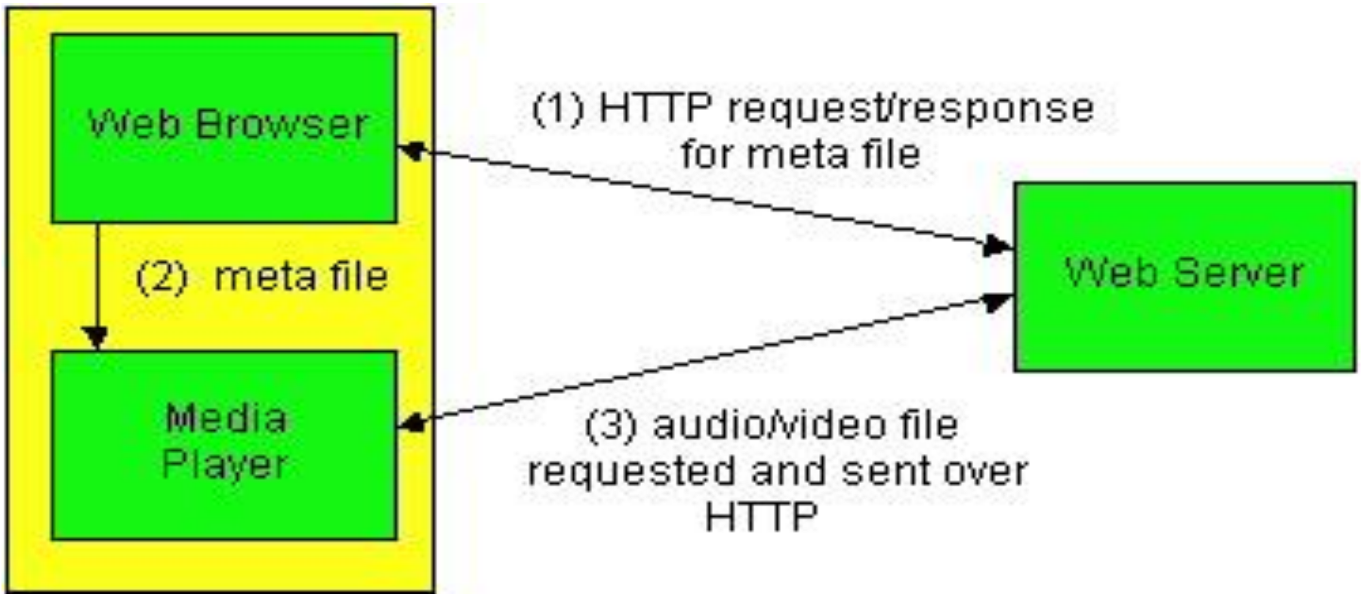


- audio or video stored in file
- files transferred as HTTP object
  - received in entirety at client
  - then passed to player

audio, video is downloaded, not streamed:

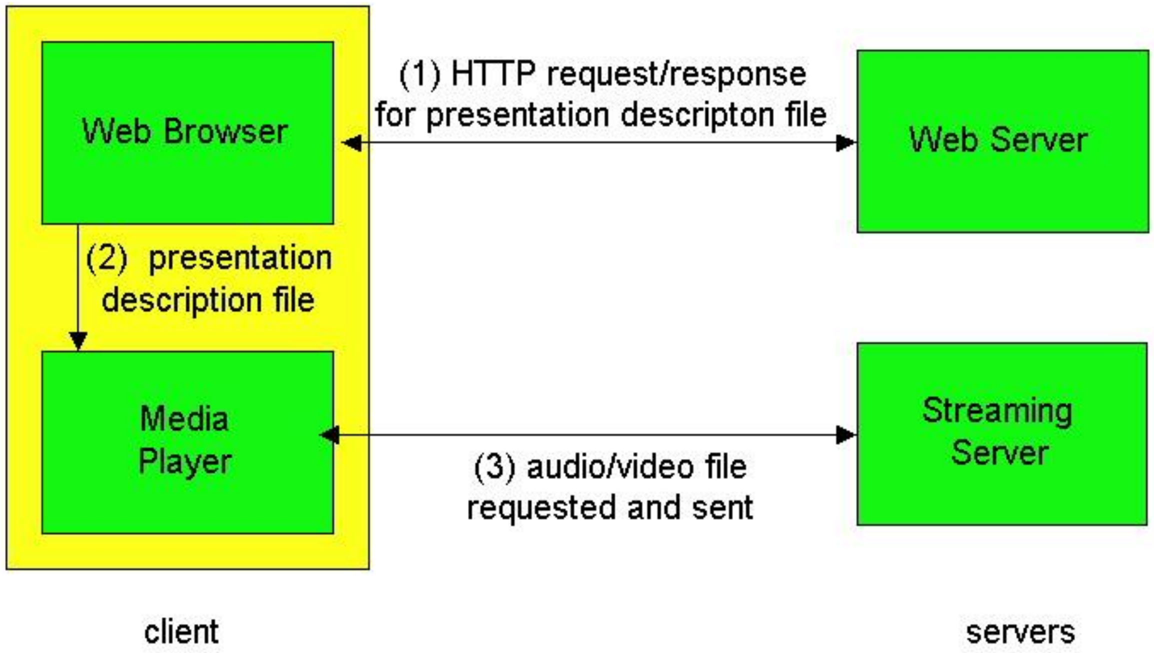
- long delays until playout, since no pipelining!

Progressive Download



- browser retrieves **metafile** using HTTP GET
- browser launches player, passing metafile to it
- media player contacts server directly
- server downloads audio/video to player

Streaming from a Streaming Server



- This architecture allows for non-HTTP protocol between server and media player
- Can also use UDP instead of TCP.

Streaming Multimedia: UDP or TCP?

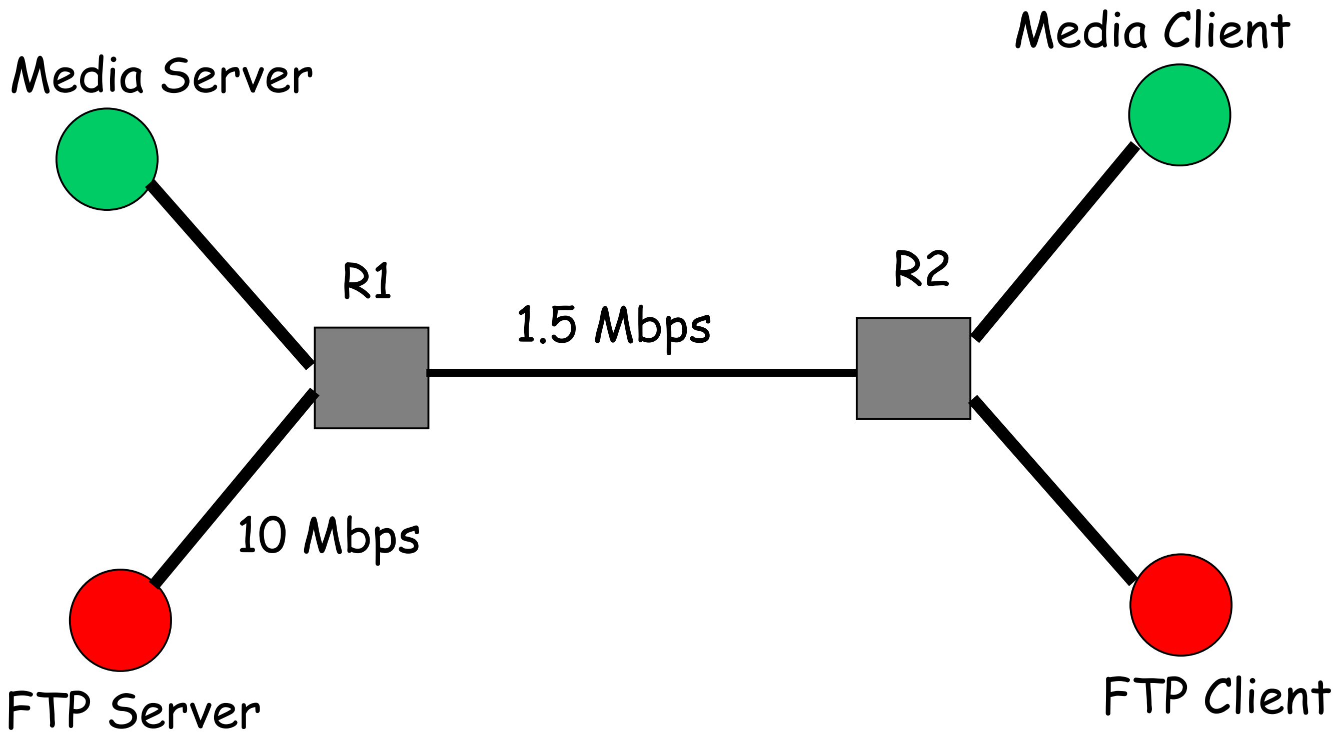
UDP

- server sends at rate appropriate for client (oblivious to network congestion !)
  - often send rate = encoding rate = constant rate
  - then, fill rate = constant rate - packet loss
- short playout delay (2-5 seconds) to compensate for network delay jitter
- error recover: time permitting

TCP

- send at maximum possible rate under TCP
- fill rate fluctuates due to TCP congestion control
- larger playout delay: smooth TCP delivery rate
- HTTP/TCP passes more easily through firewalls

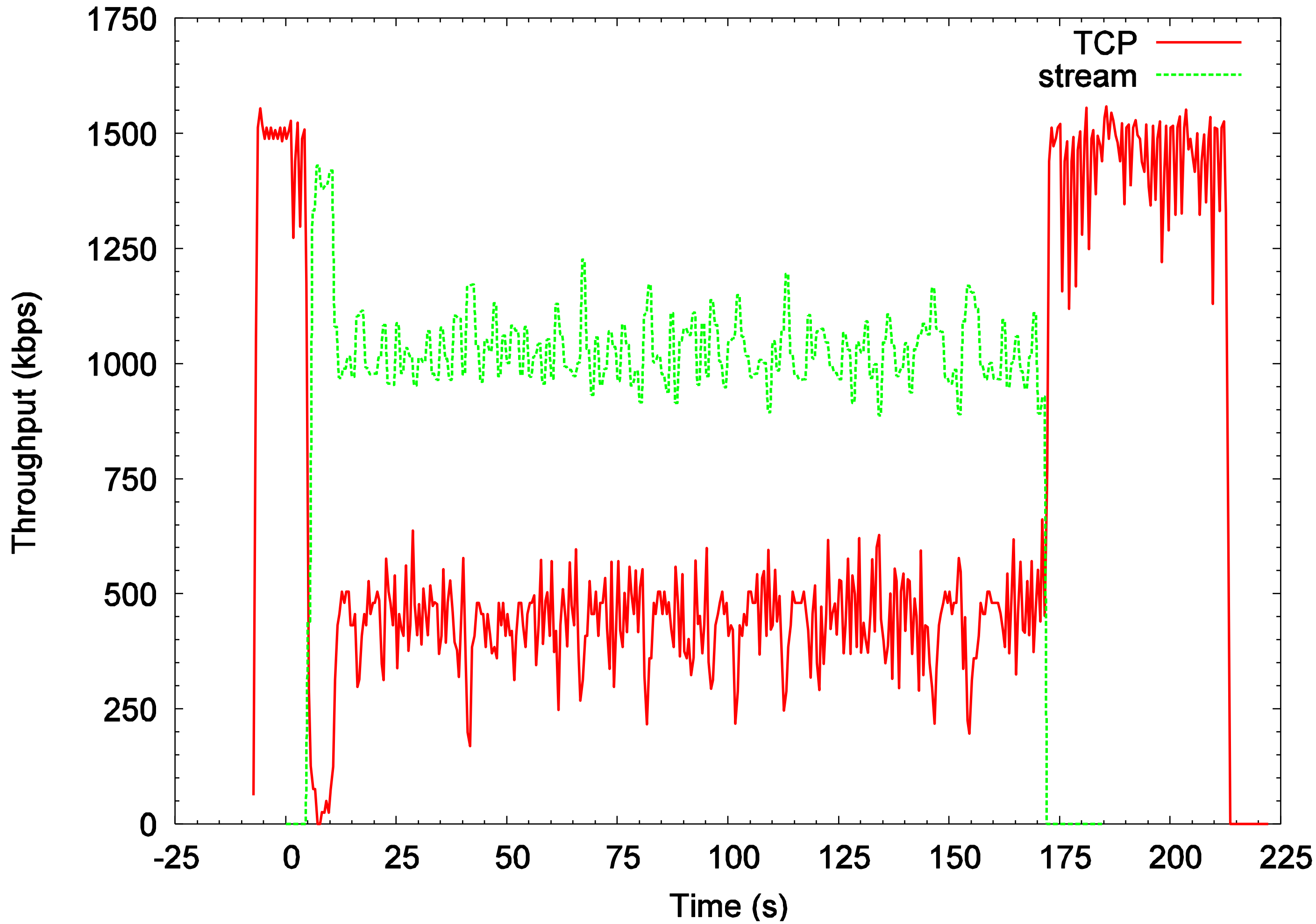
Fairness of UDP Streams (1/2)



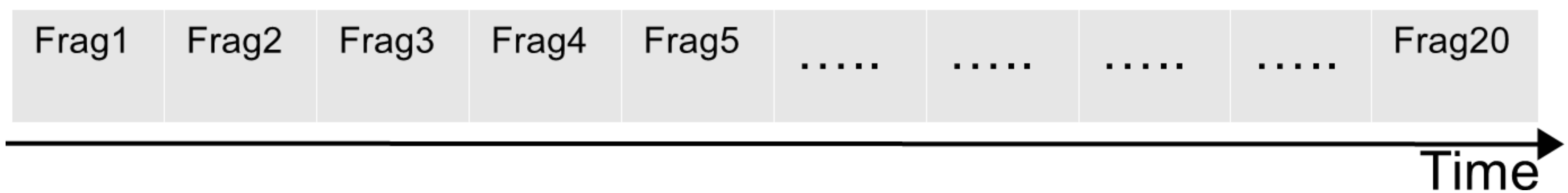
- R1-R2 is the bottleneck link
- Streaming uses UDP at the transport layer; requested media encoded at 1 Mbps
- What fraction of the bottleneck is available to FTP?

Credit: MSc thesis work by Sean Boyden (2006)

Fairness of RealVideo Streams (2/2)

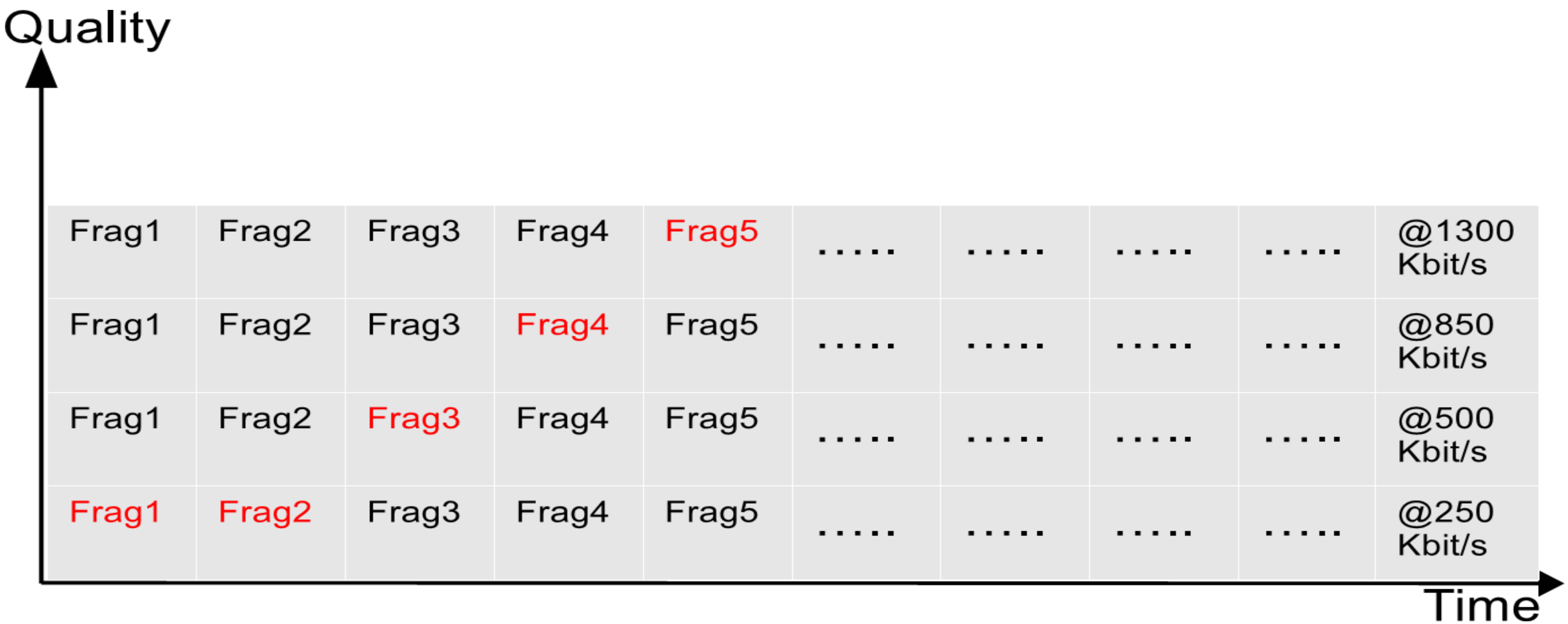


HTTP-based streaming



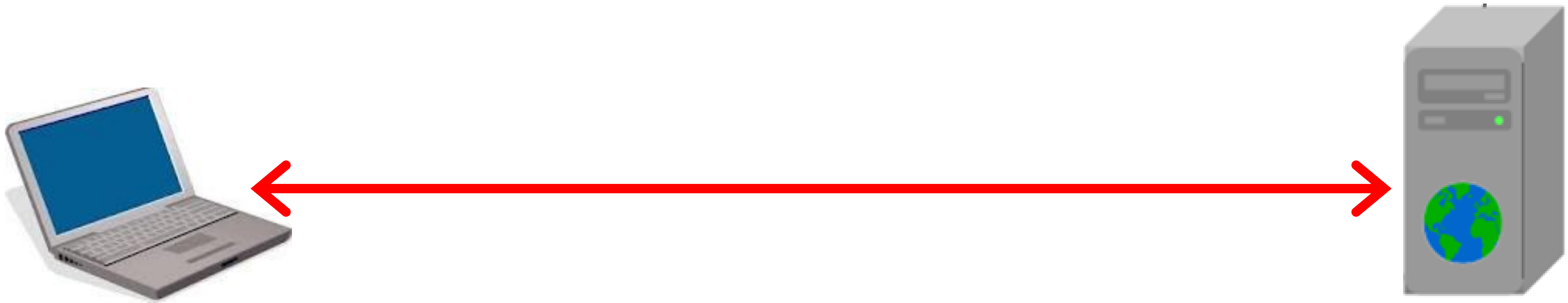
- HTTP-based streaming
  - Allows easy caching, NAT/firewall traversal, etc.
  - Use of TCP provides natural bandwidth adaptation
  - Split into fragments, download sequentially
  - Some support for interactive VoD

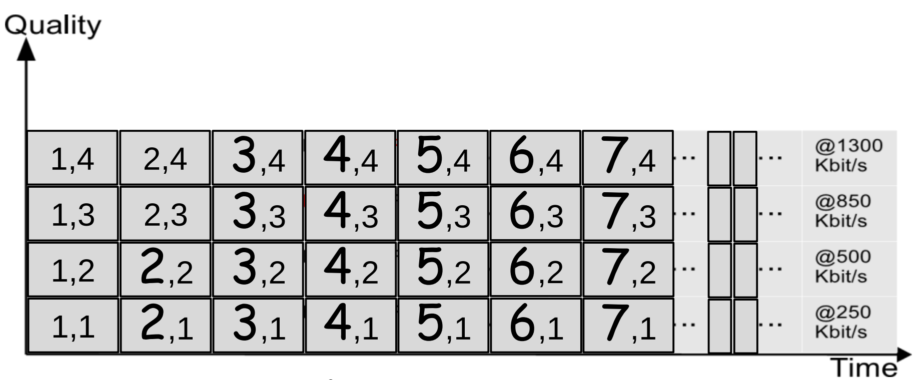
HTTP-based adaptive streaming (HAS)



- HTTP-based adaptive streaming
  - Multiple encodings of each fragment (defined in manifest file)
  - Clients adapt quality encoding based on (buffer and network) conditions

