

# TDTS04/11: Computer Networks

Instructor/examiner: Niklas Carlsson

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Office: B:476

Office Hours: TBA

Notes derived from "*Computer Networking: A Top Down Approach*", by Jim Kurose and Keith Ross, Addison-Wesley.

The slides are adapted and modified based on slides from the book's companion Web site, as well as modified slides by Anirban Mahanti and Carey Williamson.

# English

- ❑ The courses will be given in English ...
  - ❑ ... but first a few words in Swedish ... (apologies to exchange and international students)
- ❑ Many reasons, including (but not limited to):
  - Terminology mostly in English
  - Google and the literature will give you **many more** and **much better** answers ...
  - Much better lecture quality ... especially with good book + slides in English
  - (+ **many** more reasons)
- ❑ Great opportunity to practice
  - ❑ Can still ask questions in Swedish ... (please remind me to rephrase the question in English)
  - ❑ Understanding is the focus (not your language skills ...)

# People

## ❑ Examiner and lecturer

- Niklas Carlsson
- Research area: Design, modeling, and performance evaluation of distributed systems and networks

## ❑ Lab assistants TDTS04

- Vengatanathan Krishnamoorthi, PhD student
- Sergiu Rafiliu, PhD student

## ❑ Lab assistant TDTS11

- Rahul Hiran, PhD student

## ❑ Director of studies

- Patrick Lambrix

# Course Overview(s)

- ❑ Written exam
  - Grads: 'fail', 3, 4, 5.
- ❑ Five (5) or four (4) mandatory lab assignments
  - Must pass all assignments
  - Thirteen (13) or eight (8) lab opportunities
  - Register on webreg. (Deadline on Friday!!)
  - TDTS11: one (1) optional assignment
    - Up to 4 bonus marks for exam
- ❑ Fourteen (14) or ten (10) lectures
  - Ten (10) network "focus"
  - Four (4) distributed systems "focus"
  - + Last lecture with some exam preparation
- ❑ See website for more information ...

# My expectations

## ❑ Read textbook

- Good textbook, written by highly regarded researchers in the field
- Lots of content
- Not time to cover everything during lectures

## ❑ Work hard

- Pay attention during lectures
- Make sure you **understand** the material
- Start assignments early (some will take time)

## ❑ Follow deadlines and office hours

# What to expect? (What will be covered?)

- ❑ Design principles for computer networks
  - Conceptual view of Internet architecture
- ❑ Design, resource, and performance tradeoffs
  - General working knowledge of protocols/applications
  - Detailed knowledge of selected protocols/applications
  - Some practical hands on experience
- ❑ Glimpse into the future of the Internet
  - Emerging trends and technologies



# Roadmap (today's lecture)

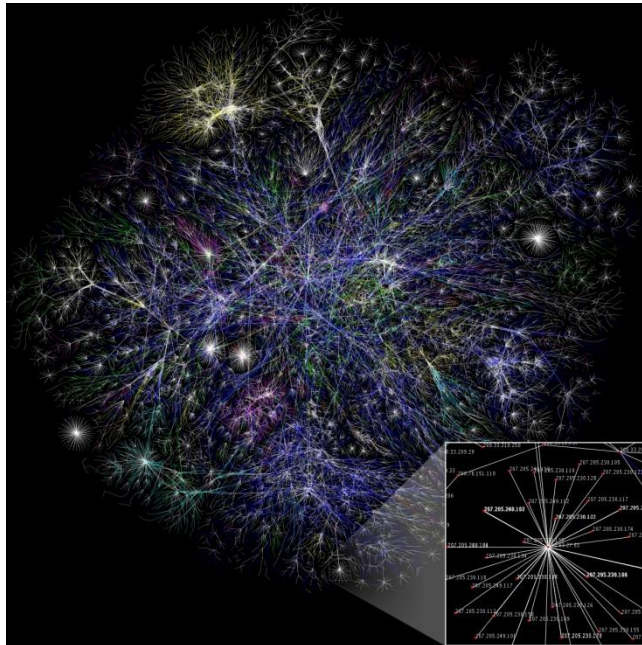
- What is a Computer Network?
- Applications of Networking
- Classification of Networks
- Layered Architecture (and Protocols)
- Network Core
- Delay & Loss in Packet-switched Networks
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- E.g., <https://www.youtube.com/watch?v=w42EsCDAhB4>

- So, what are computer networks?

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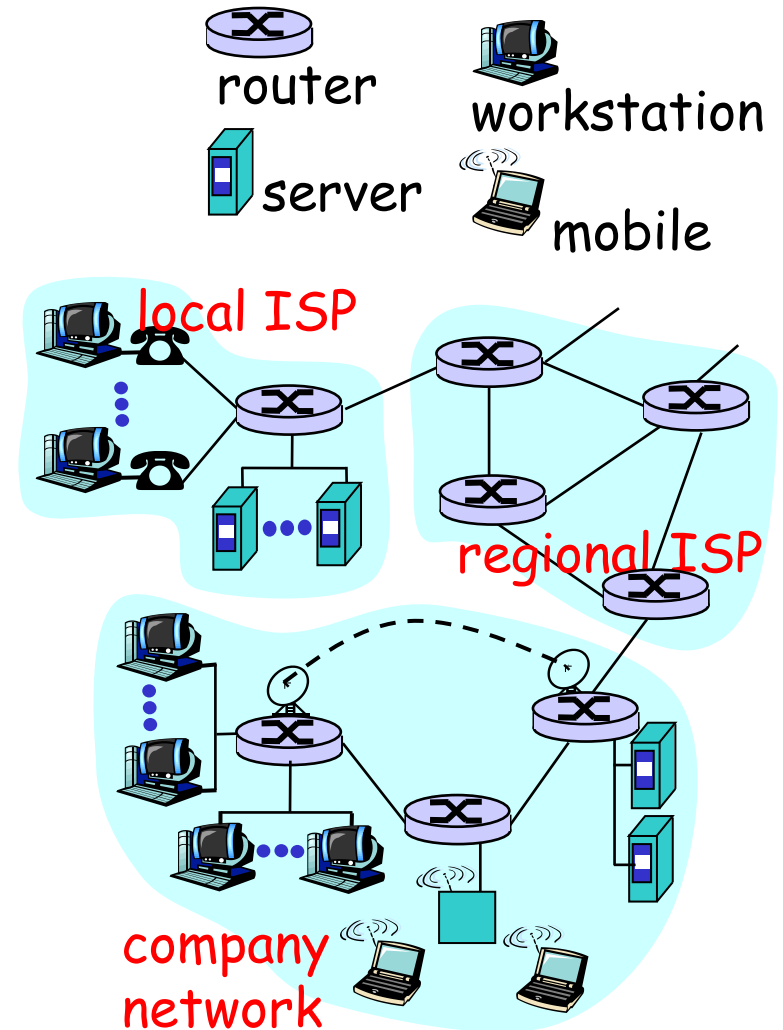


# Computer Network?



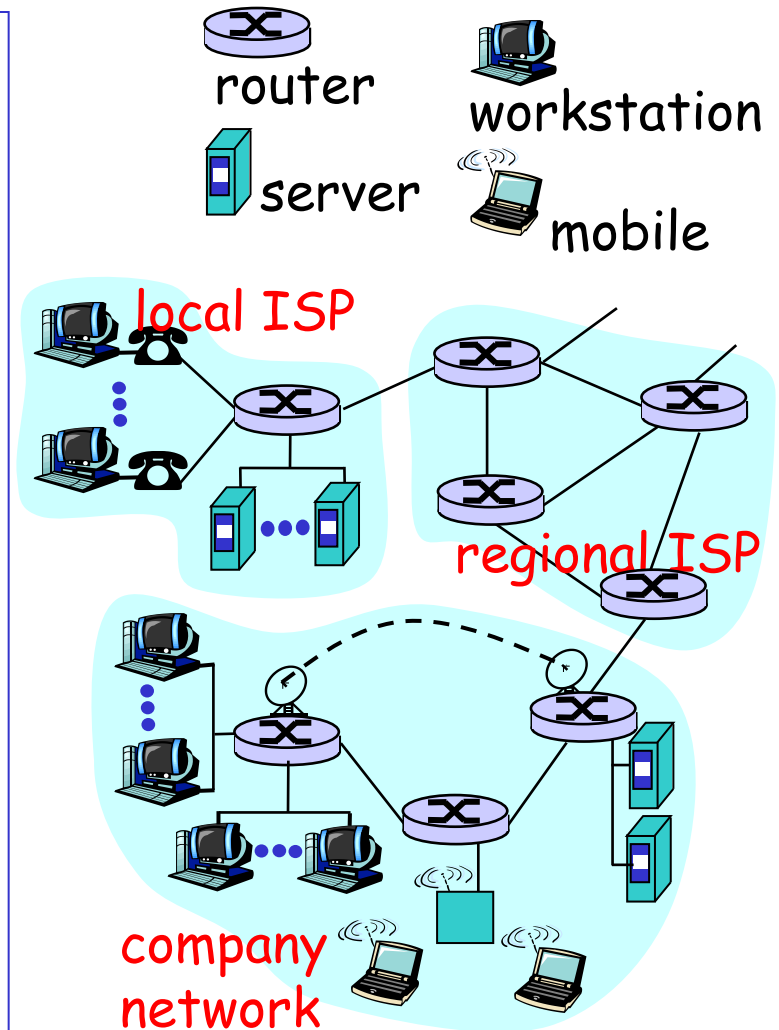
- ❑ "interconnected collection of autonomous computers connected by a communication technology"
- ❑ What is the Internet?
  - "network of networks"
  - "collection of networks interconnected by routers"
  - "a communication medium used by millions"
    - Email, chat, Web "surfing", streaming media
- ❑ Internet  $\neq$  Web

# The "nuts and bolts" view of the Internet



# The “nuts and bolts” view of the Internet

- ❑ millions of connected computing devices called *hosts or end-systems*
  - PCs, workstations, servers
  - PDAs, phones, toastersrunning *network apps*
- ❑ *communication links*
  - fiber, copper, radio, satellite
  - links have different capacities (*bandwidth*)
- ❑ *routers*: forward packets
- ❑ *packet*: piece of a message (basic unit of transfer)





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# Service/company landscape include

facebook®

Google™

skype™

NETFLIX



at&t

You Tube



Microsoft



Spotify™

ERICSSON



# Applications: Example classes

- ❑ File transfer
- ❑ Remote login (telnet, rlogin, ssh)
- ❑ World Wide Web (WWW)
- ❑ Instant Messaging (Internet chat, text messaging on cellular phones)
- ❑ Peer-to-Peer file sharing
- ❑ Internet Phone (Voice-Over-IP)
- ❑ Video-on-demand
- ❑ Distributed Games
- ❑ ... and many more to come/discuss ...

# Applications (2)

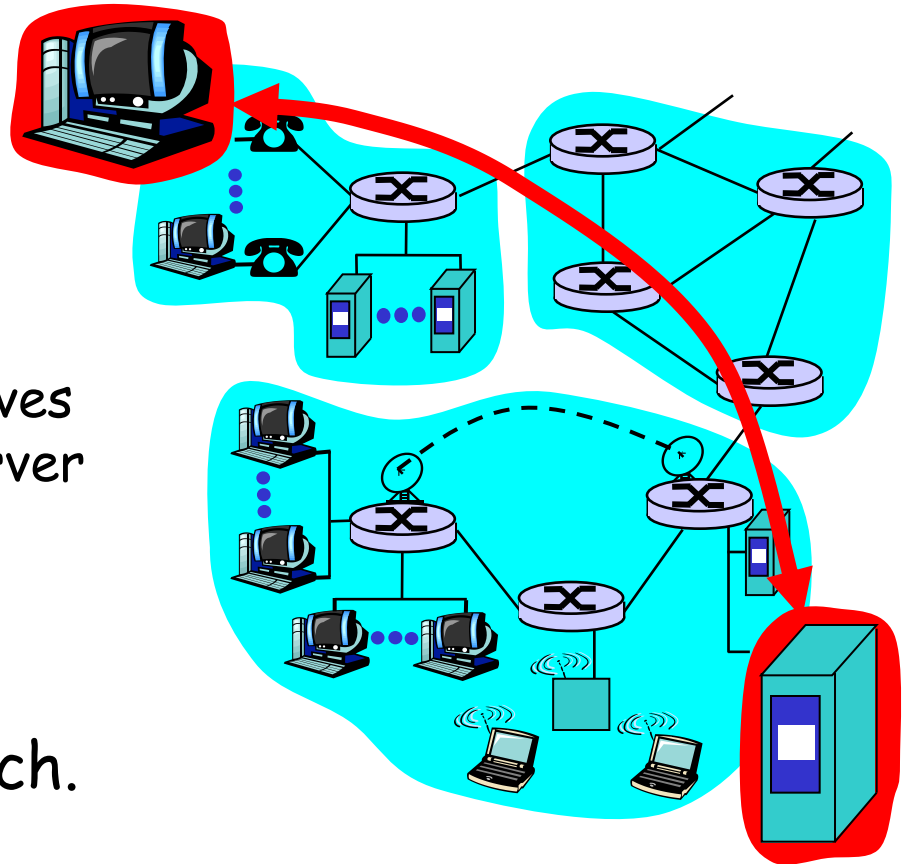
## ❑ end systems (hosts):

