### Computer Networks

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Notes derived from "*Computer Networking: A Top Down Approach"*, by Jim Kurose and Keith Ross, Addison-Wesley.

The slides are adapted and modified based on slides from the book's companion Web site, as well as modified slides by Anirban Mahanti and Carey Williamson.

### Creating a network app

#### write programs that

- run on (different) end systems
- communicate over network

No need to write software for network-core devices



## **Application architectures**

- client-server
- \* peer-to-peer (P2P)
- hybrid of client-server and P2P

## <u>App-layer protocols</u>

public-domain protocols:

- Often defined in RFCs
- allows for interoperability
- e.g., HTTP, SMTP, BitTorrent

proprietary protocols:

e.g., Skype, Spotify

### Network applications: some jargon

- Process: program running within a host.
- within same host, two processes communicate using inter-process communication (IPC, defined by OS).
- processes running on different hosts communicate with an application-layer protocol

User agent: interfaces with user "above" and network "below".

- implements user
  interface &
  application-level
  protocol
  - Web: browser
  - E-mail: mail reader
  - streaming audio/video: media player

## Processes communicating

 processes in different hosts communicate by exchanging messages

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 processes in different hosts communicate by exchanging messages

- <u>Client-server paradigm</u>
- client process: process that initiates communication
- server process: process that waits to be contacted



# process sends/receives messages to/from its socket



### Addressing processes:

- For a process to receive messages, it must have an identifier
- Identifier includes
  <u>both</u> the IP address
  and port numbers
  associated with the
  process on the host.
- Section 2 Sec
  - HTTP server: 80
  - Mail server: 25
- \* More on this later

### Addressing processes:

- For a process to receive messages, it must have an identifier
- Every host has a unique
  32-bit IP address
- Q: does the IP address of the host on which the process runs suffice for identifying the process?
- Answer: No, many processes can be running on same host

- Identifier includes
   <u>both</u> the IP address
   and port numbers
   associated with the
   process on the host.
- Example port numbers:
  - HTTP server: 80
  - Mail server: 25
- More on this later

## Remember: What's a protocol?

protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt

### Protocol: Connection oriented or not?

#### Connection oriented:

- \* Hand shaking
  - Explicit setup phase for logical connection
  - Connection release afterwards
- Establishes state information about the connection
- \* Mechanisms for
  - reliable data transfer, error control, flow control, etc.
- Guarantees that data will arrive (eventually)

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#### Connection less:

- No handshaking
- No (significant) state information (at end points or in network)
- No mechanisms for flow control etc.
- No guarantees of arrival (or when)
- Simpler (and faster?)

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Which is the best? ... It depends on (i) what it is used for, and (ii) what it is built ontop of

## Internet protocol stack



• Signaling, modulation, encoding, etc,

## Internet protocol stack



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### What transport service does an app need?

Data loss

Bandwidth

Timing

### What transport service does an app need?

#### Data loss

- some apps (e.g., file transfer, telnet) require 100% reliable data transfer
- other apps (e.g., audio) can tolerate some loss

#### Timing

 some apps (e.g., Internet telephony, interactive games) require low delay to be "effective"

#### Bandwidth

- most apps ("elastic apps") make use of whatever bandwidth they get
- other apps (e.g., multimedia) require minimum amount of bandwidth to be "effective"

### Internet transport protocols services

#### TCP service:

- connection-oriented: setup required between client and server processes
- *reliable transport* between sending and receiving process
- *flow control:* sender won't overwhelm receiver
- congestion control: throttle sender when network overloaded
- does not provide: timing, minimum throughput guarantees, security

#### <u>UDP service:</u>

- unreliable data transfer
  between sending and
  receiving process
- does not provide: connection setup, reliability, flow control, congestion control, timing, throughput guarantee, or security
- Q: why bother? Why is there a UDP?