Example





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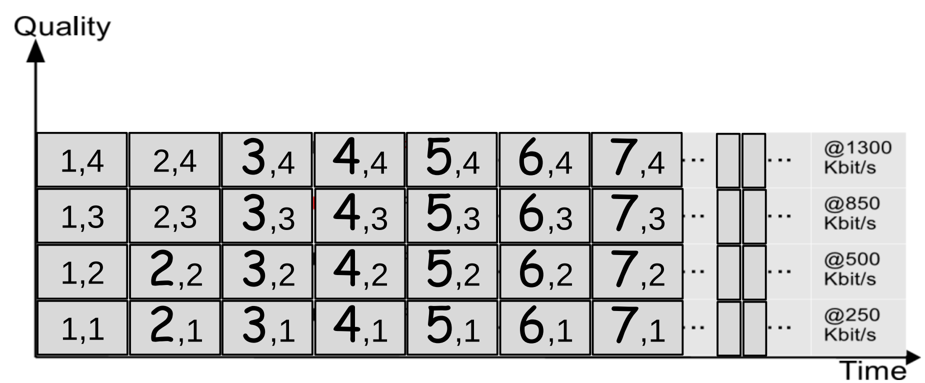
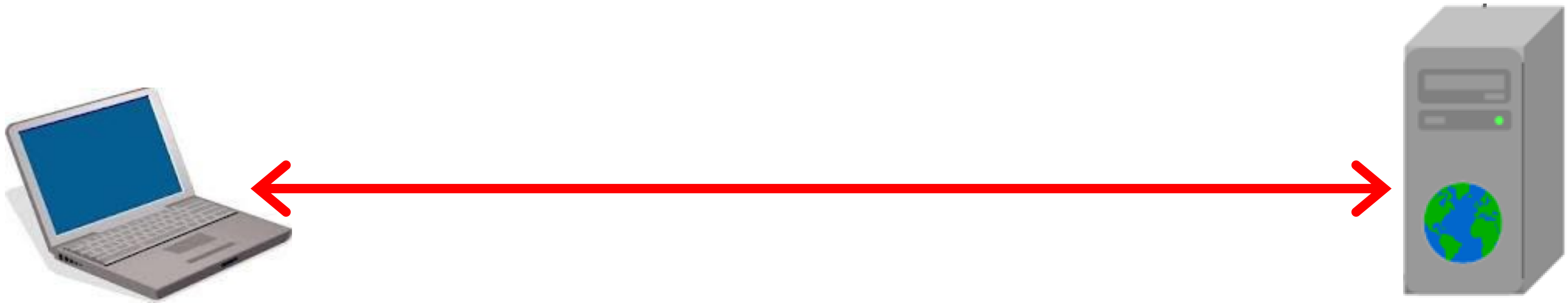
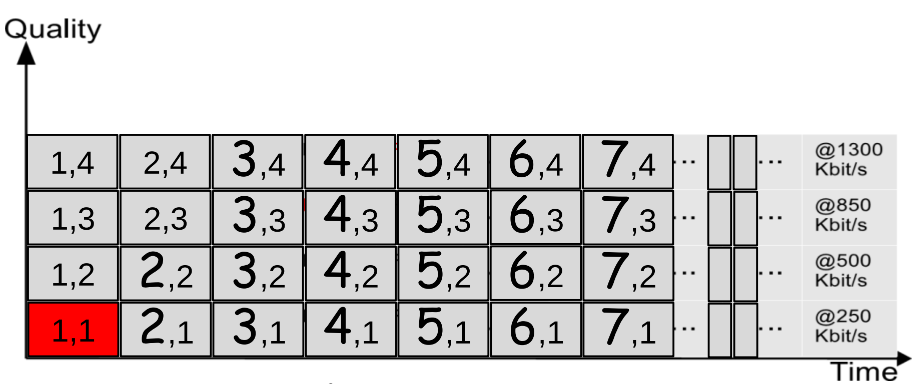


Server

Example





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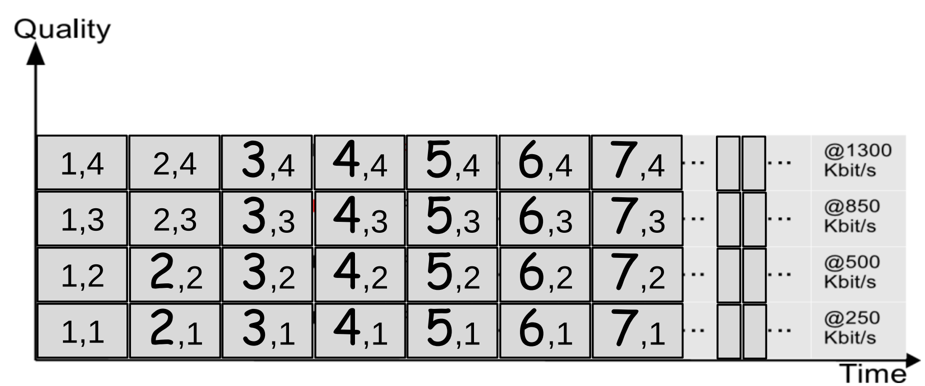

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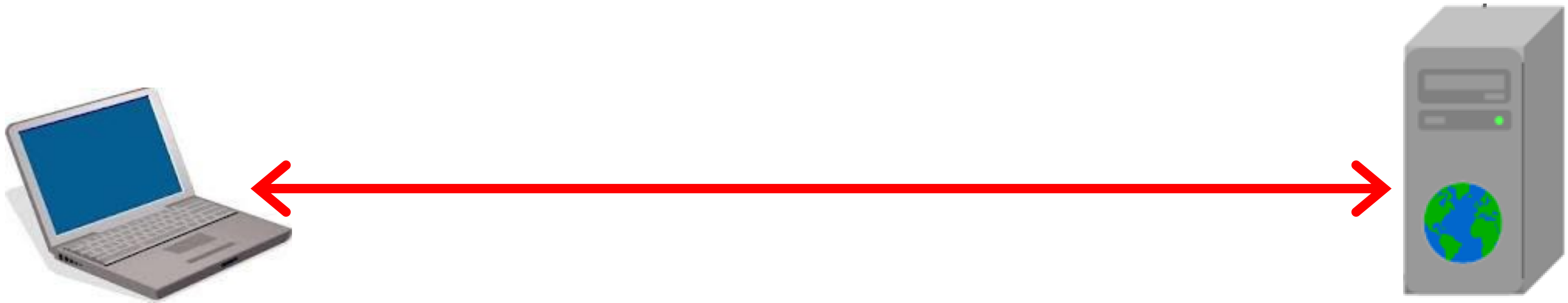

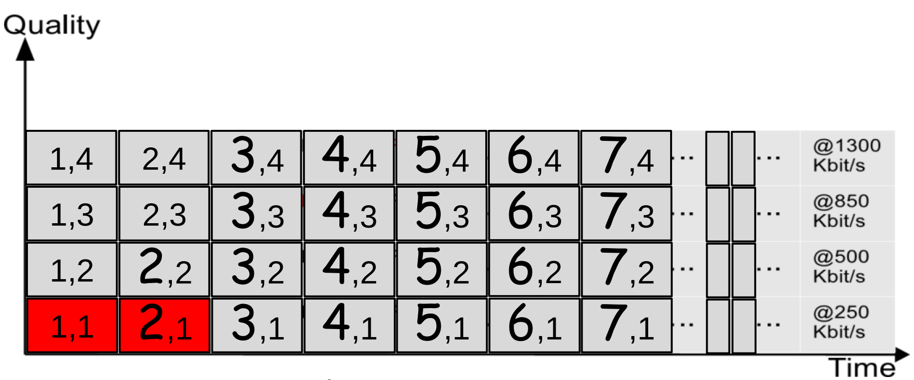

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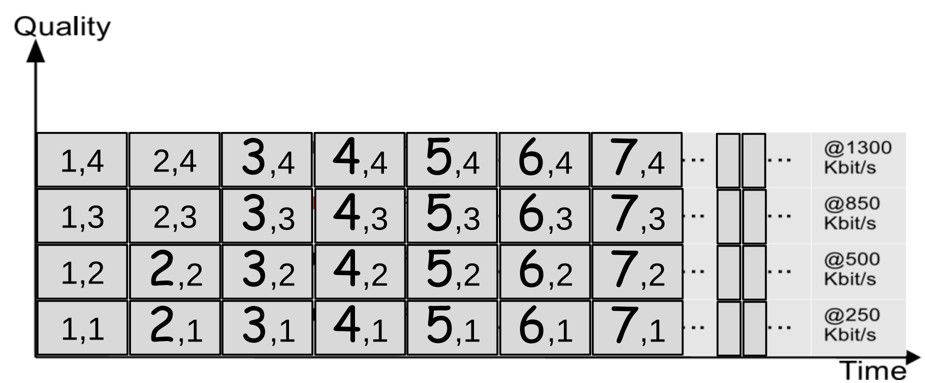

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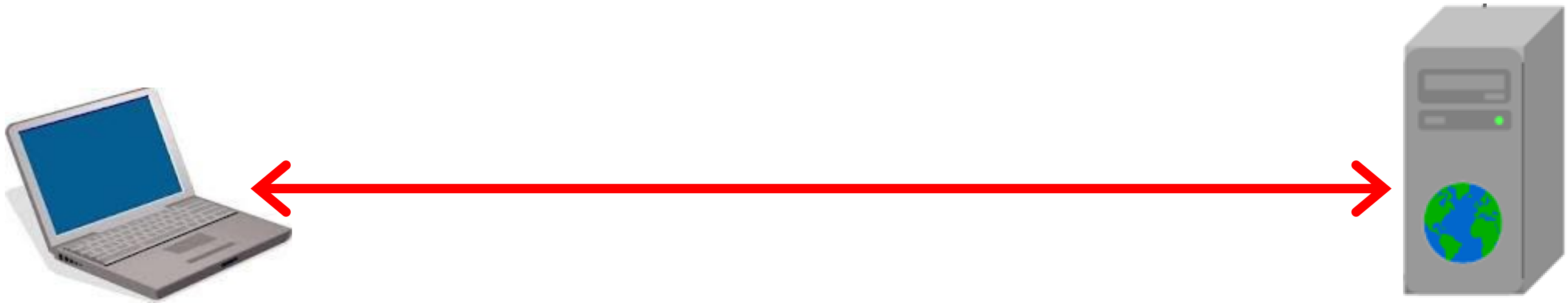

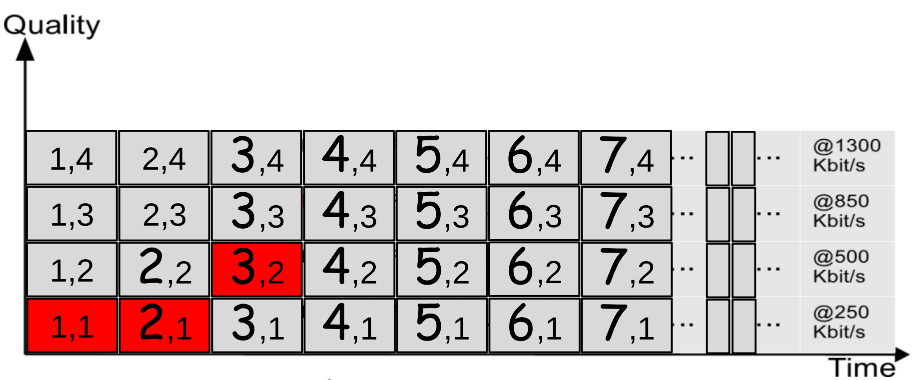

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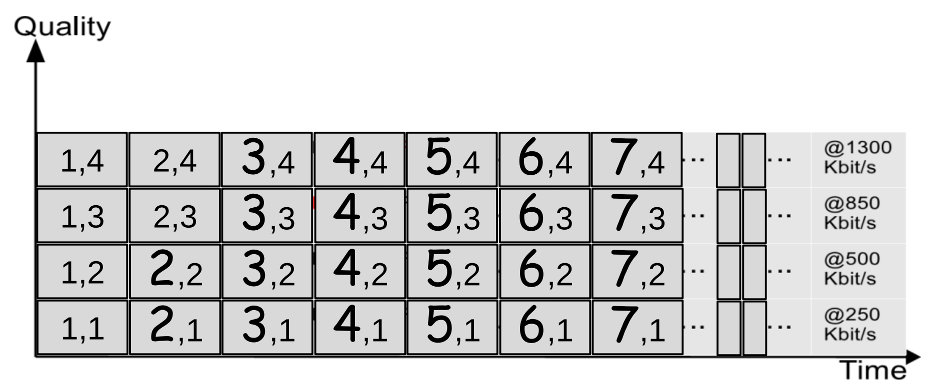

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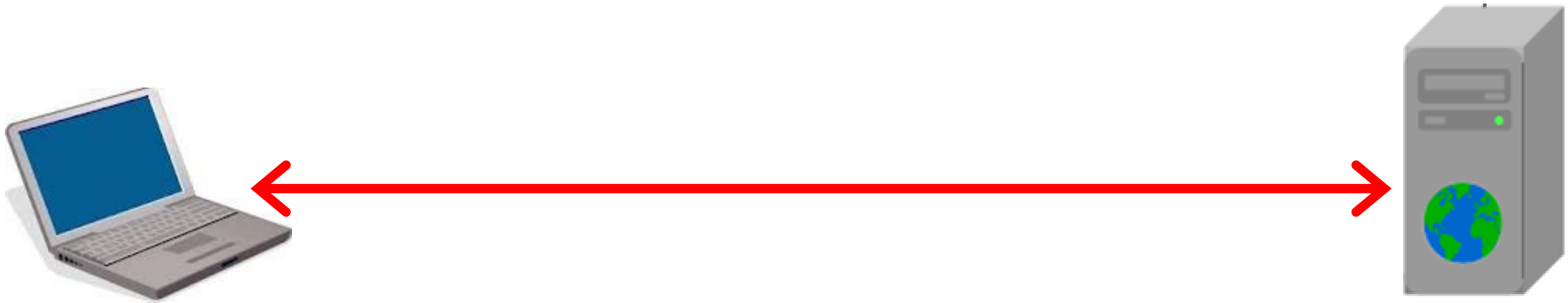

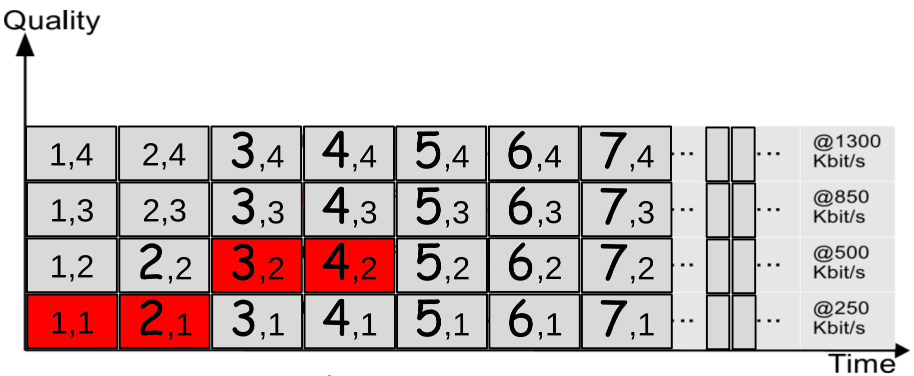

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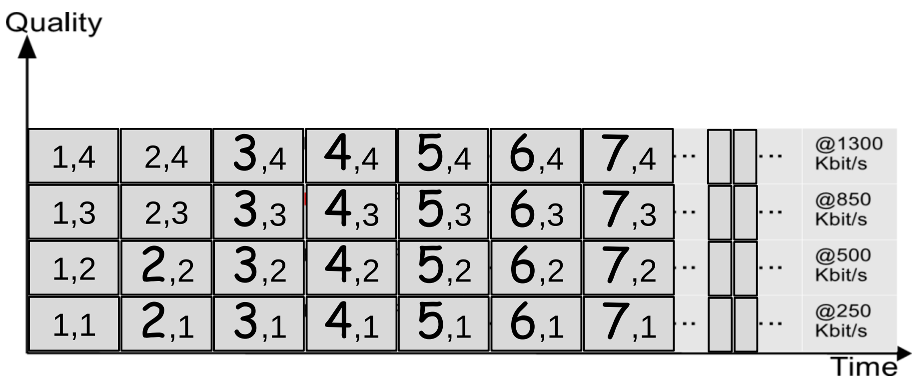

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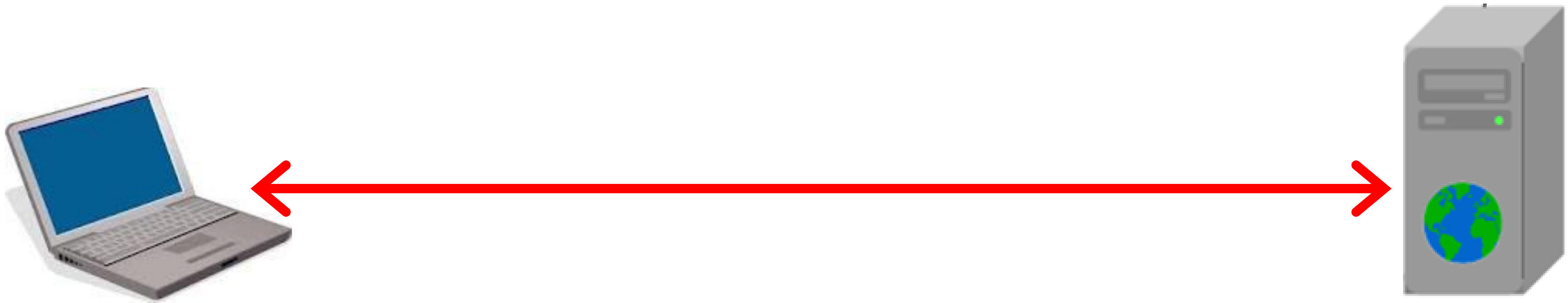

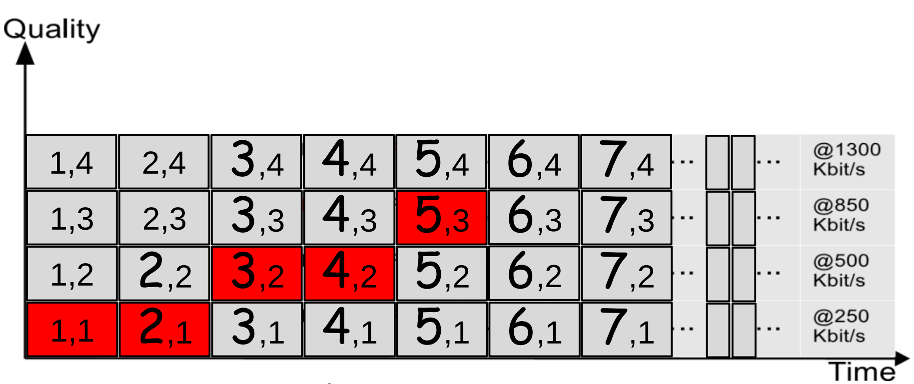

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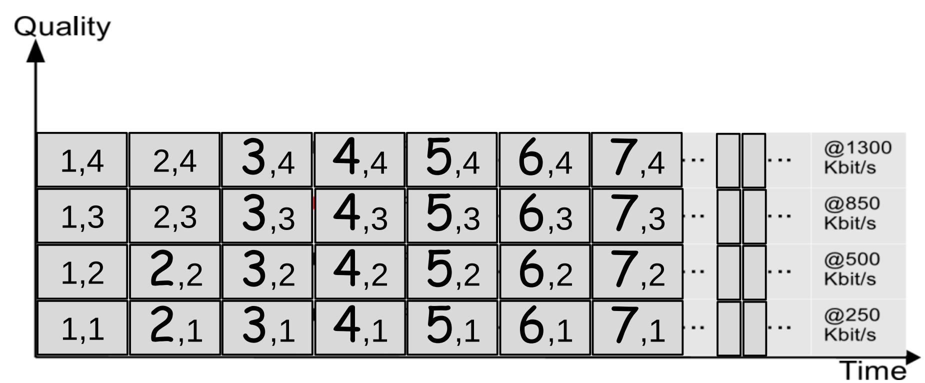

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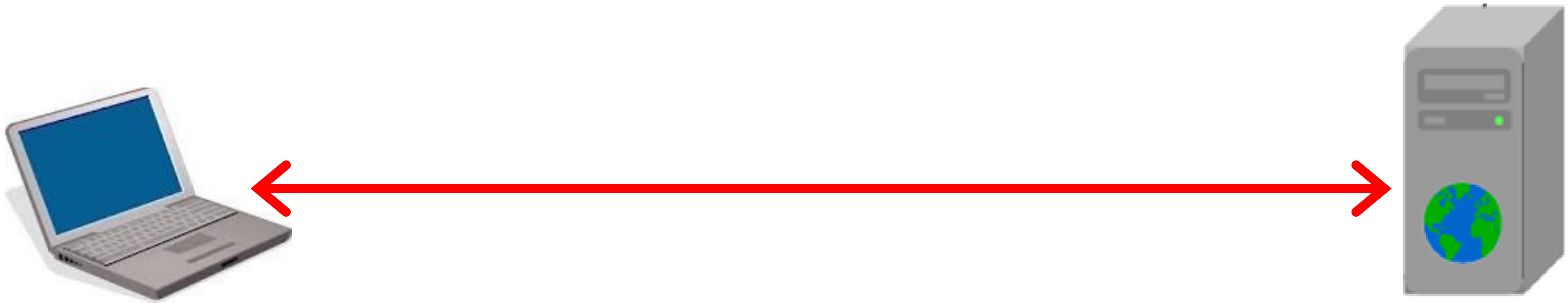