- run application programs
- e.g. Web, email, ftp
- at "edge of network"

## ❑ client/server model

- client host requests, receives service from always-on server
- e.g. Web browser/server; email client/server

## ❑ Client/server model has well-defined roles for each.

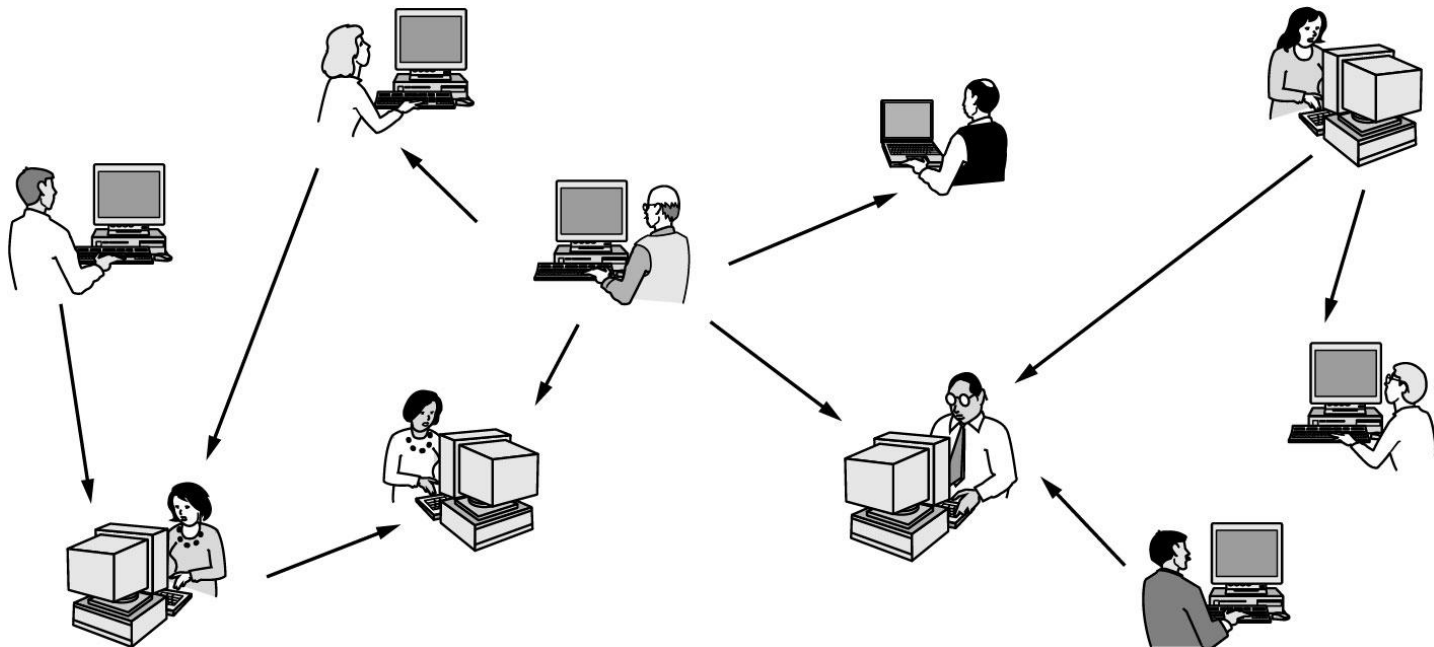


# Applications (3)

## ❑ peer-to-peer model:

- No fixed clients or servers
- Each host can act as both client and server at any time

## ❑ Examples: Napster, Gnutella, KaZaA, BitTorrent





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## ❑ Internet is an example of an internetwork.

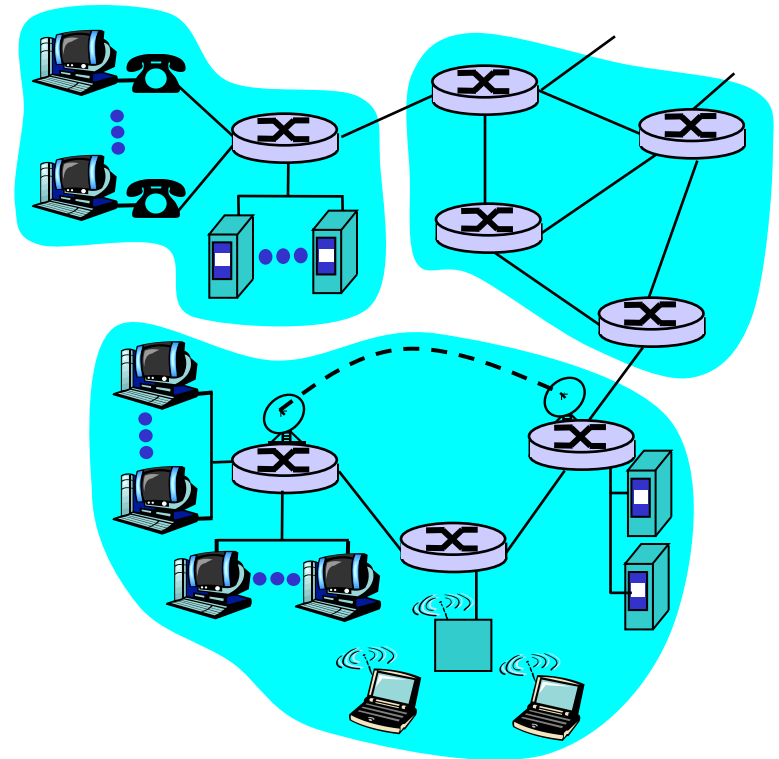
- Internetwork: interconnection of networks
- Subnetwork: a constituent of an internet
- Intermediate system: a device used to connect two networks allowing hosts of the networks to correspond with each other
  - Bridge
  - Router

# A Classification of Networks

- ❑ Wide Area Network (WAN)
- ❑ Metropolitan Area Network (MAN)
- ❑ Local Area Network (LAN)
  
- ❑ Wireless LAN (WLAN)
- ❑ Home Networks
- ❑ Personal Area Network (PAN)
- ❑ Body Area Network (BAN)
  
- ❑ ... and more (incl. sensor and ad-hoc) ...

# Wide Area Network (WAN)

- ❑ Spans a large geographic area, e.g., a country or a continent
- ❑ A WAN consists of several transmission lines and routers
- ❑ Internet is an example of a WAN