Application	Application layer protocol	Underlying transport protocol
<u> </u>		· ·
e-mail	SMTP [RFC 2821]	
remote terminal access	Telnet [RFC 854]	
Web	HTTP [RFC 2616]	
file transfer	FTP [RFC 959]	
streaming multimedia	proprietary	
	(e.g., RealNetworks,	youtube, netflix, spotify)
Internet telephony	proprietary	
	(e.g., Dialpad, skype)	)

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Application	layer protocol	transport protocol
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remote terminal access	Telnet [RFC 854]	
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streaming	multimedia	proprietary	TCP (or UDP)
		(e.g., RealNetworks, you	utube, netflix, spotify)
Interne	t telephony	proprietary	UDP or TCP
	-	(e.g., Dialpad, skype)	typically UDP

### \* WWW terminology and HTTP overview

## Some "Web" Terminology

- Web page may contain links to other pages (sometimes also called Web Objects)
- Object can be HTML file, JPEG image, Java applet, audio file,...
- Web pages are "Hypertexts"
  - One page points to another

#### Each object is addressable by a URL:



## HTTP overview

- HTTP: hypertext transfer protocol
- Web's application layer protocol

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## HTTP overview

#### HTTP: hypertext transfer protocol

- Web's application layer protocol
- client/server model
  - client: browser that requests, receives, "displays" Web objects
  - server: Web server sends objects in response to requests
- ✤ HTTP 1.0: RFC 1945
- ✤ HTTP 1.1: RFC 2616



## HTTP overview (continued)

#### Uses TCP:

- client initiates TCP connection (creates socket) to server, port 80
- server accepts TCP connection from client
- HTTP messages (applicationlayer protocol messages) exchanged between browser (HTTP client) and Web server (HTTP server)
- TCP connection closed

#### HTTP is "stateless"

 server maintains no information about past client requests
# HTTP overview (continued)

#### Uses TCP:

- client initiates TCP connection (creates socket) to server, port 80
- server accepts TCP connection from client
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### HTTP is "stateless"

- server maintains no information about past client requests
- Protocols that maintain "state" are complex!
- past history (state) must be maintained
- if server/client crashes, their views of "state" may be inconsistent, must be reconciled



- Introduction to App Layer Protocols
  Brief History of WWW
  Architecture
  HTTP Connections
  HTTP Format
- Web Performance
- Cookies

# Response time modeling

Definition of RTT: time to send a small packet to travel from client to server and back.

Response time:

- one RTT to initiate TCP connection
- one RTT for HTTP request and first few bytes of HTTP response to return
- file transmission time

total = 2\*RTT+transmit time



# HTTP connections

### Non-persistent HTTP

 At most one object is sent over a TCP connection.

#### HTTP/1.0 uses nonpersistent HTTP

### Persistent HTTP

 Multiple objects can be sent (one at a time) over single connection

- HTTP/1.1 uses
   persistent connections
   in default mode
  - Pipelined
  - Non-pipelined

## Persistent HTTP

#### Nonpersistent HTTP issues:

- requires 2 RTTs per object
- OS must work and allocate host resources for each TCP connection
- but browsers often open parallel TCP connections to fetch referenced objects

#### Persistent HTTP

- server leaves connection
   open after sending response
- subsequent HTTP messages
   between same client/server
   are sent over connection

#### Persistent without pipelining:

- client issues new request only when previous response has been received
- one RTT for each referenced object

#### Persistent with pipelining:

- default in HTTP/1.1
- client sends requests as soon as it encounters a referenced object
- as little as one RTT for all the referenced objects

# Network View: HTTP and TCP

TCP is a connection-oriented protocol



# Example Web Page

page.html

Harry Potter Movies As you all know, the new HP book will be out in June and then there will be a new movie shortly after that... "Harry Potter and the Bathtub Ring

hpface.jpg

castle.gif



The "classic" approach in HTTP/1.0 is to use one HTTP request per TCP connection, serially.





The <u>"persistent HTTP"</u> approach can re-use the same TCP connection for Multiple HTTP transfers, one after another, serially Amortizes TCP overhead, but maintains TCP state longer at server.



The "pipelining" feature in HTTP/1.1 allows requests to be issued asynchronously on a persistent connection. Requests must be processed in proper order. Can do clever packaging.