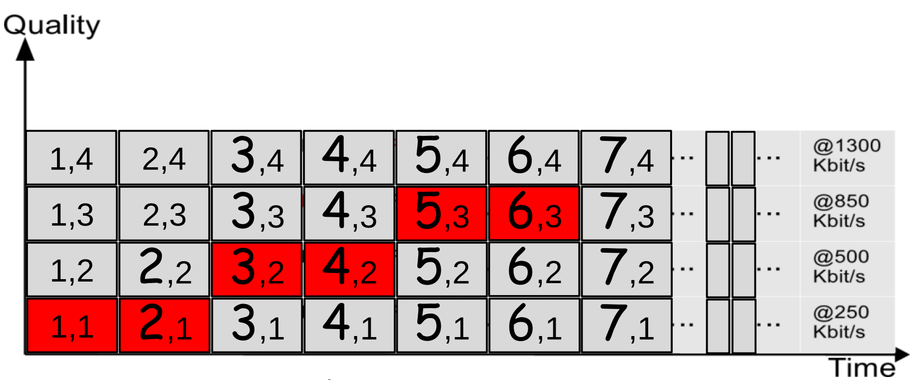

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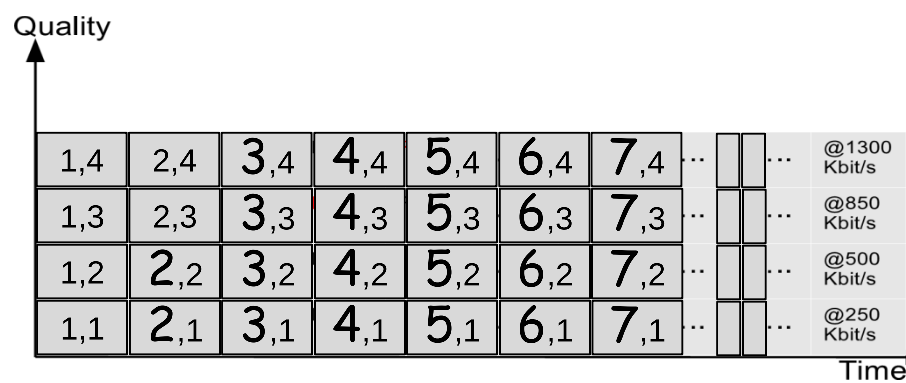

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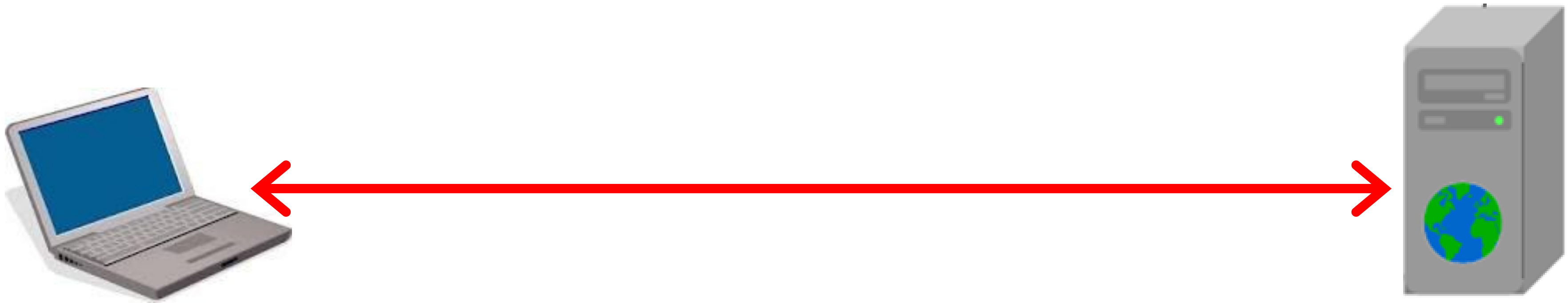

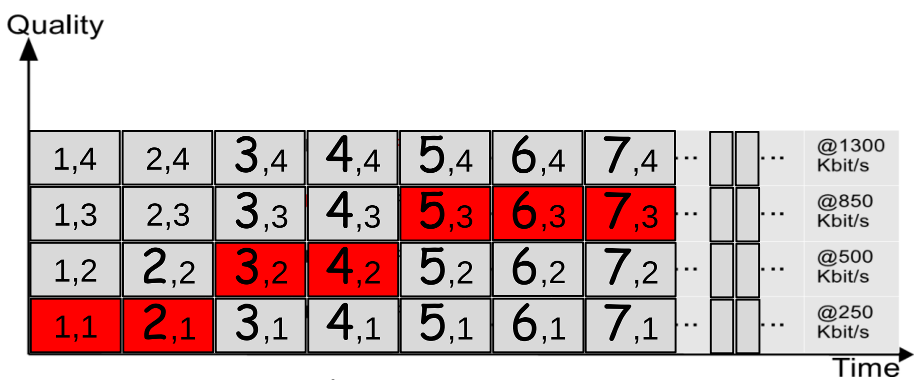

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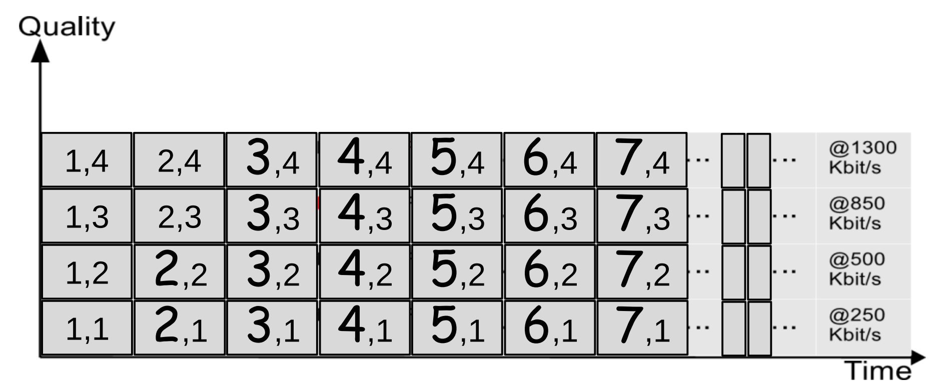

Server

Chunk-based streaming

- Chunks begin with keyframe so independent of other chunks
- Playing chunks in sequence gives seamless video
- Hybrid of streaming and progressive download:
  - Stream-like: sequence of small chunks requested as needed
  - Progressive download-like: HTTP transfer mechanism, stateless servers

HTTP-based Adaptive Streaming (HAS)

- Other terms for similar concepts: Adaptive Streaming, Smooth Streaming, HTTP Chunking, DASH
- Actually a series of small progressive downloads of chunks (or range requests)
- No standard protocol ...

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- Actually a series of small progressive downloads of chunks (or range requests)
- No standard protocol ...
  - Apple HLS: HTTP Live Streaming
  - Microsoft IIS Smooth Streaming: part of Silverlight
  - Adobe: Flash Dynamic Streaming
  - DASH: Dynamic Adaptive Streaming over HTTP

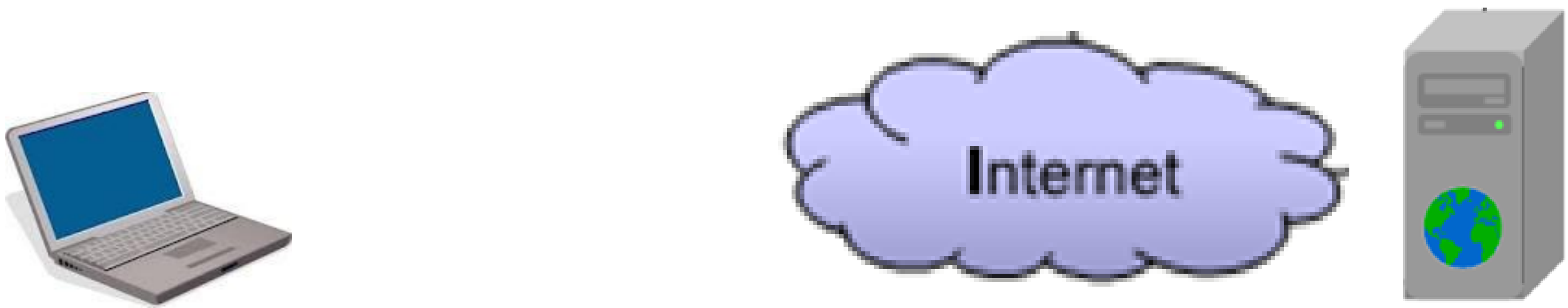
Example players

	Player	Container	Type	Open Source
Microsoft Silverlight	Microsoft Smooth Streaming	Silverlight	Chunk	✗
NETFLIX	Netflix player	Silverlight	Range	✗
	Apple HLS	QuickTime	Chunk	✗
	Adobe HDS	Flash	Chunk	✓

.. and now YouTube (Google) is also using HAS!

YouTube

Example: HAS and proxy

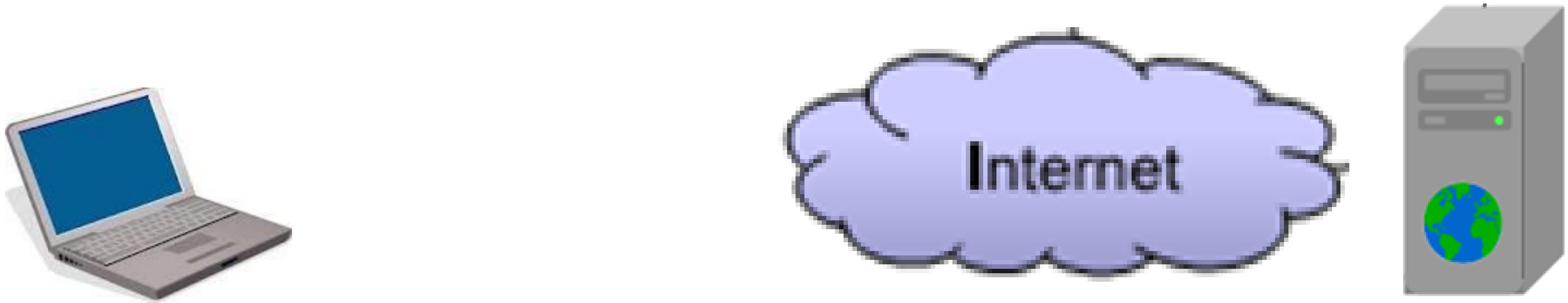


Clients' want

- High playback quality
- Small stall times
- Few buffer interruptions
- Few quality switches

Example: HAS and proxy

Slides from: V. Krishnamoorthi et al. "Helping Hand or Hidden Hurdle: Proxy-assisted HTTP-based Adaptive Streaming Performance", Proc. IEEE MASCOTS, 2013



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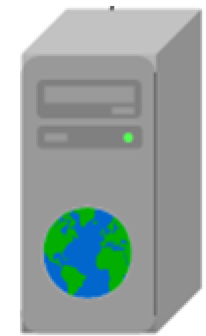
HAS is increasingly responsible for larger traffic volumes
… proxies to reduce traffic??

Example: HAS and proxy

Slides from: V. Krishnamoorthi et al. "Helping Hand or Hidden Hurdle: Proxy-assisted HTTP-based Adaptive Streaming Performance", Proc. IEEE MASCOTS, 2013



Clients' want

- High playback quality
- Small stall times
- Few buffer interruptions
- Few quality switches

Network providers' want

- High QoE of customers/clients

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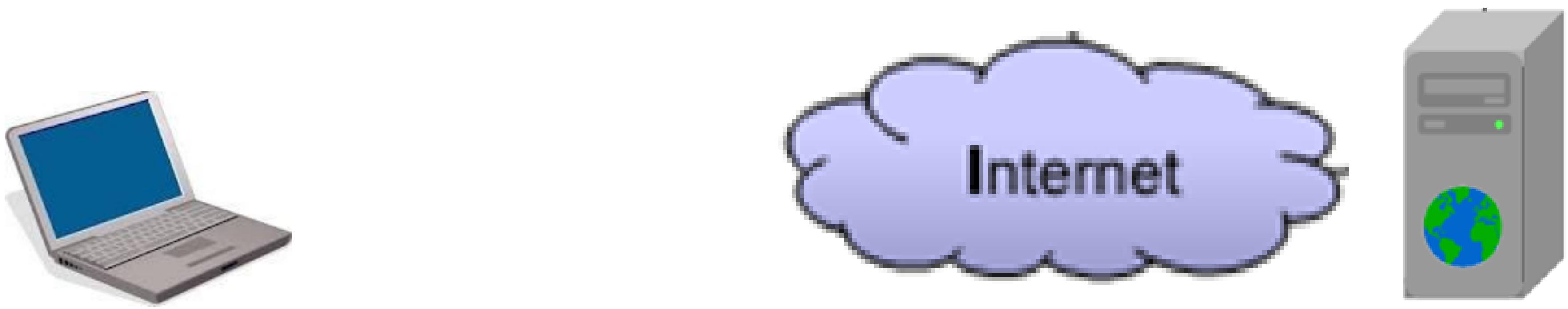

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Network providers' want

- High QoE of customers/clients
- Low bandwidth usage
- High hit rate

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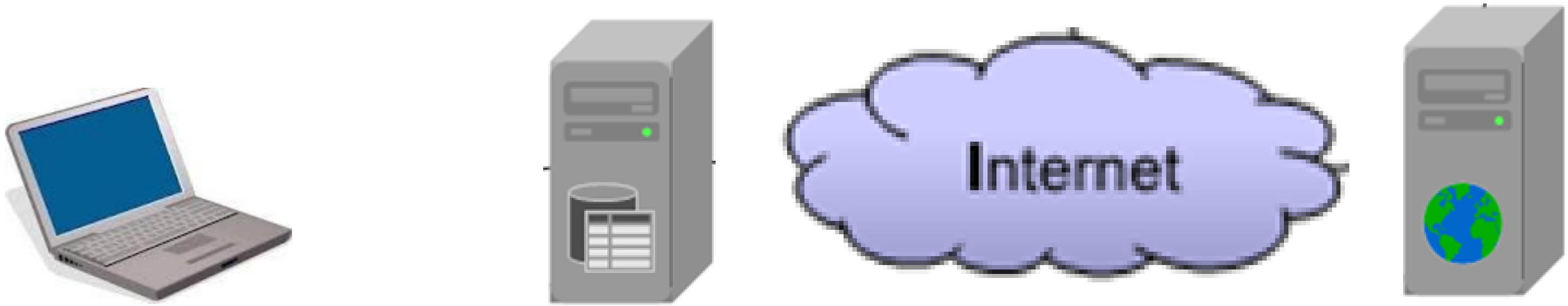

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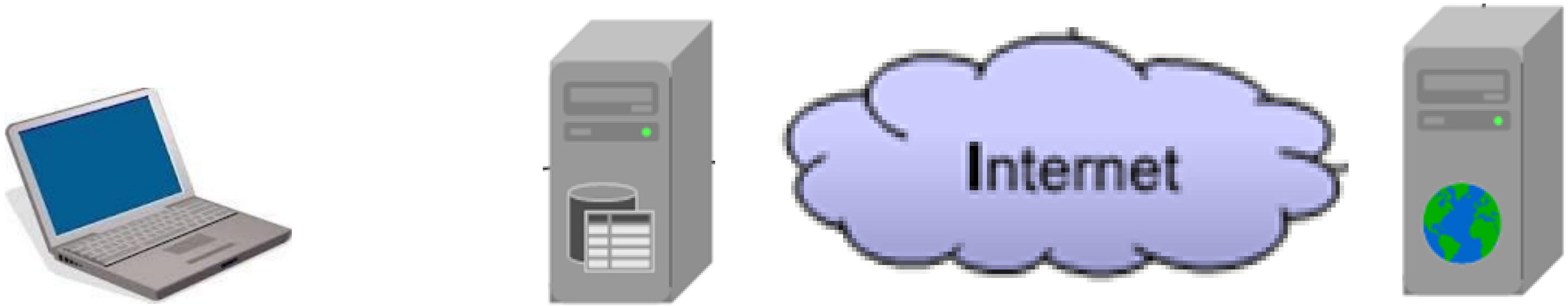
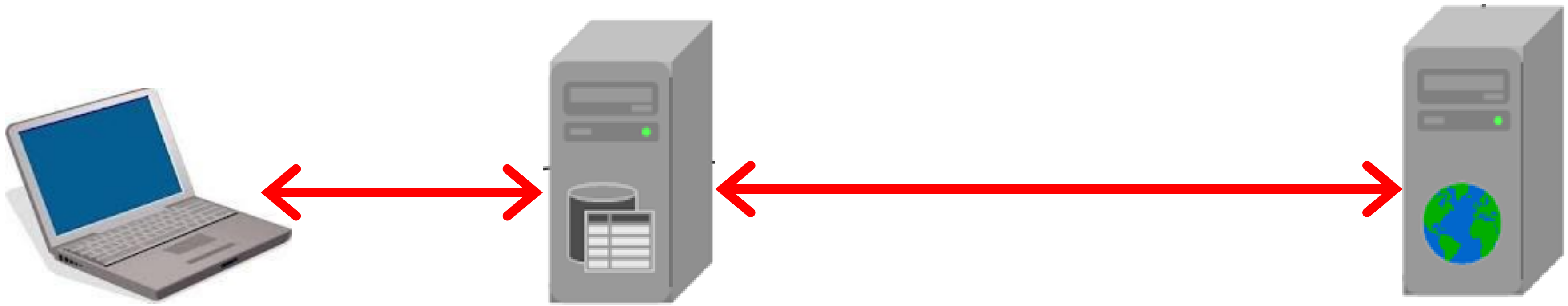


Proxy example …

Example: HAS and proxy



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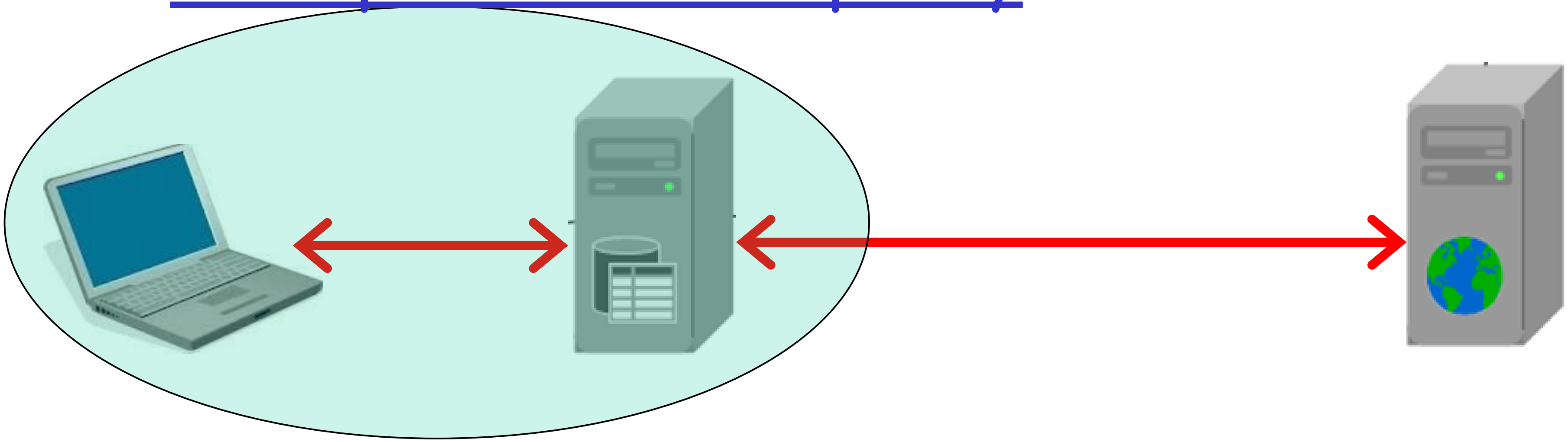