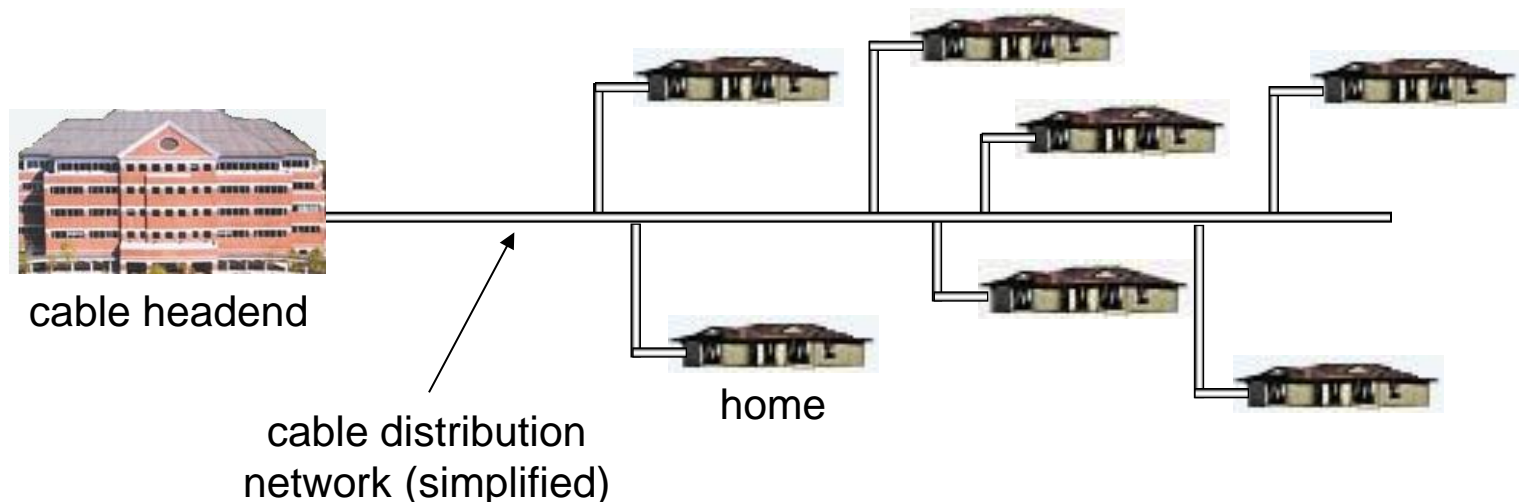


# Metropolitan Area Network (MAN)

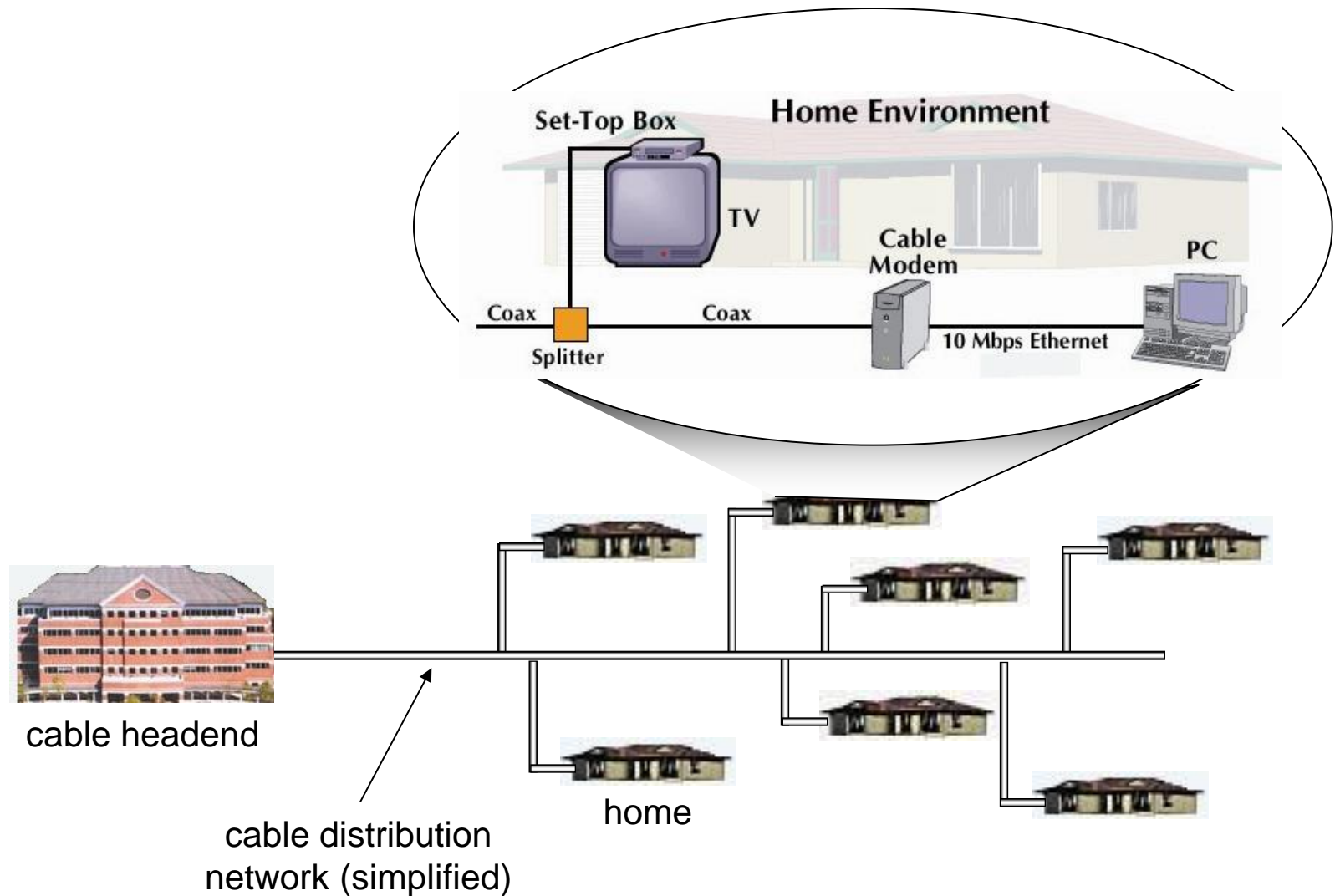
“City sized”: tens of kilometers

A Cable TV Network is an example of a MAN

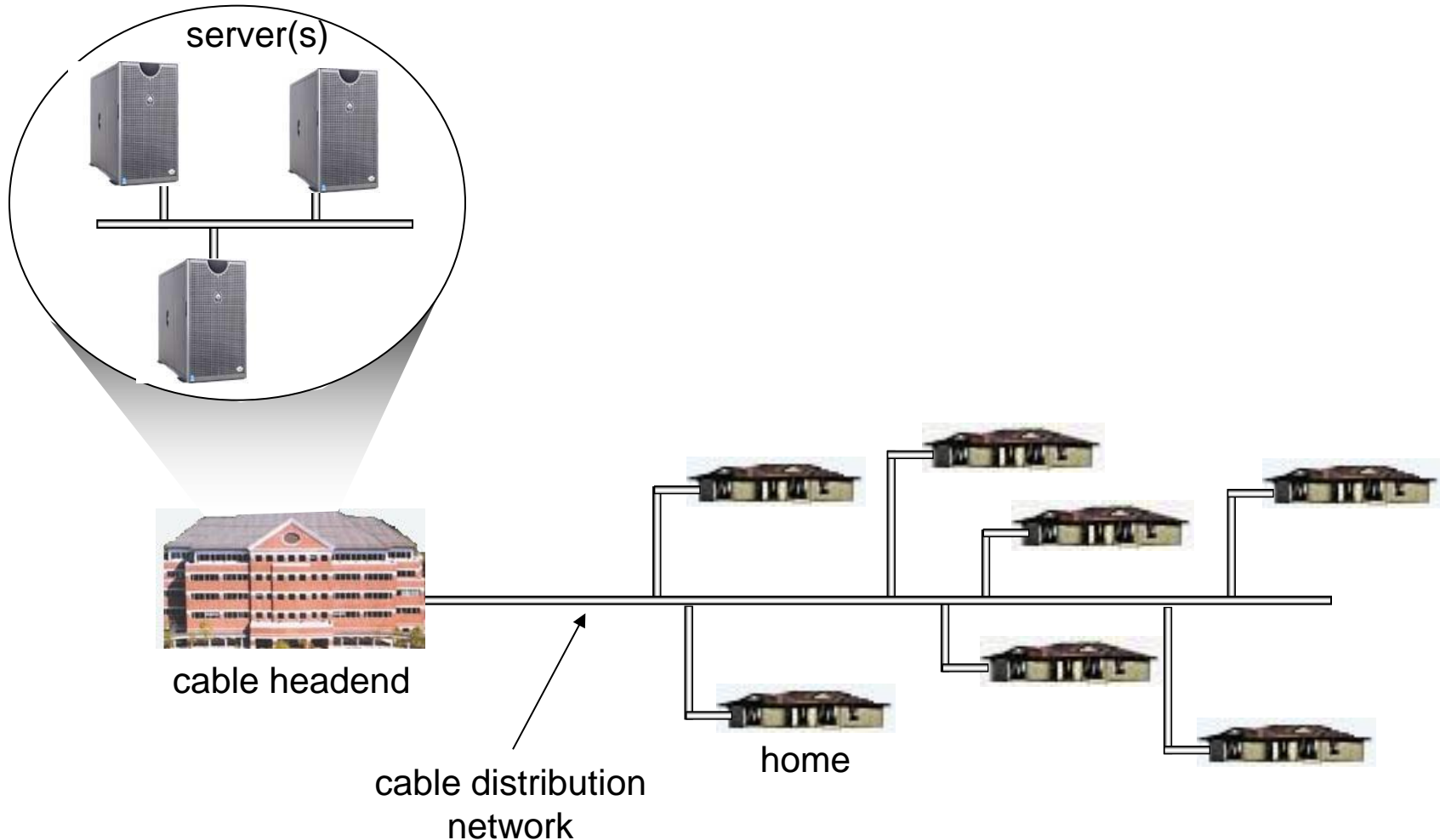
Typically 500 to 5,000 homes



# Cable Network Architecture: Overview

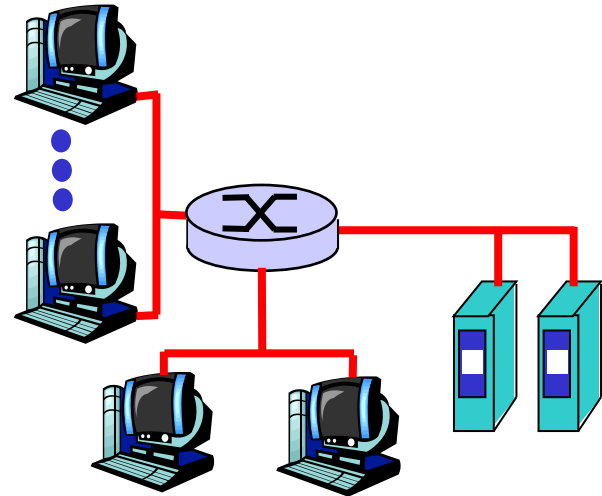


# Cable Network Architecture: Overview



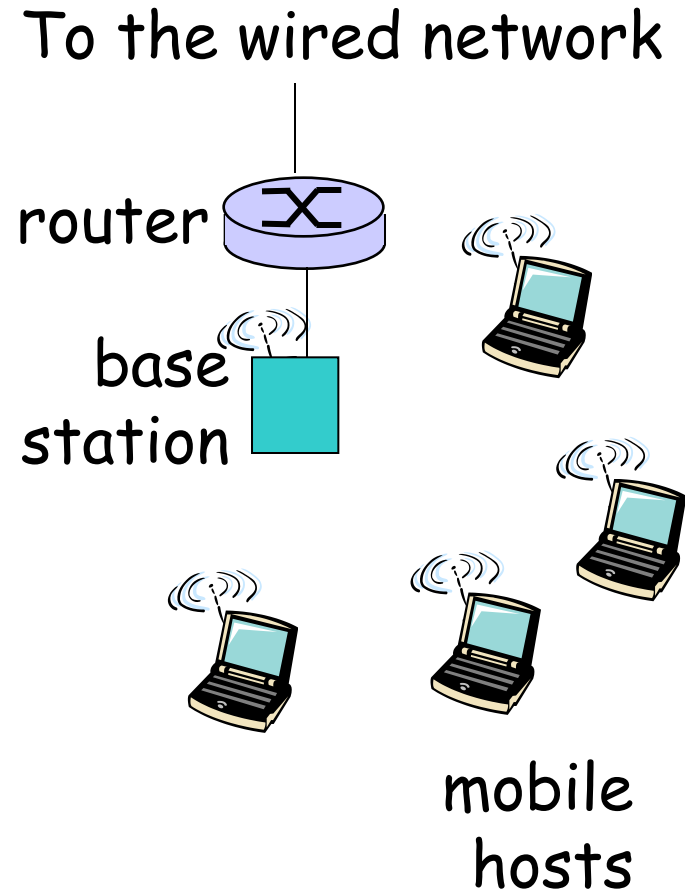
# Local Area Network (LAN)

- ❑ company/univ **local area network** (LAN) connects end system to edge router
- ❑ **Ethernet:**
  - shared or dedicated link connects end system and router (a few km)
  - 10 Mbps, 100Mbps, 1Gbps, 10Gbps ...
- ❑ **widespread deployment:** companies, univ, homeLANs
- ❑ LANs: chapter 5



# Wireless Networks (WLANs)

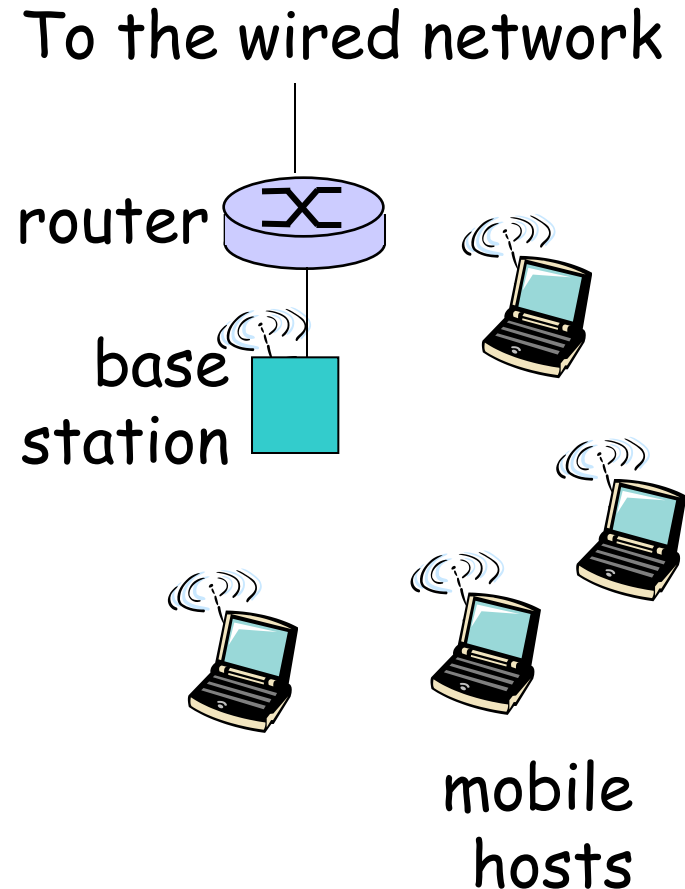
- ❑ shared *wireless* access network connects end system to router
  - via base station or “access point”
- ❑ **wireless LANs:**
  - 802.11b (WiFi)



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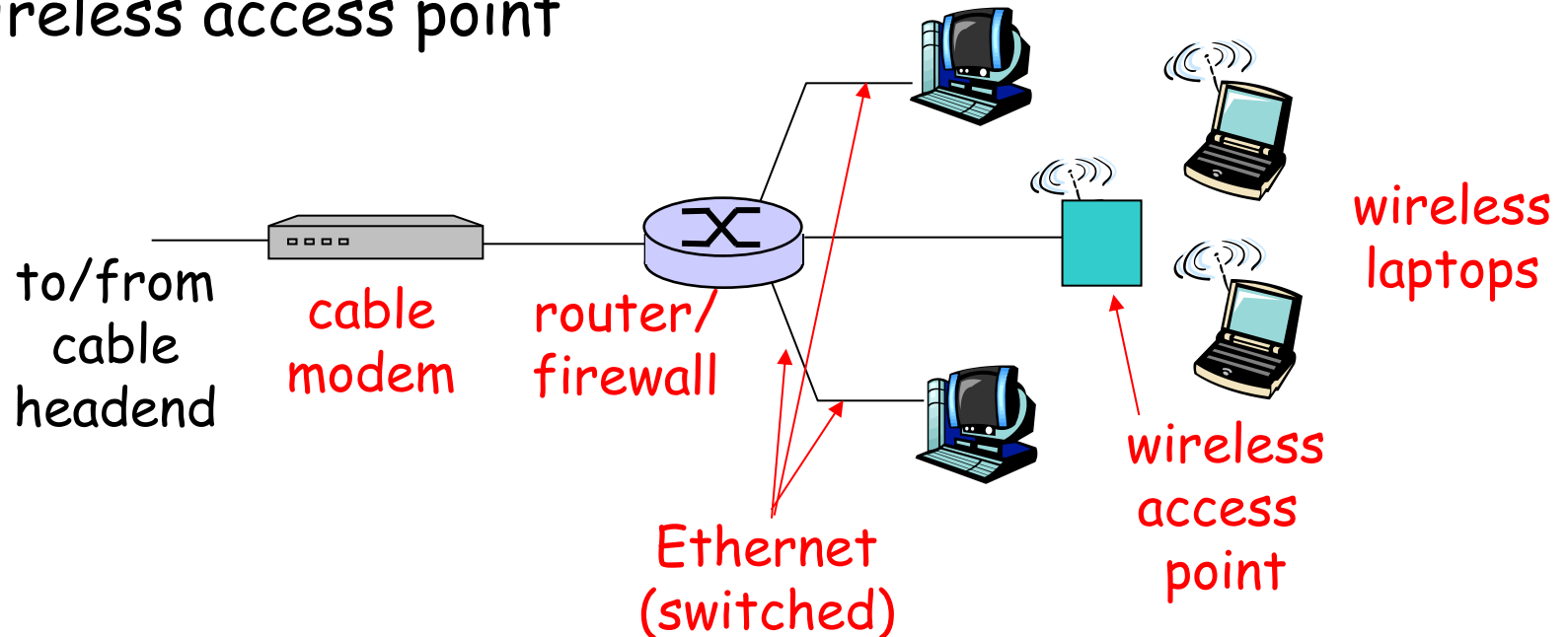
- ❑ **wider-area wireless access**
  - provided by telco operator
  - 3G, 4G, ...
  - WAP/GPRS in Europe
  - WiMax



# Home networks

## Typical home network components:

- ❑ ADSL or cable modem
- ❑ router/firewall/NAT
- ❑ Ethernet
- ❑ wireless access point





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# But first ...What's a protocol?

