

TDTS04/11: Computer Networks

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Notes derived from "*Computer Networking: A Top Down Approach*", by Jim Kurose and Keith Ross, Addison-Wesley.

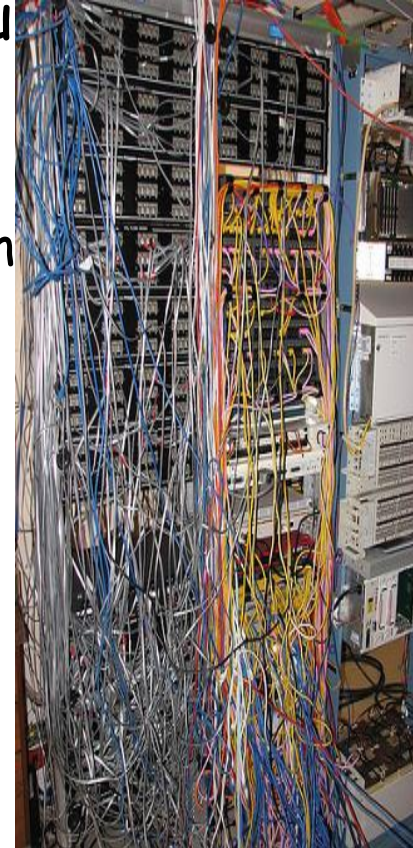
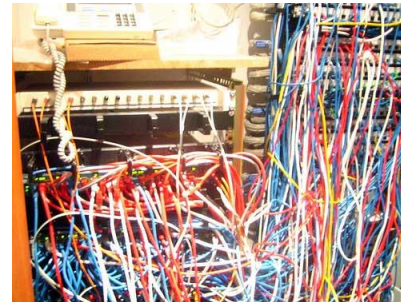
The slides are adapted and modified based on slides from the book's companion Web site, as well as modified slides by Anirban Mahanti and Carey Williamson.

What is Wireless Networking?

- ❖ The use of infra-red (IR) or radio frequency (RF) signals to share information and resources between devices
- ❖ Promises *anytime, anywhere* connectivity
 - Laptops, palmtops, PDAs, Internet-enabled phone promise anytime *untethered* Internet access
- ❖ No wires!

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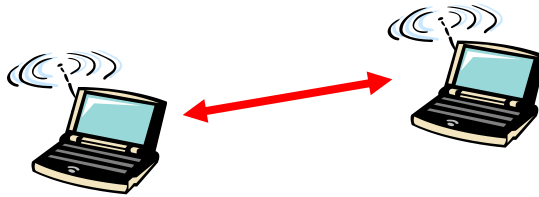


Two important (but different!) challenges

- ❖ Communication over wireless link
- ❖ Handling mobile user who changes point of attachment to network

Two important (but different!) challenges

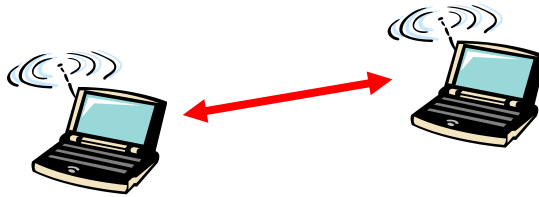
- ❖ Communication over wireless link



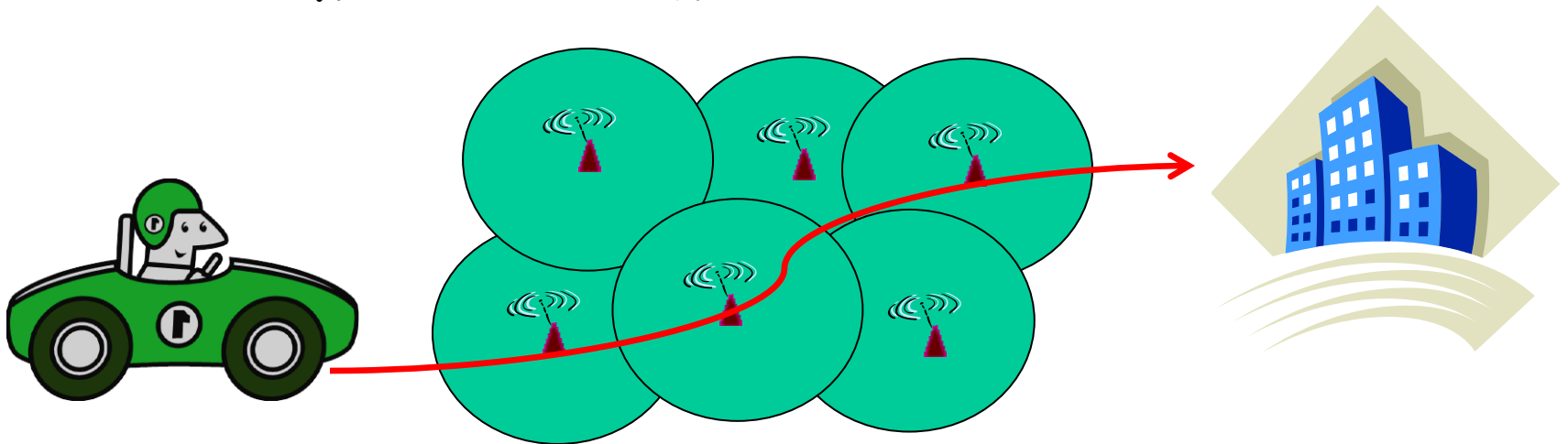
- ❖ Handling mobile user who changes point of attachment to network

Two important (but different!) challenges

- ❖ Communication over wireless link



- ❖ Handling mobile user who changes point of attachment to network



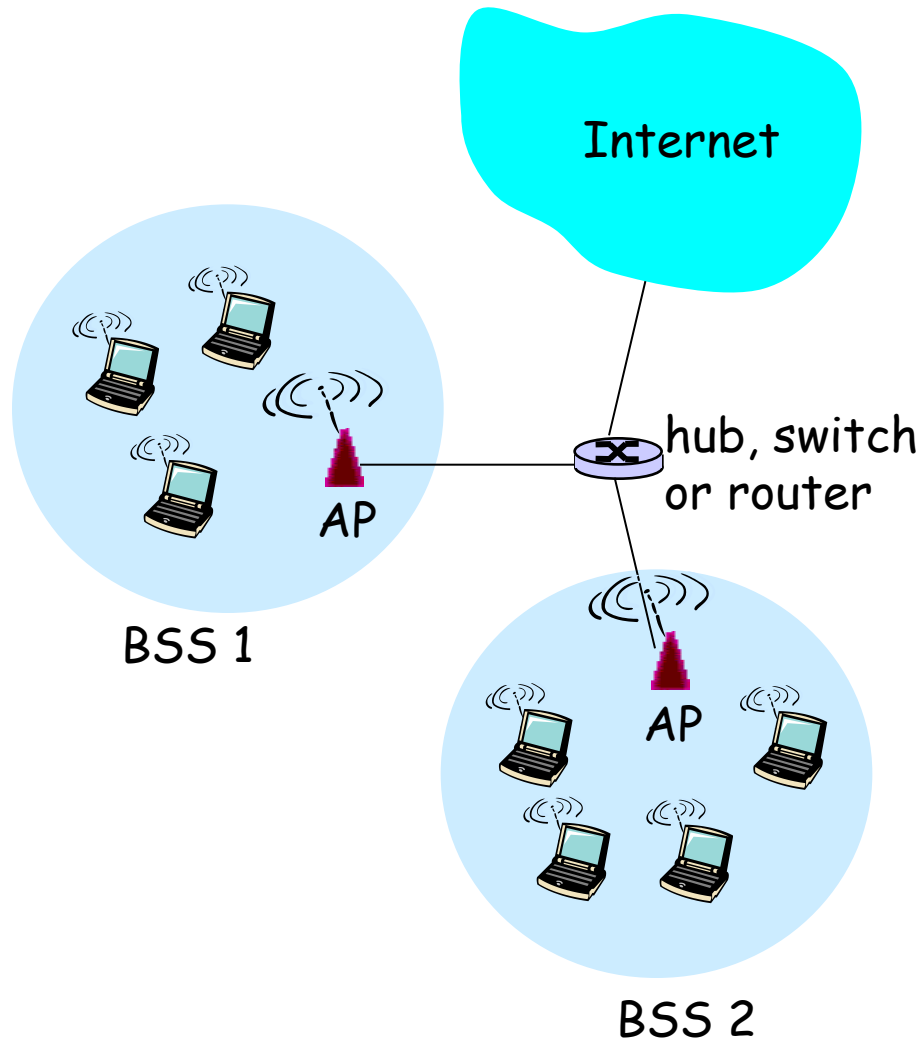
Wireless Link Characteristics

Differences from wired link

- **Decreasing signal strength:** radio signal attenuates as it propagates through matter (path loss)
- **Interference from other sources:** standardized wireless network frequencies (e.g., 2.4 GHz) shared by other devices (e.g., phone); devices (motors) interfere as well
- **Multi-path propagation:** radio signal reflects off objects ground, arriving at destination at slightly different times

.... make communication across (even a point to point) wireless link much more "difficult"

802.11 LAN architecture

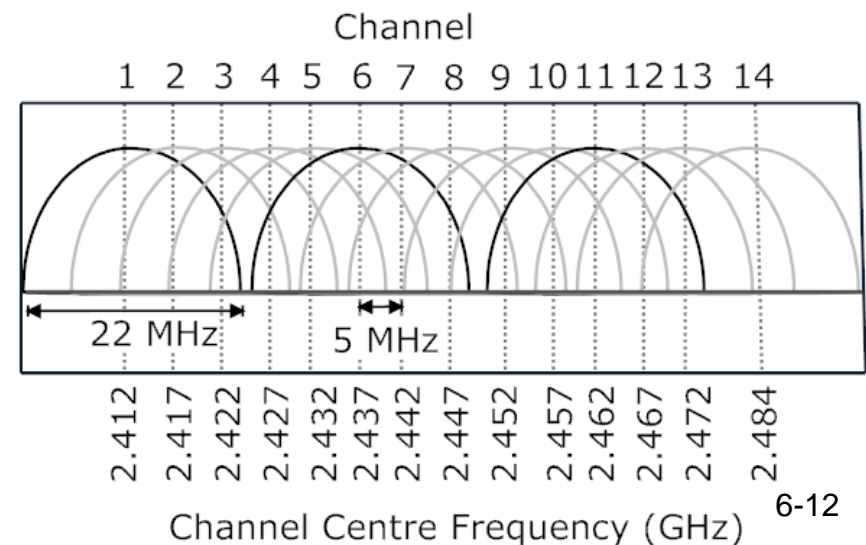


- ❑ Wireless host communicates with base station
 - base station = access point (AP)
- ❑ Basic Service Set (BSS) (aka "cell") in infrastructure mode contains:
 - wireless hosts
 - access point (AP)
 - ad hoc mode: hosts only