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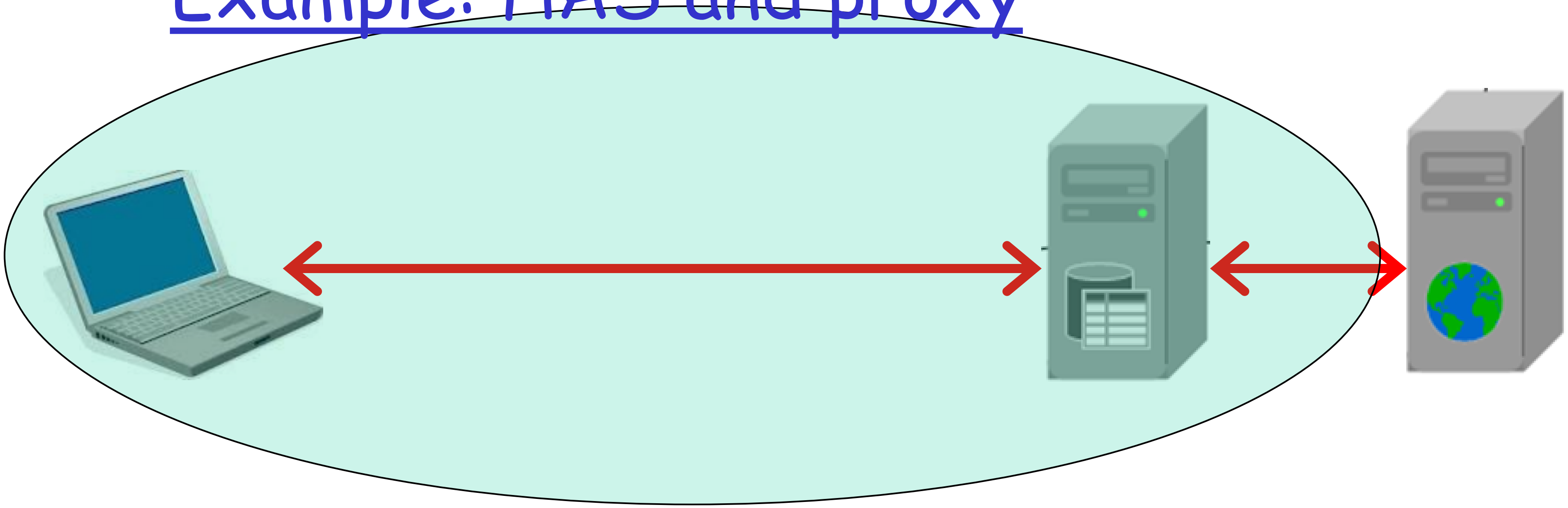

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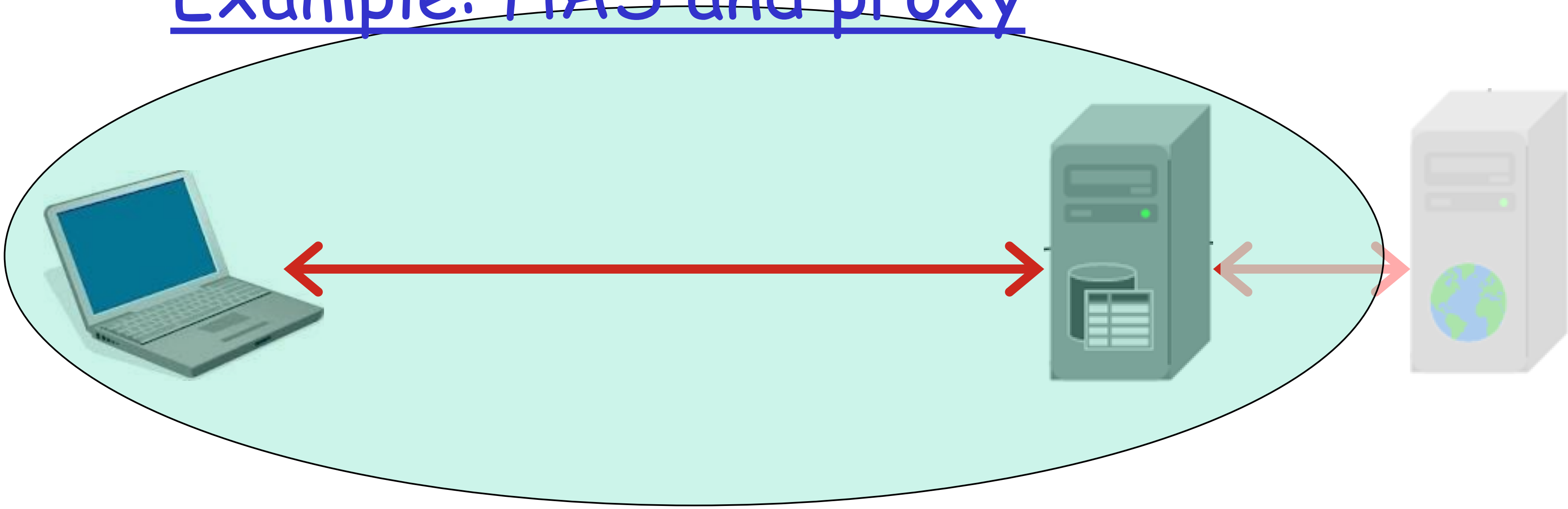

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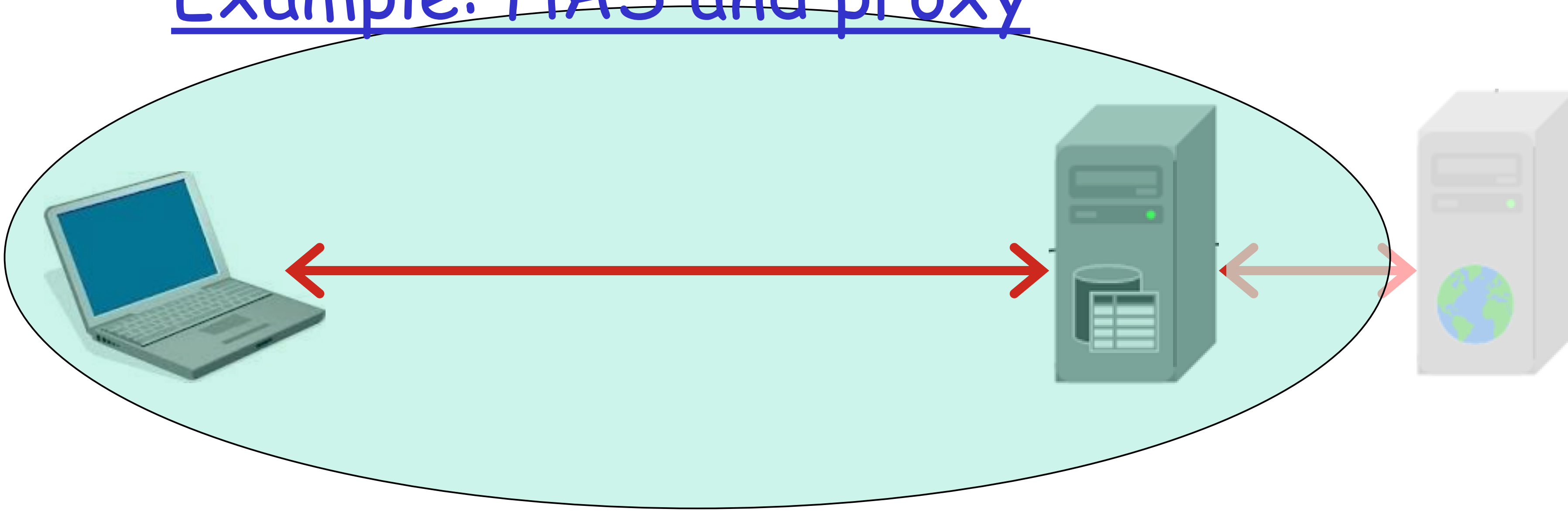
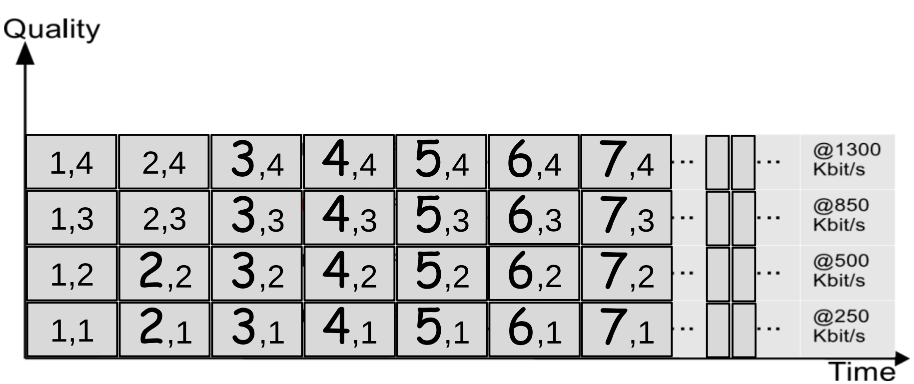
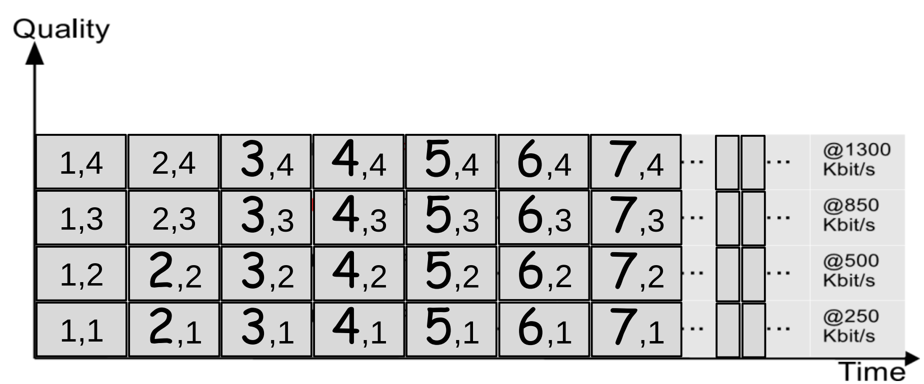


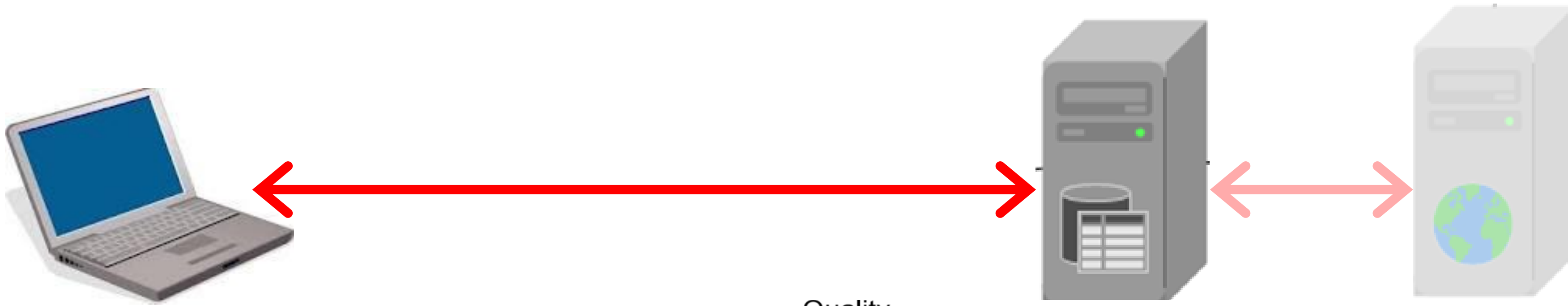


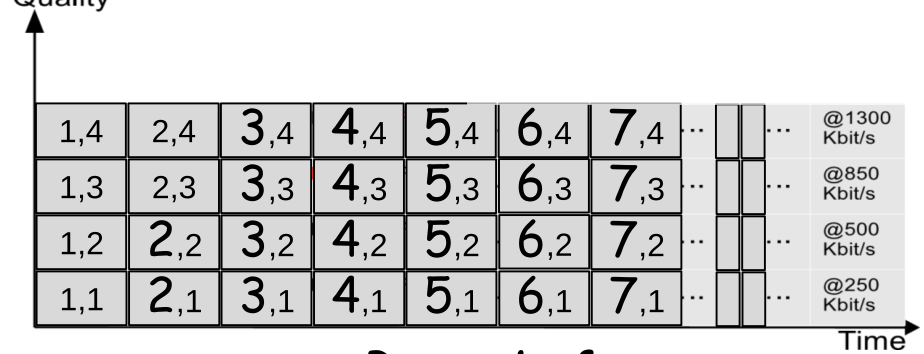
Client



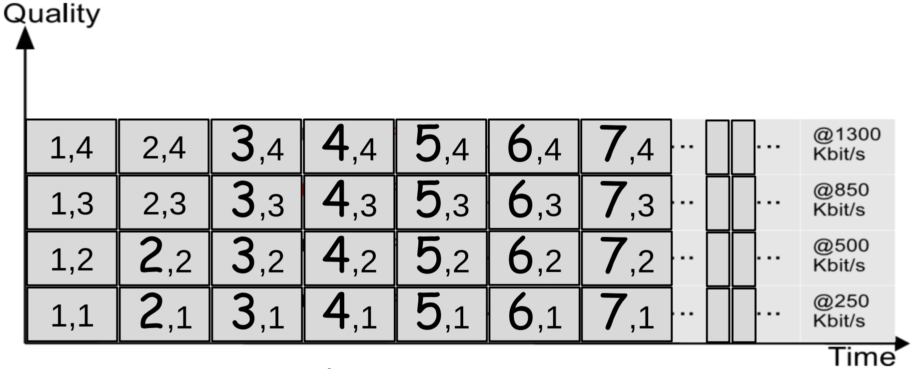
Proxy

Example: HAS and proxy

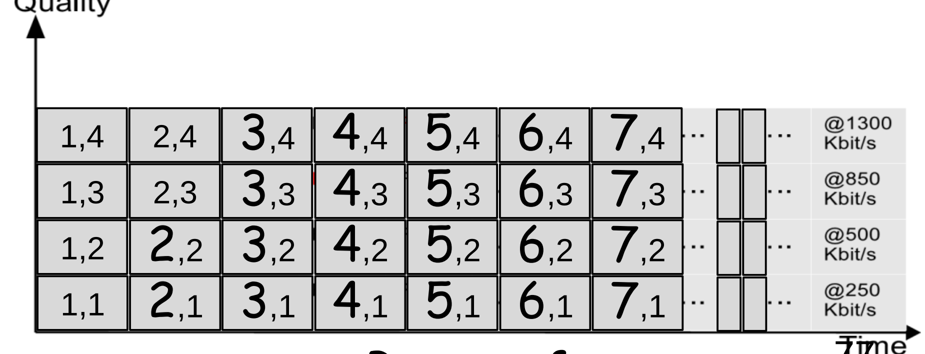




Proxy before

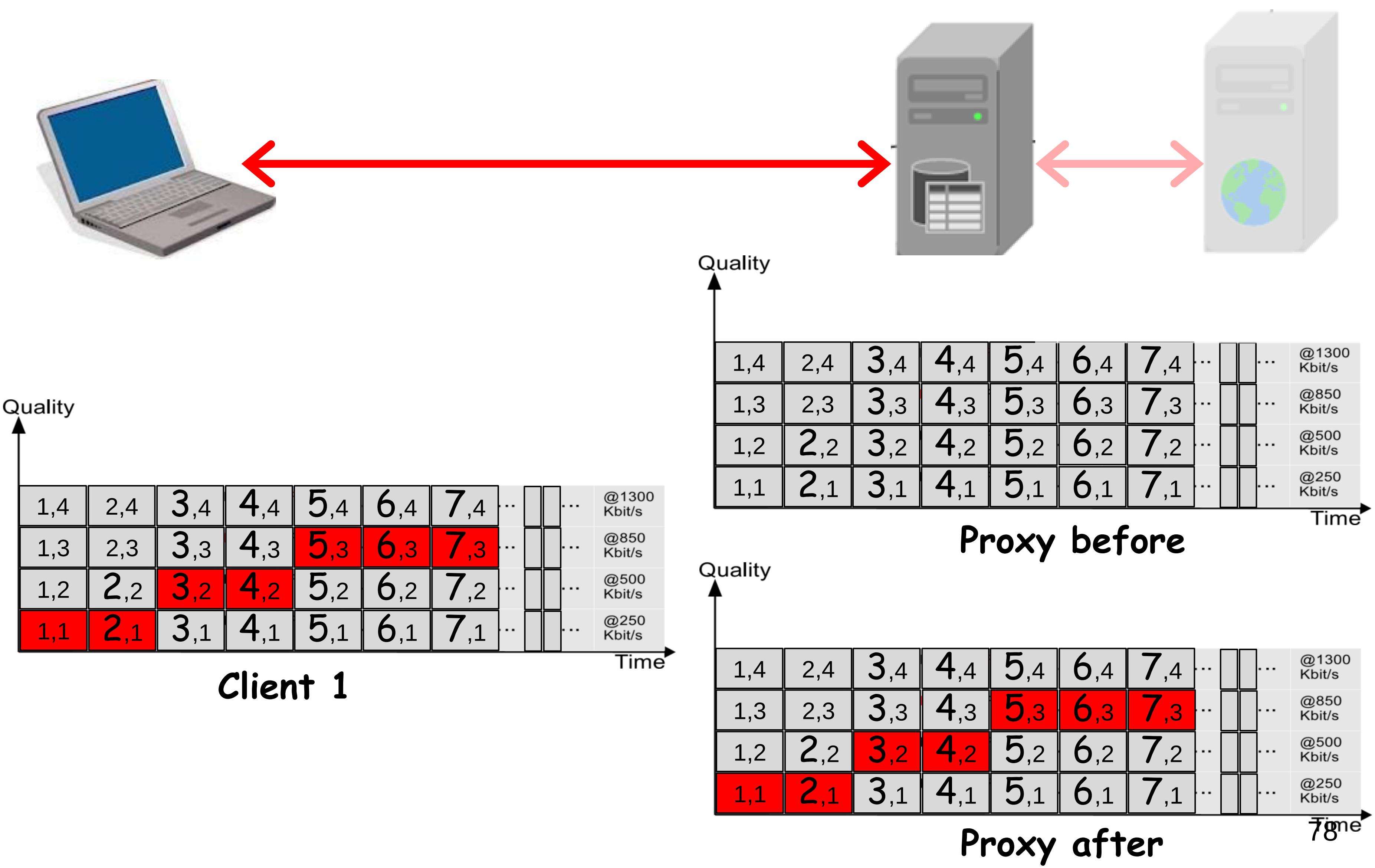


Client 1



Proxy after

Example: HAS and proxy



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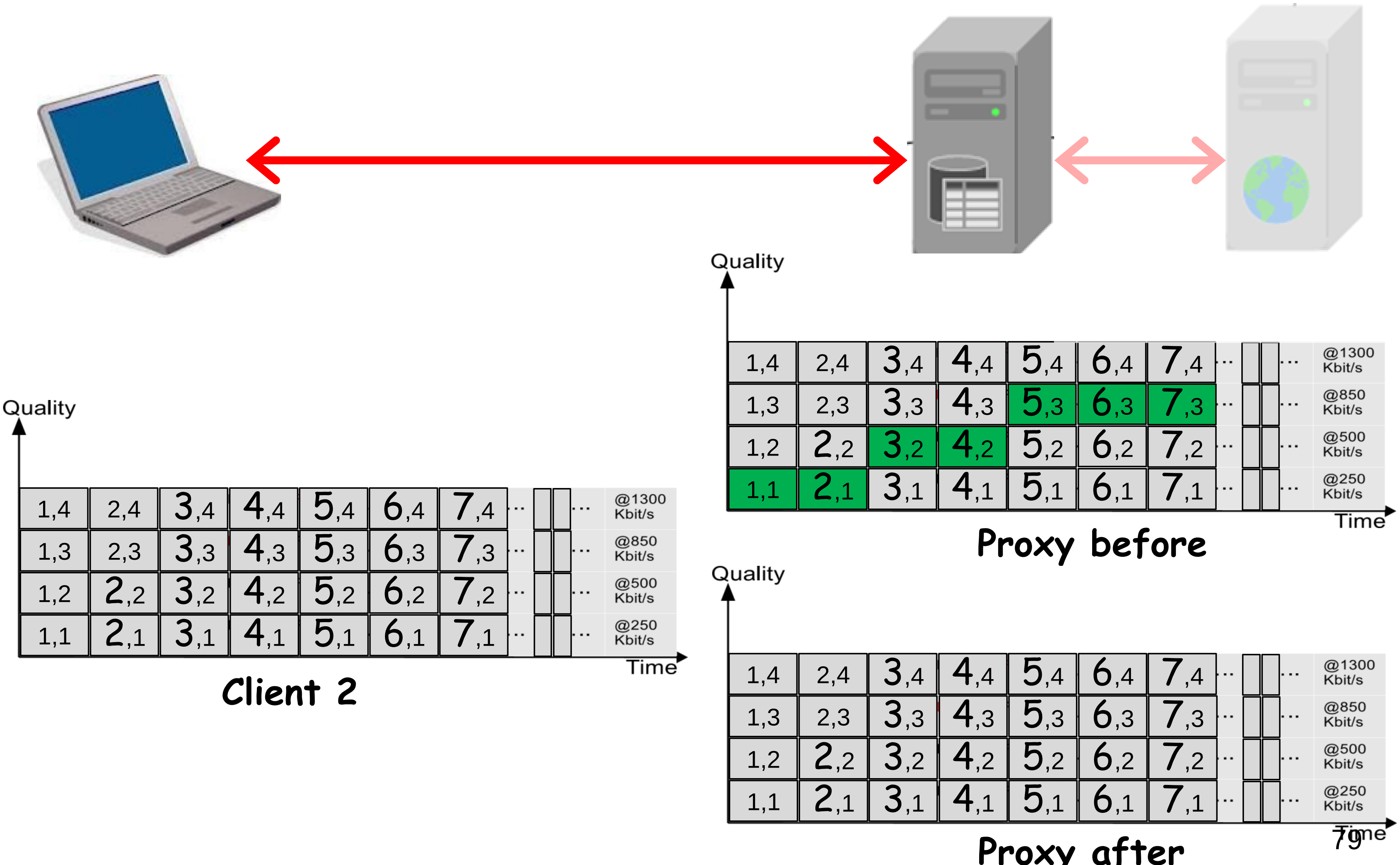