- Introduction to App Layer Protocols
- Srief History of WWW
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## <u>HTTP request message</u>

- HTTP request message:
  - ASCII (human-readable format)



## HTTP request message: general format



# HTTP Methods

- \* GET: retrieve a file (95% of requests)
- HEAD: just get meta-data (e.g., mod time)
- POST: submitting a form to a server
- PUT: store enclosed document as URI
- DELETE: removed named resource
- LINK/UNLINK: in 1.0, gone in 1.1
- TRACE: http "echo" for debugging (added in 1.1)
- CONNECT: used by proxies for tunneling (1.1)
- OPTIONS: request for server/proxy options (1.1)

## Trying out HTTP (client side) for yourself

- 1. Telnet to your favorite Web server:
- telnet www.eurecom.fr 80 Opens TCP connection to port 80 (default HTTP server port) at www.eurecom.fr. Anything typed in sent to port 80 at www.eurecom.fr
- 2. Type in a GET HTTP request:

GET /~ross/index.html HTTP/1.0

By typing this in (hit carriage return twice), you send this minimal (but complete) GET request to HTTP server

3. Look at response message sent by HTTP server!

## HTTP response message

status line (protocol <u>status code</u> status phrase)

> header lines

HTTP/1.1 200 OK Connection: close Date: Thu, 06 Aug 1998 12:00:15 GMT Server: Apache/1.3.0 (Unix) Last-Modified: Mon, 22 Jun 1998 ..... Content-Length: 6821 Content-Type: text/html

data, e.g., – requested HTML file

data data data data ...

## HTTP Response Status Codes

- Alx: Informational (def'd in 1.0, used in 1.1) 100 Continue, 101 Switching Protocols
- 2XX: Success

200 OK, 206 Partial Content

3XX: Redirection

301 Moved Permanently, 304 Not Modified

✤ 4XX: Client error

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Introduction to App Layer Protocols Srief History of WWW Architecture HTTP Connections HTTP Format Web Performance \* Cookies

## Web caches (proxy server)

Goal: satisfy client request without involving origin server

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- Use a proxy cache
  - Acts as both client and server
- Typically cache is installed by ISP (university, company, residential ISP)

## Web caches (proxy server)

Goal: satisfy client request without involving origin server

- user sets browser: Web accesses via cache
- browser sends all HTTP requests to cache
  - If object in cache: cache returns object
  - Else: cache requests object from origin server, then returns object to client



## Caching example









## Caching example (cont)

### Why Web caching?

- Reduce response time for client request
- Reduce traffic on an institution's access link
- Offloads server



## "Typical" hit rates?



 Traces suggests 40-50% hit rate for objects and 20-25% for bytes (across geographies and over time)

Example references

- P Gill, M. Arlitt, N. Carlsson, A. Mahanti, C. Williamson, "Characterizing Organizational Use of Web-based Services: Methodology, Challenges, Observations, and Insights", ACM Transactions on the Web (ACM TWEB), Vol. 5, No. 4 (Oct. 2011), pp. 19:1--19:23.
- A. Wolman, G. Voelker, N. Sharma, N. Cardwell, A. Karlin, and H. Levy, "On the scale and performance of cooperative Web proxy caching". Proc. ACM Symposium on Operating Systems Principles (ACM SOSP). Kiawah Island, SC, Dec. 1999, pp. 16–31.



## Some Issues

Not all objects can be cached
Cache Replacement Policies

Prefetch?

Cache consistency

# Some Issues

- \* Not all objects can be cached
  - E.g., dynamic objects, copyrighted material
- Cache Replacement Policies
  - Variable size objects
  - Varying cost of not finding an object (a "miss") in the cache
- Prefetch?
  - A large fraction of the requests are one-timers
- Cache consistency
  - strong
  - weak
## Weak Consistency

- Search cached copy has a TTL beyond which it must be validated with the origin server
- Age Penalty?

### <u>Conditional GET: client-side caching</u>

- Goal: don't send object if client has up-to-date cached version
- client: specify date of cached copy in HTTP request
- server: response contains no object if cached copy is upto-date.