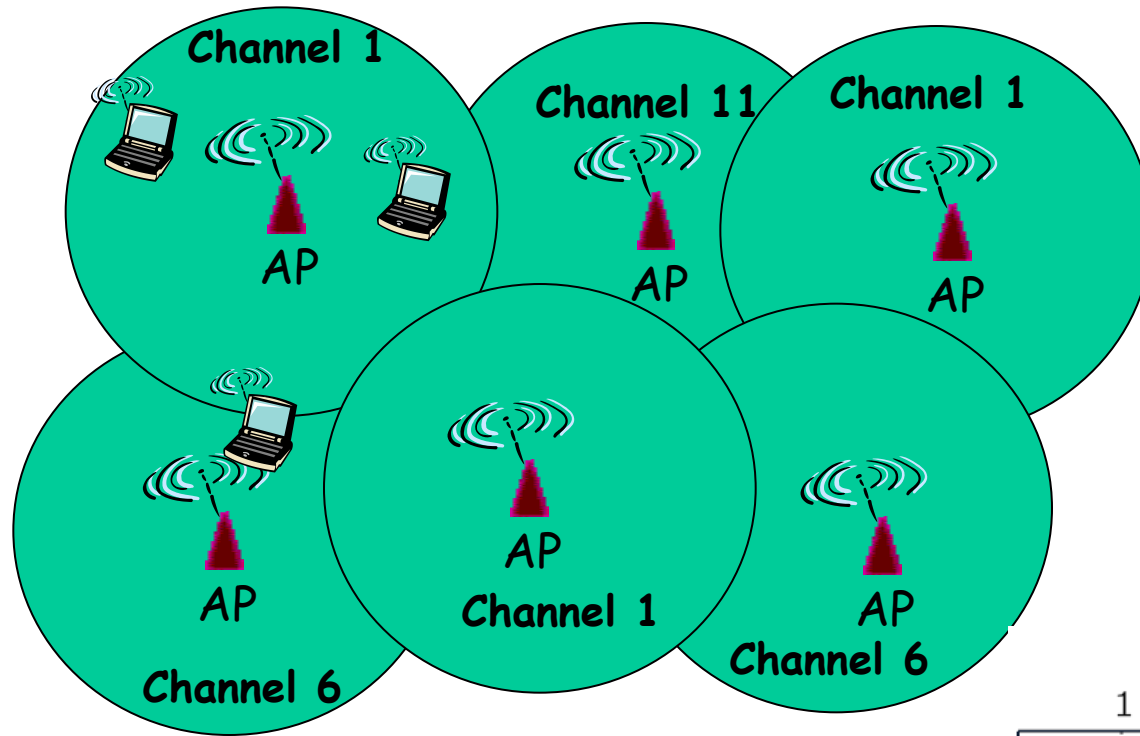
Wireless Cells

- ❖ 802.11b has 11 channels
- ❖ Channels 1, 6, and 11 are non-overlapping

- ❖ Admin chooses frequency for AP

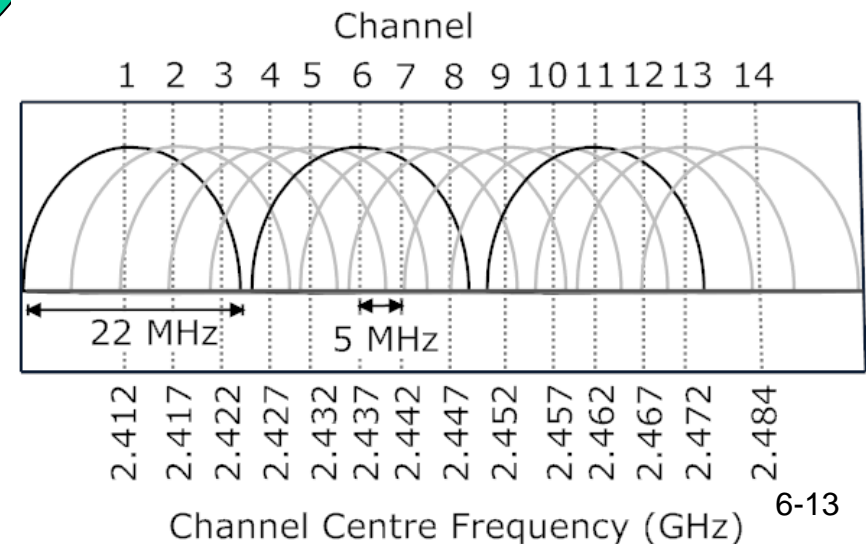


Wireless Cells



- ❖ 802.11b has 11 channels
- ❖ Channels 1, 6, and 11 are non-overlapping
- ❖ Each AP coverage area is called a "cell"
- ❖ Wireless nodes can roam between cells

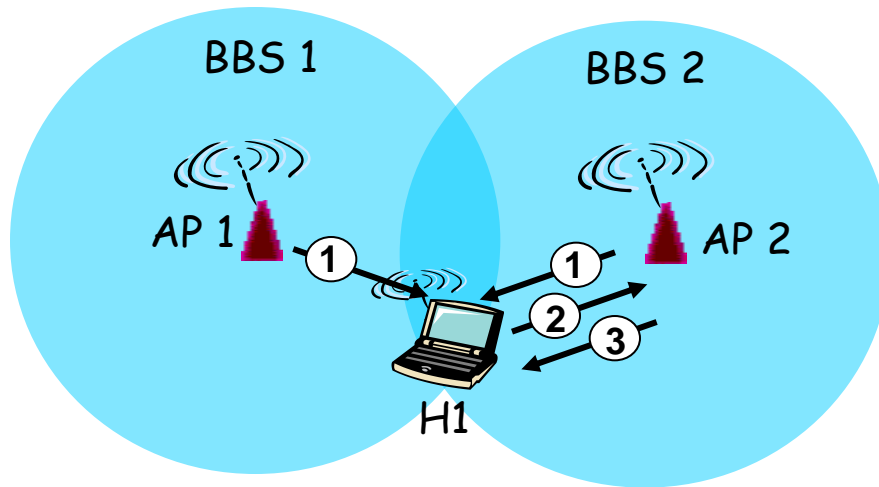
- ❖ Admin chooses frequency for AP
- ❖ Interference possible: channel can be same as that chosen by neighboring AP!



802.11: Channels, association

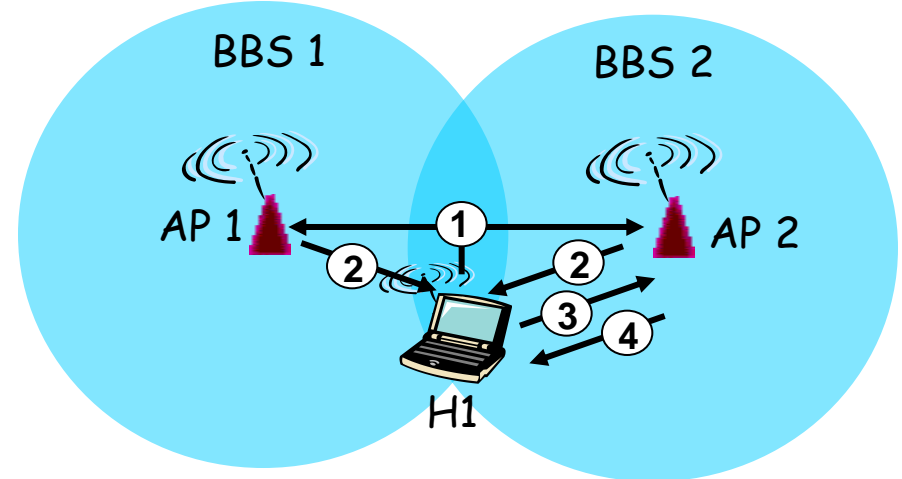
- ❖ host: must *associate* with an AP
 - scans channels, listening for *beacon frames* containing AP's name (SSID) and MAC address
 - selects AP to associate with
 - may perform authentication
 - will typically run DHCP to get IP address in AP's subnet

802.11: passive/active scanning



Passive Scanning:

- (1) beacon frames sent from APs
- (2) association Request frame sent:
H1 to selected AP
- (3) association Response frame sent:
H1 to selected AP



Active Scanning:

- (1) Probe Request frame broadcast
from H1
- (2) Probes response frame sent
from APs
- (3) Association Request frame
sent: H1 to selected AP
- (4) Association Response frame
sent: H1 to selected AP

IEEE 802.11: multiple access

- ❖ avoid collisions: 2+ nodes transmitting at same time
- ❖ 802.11: CSMA - sense before transmitting
 - don't collide with ongoing transmission by other node
- ❖ 802.11: *no* collision detection!
 - difficult to receive (sense collisions) when transmitting due to weak received signals (fading)
 - can't sense all collisions in any case: hidden terminal, fading
 - goal: *avoid collisions*: CSMA/C(ollision)A(voidance)

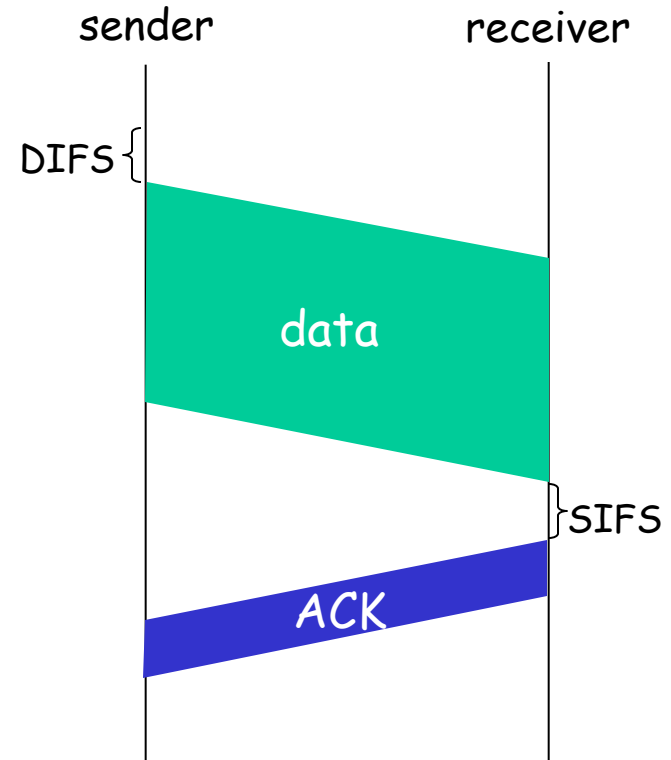
IEEE 802.11 MAC Protocol: CSMA/CA

802.11 sender

- 1 if sense channel idle for **DIFS** then
transmit entire frame (no CD)
- 2 if sense channel busy then
start random backoff time
timer counts down while channel idle
transmit when timer expires
- 3 if no **ACK** then increase random backoff
interval, repeat step 2

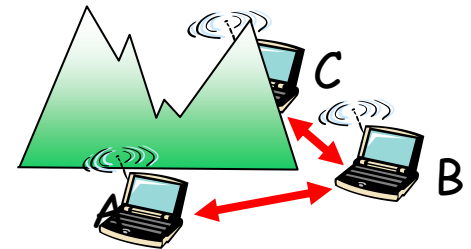
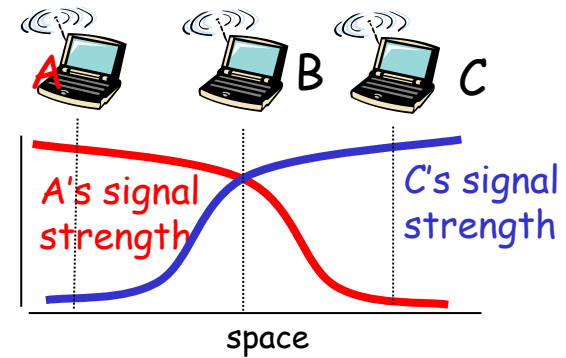
802.11 receiver

- if frame received OK
return **ACK** after **SIFS**
(service model is connectionless, acked)