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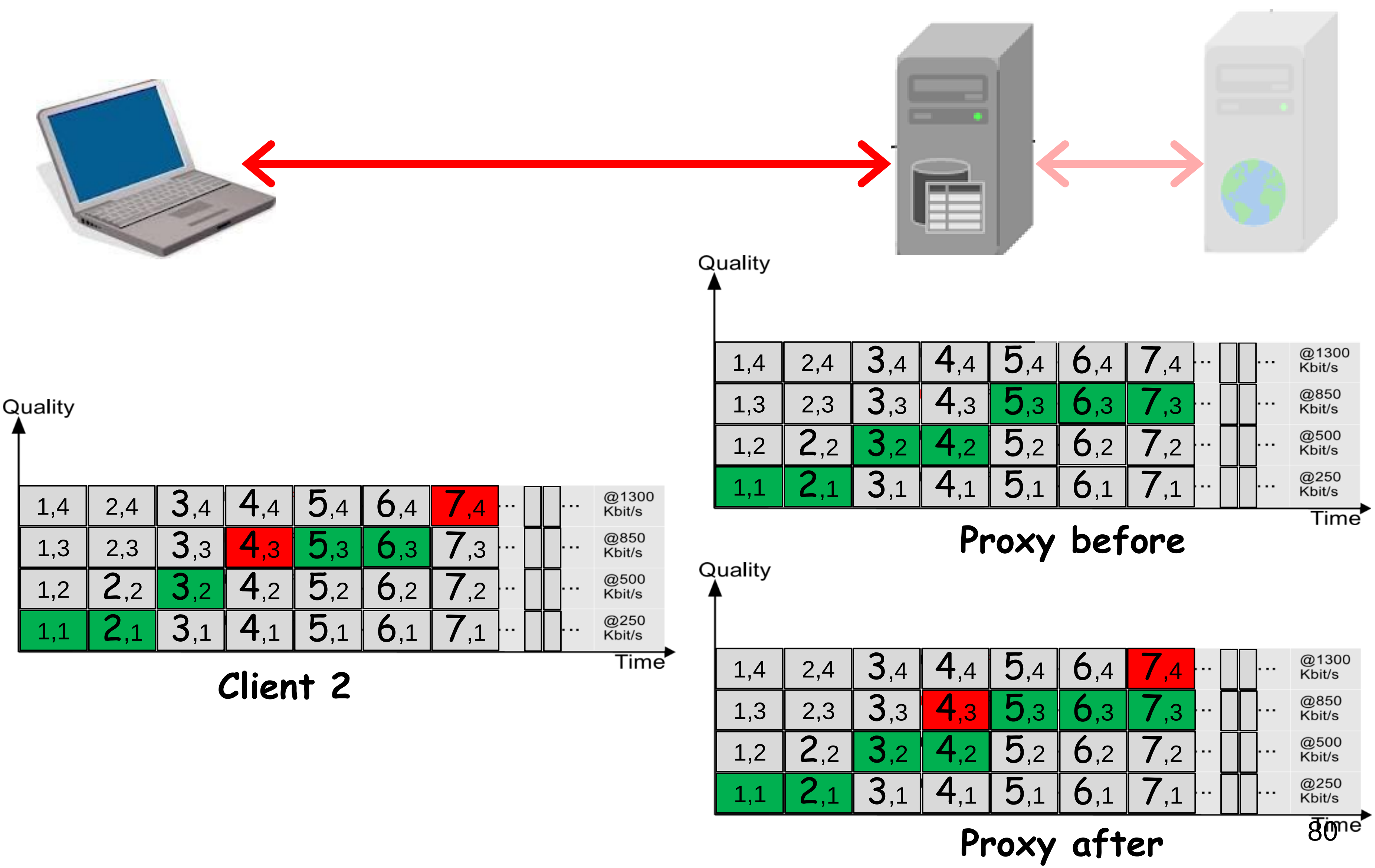

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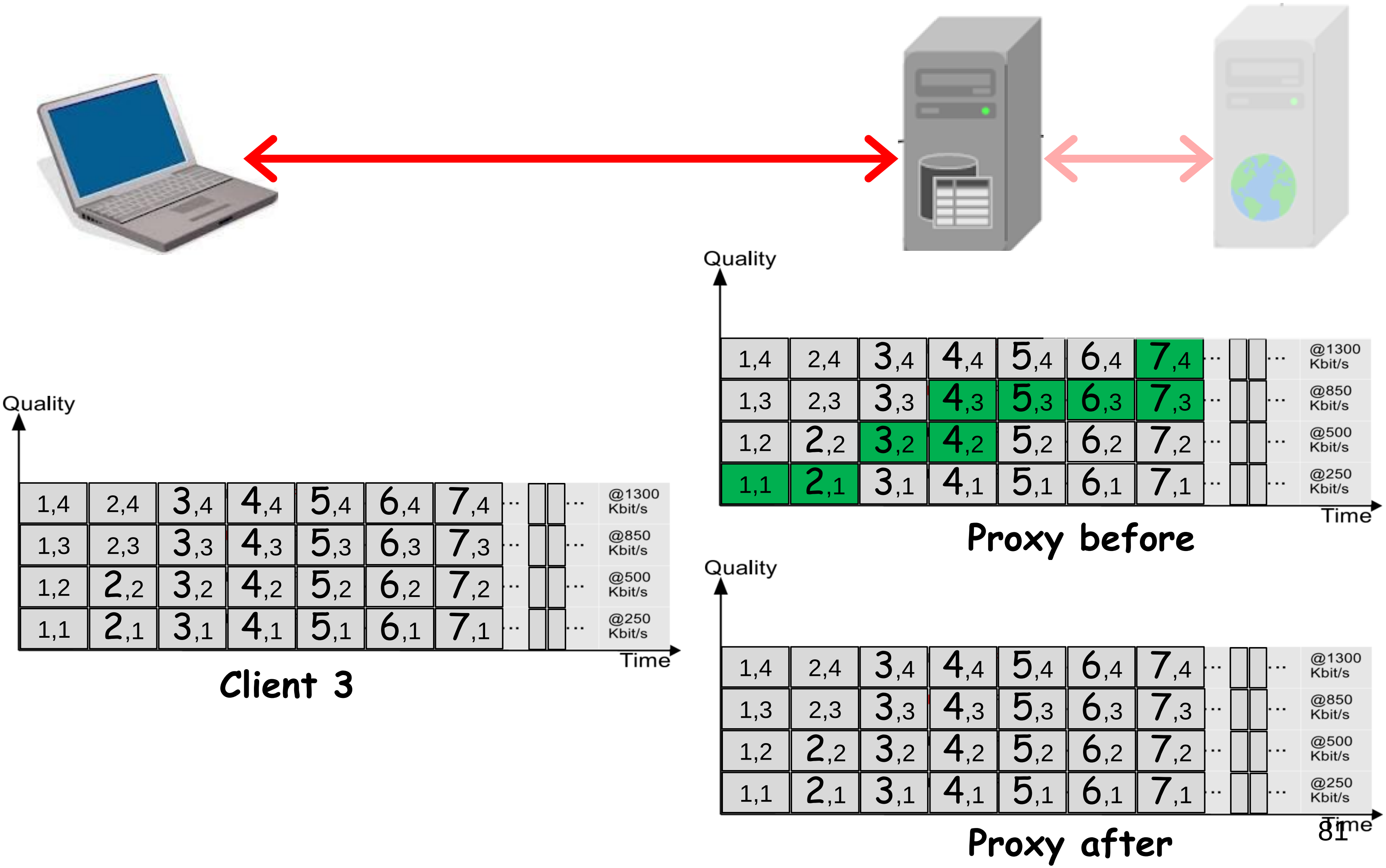

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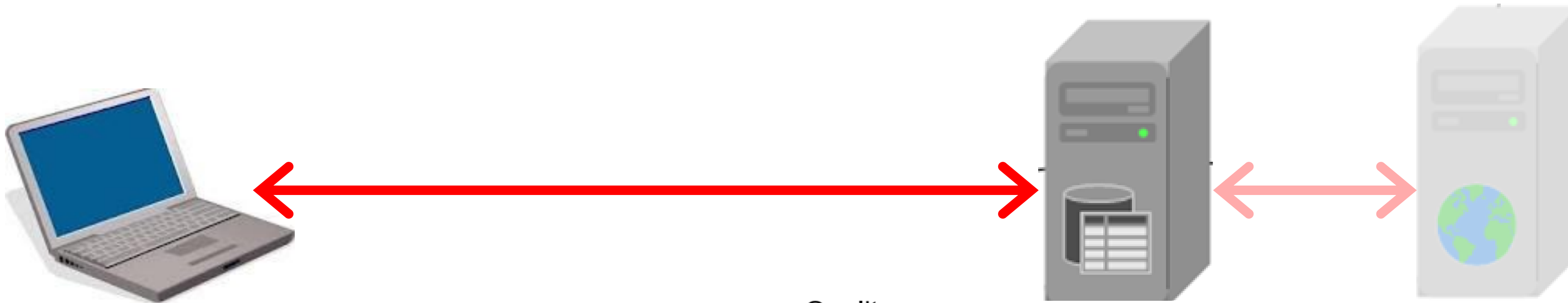
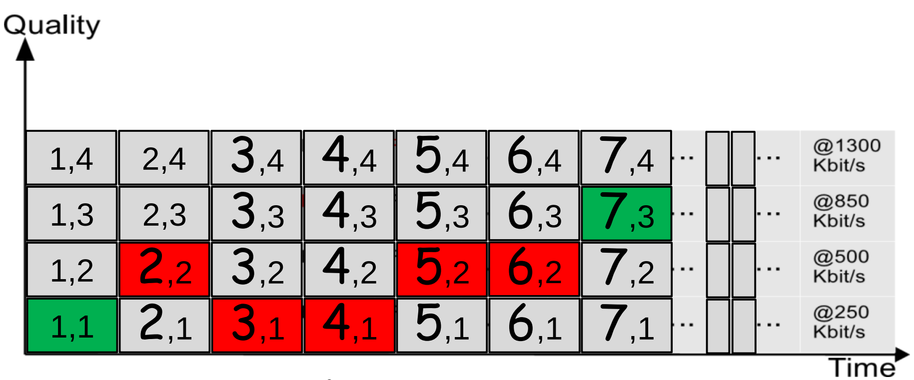
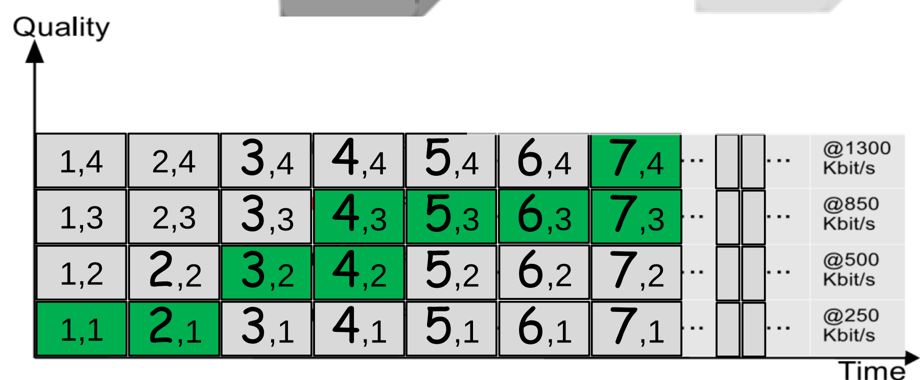
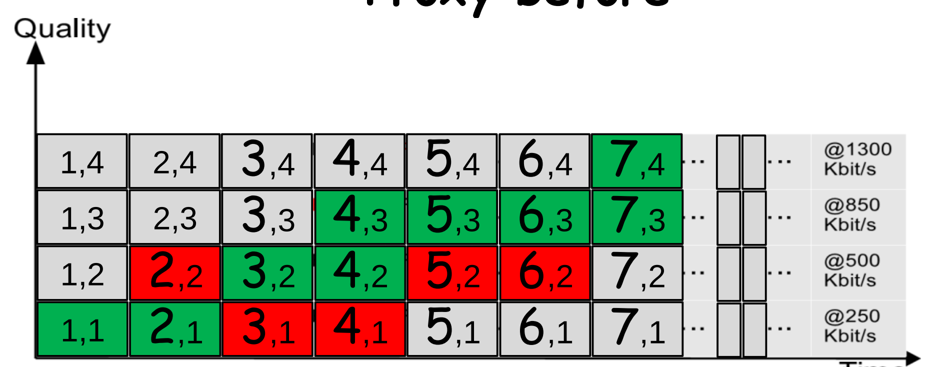




Client 3



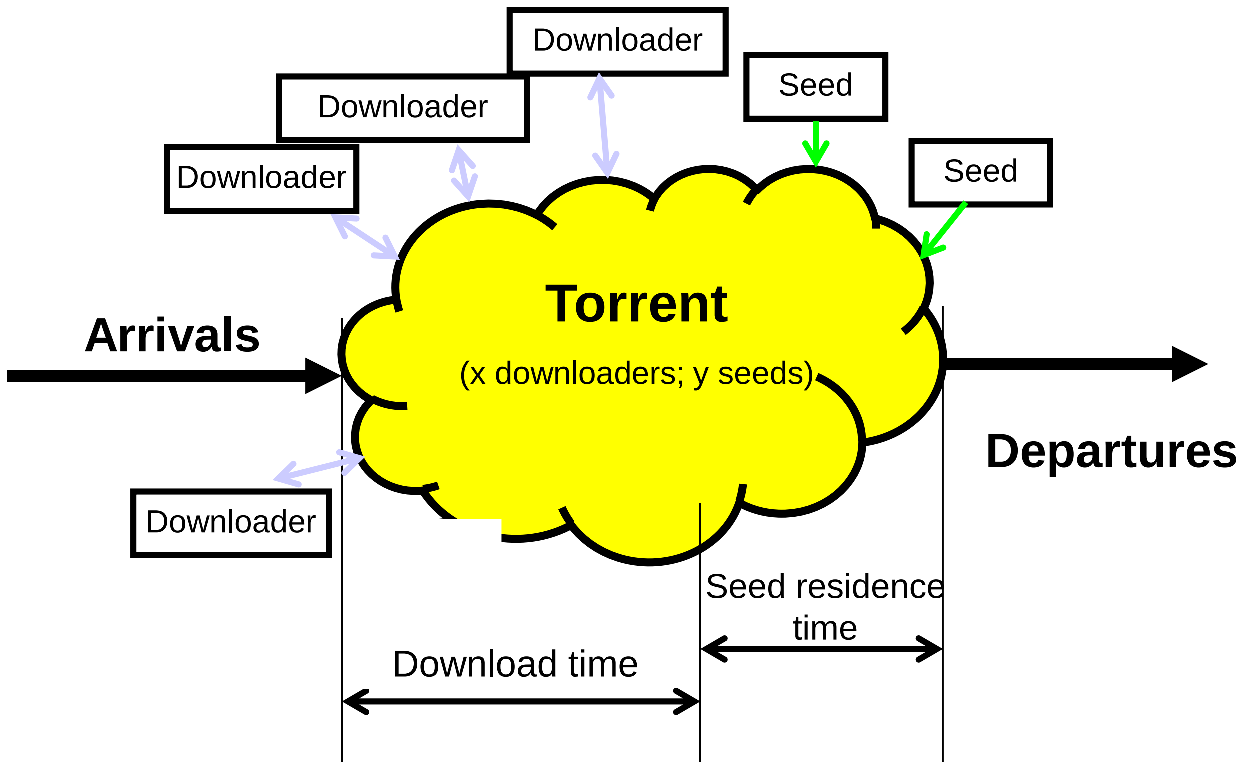
Proxy before



Proxy after

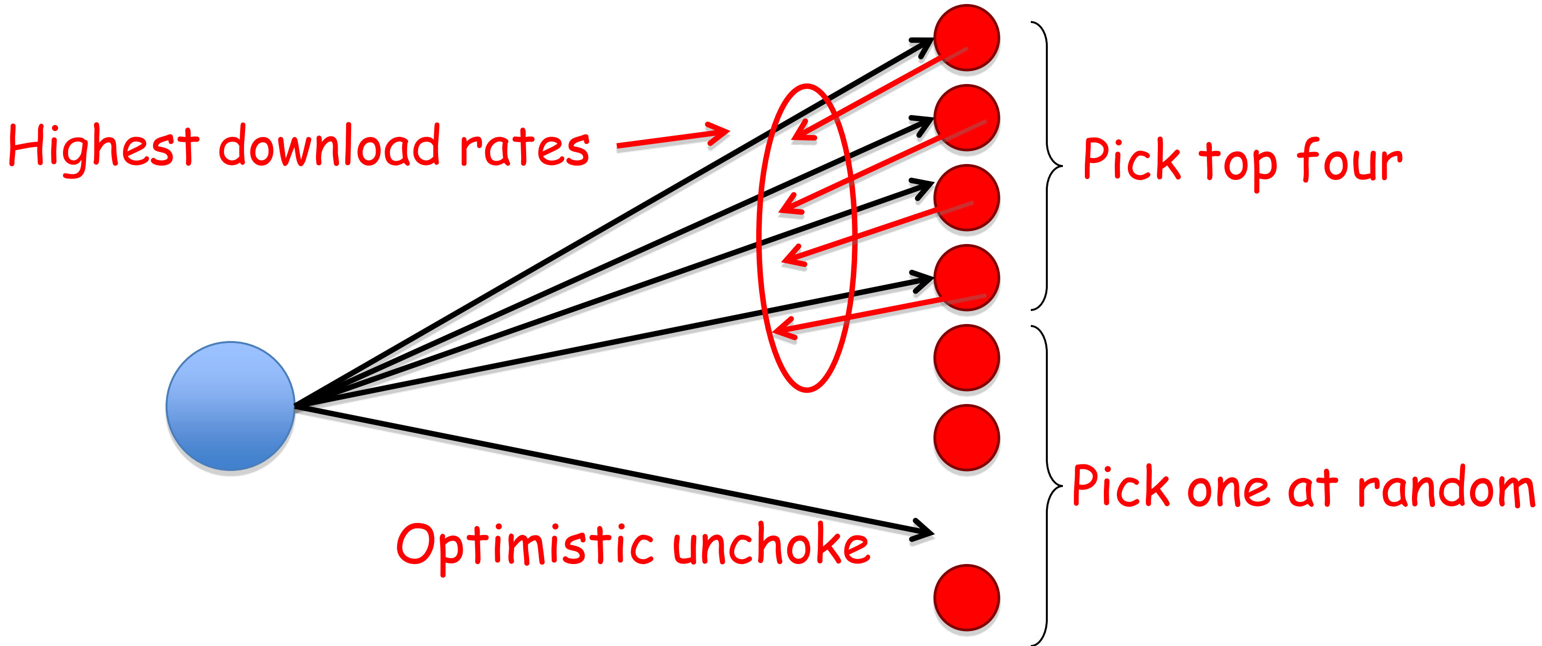
BitTorrent-like systems

- File split into many smaller pieces
- Pieces are downloaded from both seeds and downloaders
- Distribution paths are dynamically determined
    - Based on data availability