### Content distribution networks (CDNs)

### **Content replication**

- replicate content at hundreds of servers throughout Internet (often in edge/access network)
- content "close" to user reduce impairments (loss, delay) of sending content over long paths



### Content distribution networks (CDNs)

### **Content replication**

- CDN (e.g., Akamai, Limewire) customer is the content provider (e.g., CNN)
- Other companies build their own CDN (e.g., Google)
- CDN replicates customers' content in CDN servers.
- When provider updates content, CDN updates servers



# Cookies: keeping "state"

Many major Web sites use cookies

### Four components:

- 1) cookie header line in the HTTP response message
- 2) cookie header line in HTTP request message
- cookie file kept on user's host and managed by user's browser
- 4) back-end database at Web site

### Example:

 User visits a specific ecommerce site ...









# <u>Cookies (continued)</u>

### What cookies can bring:

- authorization
- shopping carts
- recommendations
- user session state(Web e-mail)

### <u>Cookies and privacy:</u>

- cookies permit sites to learn a lot about you
- you may supply name and e-mail to sites
- search engines use redirection & cookies to learn yet more
- advertising companies obtain info across sites

# Web & HTTP

The major application on the Internet

- A large fraction of traffic is HTTP
- Client/server model:
  - Clients make requests, servers respond to them
  - Done mostly in ASCII text (helps debugging!)
- Various headers and commands
- Web Caching & Performance
- Content Distribution Networks



# Introduction to HTTP





Laptop w/ Netscape

http response



Server w/ Apache

http request

http response



Desktop w/ Explorer

- HTTP: HyperText Transfer Protocol
  - Communication protocol between clients and servers
  - Application layer protocol for WWW
- Client/Server model:
  - Client: browser that requests, receives, displays object
  - Server: receives requests and responds to them
- Protocol consists of various operations
  - Few for HTTP 1.0 (RFC 1945, 1996)
  - Many more in HTTP 1.1 (RFC 2616, 1999)

## Request Generation

- User clicks on something
- Uniform Resource Locator (URL):
  - http://www.cnn.com
  - http://www.cpsc.ucalgary.ca
  - https://www.paymybills.com
  - ftp://ftp.kernel.org
- Different URL schemes map to different services
- Hostname is converted from a name to a 32-bit IP address (DNS lookup, if needed)
- Connection is established to server (TCP)

# What Happens Next?

- Client downloads HTML document
  - Sometimes called "container page"
  - Typically in text format (ASCII)
  - Contains instructions for rendering (e.g., background color, frames)
  - Links to other pages
- Many have embedded objects:
  - Images: GIF, JPG (logos, banner ads)
  - Usually automatically retrieved
    - I.e., without user involvement
    - can control sometimes
       (e.g. browser options, junkbusters)

<html> <head> <meta name="Author" content="Erich Nahum"> <title> Linux Web Server Performance </title> </head><body text="#00000"> <img width=31 height=11 src="ibmlogo.gif"> <img src="images/new.gif> <h1>Hi There!</h1> Here's lots of cool linux stuff! <a href="more.html"> Click here</a> for more!

</sample html file

### Web Server Role

- Respond to client requests, typically a browser
  - Can be a proxy, which aggregates client requests
  - Could be search engine spider or robot
- May have work to do on client's behalf:
  - Is the client's cached copy still good?
  - Is client authorized to get this document?
- Hundreds or thousands of simultaneous clients
- Hard to predict how many will show up on some day (e.g., "flash crowds", diurnal cycle, global presence)
- Many requests are in progress concurrently

## HTTP Request Format

```
GET /images/penguin.gif HTTP/1.0
User-Agent: Mozilla/0.9.4 (Linux 2.2.19)
Host: www.kernel.org
Accept: text/html, image/gif, image/jpeg
Accept-Encoding: gzip
Accept-Language: en
Accept-Charset: iso-8859-1,*,utf-8
Cookie: B=xh203jfsf; Y=3sdkfjej
```

<cr><lf>

- Messages are in ASCII (human-readable)
- Carriage-return and line-feed indicate end of headers
- Headers may communicate private information (browser, OS, cookie information, etc.)