Hidden Terminal Problem

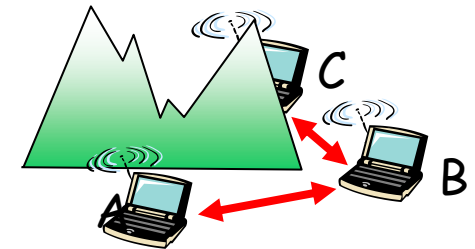
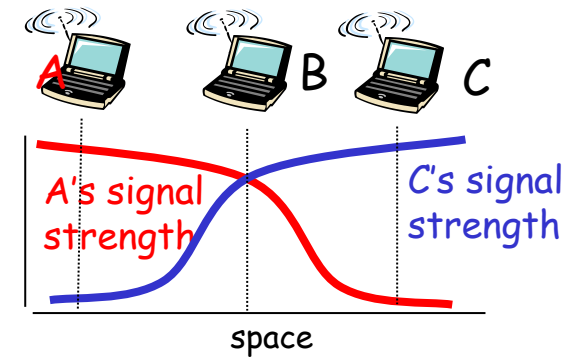
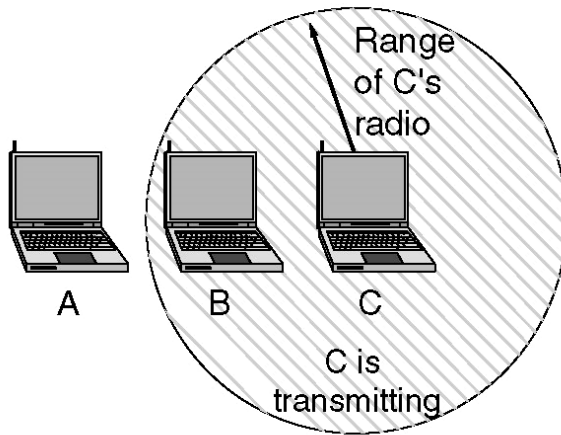
Hidden terminal problem (ad-hoc and WLAN)



Hidden Terminal Problem

Hidden terminal problem (ad-hoc and WLAN)

A wants to send to B
but cannot hear that
B is busy

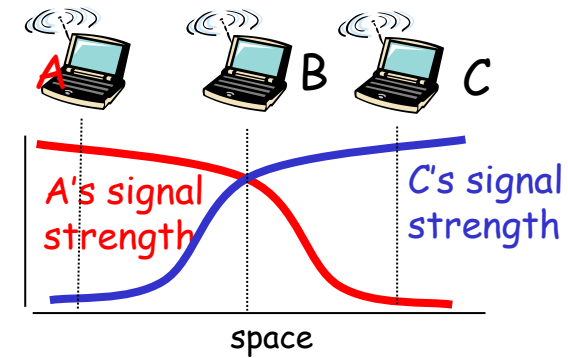


Hidden Terminal Problem

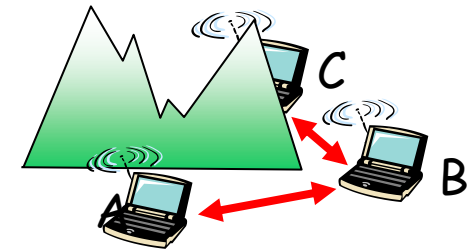
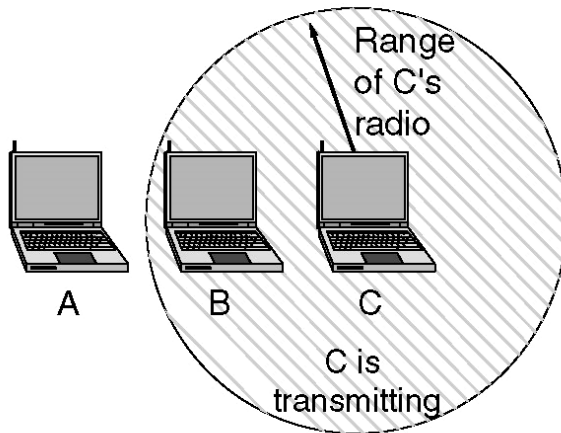
Hidden terminal problem (ad-hoc and WLAN)

- medium free near the transmitter
- medium not free near the receiver

=> *Packet collision*



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Hidden Terminal Problem

Hidden terminal problem (ad-hoc and WLAN)

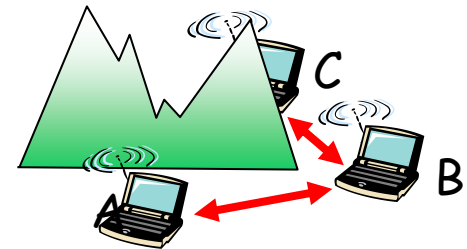
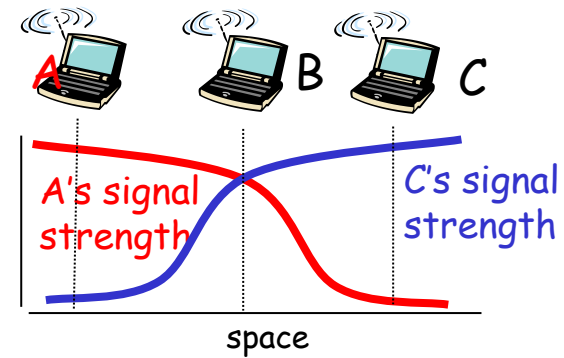
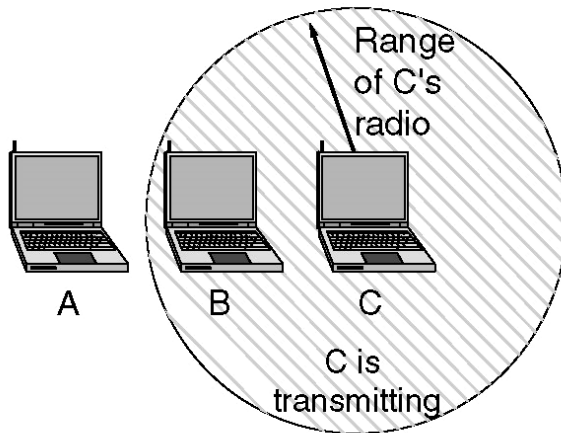
- medium free near the transmitter
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=> *Packet collision*

Possible solution:

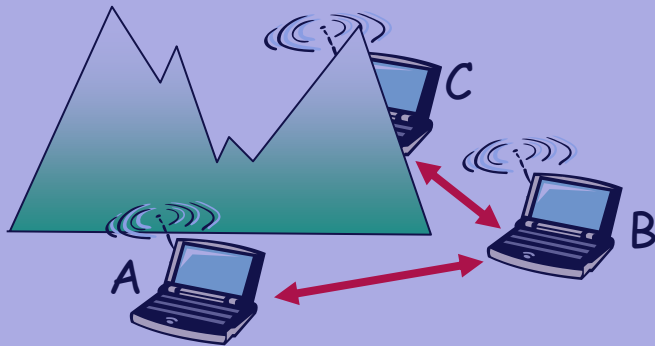
- MAC scheme using RTS-CTS scheme

A wants to send to B
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B is busy



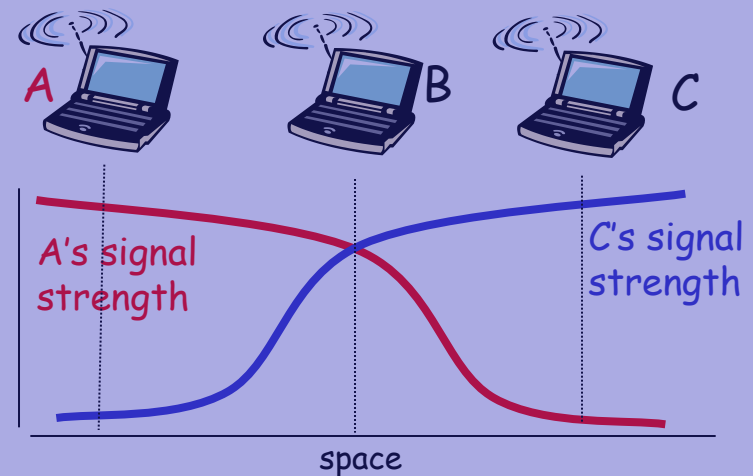
Wireless Network Characteristics

Multiple wireless senders and receivers create additional problems (beyond multiple access):



Hidden terminal problem

- ☐ A and B can hear each other
- ☐ B and C can hear each other
- ☐ A and C can't hear each other
- ☐ thus A and C are unaware of their interference at B



Signal fading:

- ☐ A and B hear each other
- ☐ B and C hear each other
- ☐ A and C can't hear each other interfering at B

Avoiding collisions (more)

idea: allow sender to “reserve” channel rather than random access of data frames: avoid collisions of long data frames

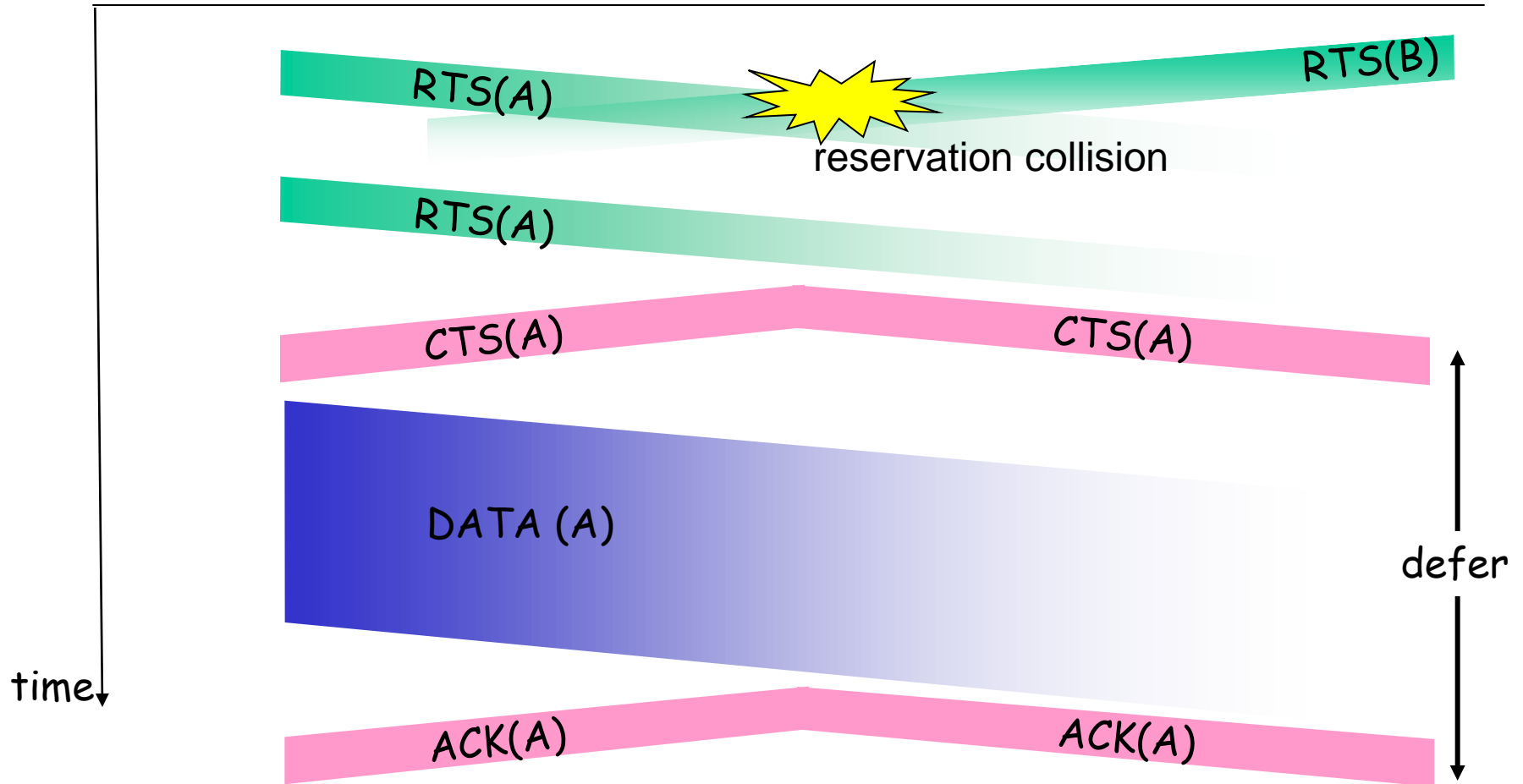
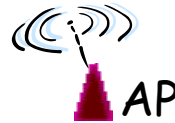
Avoiding collisions (more)

idea: allow sender to “reserve” channel rather than random access of data frames: avoid collisions of long data frames

- ❖ sender first transmits *small* request-to-send (RTS) packets to base station using CSMA
 - RTS may still collide with each other (but they're short)
- ❖ BS broadcasts clear-to-send CTS to host in response to RTS
- ❖ RTS heard by all nodes because of broadcast property
 - sender transmits (large) data frame
 - other stations defer transmissions until it is done

Avoid data frame collisions completely
using small reservation packets!