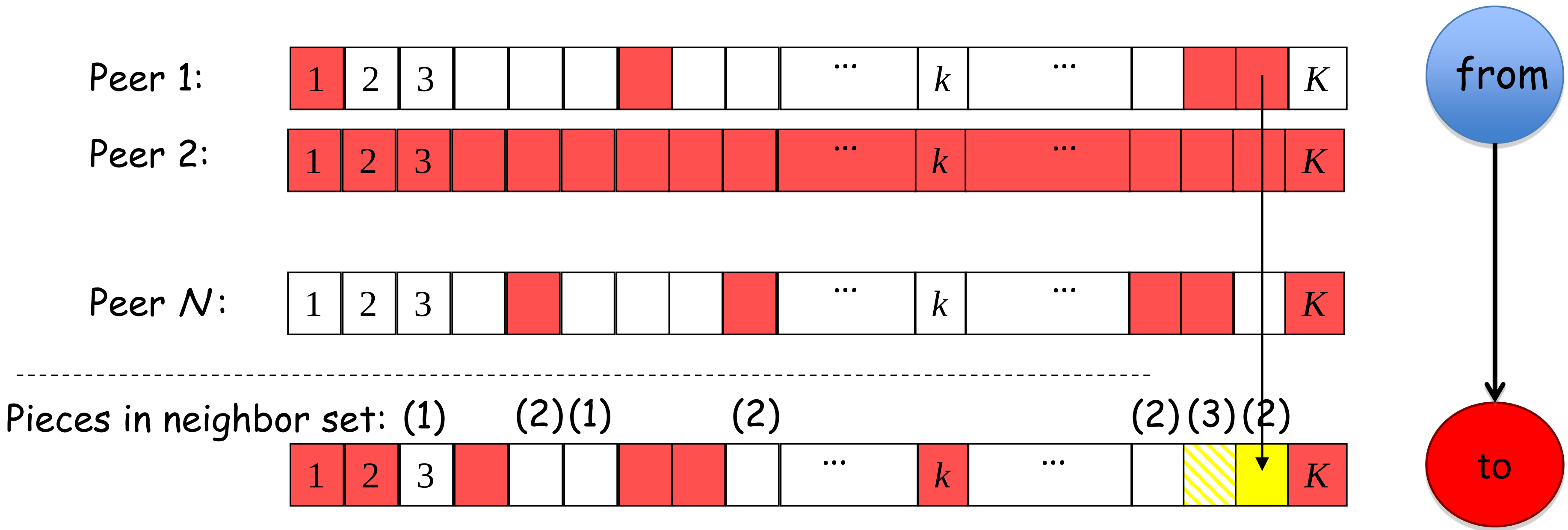


Download using BitTorrent
Background: Incentive mechanism

- Establish connections to large set of peers
  - At each time, only upload to a small (changing) set of peers
- Rate-based tit-for-tat policy
  - Downloaders give upload preference to the downloaders that provide the highest download rates

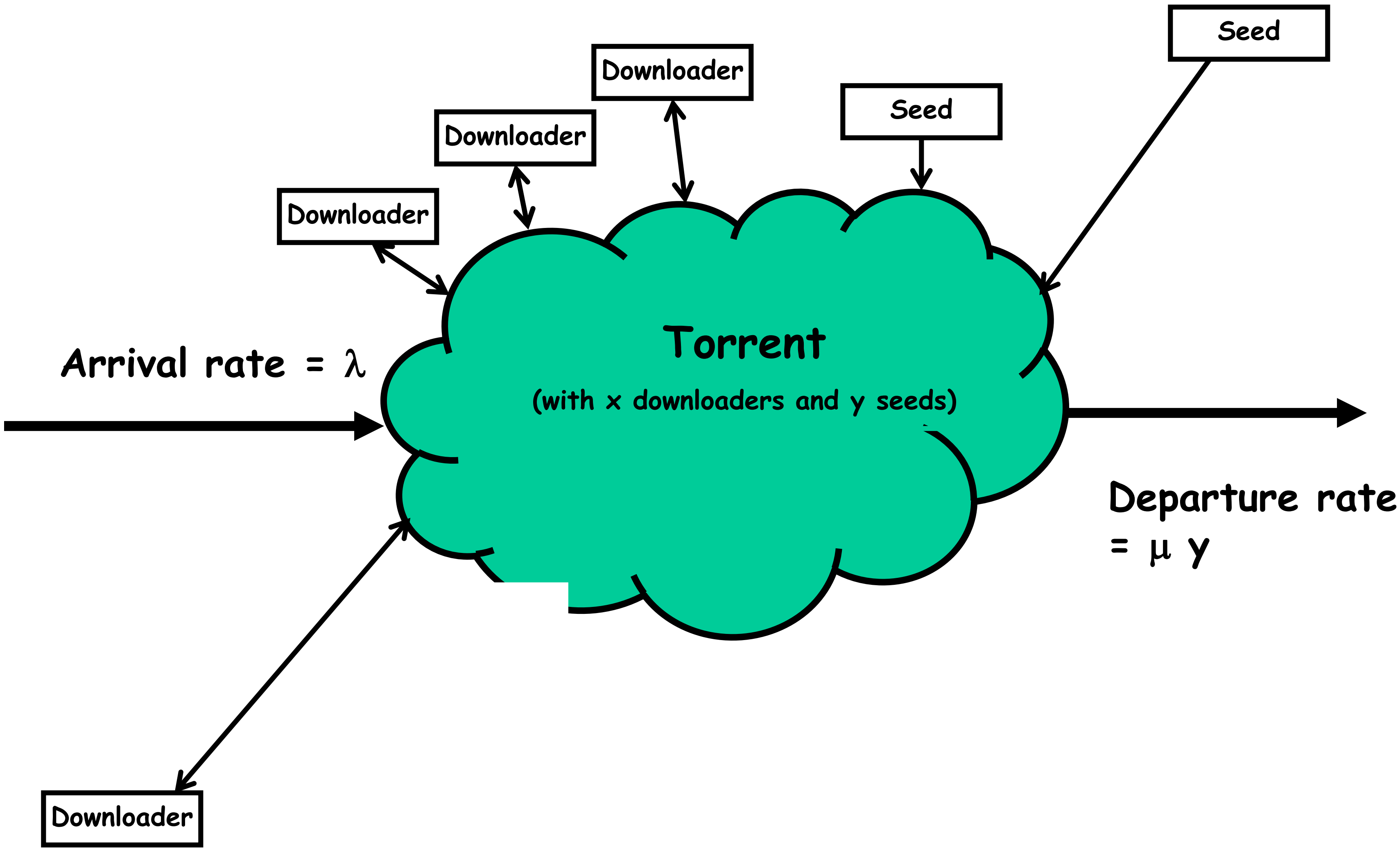


Download using BitTorrent Background: Piece selection

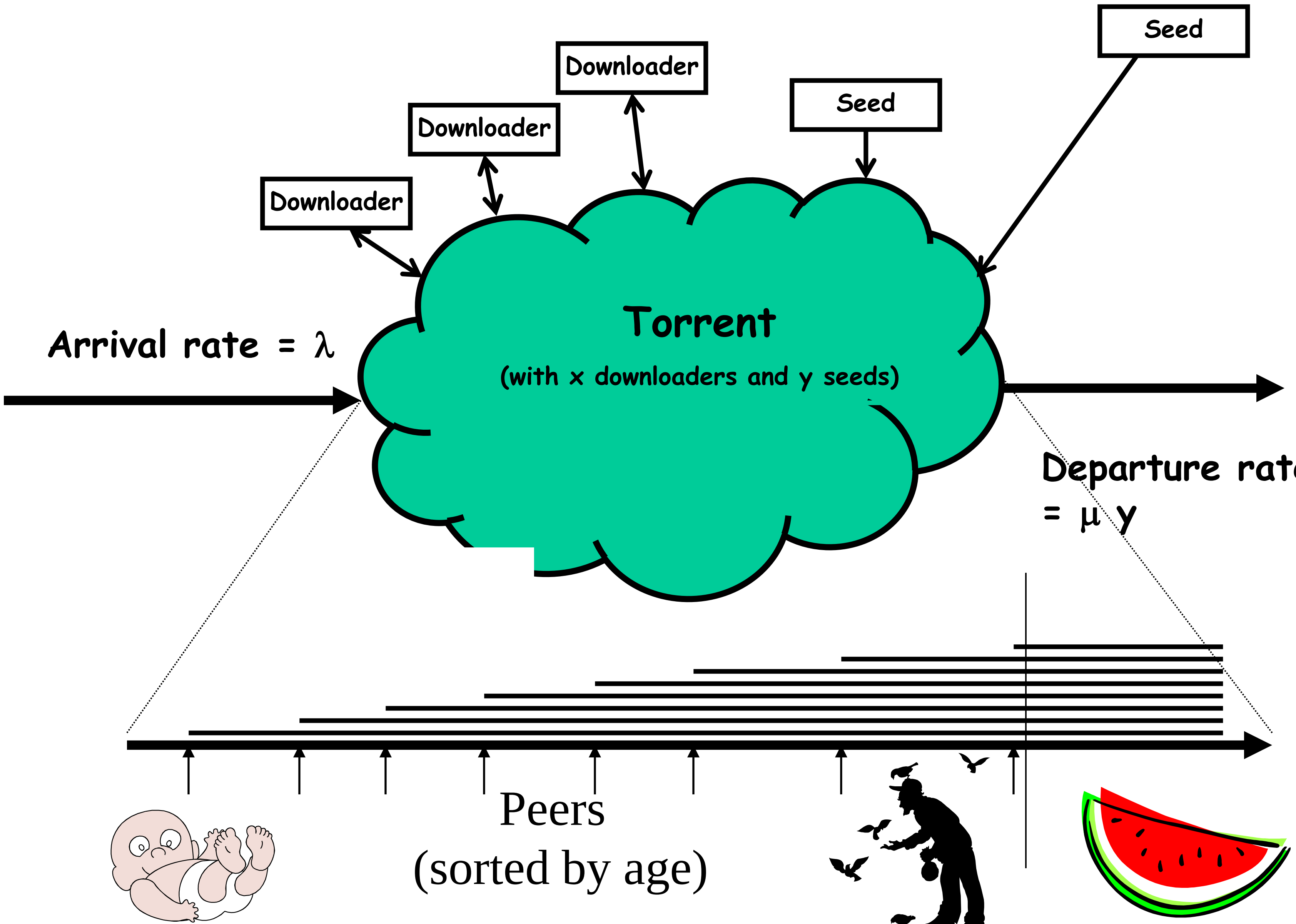


- Rarest first piece selection policy
  - Achieves high piece diversity
- Request pieces that
  - the uploader has;
  - the downloader is interested (wants); and
  - is the rarest among this set of pieces

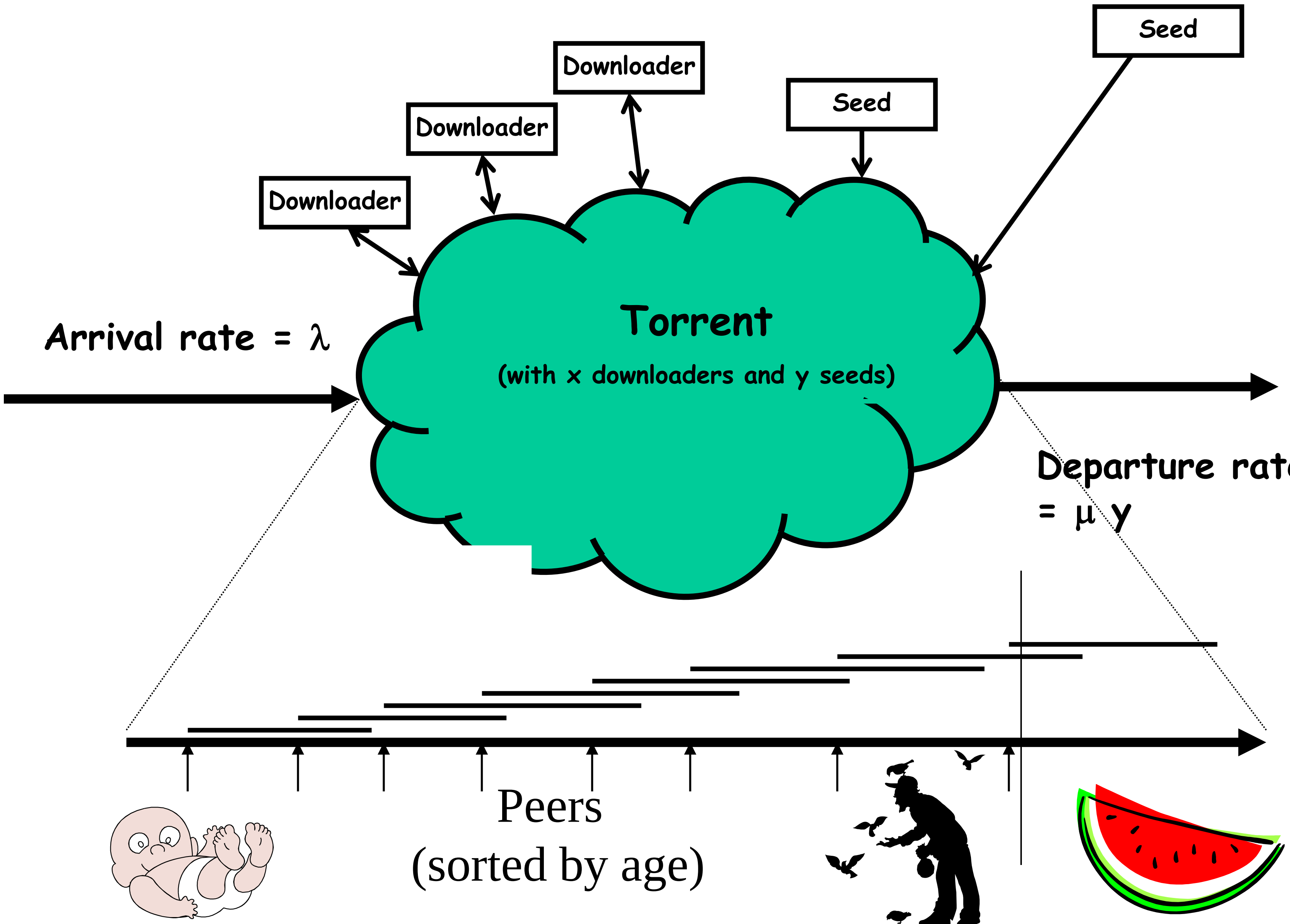
BitTorrent Model



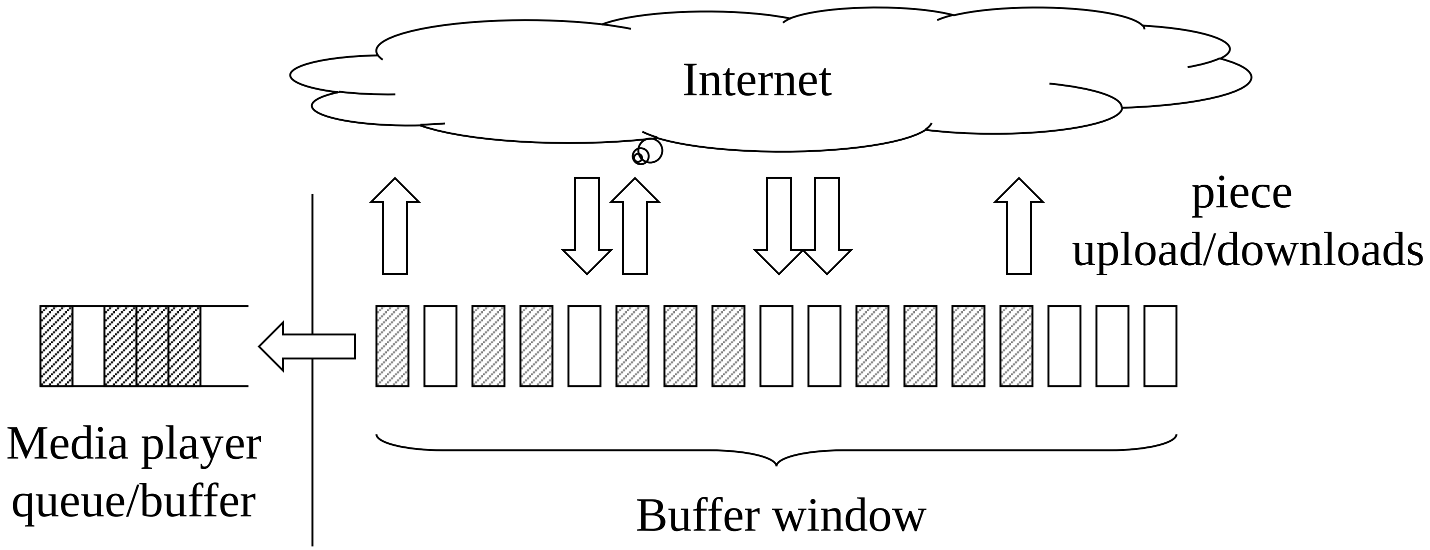
BitTorrent Model (random)



BitTorrent Model (chaining)



Live Streaming using Bittorrent-like systems



- Live streaming (e.g., CoolStreaming)
  - All peers at roughly the same play/download position
    - High bandwidth peers can easily contribute more …
  - (relatively) Small buffer window
    - Within which pieces are exchanged

Peer-assisted VoD streaming

- Can BitTorrent-like protocols provide scalable on-demand streaming?
- How sensitive is the performance to the application configuration parameters?
  - Piece selection policy (rarest vs. in-order tradeoff)
  - Peer selection policy
  - Upload/download bandwidth
- What is the user-perceived performance?
  - Start-up delay
  - Probability of disrupted playback

ACM SIGMETRICS 2008; IFIP Networking 2007/2009, IEEE/ACM ToN 2012

A protocol family for streaming

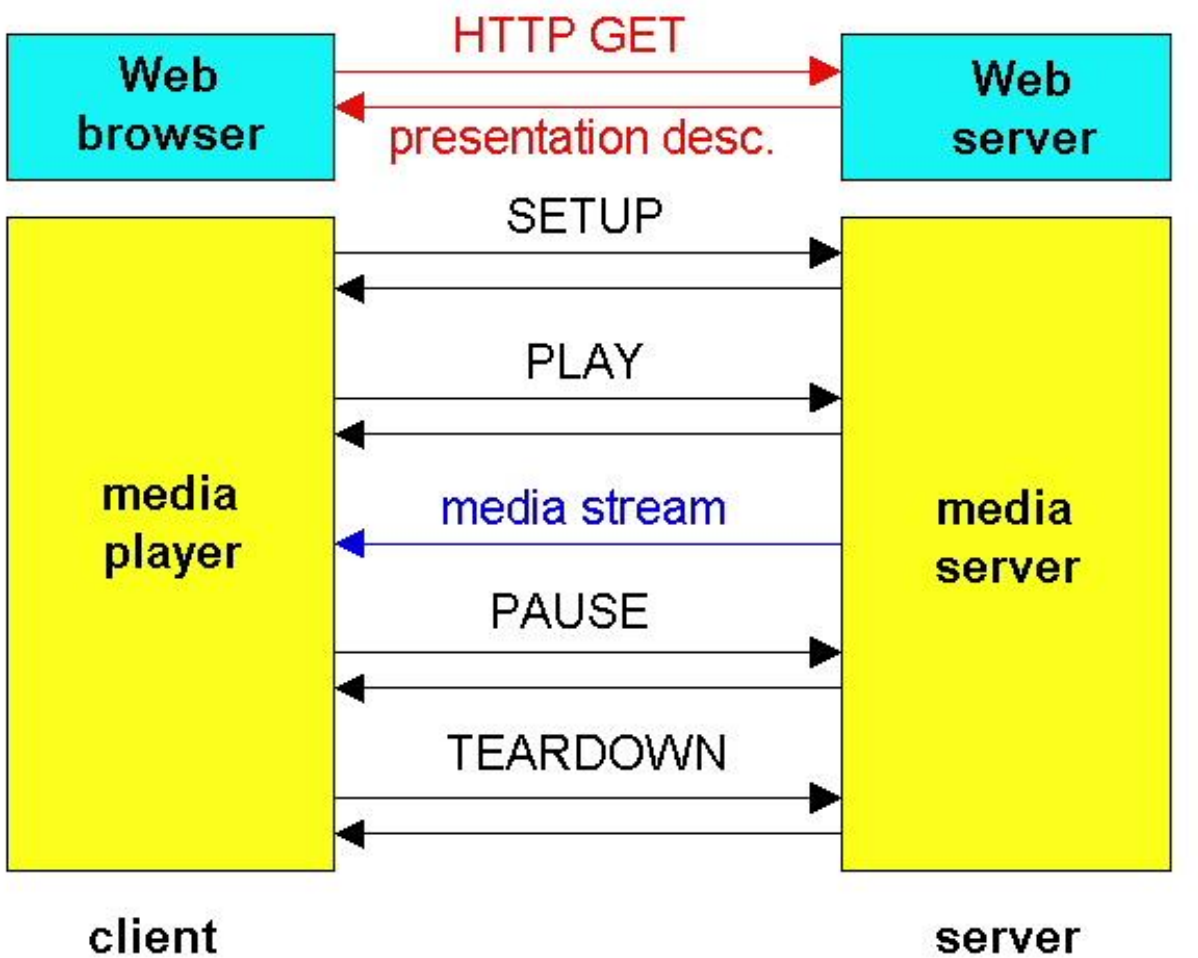
- RTSP
- RTP
- RTCP

RTSP Example

Scenario:

- metafile communicated to web browser
- browser launches player
- player sets up an RTSP control connection, data connection to streaming server

RTSP Operation



Control messages "out-of-band"

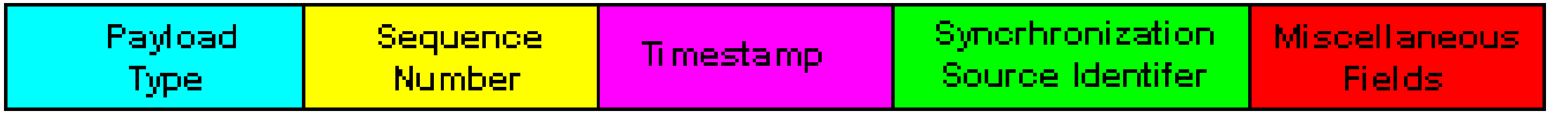
- port 554

Media stream "in-band".