## Protocols:

- ❖ The rules used for communication
- ❖ Proper, accepted, and expected behavior

# But first ... What's a protocol?

## human protocols:

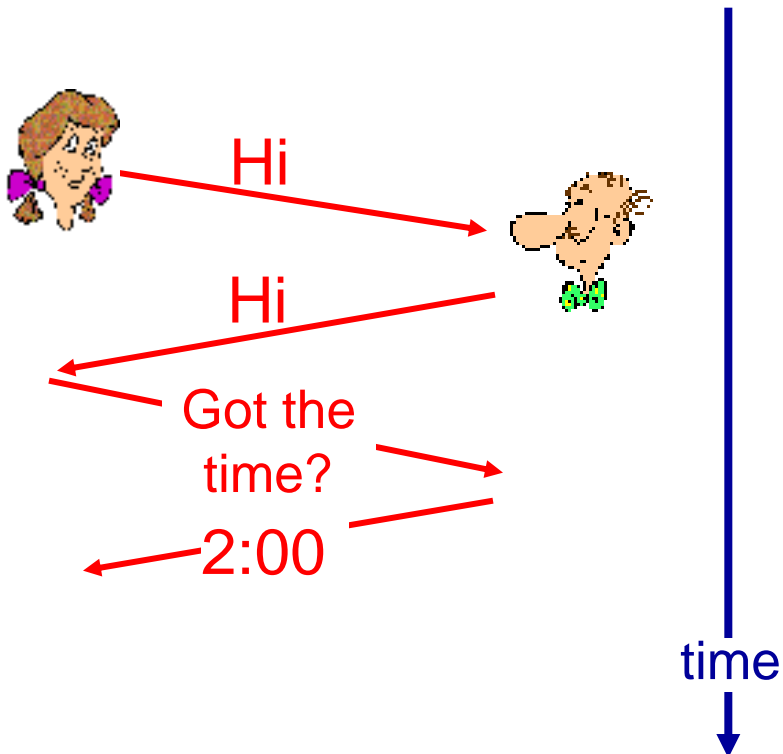
- ❖ "What's the time?"
- ❖ "I have a question"
- ❖ Introductions



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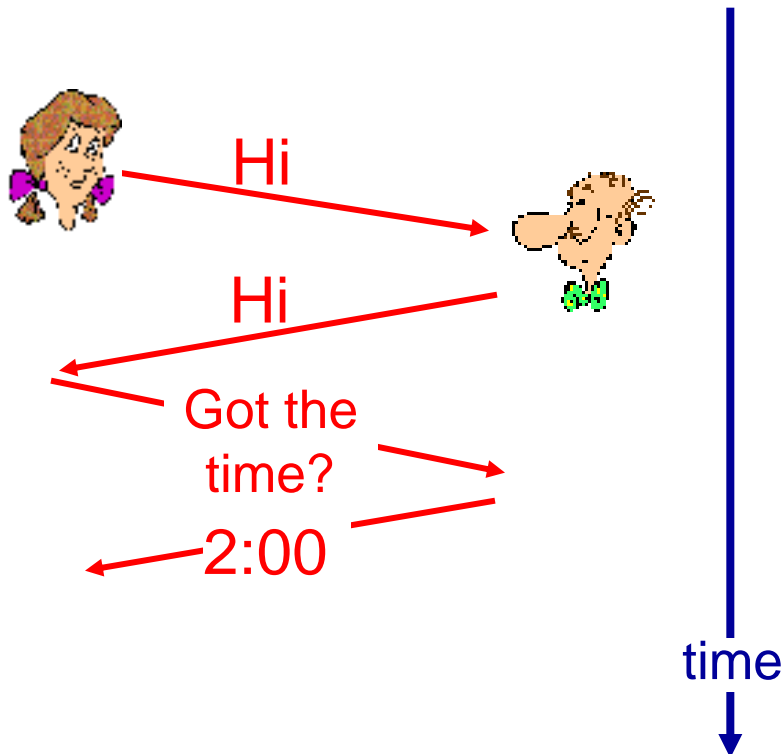
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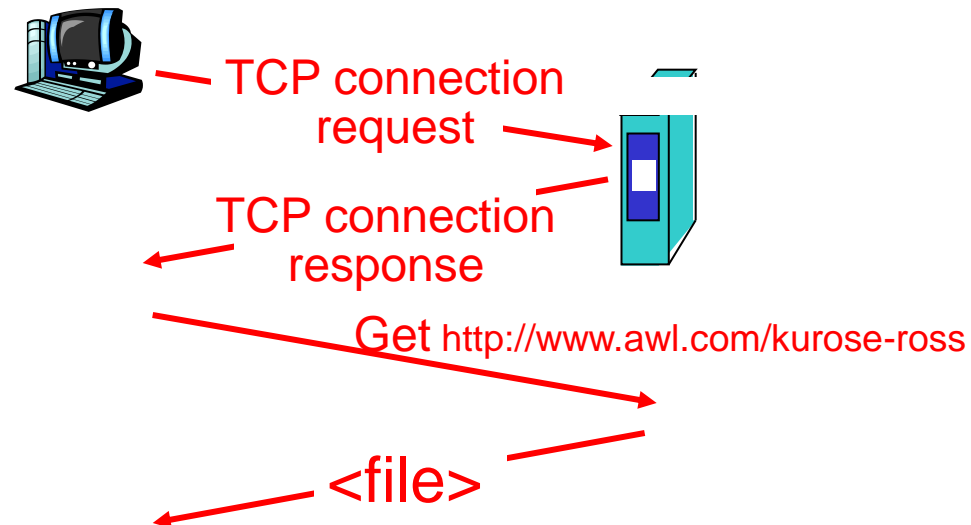
## human protocols:

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- ❖ "I have a question"
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## network protocols:

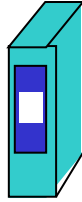
- ❖ Machines rather than humans
- ❖ All communication activity in Internet governed by protocols



# But first ...What's a protocol?



messages



[actions on  
events]

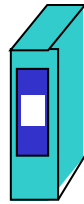
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Need:

# But first ...What's a protocol?



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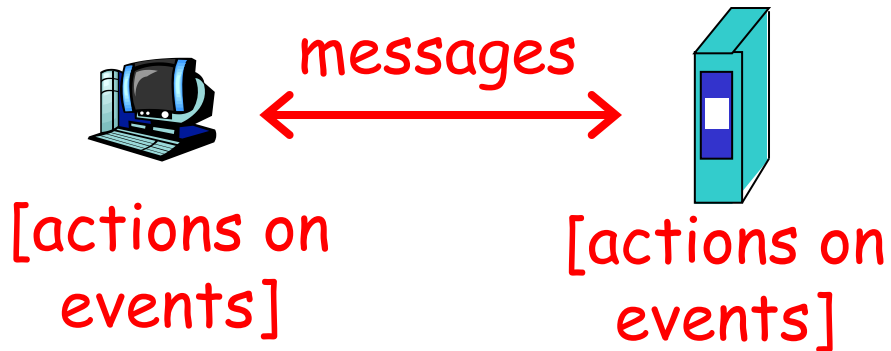
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... specific msgs sent

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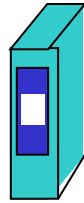
Need:

- ... specific msgs sent
- ... specific actions taken  
when msgs received,  
or other events

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messages



[actions on  
events]

[actions on  
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Need:

- ... specific msgs sent
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## Network protocols:

- ❖ Define the order and format of messages exchanged
- ❖ Defines the actions to take in response to events (e.g., message arrivals, transmissions, losses, and timeouts)



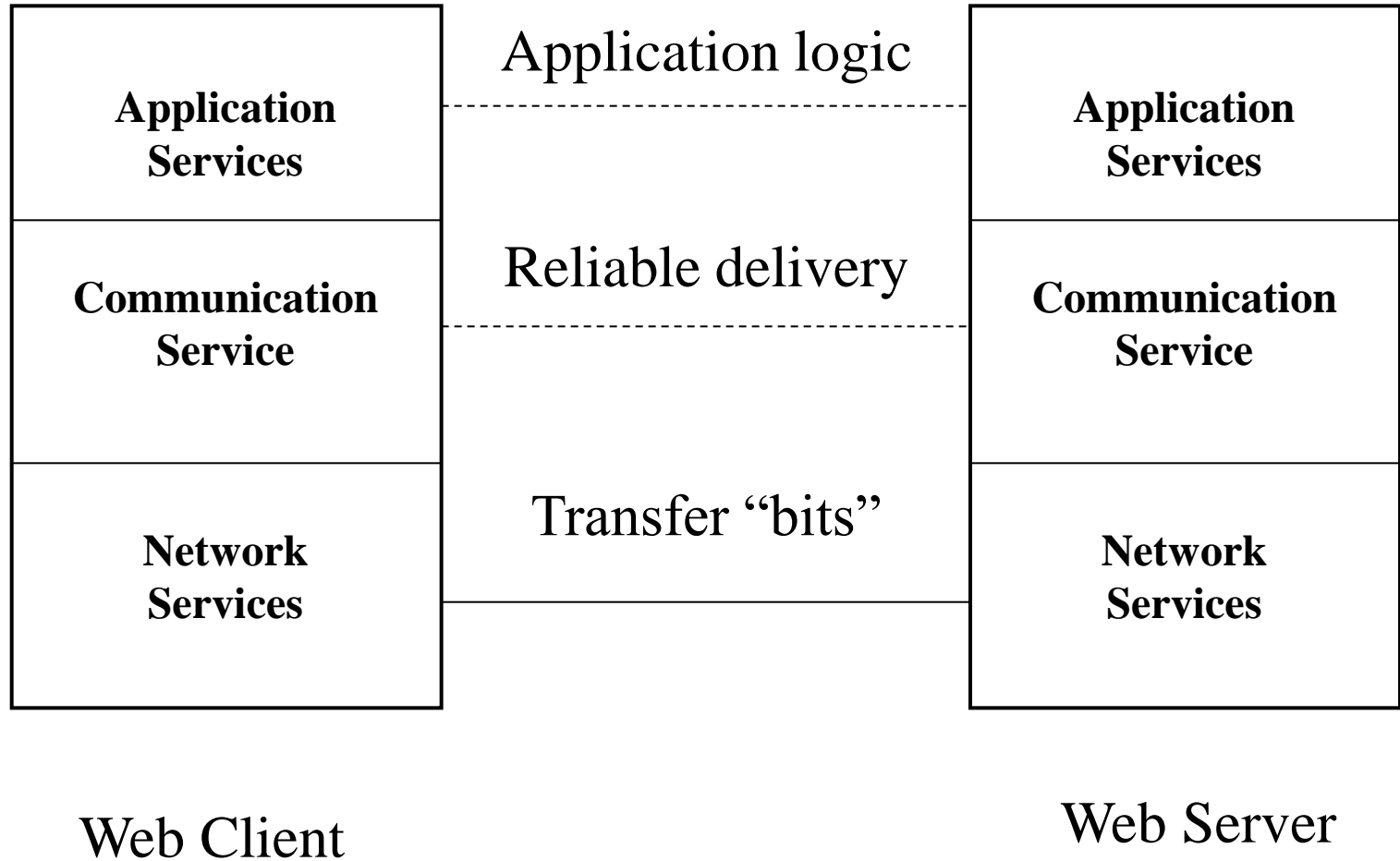
# Layered Architecture: Why?

- ❑ Networks are complex with many pieces
  - Hosts, routers, links, applications, protocols, hardware, software
- ❑ Can we organize it, somehow?

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- ❑ Networks are complex with many pieces
  - Hosts, routers, links, applications, protocols, hardware, software
- ❑ Can we organize it, somehow?
- ❑ Let's consider a Web page request ...