# Request Types

Called Methods:

- ✤ GET: retrieve a file (95% of requests)
- HEAD: just get meta-data (e.g., mod time)
- POST: submitting a form to a server
- PUT: store enclosed document as URI
- DELETE: removed named resource
- LINK/UNLINK: in 1.0, gone in 1.1
- TRACE: http "echo" for debugging (added in 1.1)
- CONNECT: used by proxies for tunneling (1.1)
- OPTIONS: request for server/proxy options (1.1)

### <u>Response Format</u>

```
HTTP/1.0 200 OK
Server: Tux 2.0
Content-Type: image/gif
Content-Length: 43
Last-Modified: Fri, 15 Apr 1994 02:36:21 GMT
Expires: Wed, 20 Feb 2002 18:54:46 GMT
Date: Mon, 12 Nov 2001 14:29:48 GMT
Cache-Control: no-cache
Pragma: no-cache
Connection: close
Set-Cookie: PA=wefj2we0-jfjf
```

<cr><lf> Similar format to requests (i.e., ASCII)

### Response Types

- \* 1XX: Informational (def'd in 1.0, used in 1.1) 100 Continue, 101 Switching Protocols
- 2XX: Success
  - 200 OK, 206 Partial Content
- 3XX: Redirection
  - 301 Moved Permanently, 304 Not Modified
- ✤ 4XX: Client error

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# **Outline of an HTTP Transaction**

- This section describes the basics of servicing an HTTP GET request from user space
- Assume a single process running in user space, similar to Apache 1.3
- We'll mention relevant socket operations along the way

initialize; forever do { get request; process; send response; log request;

> server in a nutshell

### Readying a Server

```
s = socket(); /* allocate listen socket */
bind(s, 80); /* bind to TCP port 80  */
listen(s); /* indicate willingness to accept */
while (1) {
    newconn = accept(s); /* accept new connection */b
```

- First thing a server does is notify the OS it is interested in WWW server requests; these are typically on TCP port 80.
   Other services use different ports (e.g., SSL is on 443)
- Allocate a socket and bind()'s it to the address (port 80)
- Server calls listen() on the socket to indicate willingness to receive requests
- Calls accept() to wait for a request to come in (and blocks)
- When the accept() returns, we have a new socket which represents a new connection to a client

### Processing a Request

```
remoteIP = getsockname(newconn);
remoteHost = gethostbyname(remoteIP);
gettimeofday(currentTime);
read(newconn, reqBuffer, sizeof(reqBuffer));
reqInfo = serverParse(reqBuffer);
```

- setsockname() called to get the remote host name
  - for logging purposes (optional, but done by most)
- sethostbyname() called to get name of other end
  - again for logging purposes
- settimeofday() is called to get time of request
  - both for Date header and for logging
- read() is called on new socket to retrieve request
- request is determined by parsing the data
  - "GET /images/jul4/flag.gif"

### Processing a Request (cont)

```
fileName = parseOutFileName(requestBuffer);
fileAttr = stat(fileName);
serverCheckFileStuff(fileName, fileAttr);
open(fileName);
```

- stat() called to test file path
  - to see if file exists/is accessible
  - may not be there, may only be available to certain people
  - "/microsoft/top-secret/plans-for-world-domination.html"
- stat() also used for file meta-data
  - e.g., size of file, last modified time
  - "Has file changed since last time I checked?"
- might have to stat() multiple files and directories
- Assuming all is OK, open() called to open the file

### Responding to a Request

```
read(fileName, fileBuffer);
headerBuffer = serverFigureHeaders(fileName, reqInfo);
write(newSock, headerBuffer);
write(newSock, fileBuffer);
close(newSock);
close(fileName);
write(logFile, requestInfo);
```

- read() called to read the file into user space
- write() is called to send HTTP headers on socket (early servers called write() for *each header!*)
- \* write() is called to write the file on the socket
- \$ close() is called to close the socket
- close() is called to close the open file descriptor
- write() is called on the log file

## Summary of Web and HTTP

- The major application on the Internet
  - Majority of traffic is HTTP (or HTTP-related)
- Client/server model:
  - Clients make requests, servers respond to them
  - Done mostly in ASCII text (helps debugging!)
- Various headers and commands
  - Too many to go into detail here
  - Many web books/tutorials exist (e.g., Krishnamurthy & Rexford 2001)