Real-Time Protocol (RTP)

- RTP specifies packet structure for packets carrying audio, video data
- RFC 3550
- RTP runs in end systems
- RTP packets encapsulated in UDP segments

RTP Header



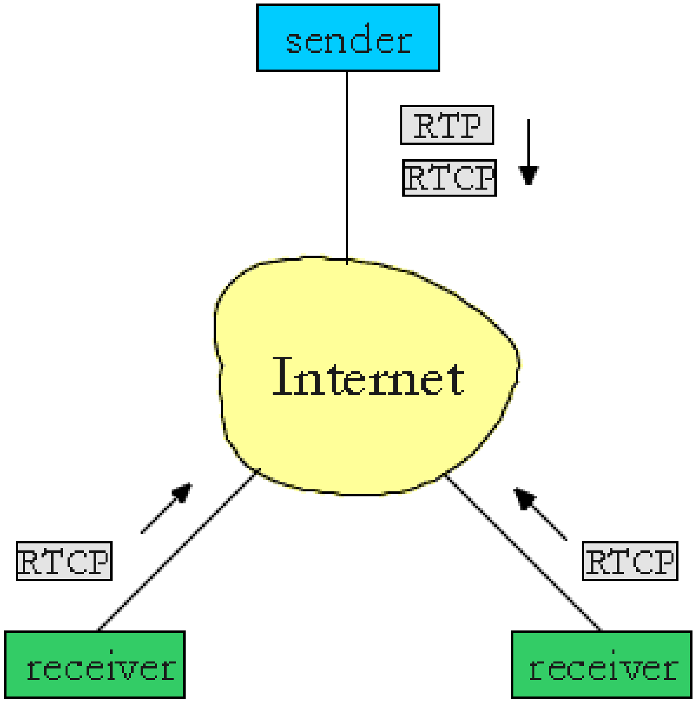
RTP Header

Payload Type (7 bits): Indicates type of encoding currently being used. If sender changes encoding in middle of conference, sender informs receiver via payload type field.

- Payload type 0: PCM mu-law, 64 kbps
- Payload type 3, GSM, 13 kbps
- Payload type 7, LPC, 2.4 kbps
- Payload type 26, Motion JPEG
- Payload type 31. H.261
- Payload type 33, MPEG2 video

Sequence Number (16 bits): Increments by one for each RTP packet sent, and may be used to detect packet loss and to restore packet sequence.

Real-time Control Protocol (RTCP)



Receiver report packets:

- fraction of packets lost, last sequence number, average interarrival jitter

Sender report packets:

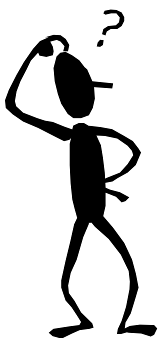
- SSRC of RTP stream, current time, number of packets sent, number of bytes sent

- feedback can be used to control performance

- RTCP attempts to limit its traffic to 5% of session bandwidth

Multimedia Over "Best Effort" Internet

- TCP/UDP/IP: *no* guarantees on delay, loss



But you said multimedia apps requires QoS and level of performance to be effective!

Today's multimedia applications implement functionality at the app. layer to mitigate (as best possible) effects of delay, loss

Packet Loss

- network loss: IP datagram lost due to network congestion (router buffer overflow) or losses at wireless link(s)
- delay loss: IP datagram arrives too late for playout at receiver (effectively the same as if it was lost)
  - delays: processing, queueing in network; end-system (sender, receiver) delays
  - Tolerable delay depends on the application
- How can packet loss be handled?
  - We will discuss this next …

Receiver-based Packet Loss Recovery

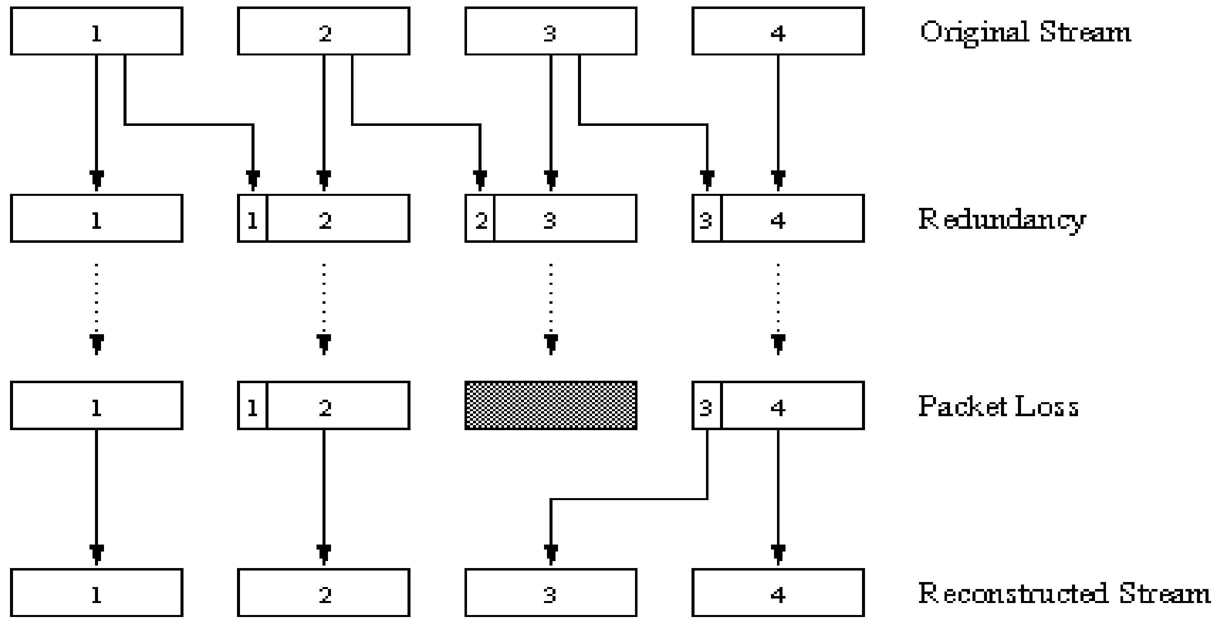
- Generate replacement packet
  - Packet repetition
  - Interpolation
  - Other sophisticated schemes
- Works when audio/video streams exhibit short-term correlations (e.g., self-similarity)
- Works for relatively low loss rates (e.g., < 5%)
- Typically, breaks down on "bursty" losses

Forward Error Correction (FEC)

- For every group of n actual media packets, generate k additional redundant packets
- Send out n+k packets, which increases the bandwidth consumption by factor k/n.
- Receiver can reconstruct the original n media packets provided at most k packets are lost from the group
- Works well at high loss rates (for a proper choice of k)
- Handles "bursty" packet losses
- Cost: increase in transmission cost (bandwidth)

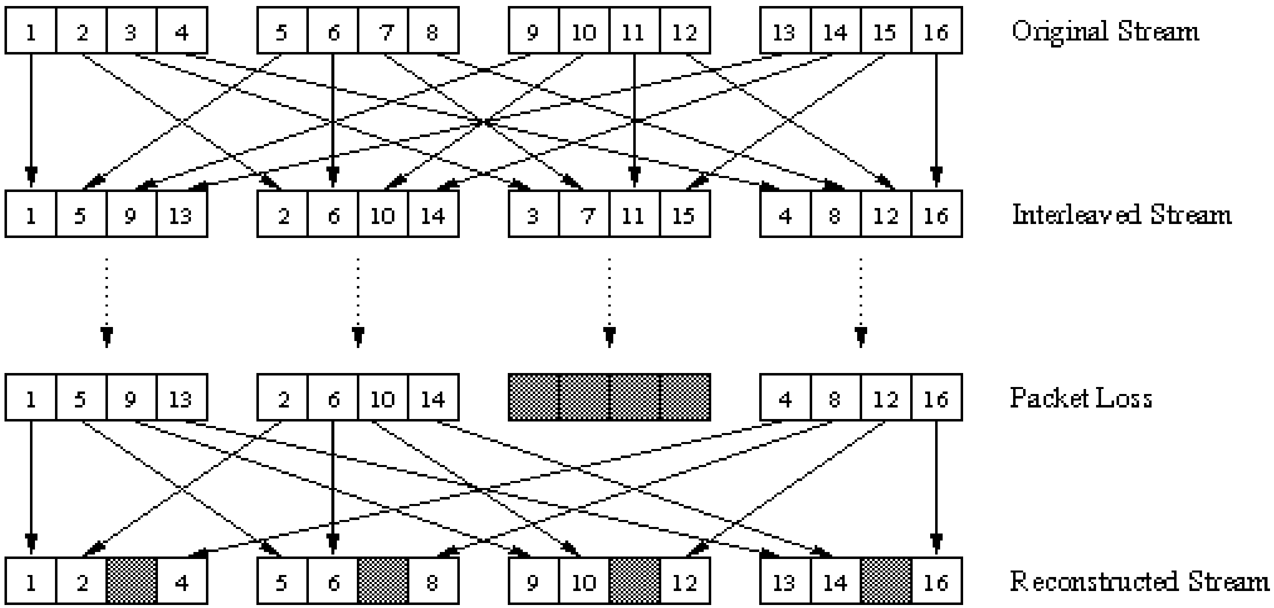
Another FEC Example

- "piggyback lower quality stream"



- Whenever there is non-consecutive loss, the receiver can conceal the loss.
- Can also append (n-1)st and (n-2)nd low-bit rate chunk

Interleaving: Recovery from packet loss



Interleaving

- Intentionally alter the sequence of packets before transmission
- Better robustness against "burst" losses of packets
- Results in increased playout delay from inter-leaving

More slides ….

Outline

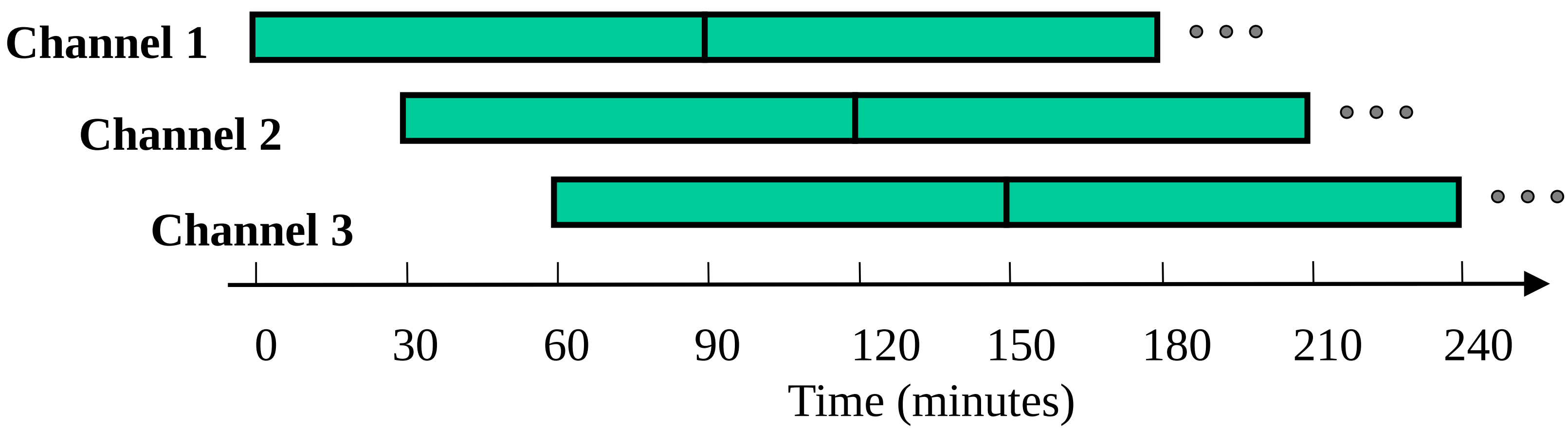
- Multimedia Networking Applications
- Streaming stored audio and video
- Scalable Streaming Techniques
- Content Distribution Networks
- Beyond Best Effort

Streaming Popular Content

- Consider a popular media file
  - Playback rate: 1 Mbps
  - Duration: 90 minutes
  - Request rate: once every minute
- How can a video server handle such high loads?
  - Approach 1: Start a new "stream" for each request
  - Allocate server and disk I/O bandwidth for each request
  - Bandwidth required at server= 1 Mbps x 90

Streaming Popular Content using Batching

- Approach 2: Leverage the multipoint delivery capability of modern networks
- Playback rate = 1 Mbps, duration = 90 minutes
- Group requests in non-overlapping intervals of 30 minutes:
  - Max. start-up delay = 30 minutes
  - Bandwidth required = 3 channels = 3 Mbps



Batching Issues

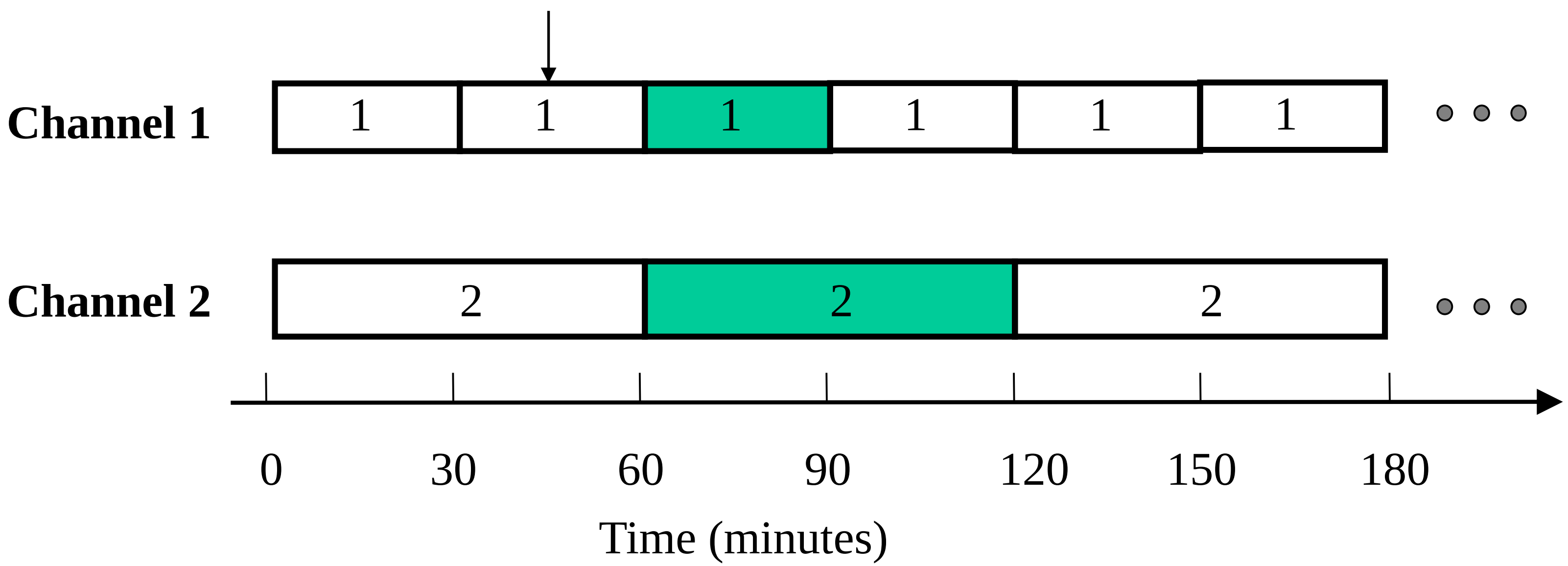
- Bandwidth increases linearly with decrease in start-up delays

- Can we reduce or eliminate "start-up" delays?
  - Periodic Broadcast Protocols

  - Stream Merging Protocols

Periodic Broadcast Example

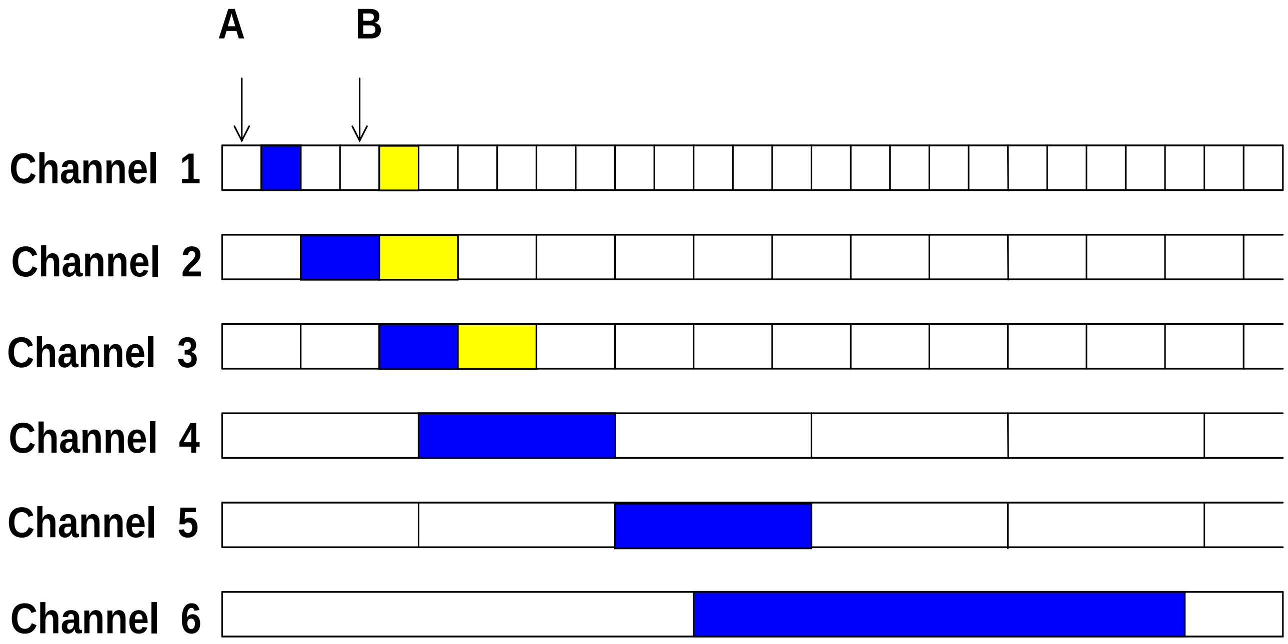
- Partition the media file into 2 segments with relative sizes {1, 2}. For a 90 min. movie:
  - Segment 1 = 30 minutes, Segment 2 = 60 minutes
- Advantage:
  - Max. start-up delay = 30 minutes
  - Bandwidth required = 2 channels = 2 Mbps
- Disadvantage: Requires increased client capabilities



Skyscraper Broadcasts (SB)

[Hua & Sheu 1997]

- Divide the file into *K* segments of increasing size
  - Segment size progression: 1, 2, 2, 5, 5, 12, 12, 25, …
- Multicast each segment on a separate channel at the playback rate
- Aggregate rate to clients: *2 x playback rate*



Comparing Batching and SB

Server Bandwidth	Start-up Delay	
	Batching	SB
1 Mbps	90 minutes	90 minutes
2 Mbps	45 minutes	30 minutes
6 Mbps	15 minutes	3 minutes
10 Mbps	9 minutes	30 seconds