# Motivation Continued ...



# Motivation Continued ...

Dealing with complex systems:

- ❑ explicit structure allows identification, relationship of complex system's pieces
  - layered **reference model** for discussion
- ❑ modularization eases maintenance, updating of system
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in network technology doesn't affect rest of system
- ❑ layering considered harmful? (design vs implementation)

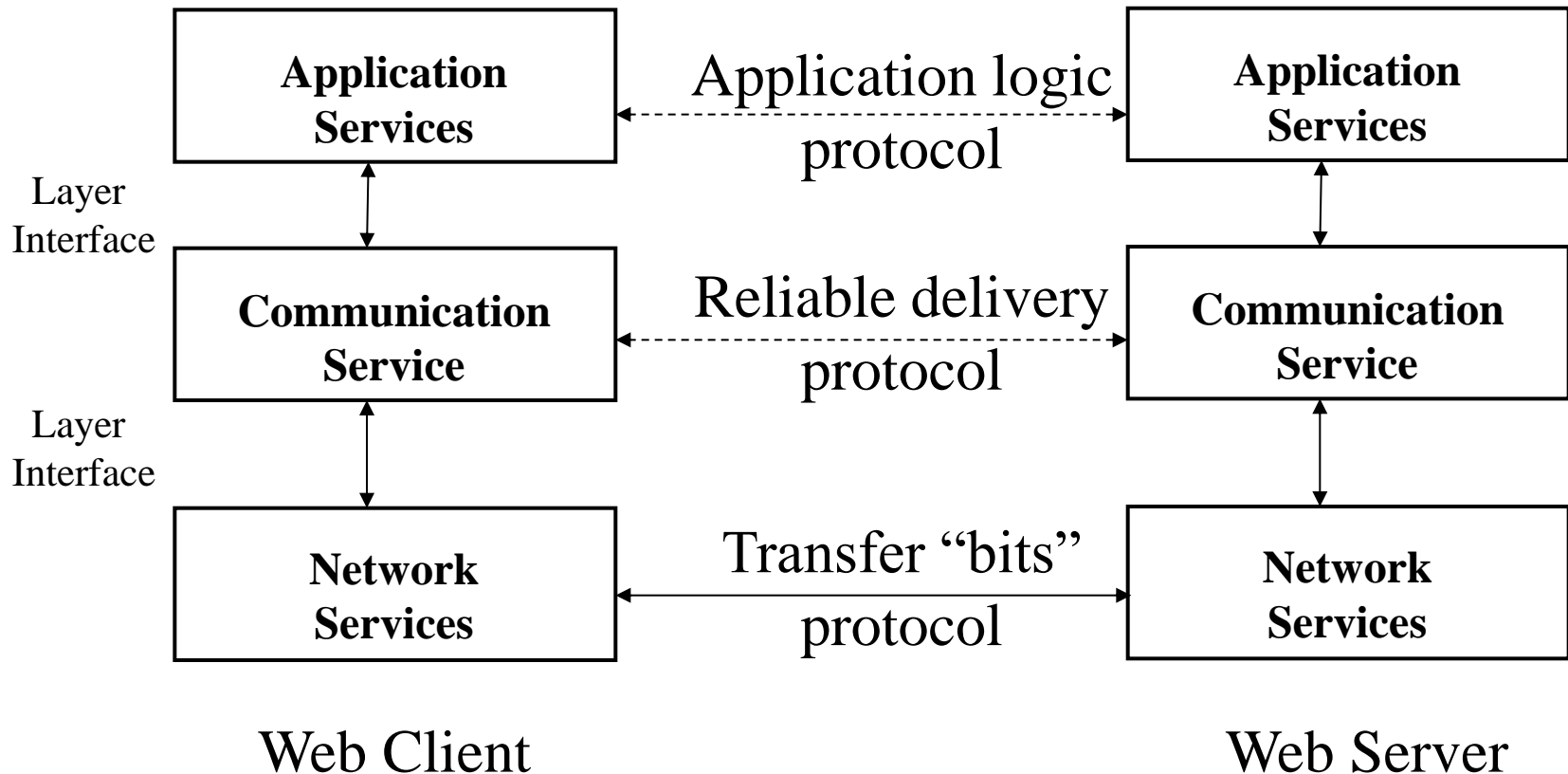


# Layers, Protocols, Interfaces

# Layers, Protocols, Interfaces

- ❑ Networks organized as a **stack of layers**
  - Offer services to the layer above it using a well-defined **interface**
    - programming language analogy: libraries hide details while providing a service)
  - Reduces design complexity
- ❑ **Protocols:** Logical “horizontal” conversations at any layer (between peers)
- ❑ **Data Transfer:** each layer passes data & control information over the interfaces (between neighboring layers)

# Layers, Protocols, Interfaces



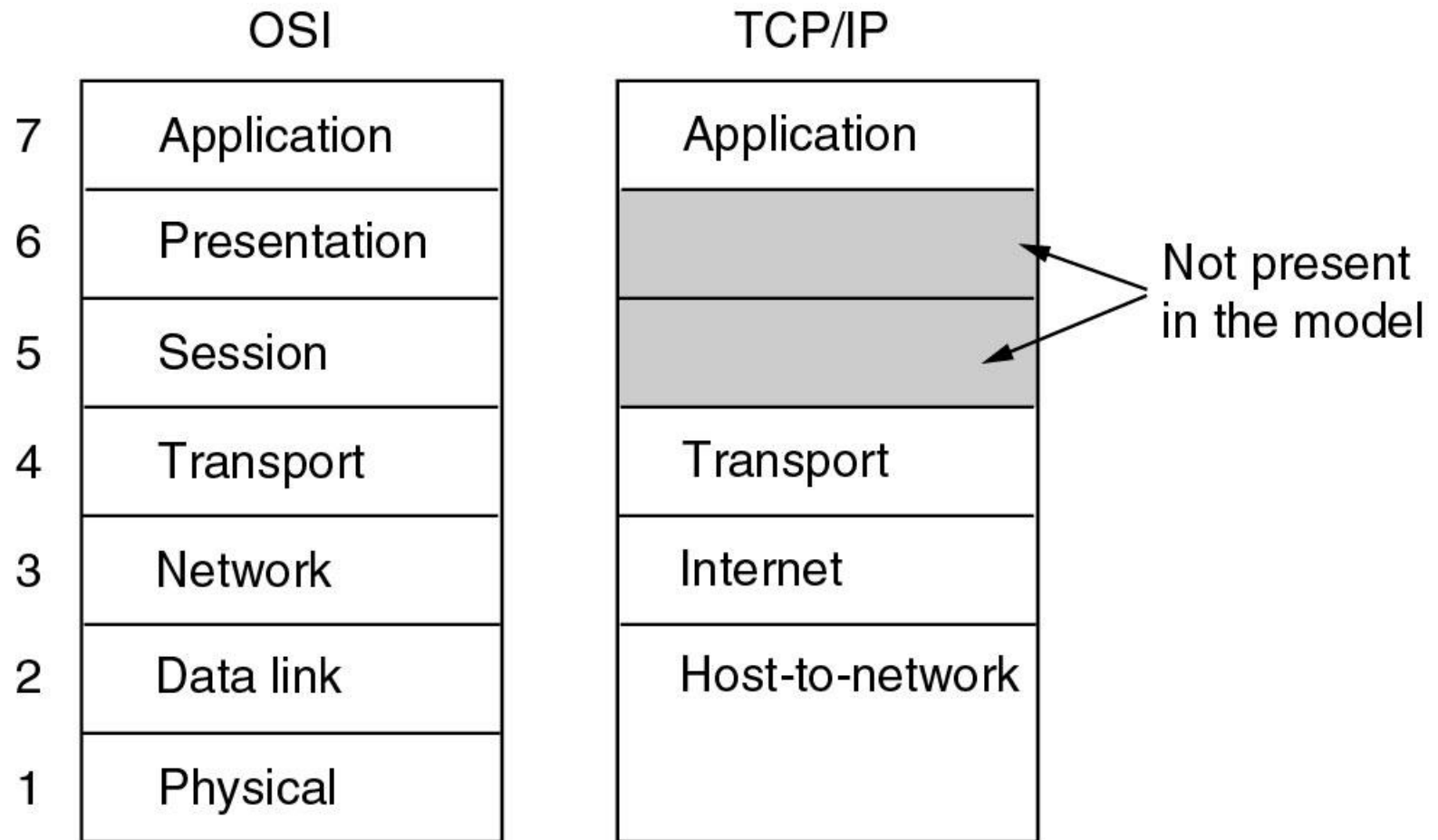
# Layered Architecture (cont'd)

- ❑ A set of layers & protocols is called a Network Architecture.
- ❑ These specifications enable hardware/software developers to build systems compliant with a particular architecture.
  - E.g., TCP/IP, OSI

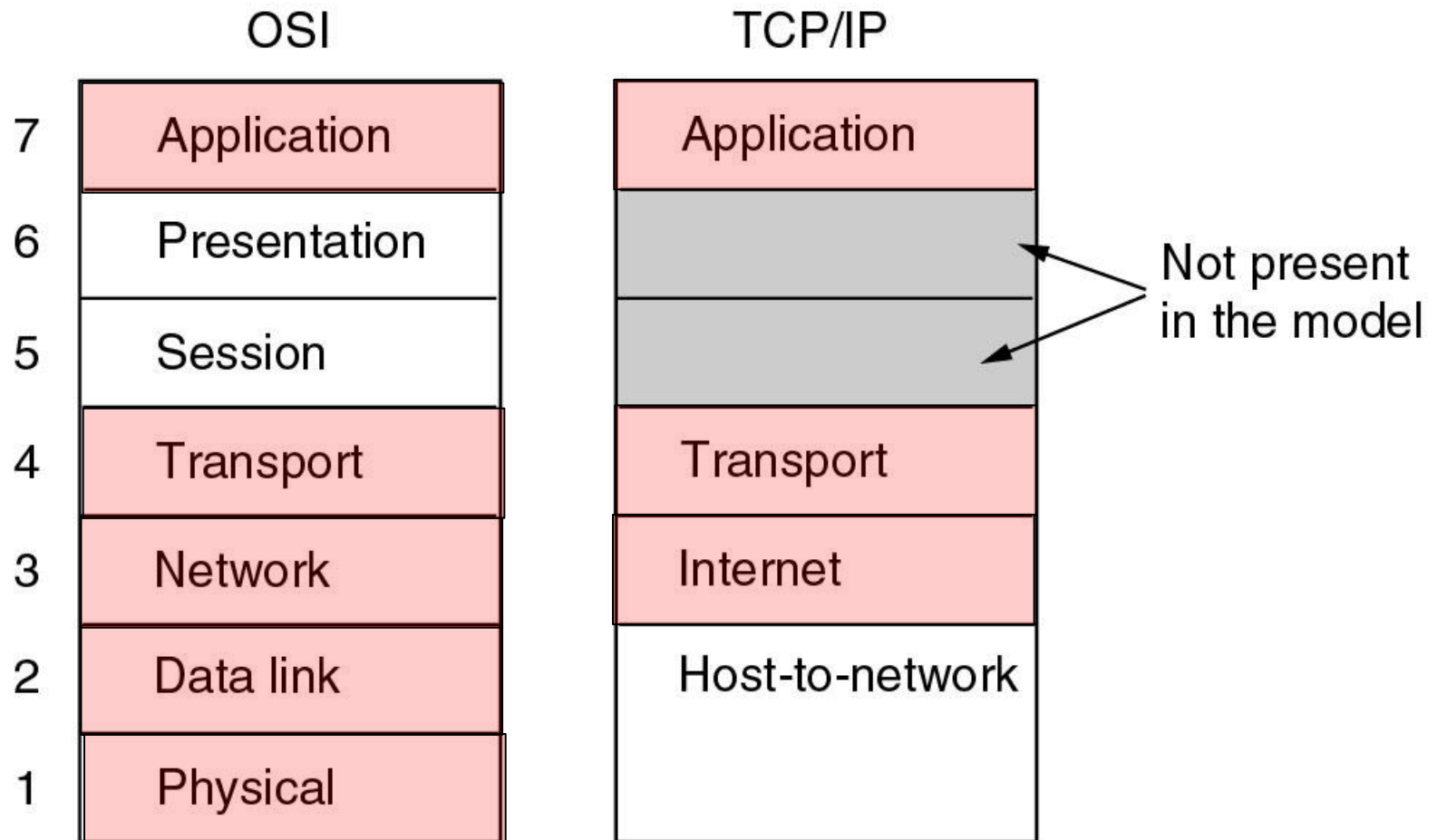
# Layering: Design Issues

- ❑ How many layers? What do they each do?
- ❑ How to identify senders/receivers?
  - Addressing
- ❑ Unreliable physical communication medium?
  - Error detection
  - Error control
  - Message reordering
- ❑ Sender can swamp the receiver?
  - Flow control
- ❑ Multiplexing/Demultiplexing