Collision Avoidance: RTS-CTS exchange



Exposed Terminal Problems

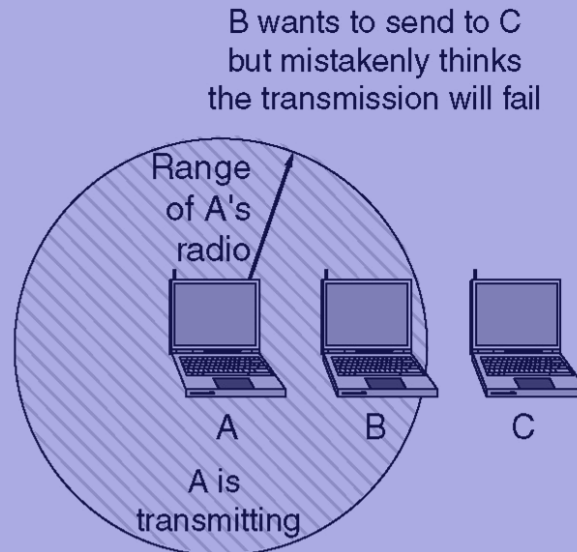
Exposed terminal problem - ad-hoc and WLAN

- medium free near the receiver
- medium busy near the transmitter

=> *Waist of bandwidth*

Possible solutions:

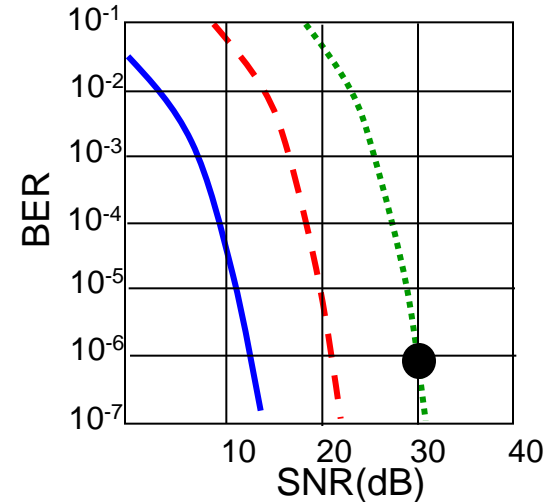
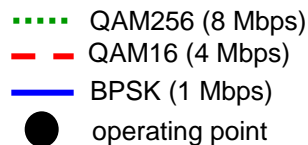
- directional antennas
- separate channels for control and data



802.11: advanced capabilities

Rate Adaptation

- ❖ base station, mobile dynamically change transmission rate (physical layer modulation technique) as mobile moves, SNR varies



1. SNR decreases, BER increase as node moves away from base station
2. When BER becomes too high, switch to lower transmission rate but with lower BER

802.11: advanced capabilities

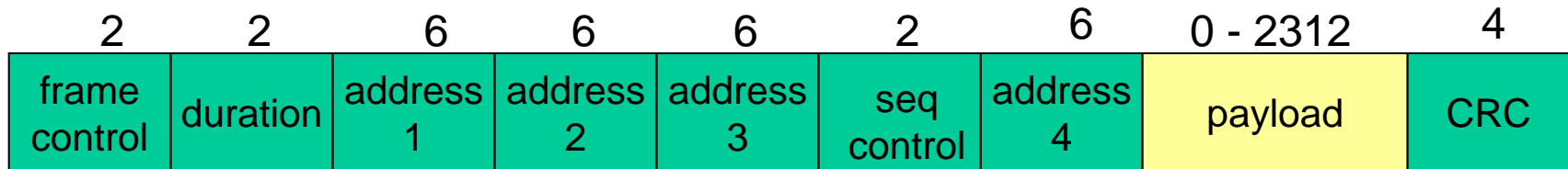
Power Management

- ❑ node-to-AP: "I am going to sleep until next beacon frame"
 - AP knows not to transmit frames to this node
 - node wakes up before next beacon frame
- ❑ beacon frame: contains list of mobiles with AP-to-mobile frames waiting to be sent
 - Every 100ms (250 μ s wakeup time)
 - node will stay awake if AP-to-mobile frames to be sent; otherwise sleep again until next beacon frame
 - Explicit pull request

Note: Nodes with nothing to send/receive can save 99% of energy

More slides ...

802.11 frame: addressing



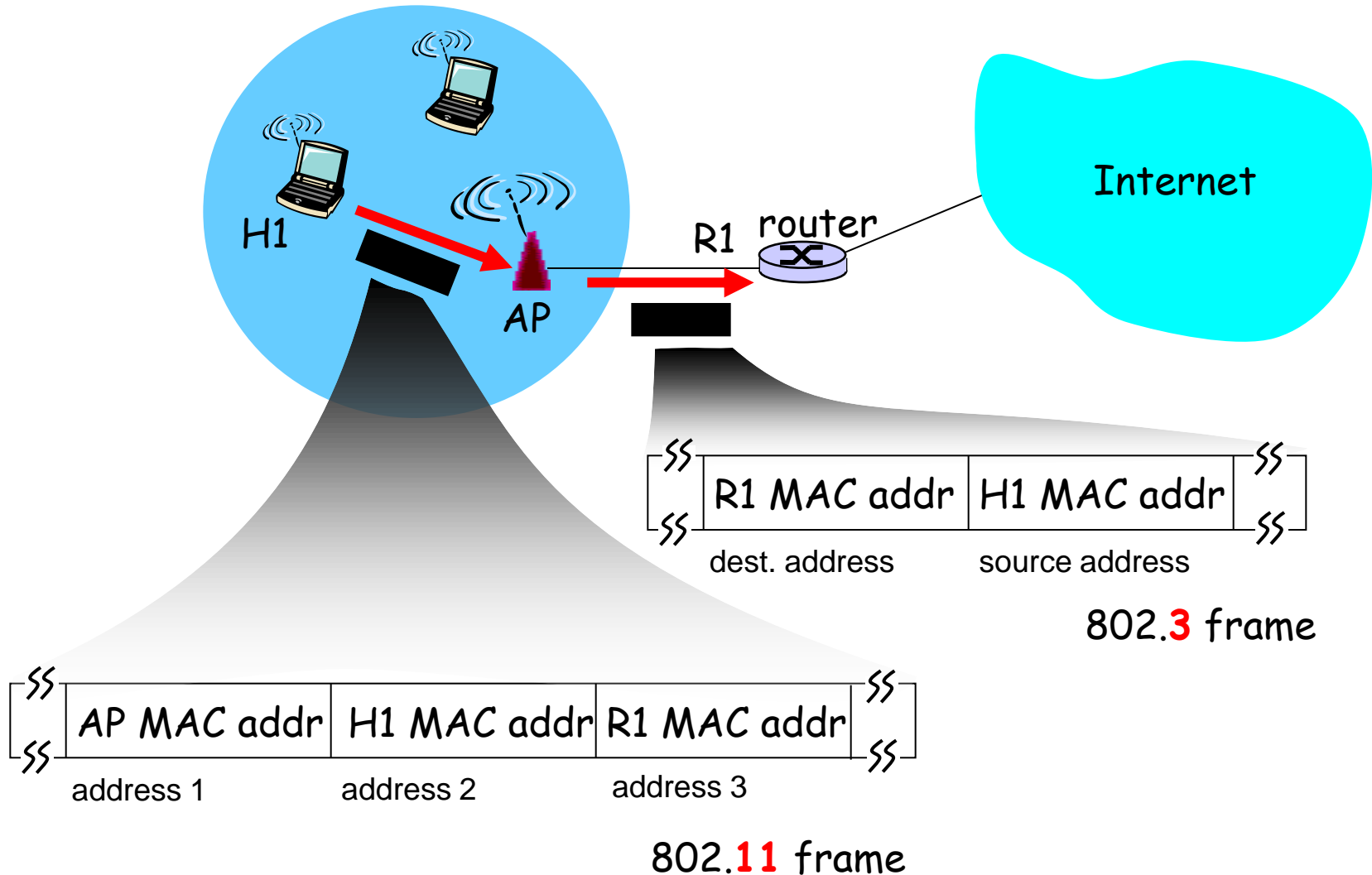
Address 1: MAC address of wireless host or AP to receive this frame