- Playback rate = 1 Mbps, duration = 90 minutes

- Limitations of Skyscraper:
  - Ad hoc segment size progress
  - Does not work for low client data rates

Reliable Periodic Broadcasts (RPB)

[Mahanti *et al*. 2001, 2003, 2004]

- Optimized PB protocols (no packet loss recovery)
  - client fully downloads each segment before playing
  - required server bandwidth near minimal
  - Segment size progression is *not* ad hoc
  - Works for client data rates < 2 x playback rate
- extend for packet loss recovery
- extend for "bursty" packet loss
- extend for client heterogeneity

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Optimized Periodic Broadcasts

- r = segment streaming rate = 1
- s = maximum # streams client listens to concurrently = 2
- b = client data rate = s x r = 2
- length of first s segments: $\frac{1}{r} l_k = \frac{1}{r} l_1 + \sum_{j=1}^{k-1} l_j$
- length of segment k > s: $\frac{1}{r} l_k = \sum_{j=k-s}^{k-1} l_j$

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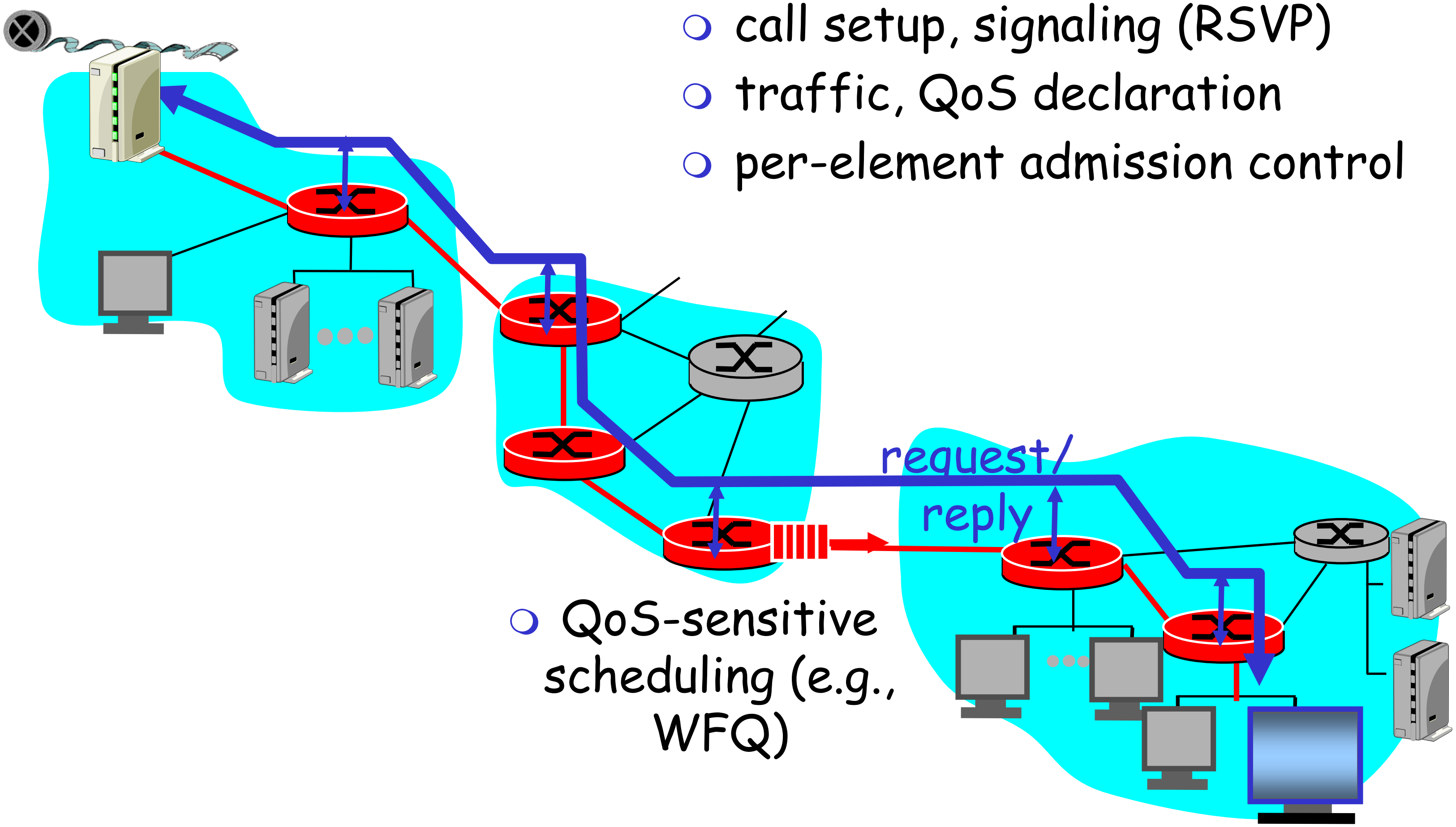
Integrated Services (IntServ) Architecture

- architecture for providing QOS guarantees in IP networks for individual flows
- flow: a distinguishable stream of distinct IP datagrams
  - Unidirectional
  - Multiple recipient
- Components of this architecture:
  - Admission control
  - Reservation protocol
  - Routing protocol
  - Classifier and route selection
  - Packet scheduler

Intserv: QoS guarantee scenario

- Resource reservation
  - call setup, signaling (RSVP)
  - traffic, QoS declaration
  - per-element admission control



- QoS-sensitive scheduling (e.g., WFQ)

Call Admission

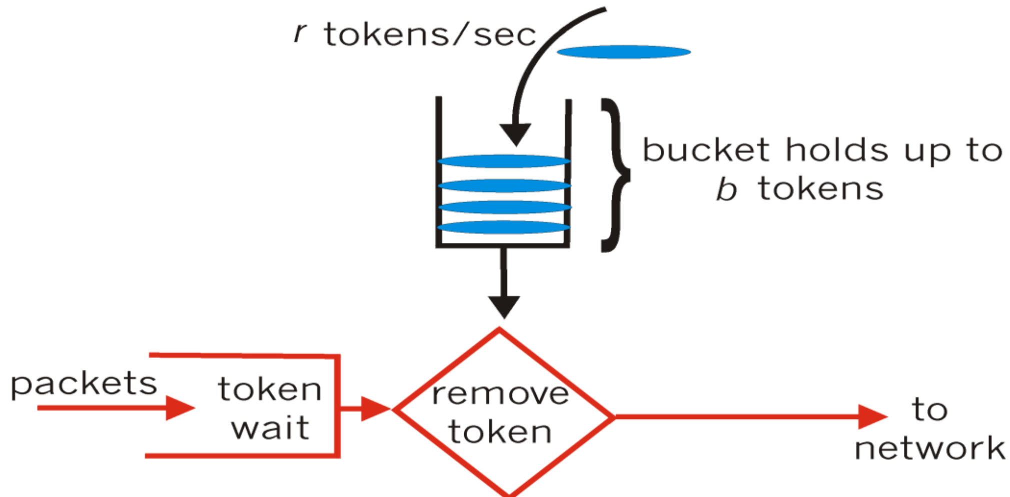
Arriving session must :

- declare its QoS requirement
  - R-spec: defines the QoS being requested
- characterize traffic it will send into network
  - T-spec: defines traffic characteristics
- signaling protocol: needed to carry R-spec and T-spec to routers (where reservation is required)
  - RSVP

Need Scheduling and Policing Policies to provide QoS

Policing: Token Bucket

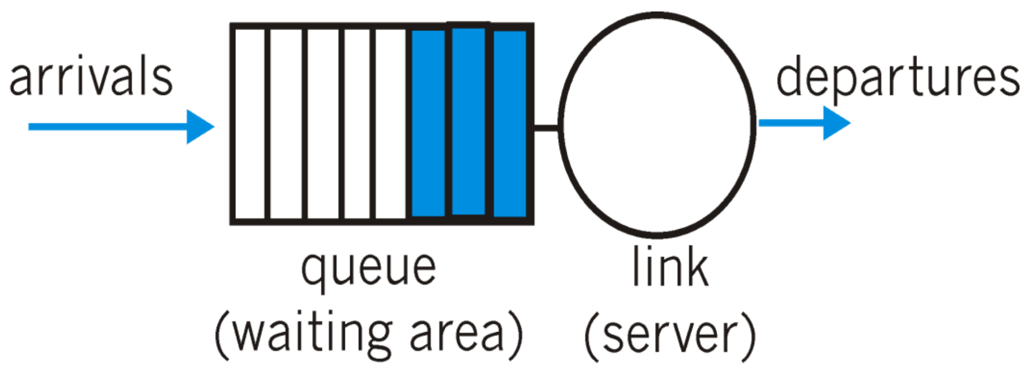
Token Bucket: limit input to specified Burst Size and Average Rate.



- bucket can hold b tokens
- tokens generated at rate *r token/sec* unless bucket full
- *over interval of length t: number of packets admitted less than or equal to (r t + b).*

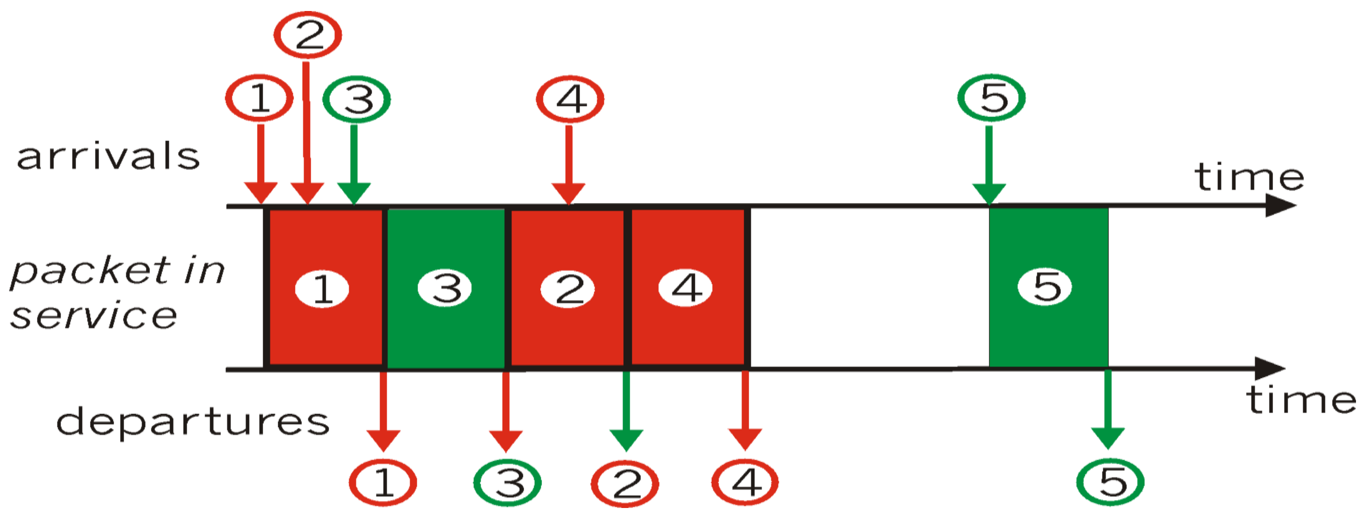
Link Scheduling

- scheduling: choose next packet to send on link
- FIFO (First In First Out) scheduling: send in order of arrival to queue
  - discard policy: if packet arrives to full queue: who to discard?
    - DropTail: drop arriving packet
    - Priority: drop/remove on priority basis
    - Random: drop/remove randomly (e.g., RED)



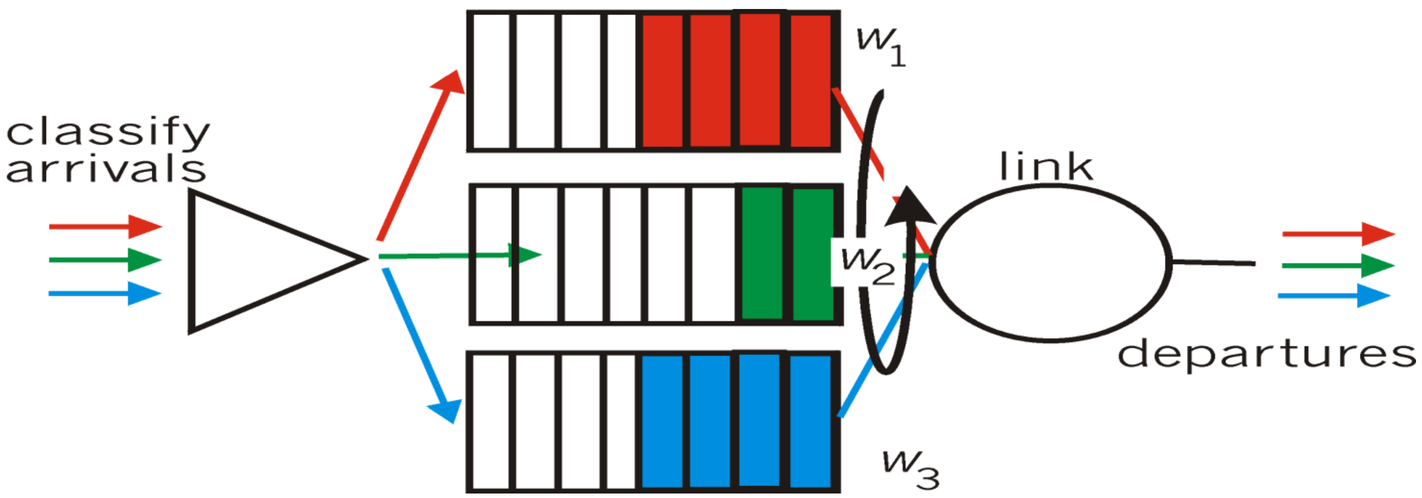
Round Robin

- multiple classes
- cyclically scan class queues, serving one from each class (if available)



Weighted Fair Queuing

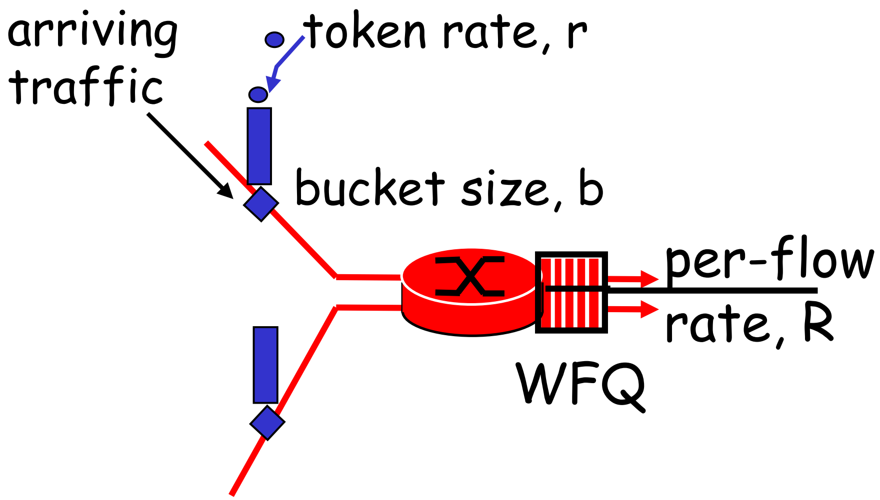
- generalized Round Robin
- each class gets weighted amount of service in each cycle



IntServ QoS: Service models [rfc2211, rfc 2212]

Guaranteed service:

- Assured data rate
- A specified upper bound on queuing delay



$D_{max} = b/R$

Controlled load service:

- "a quality of service closely approximating the QoS that same flow would receive from an unloaded network element."
- Similar to behavior best effort service in an unloaded network

Differentiated Services

Concerns with IntServ:

- Scalability: signaling, maintaining per-flow router state difficult with large number of flows
- Flexible Service Models: Intserv has only two classes. Desire "qualitative" service classes
  - E.g., Courier, xPress, and normal mail
  - E.g., First, business, and cattle class ☺

DiffServ approach:

- simple functions in network core, relatively complex functions at edge routers (or hosts)
- Don't define service classes, just provide functional components to build service classes

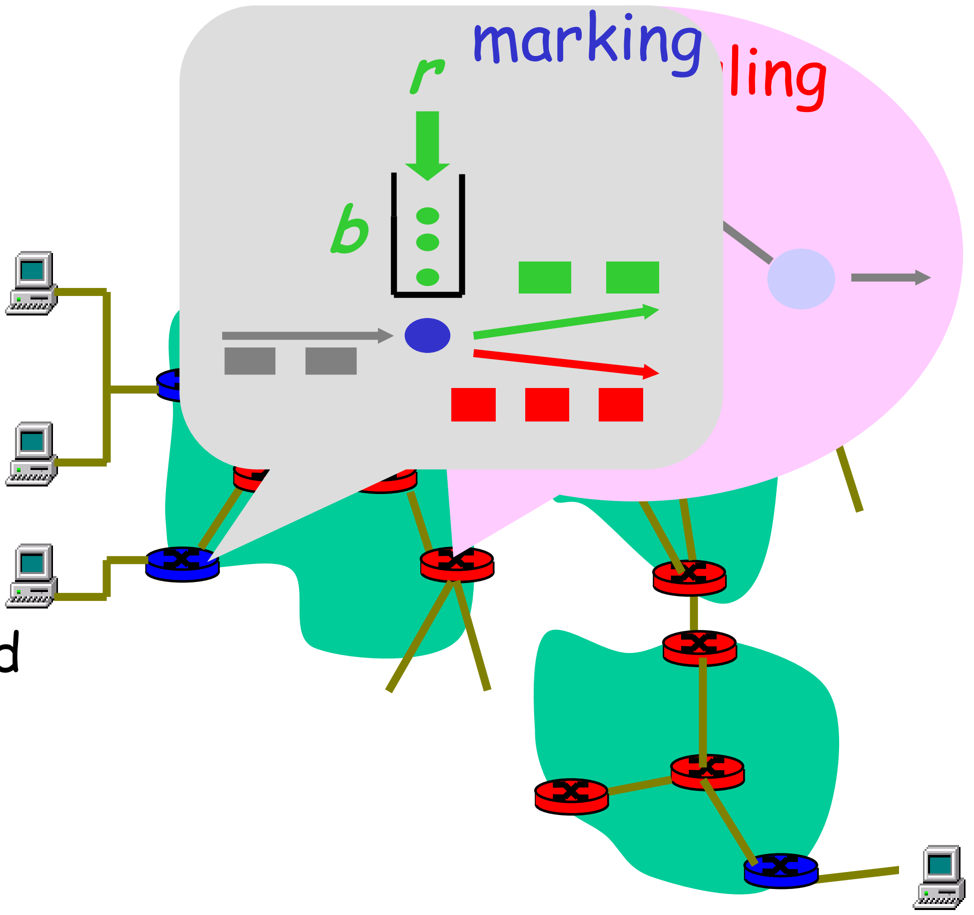
DiffServ Architecture

Edge router:

- per-flow traffic management
- Set the DS field; value determines type of service (PHB: Per-Hop Behavior)

Core router:

- buffering and scheduling based on marking at edge
- per-class traffic management



Traffic Classification/Conditioning

- How can packet markings be carried in IPv4 datagrams?
- Sender may agree to conform to a "traffic profile", thus a leaky bucket policer may be used at the network edge to enforce
  - Peak rate
  - Average rate
  - Burst size
- What happens when traffic profile is violated?
  - Employ traffic shaping?

Deployment Issues

- Single administrative domain
- Incremental deployment
- Traffic policing/shaping complexity
- Charging models

Signaling in the Internet

connectionless (stateless) forwarding by IP routers + best effort service = no network signaling protocols in initial IP design

- New requirement: reserve resources along end-to-end path (end system, routers) for QoS for multimedia applications
- RSVP: Resource reSerVation Protocol [RFC 2205]
  - " ... allow users to communicate requirements to network in robust and efficient way." i.e., signaling !
- earlier Internet Signaling protocol: ST-II [RFC 1819]

RSVP Design Goals

1. accommodate heterogeneous receivers (different bandwidth along paths)
2. accommodate different applications with different resource requirements
3. make multicast a first class service, with adaptation to multicast group membership
4. leverage existing multicast/unicast routing, with adaptation to changes in underlying unicast, multicast routes
5. control protocol overhead to grow (at worst) linear in # receivers
6. modular design for heterogeneous underlying technologies

RSVP: does not…

- specify how resources are to be reserved
  - rather: a mechanism for communicating needs
- determine routes packets will take
  - that's the job of routing protocols
  - signaling decoupled from routing
- interact with forwarding of packets
  - separation of control (signaling) and data (forwarding) planes

Multimedia Networking: Summary

- multimedia applications and requirements
- making the best of today's "best effort" service
- scheduling and policing mechanisms
- next generation Internet: IntServ, RSVP, DiffServ, IPv6, IP-QoS