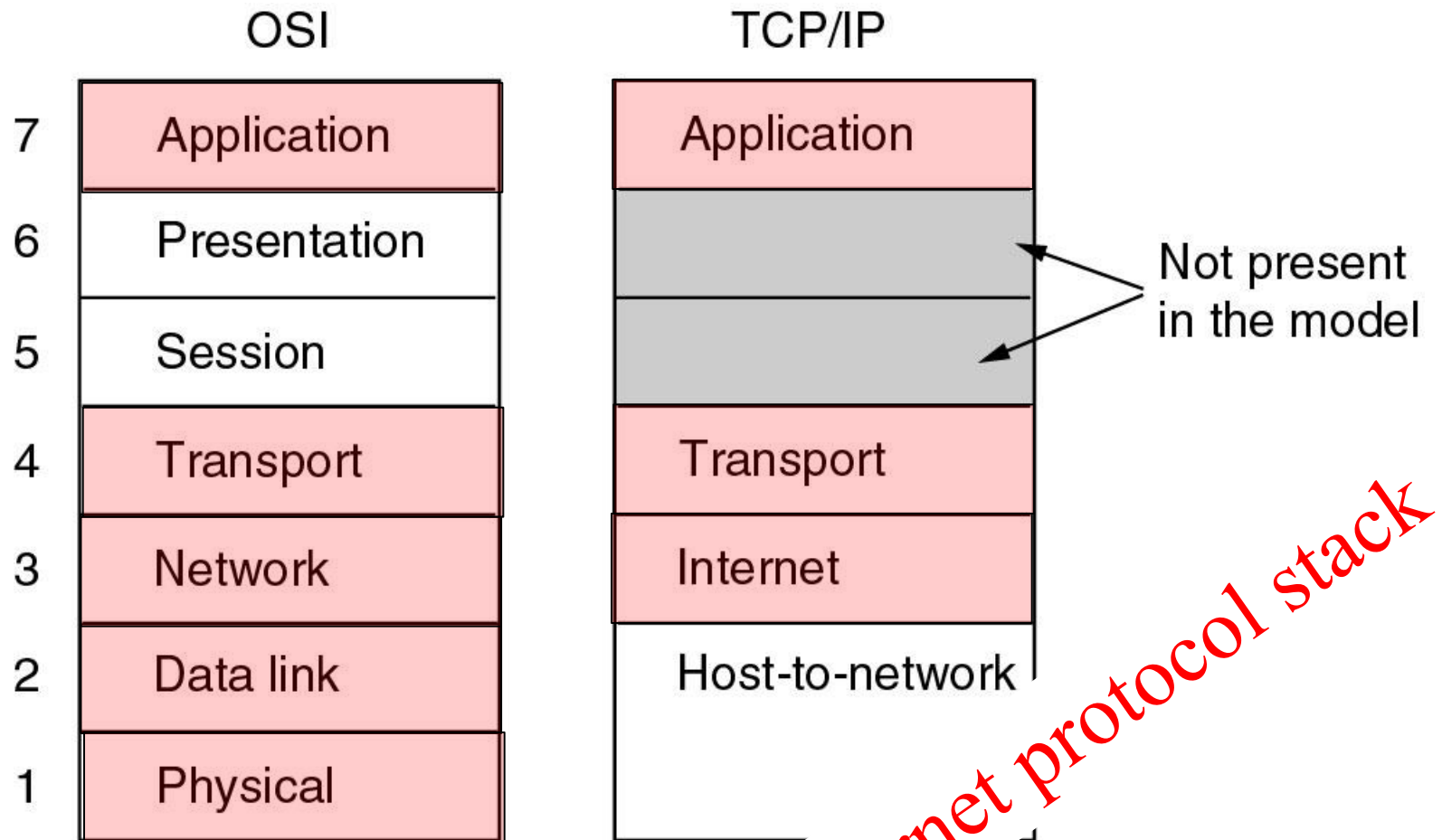
# Reference Models



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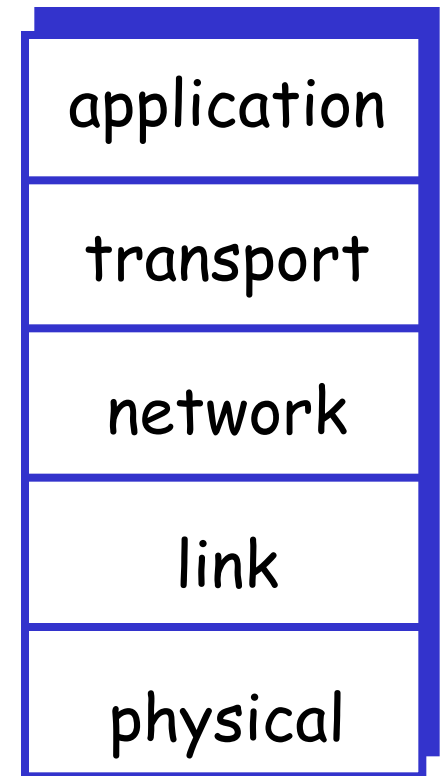


# Reference Models



# Internet protocol stack

- ❑ **application:** supporting network applications
  - FTP, SMTP, STTP
- ❑ **transport:** host-host data transfer
  - TCP, UDP
- ❑ **network:** routing of datagrams from source to destination
  - IP, routing protocols
- ❑ **link:** data transfer between neighboring network elements
  - PPP, Ethernet
- ❑ **physical:** bits “on the wire”



# The Application Layer

- ❑ Residence of network applications and their application control logic
- ❑ Examples include:
  - HTTP (Hyper-Text Transfer Protocol)
  - FTP (File Transfer Protocol)
  - Telnet
  - SMTP (Simple Mail Transfer Protocol)
  - DNS (Domain Name Service)

# The Transport Layer

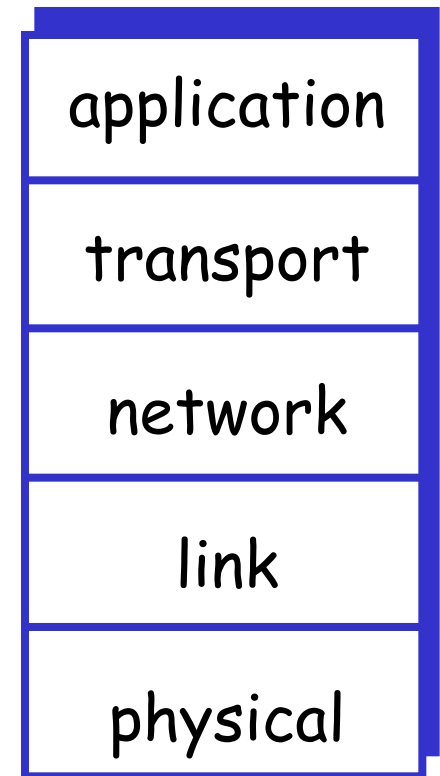
- ❑ Concerned with end-to-end data transfer between end systems (hosts)
- ❑ Transmission unit is called segment
- ❑ TCP/IP networks such as the Internet provides two types of services to applications
  - "connection-oriented" service - Transmission Control Protocol (TCP)
  - "connectionless" service - User Datagram Protocol (UDP)

# The Network Layer

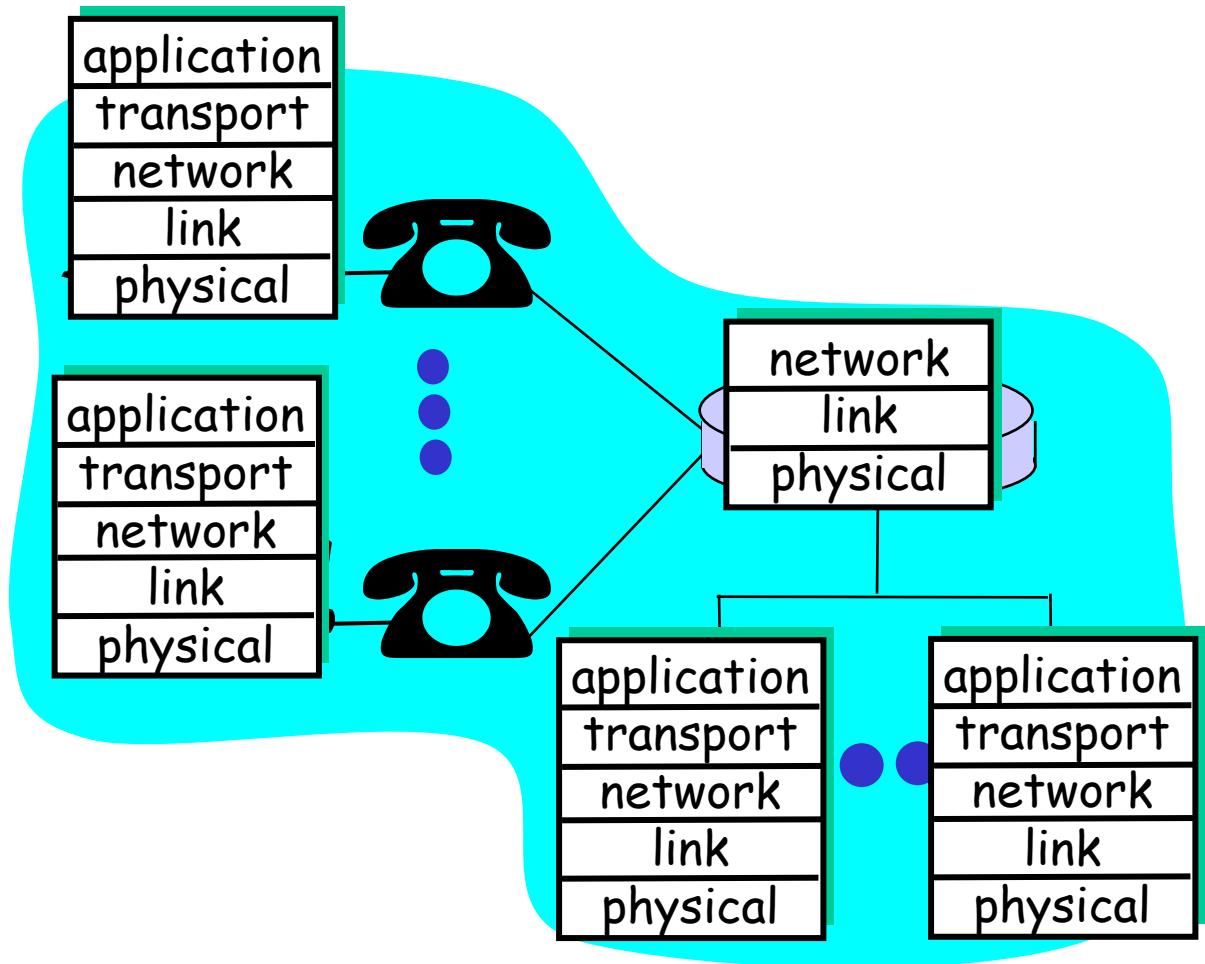
- ❑ End systems inject datagrams in the networks
- ❑ A transmission path is determined for each packet (routing)
- ❑ A “best effort” service
  - Datagrams might be lost
  - Datagrams might arrive out of order
- ❑ Analogy: Postal system