Address 2: MAC address of wireless host or AP transmitting this frame

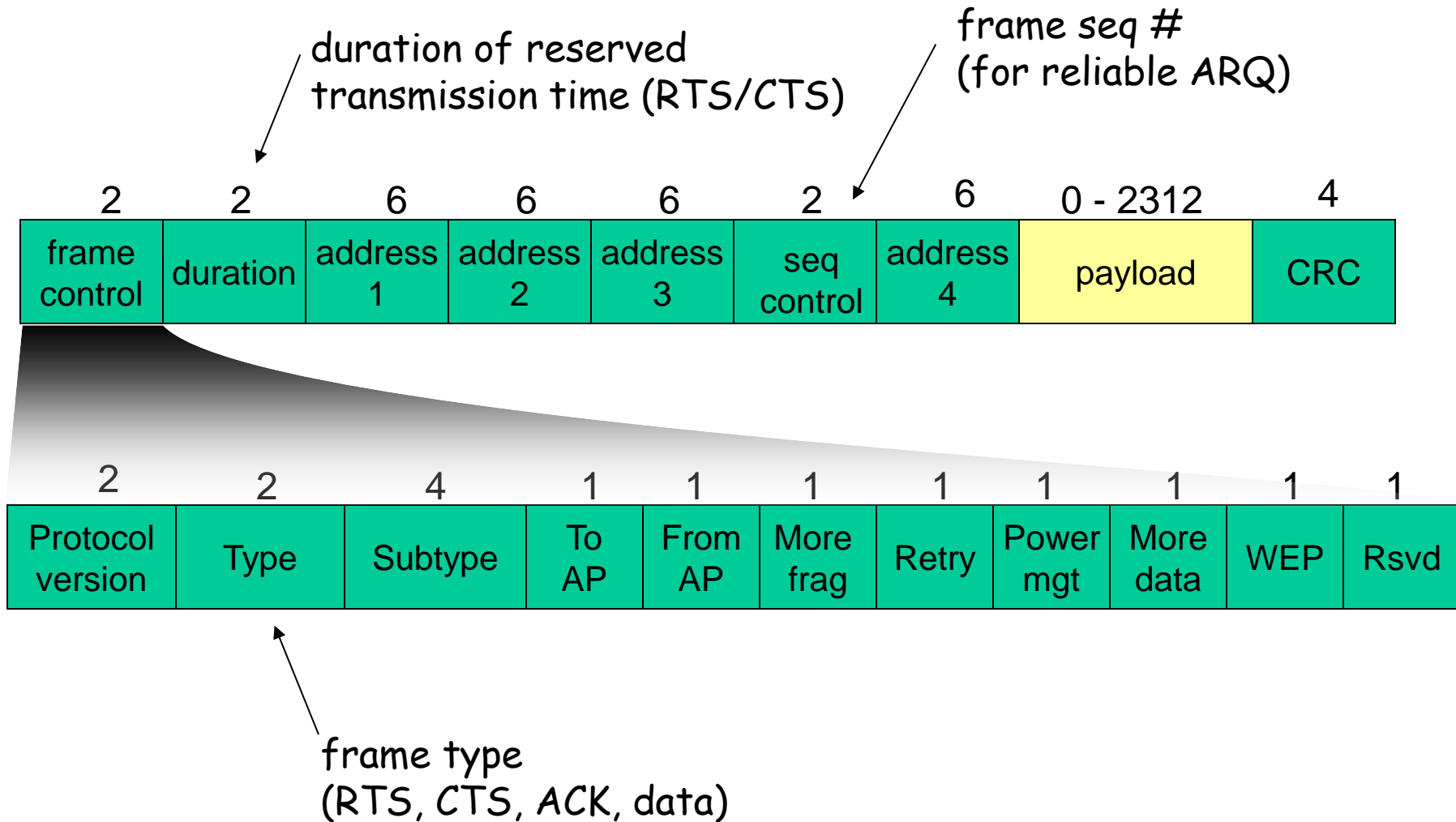
Address 3: MAC address of router interface to which AP is attached

Address 3: used only in ad hoc mode

802.11 frame: addressing

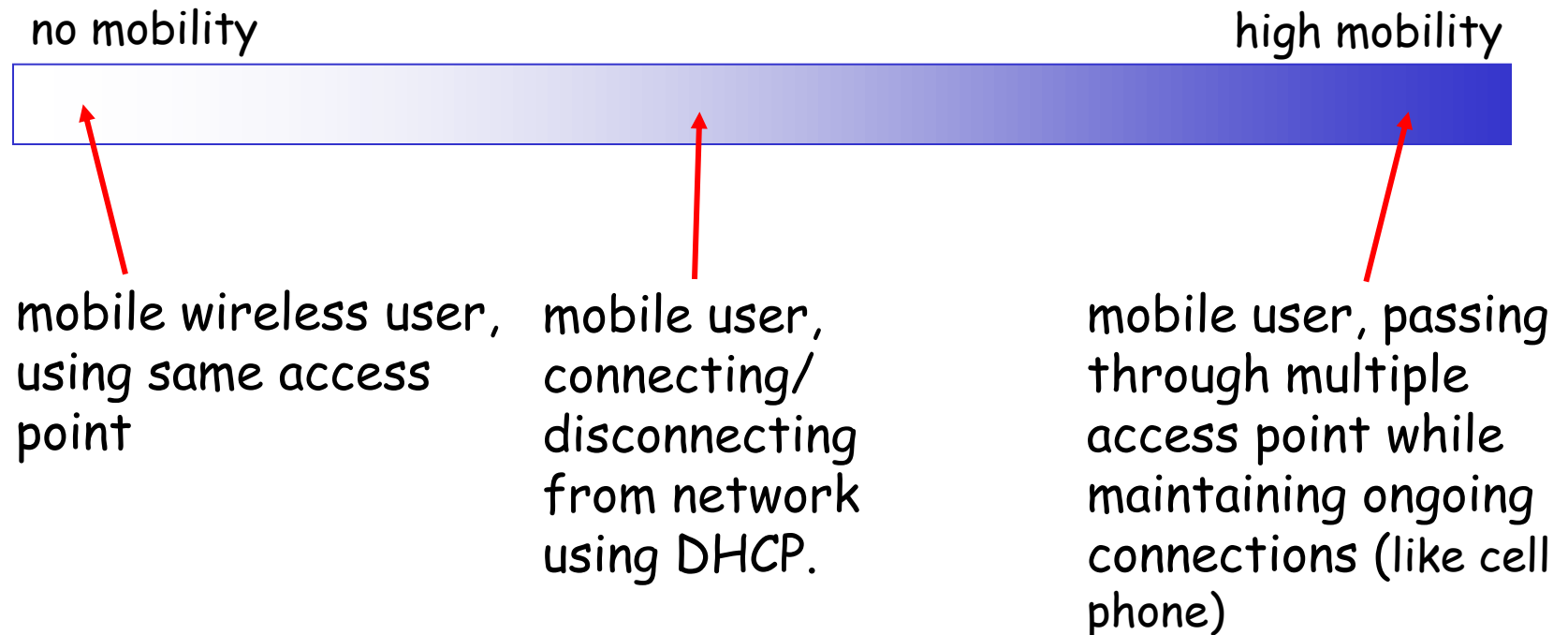


802.11 frame: more



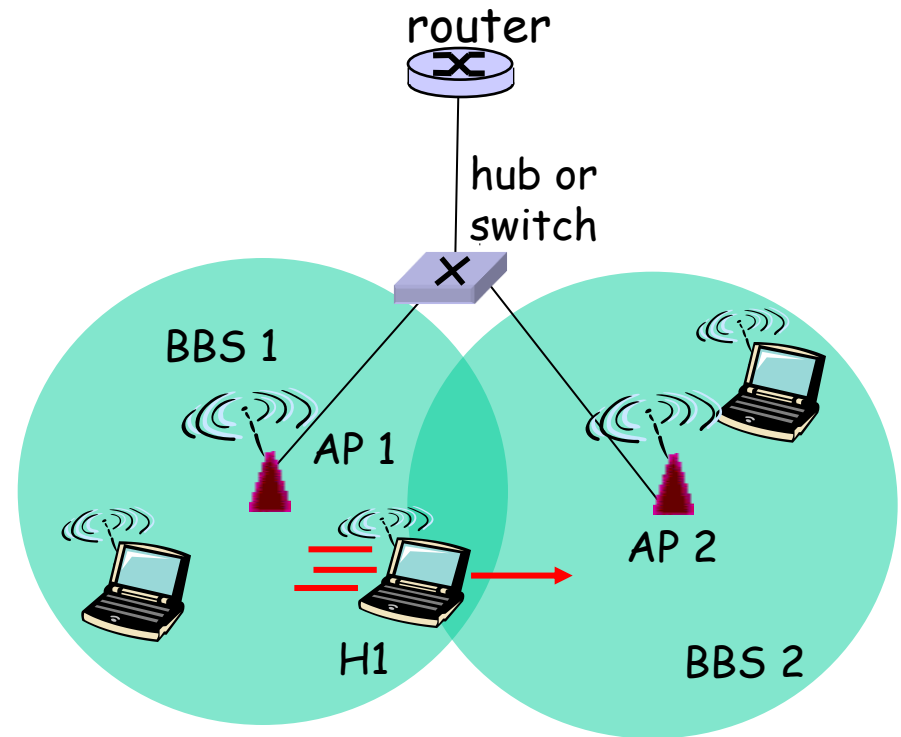
What is mobility?

❖ spectrum of mobility, from the *network* perspective:



802.11: mobility within same subnet

- ❖ H1 remains in same IP subnet: IP address can remain same
- ❖ switch: which AP is associated with H1?
 - self-learning (Ch. 5): switch will see frame from H1 and "remember" which switch port can be used to reach H1

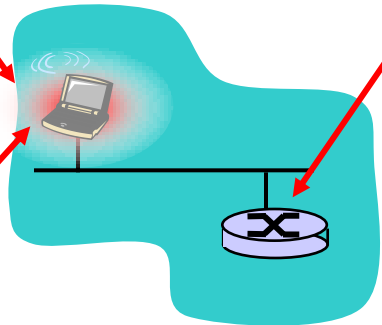


Mobility: Vocabulary

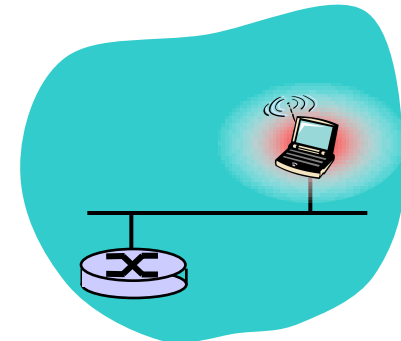
home network: permanent
"home" of mobile
(e.g., 128.119.40/24)

home agent: entity that will
perform mobility functions on
behalf of mobile, when mobile
is remote

permanent address:
address in home
network, *can always* be
used to reach mobile
e.g., 128.119.40.186



wide area
network




correspondent

Mobility: more vocabulary

permanent address: remains constant (e.g., 128.119.40.186)

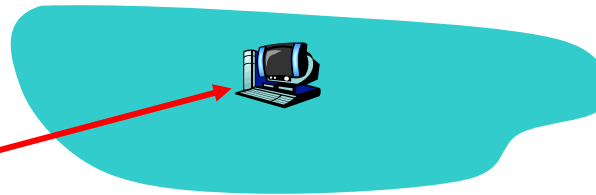
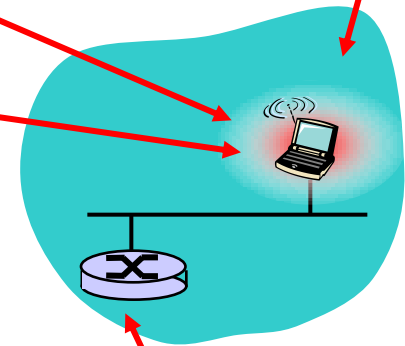
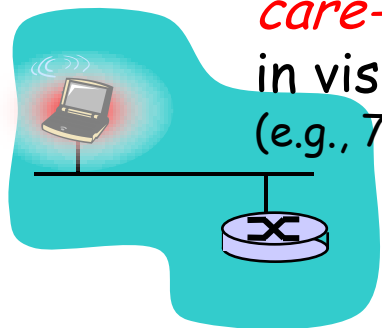
visited network: network in which mobile currently resides (e.g., 79.129.13/24)

care-of-address: address in visited network. (e.g., 79.129.13.2)