# Internet protocol stack

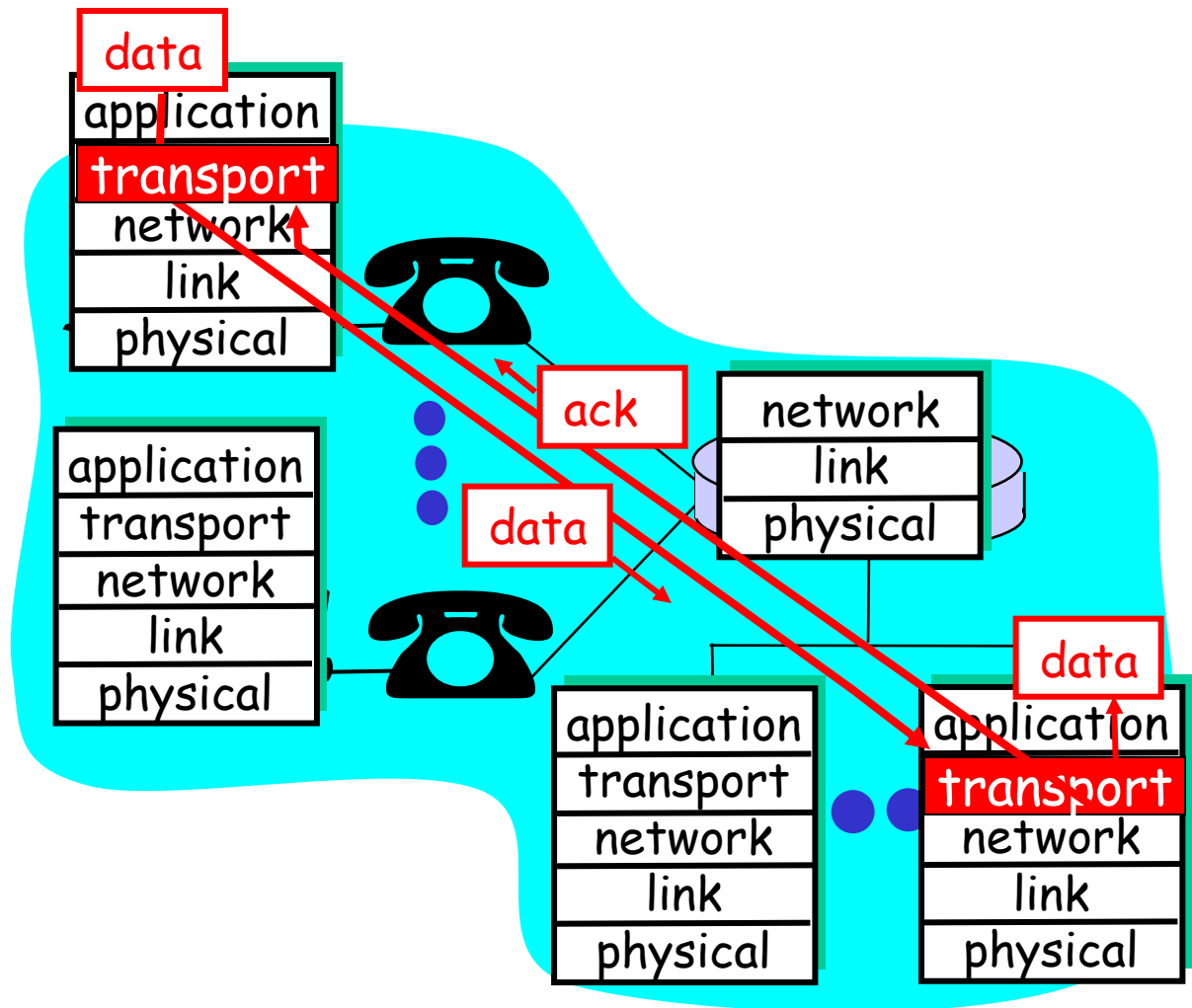
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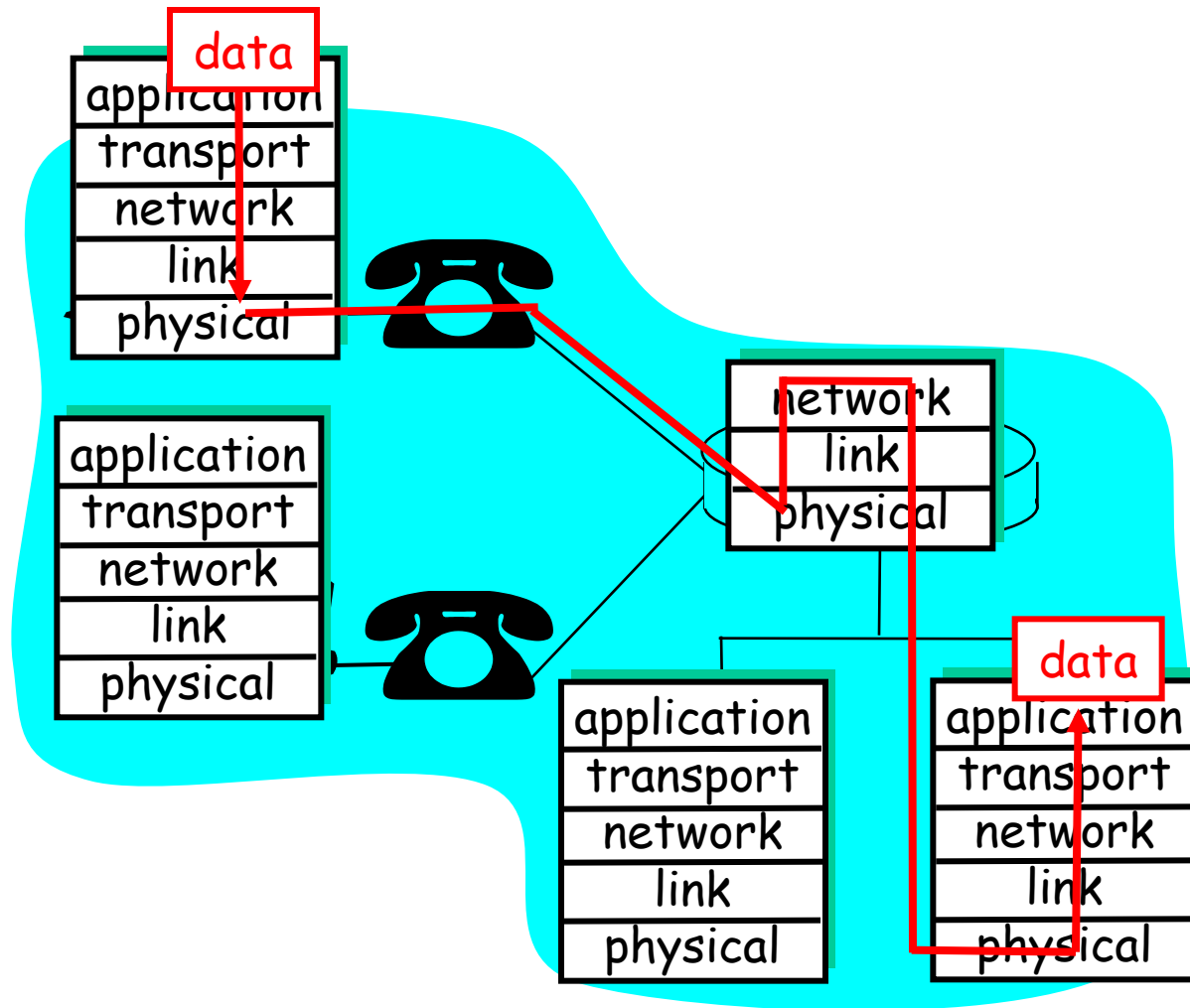
# Layering: logical communication



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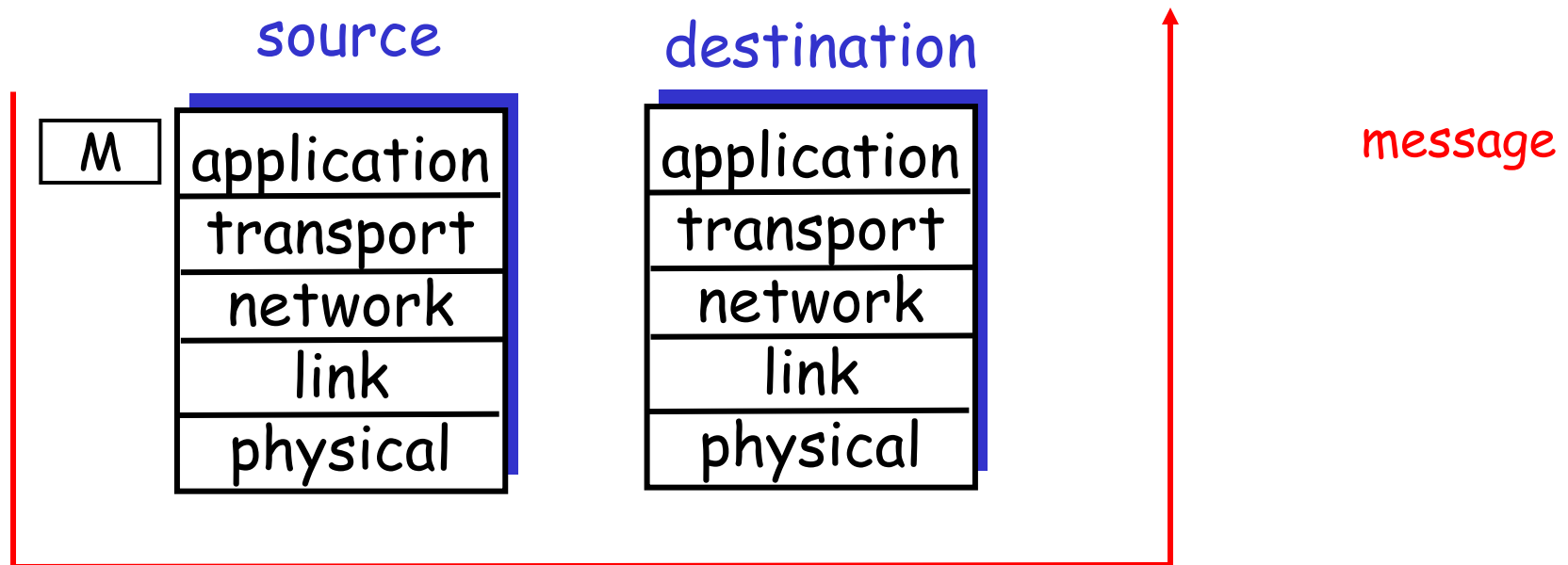
# Layering: *physical* communication



# Encapsulation: Layering and data

Each layer takes data from above

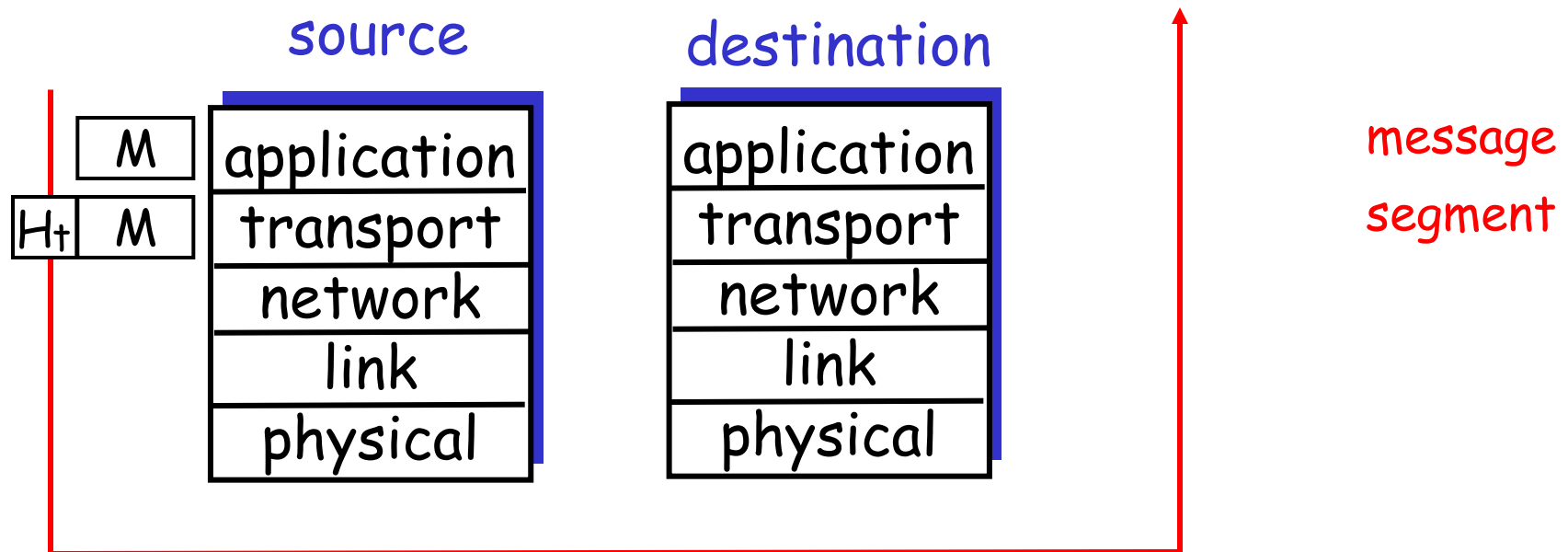
- ❑ adds header information to create new data unit
- ❑ passes new data unit to layer below