wide area network

correspondent: wants to communicate with mobile

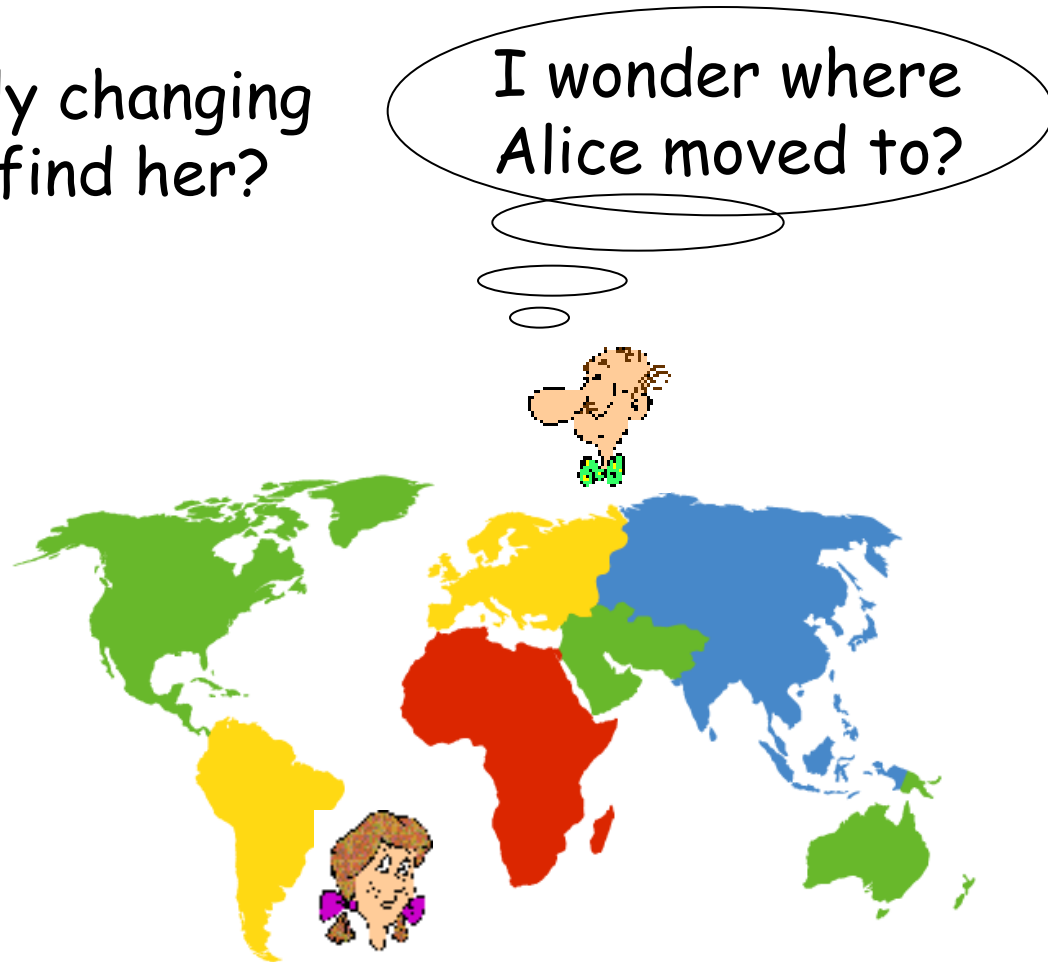
foreign agent: entity in visited network that performs mobility functions on behalf of mobile.



How do *you* contact a mobile friend:

Consider friend frequently changing addresses, how do you find her?

- ❖ search all phone books?
- ❖ call her parents?
- ❖ expect her to let you know where he/she is?



Mobility: approaches

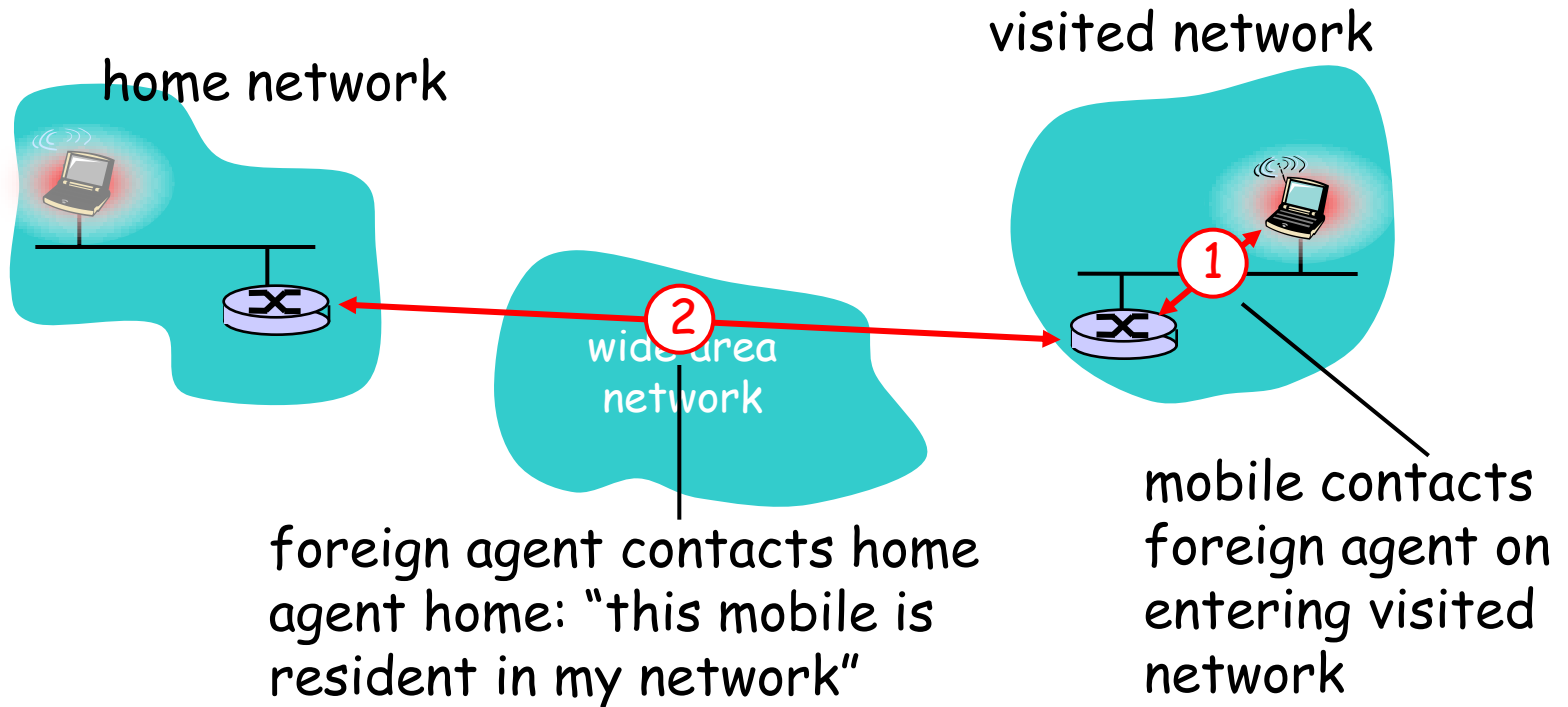
- ❖ *Let routing handle it:* routers advertise permanent address of mobile-nodes-in-residence via usual routing table exchange.
 - routing tables indicate where each mobile located
 - no changes to end-systems
- ❖ *Let end-systems handle it:*
 - *indirect routing:* communication from correspondent to mobile goes through home agent, then forwarded to remote
 - *direct routing:* correspondent gets foreign address of mobile, sends directly to mobile

Mobility: approaches

- ❖ *Let routing handle it:* routers advertise permanent address of mobile, mobile residence via usual routing table entries
 - routing table entry for each mobile location
 - no changes to end systems
- ❖ *let end-systems handle it:*
 - *indirect routing:* communication from correspondent to mobile goes through home agent, then forwarded to remote
 - *direct routing:* correspondent gets foreign address of mobile, sends directly to mobile

not
scalable
to millions of
mobiles

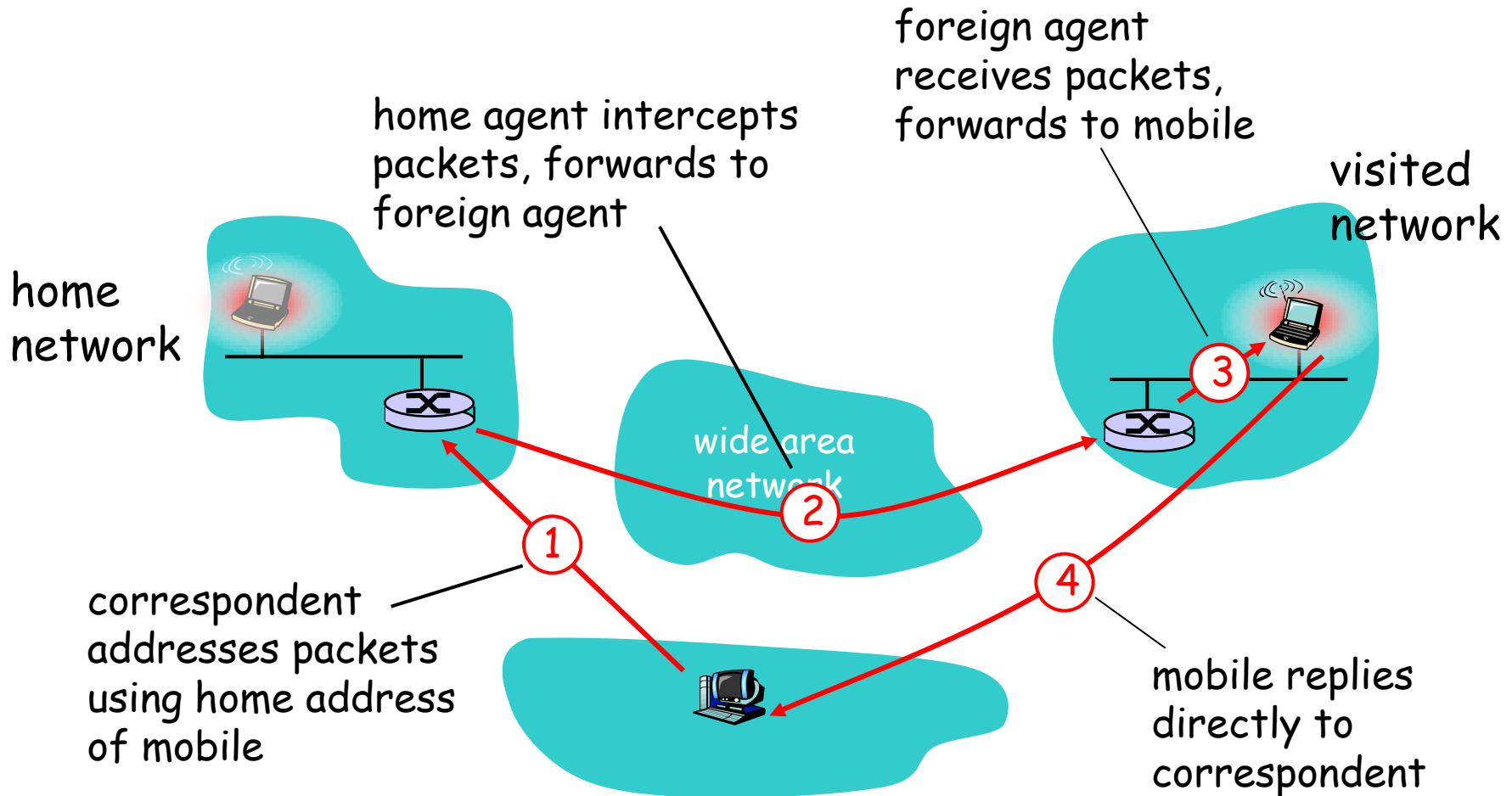
Mobility: registration



End result:

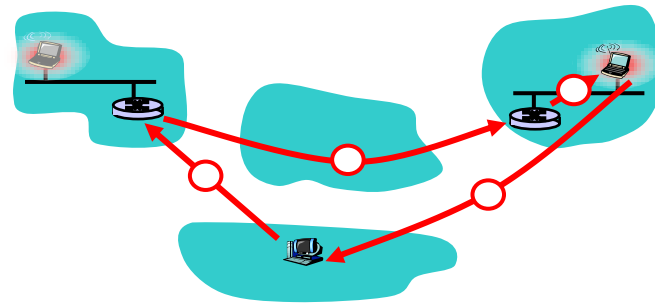
- ❖ Foreign Agent (FA) knows about mobile
- ❖ Home Agent (HA) knows location of mobile

Mobility via Indirect Routing



Indirect Routing: comments

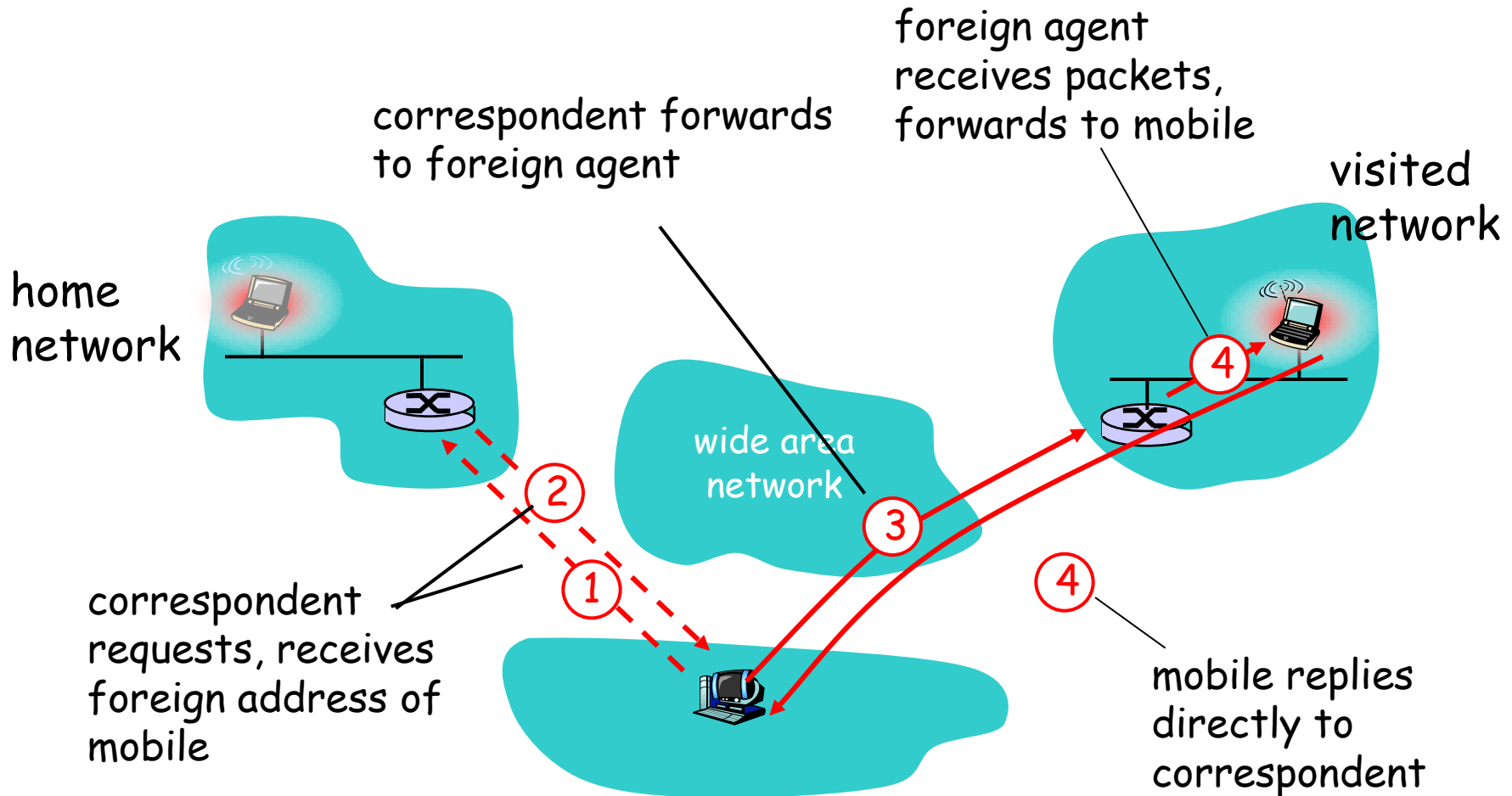
- ❖ Mobile uses two addresses:
 - **permanent address**: used by correspondent (hence mobile location is *transparent* to correspondent)
 - **care-of-address**: used by home agent to forward datagrams to mobile
- ❖ foreign agent functions may be done by mobile itself
- ❖ **triangle routing**: correspondent-home-network-mobile
 - inefficient when correspondent, mobile are in same network



Indirect Routing: moving between networks

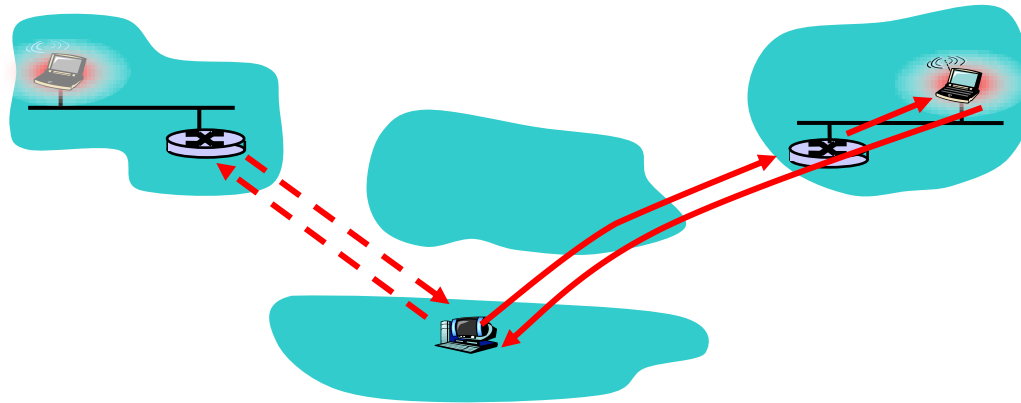
- ❖ suppose mobile user moves to another network
 - registers with new foreign agent
 - new foreign agent registers with home agent
 - home agent update care-of-address for mobile
 - packets continue to be forwarded to mobile (but with new care-of-address)
- ❖ mobility, changing foreign networks
transparent: *ongoing connections can be maintained!*

Mobility via Direct Routing



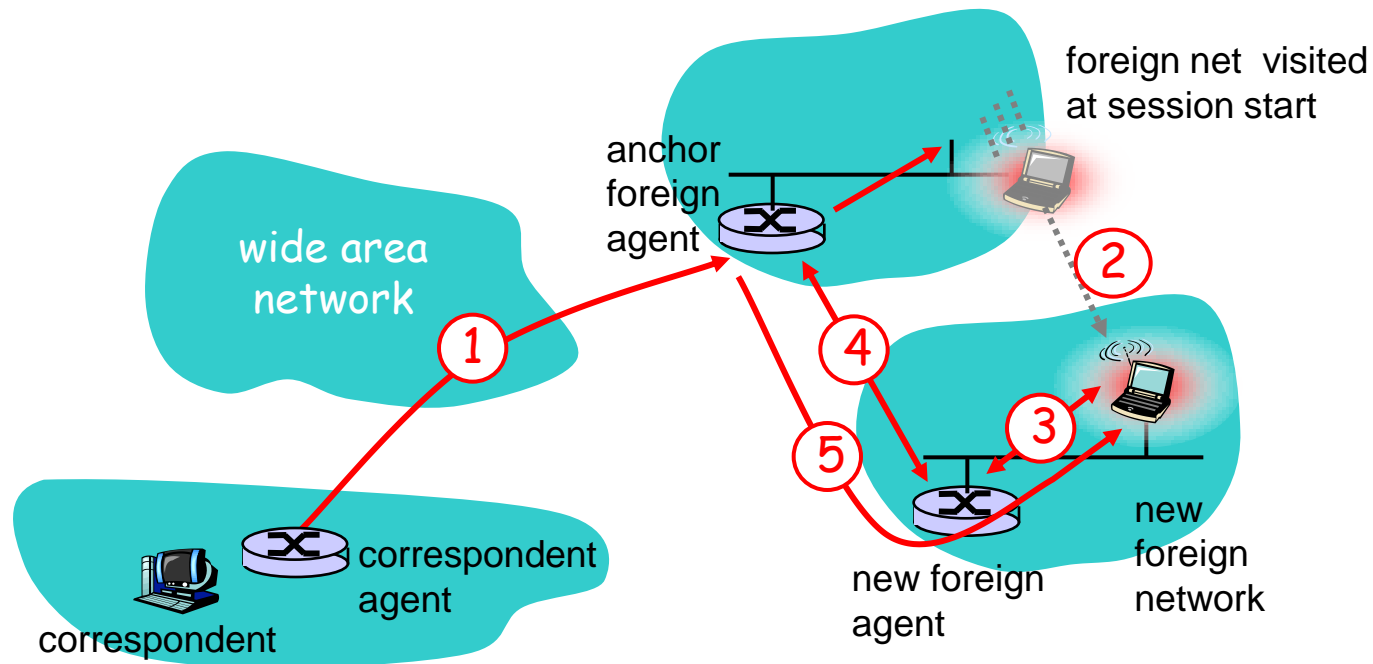
Mobility via Direct Routing: comments

- ❖ overcome triangle routing problem
- ❖ **non-transparent to correspondent:**
correspondent must get care-of-address from home agent
 - what if mobile changes visited network?



Accommodating mobility with direct routing

- ❖ anchor foreign agent: FA in first visited network
- ❖ data always routed first to anchor FA
- ❖ when mobile moves: new FA arranges to have data forwarded from old FA (chaining)



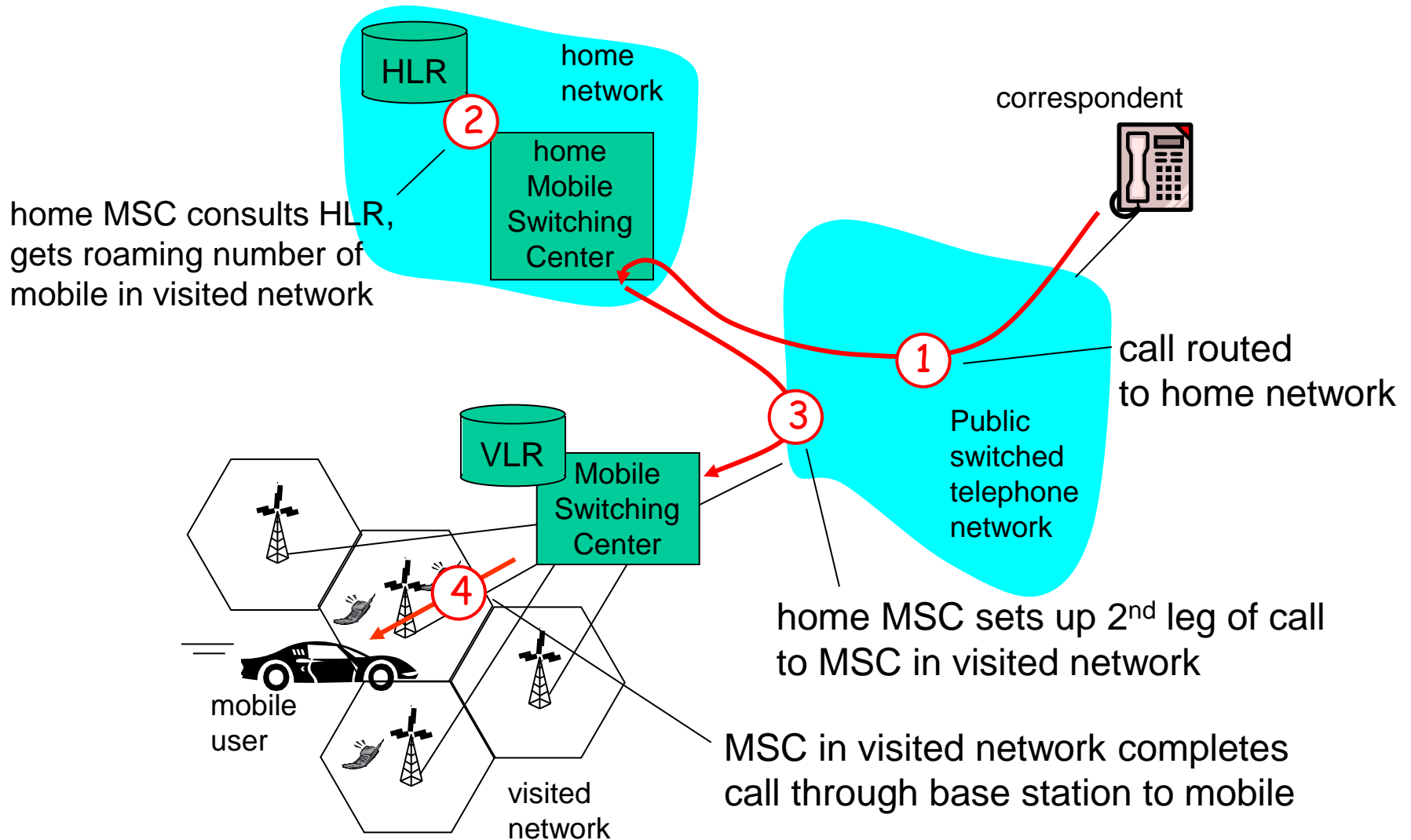
Mobile IP

- ❖ RFC 3220
- ❖ has many features we've seen:
 - home agents, foreign agents, foreign-agent registration, care-of-addresses, encapsulation (packet-within-a-packet)
- ❖ three components to standard:
 - indirect routing of datagrams
 - agent discovery
 - registration with home agent

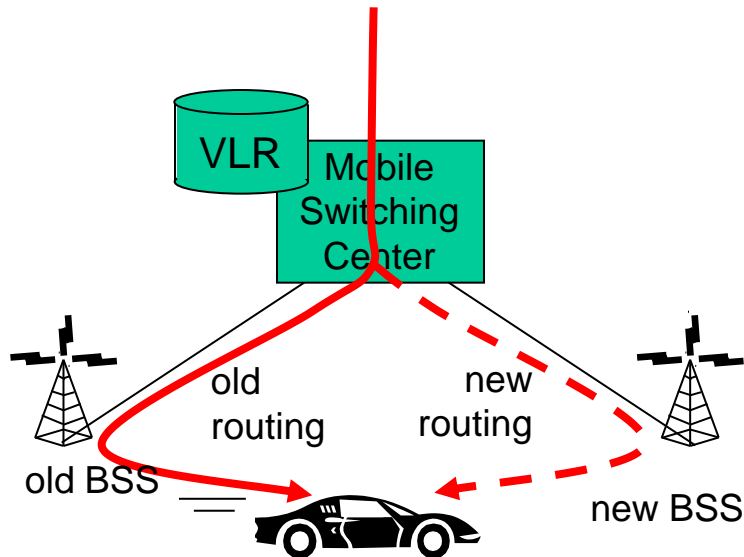
Handling mobility in cellular networks

- ❖ *home network*: network of cellular provider you subscribe to (e.g., Sprint PCS, Verizon)
 - *home location register (HLR)*: database in home network containing permanent cell phone #, profile information (services, preferences, billing), information about current location (could be in another network)
- ❖ *visited network*: network in which mobile currently resides
 - *visitor location register (VLR)*: database with entry for each user currently in network
 - could be home network

GSM: indirect routing to mobile

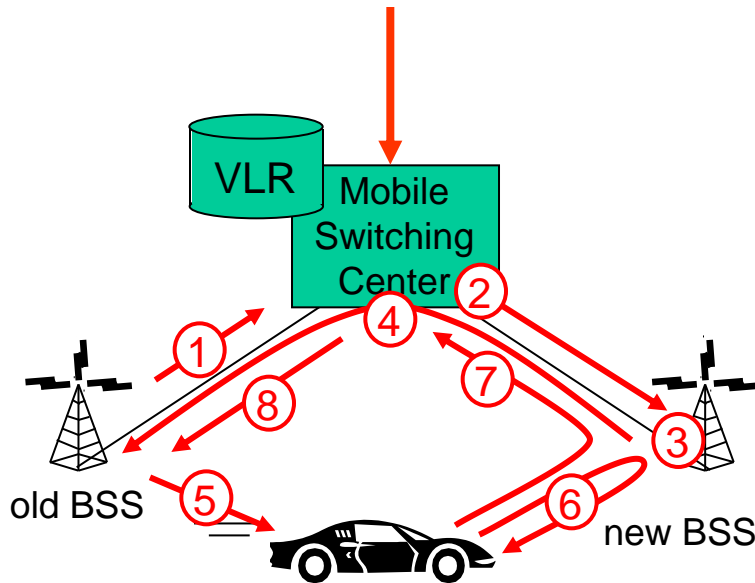


GSM: handoff with common MSC



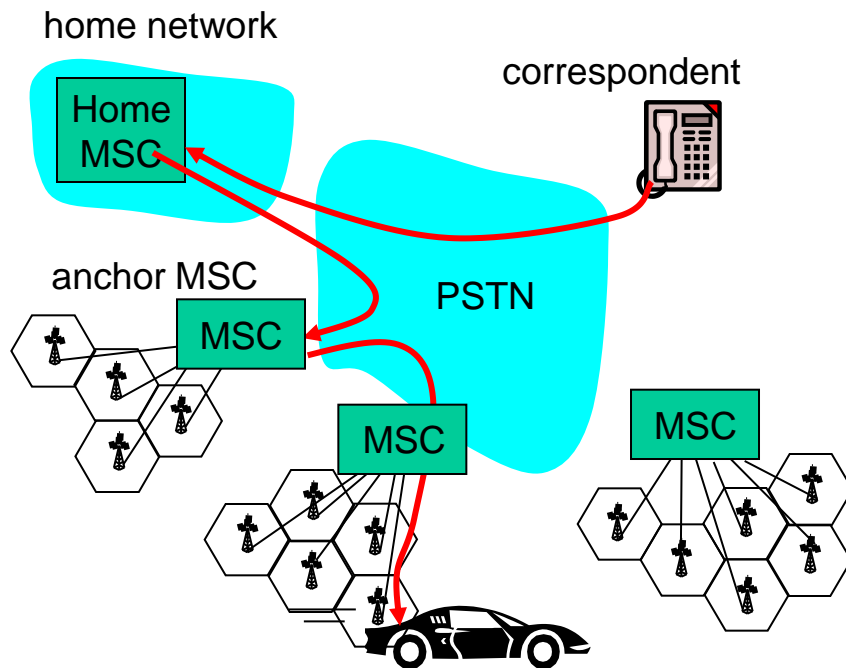
- ❖ Handoff goal: route call via new base station (without interruption)
- ❖ reasons for handoff:
 - stronger signal to/from new BSS (continuing connectivity, less battery drain)
 - load balance: free up channel in current BSS
 - GSM doesn't mandate why to perform handoff (policy), only how (mechanism)
- ❖ handoff initiated by old BSS

GSM: handoff with common MSC



1. old BSS informs MSC of impending handoff, provides list of 1+ new BSSs
2. MSC sets up path (allocates resources) to new BSS
3. new BSS allocates radio channel for use by mobile
4. new BSS signals MSC, old BSS: ready
5. old BSS tells mobile: perform handoff to new BSS
6. mobile, new BSS signal to activate new channel
7. mobile signals via new BSS to MSC: handoff complete. MSC reroutes call
8. MSC-old-BSS resources released

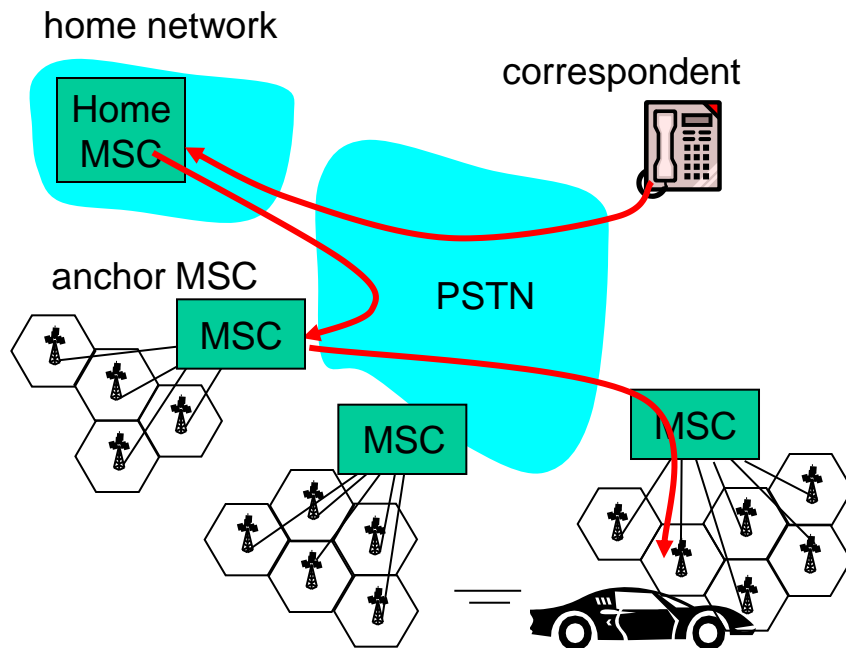
GSM: handoff between MSCs



(a) before handoff

- ❖ *anchor MSC*: first MSC visited during cal
 - call remains routed through anchor MSC
- ❖ new MSCs add on to end of MSC chain as mobile moves to new MSC
- ❖ IS-41 allows optional path minimization step to shorten multi-MSC chain

GSM: handoff between MSCs



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More slides ...

Wireless Networking Technologies

- ❖ Mobile devices - laptop, PDA, cellular phone, wearable computer, sensors, ...
- ❖ Operating modes
 - Infrastructure mode (Access Point (AP))
 - Ad hoc mode
- ❖ Access technology
 - Bluetooth (1 Mbps, up to 3m)
 - IEEE 802.11 (up to 54 Mbps, 20 - 100m)

Chapter 6: Wireless and Mobile Networks

Background:

- ❖ # wireless (mobile) phone subscribers now exceeds # wired phone subscribers!
- ❖ # wireless Internet-connected devices soon to exceed # wireline Internet-connected devices