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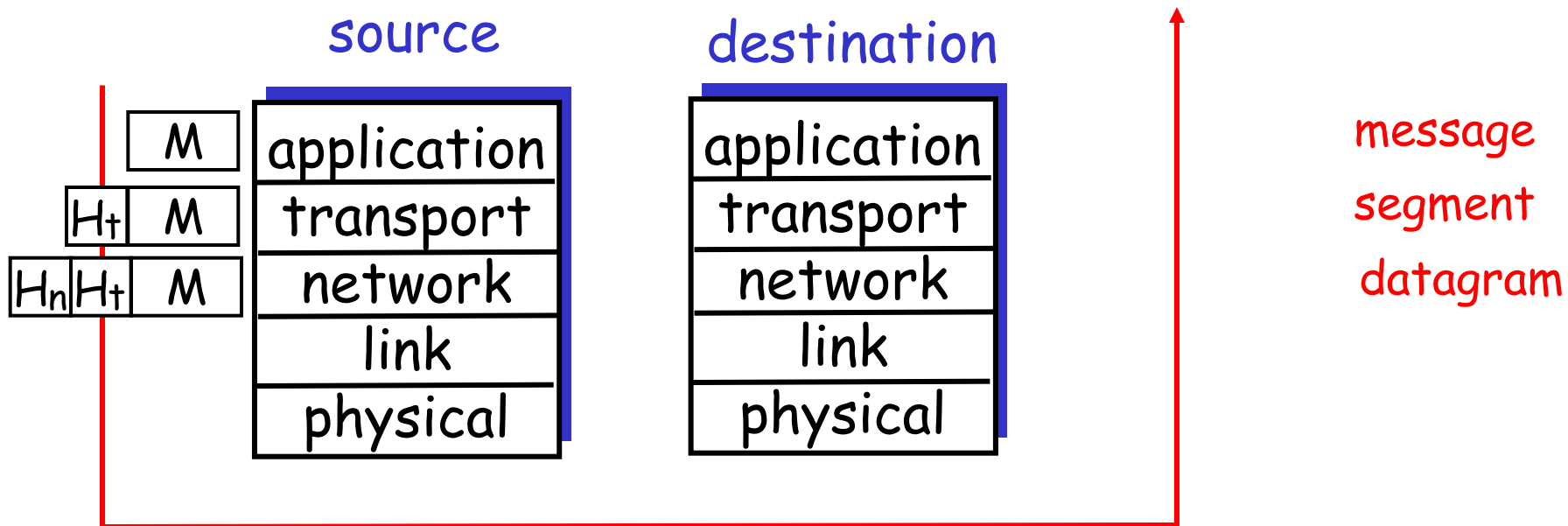
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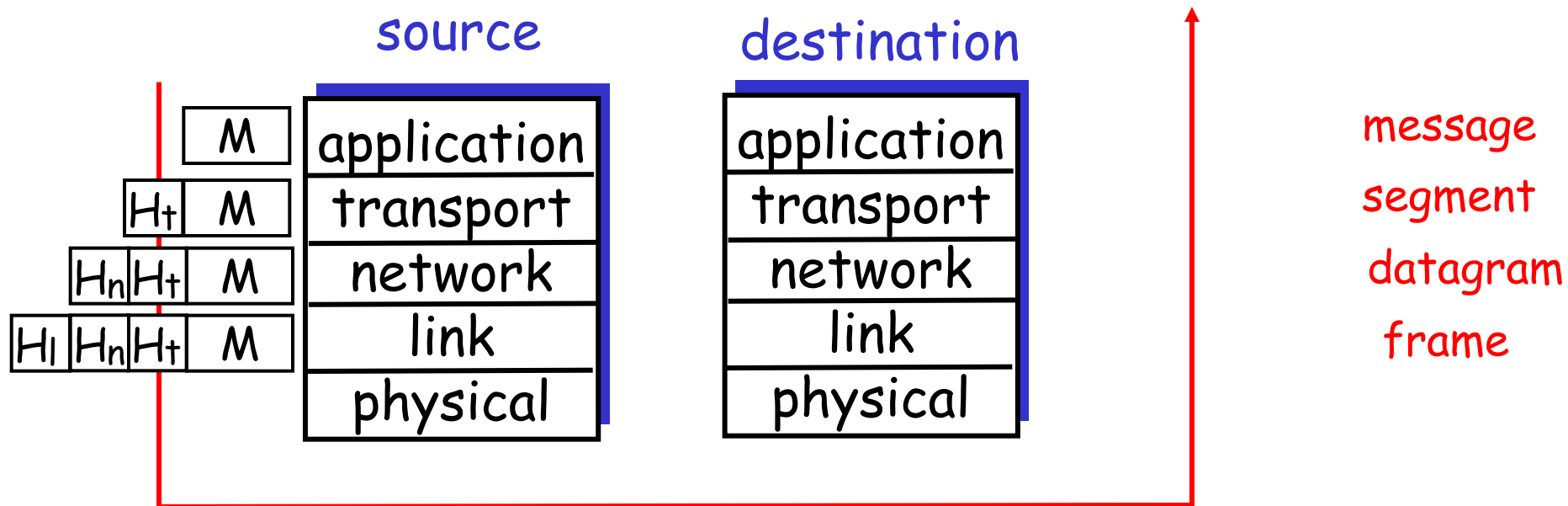
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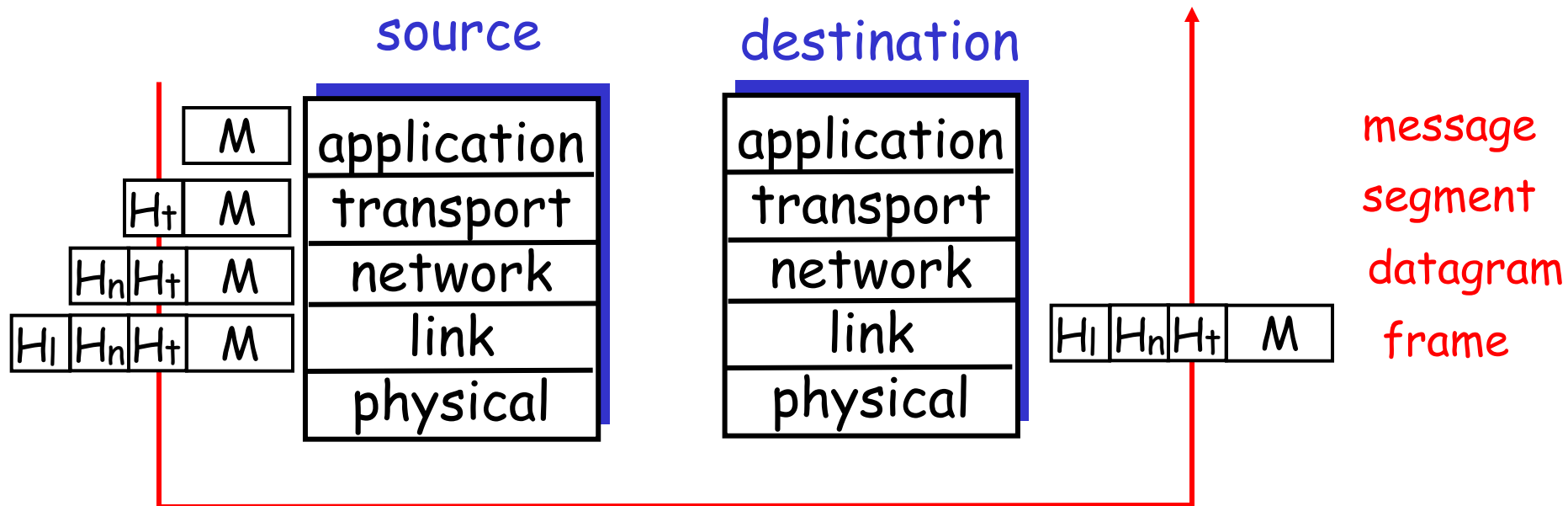
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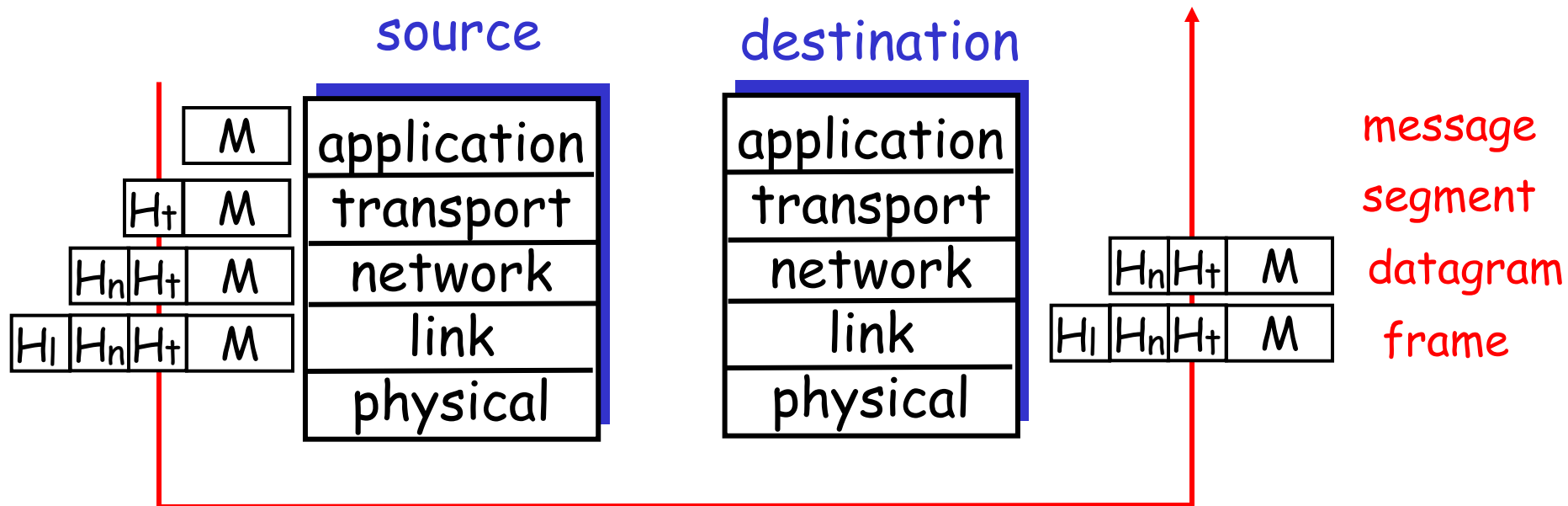
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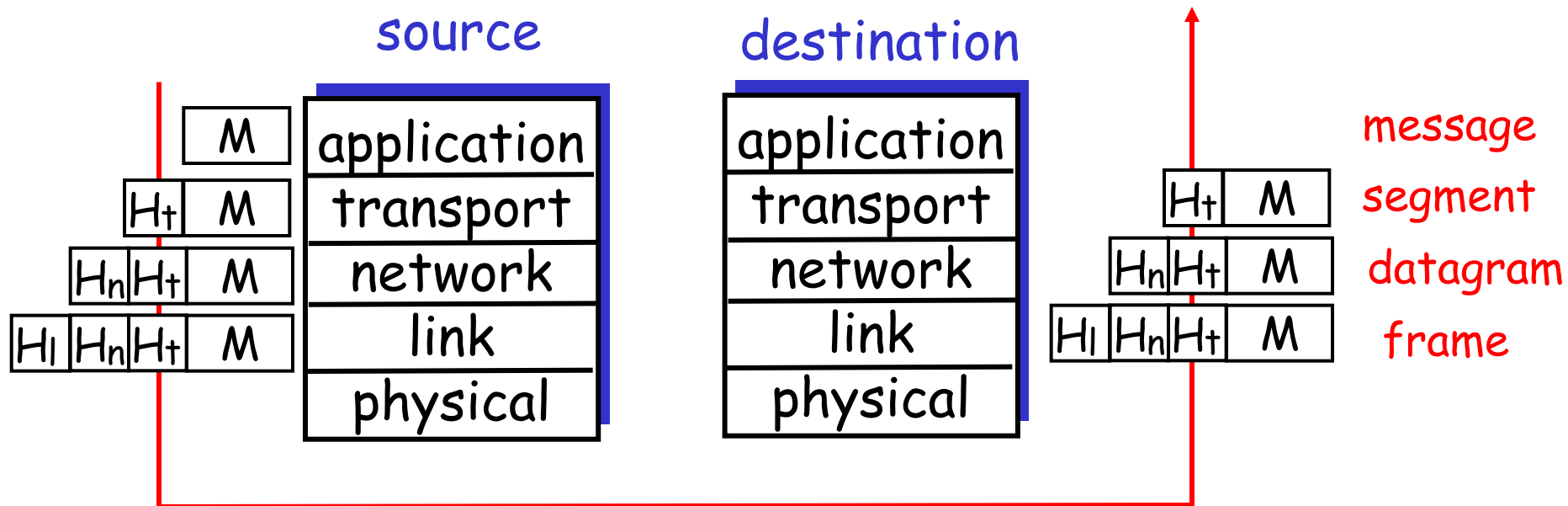
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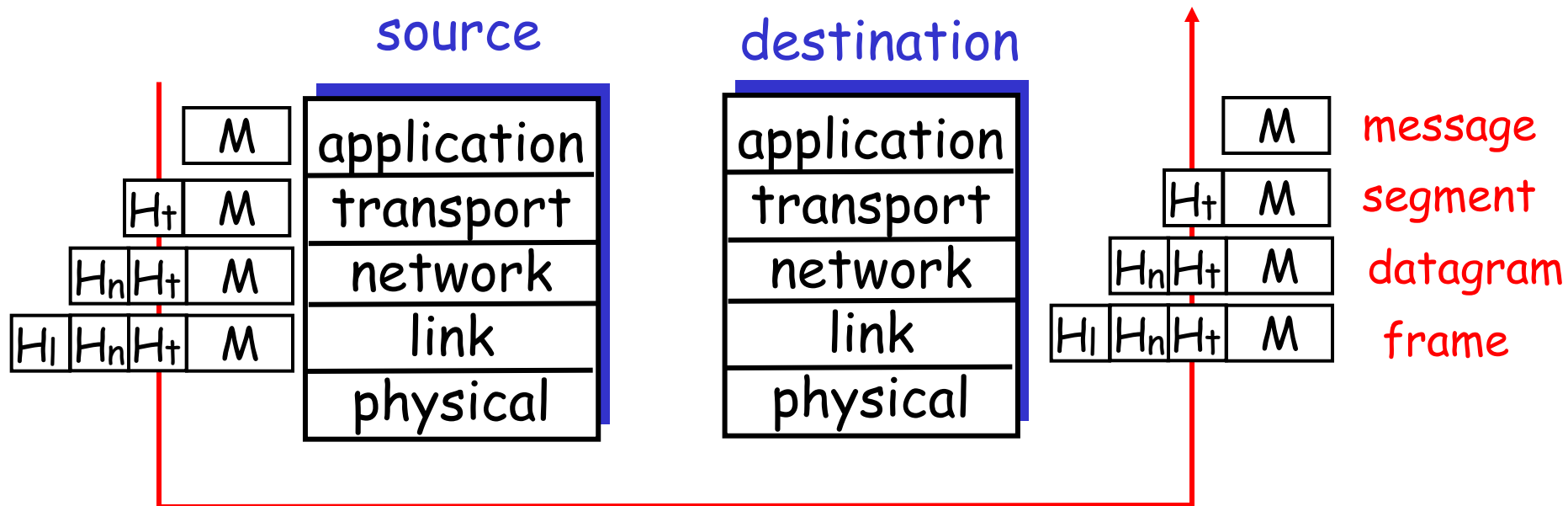
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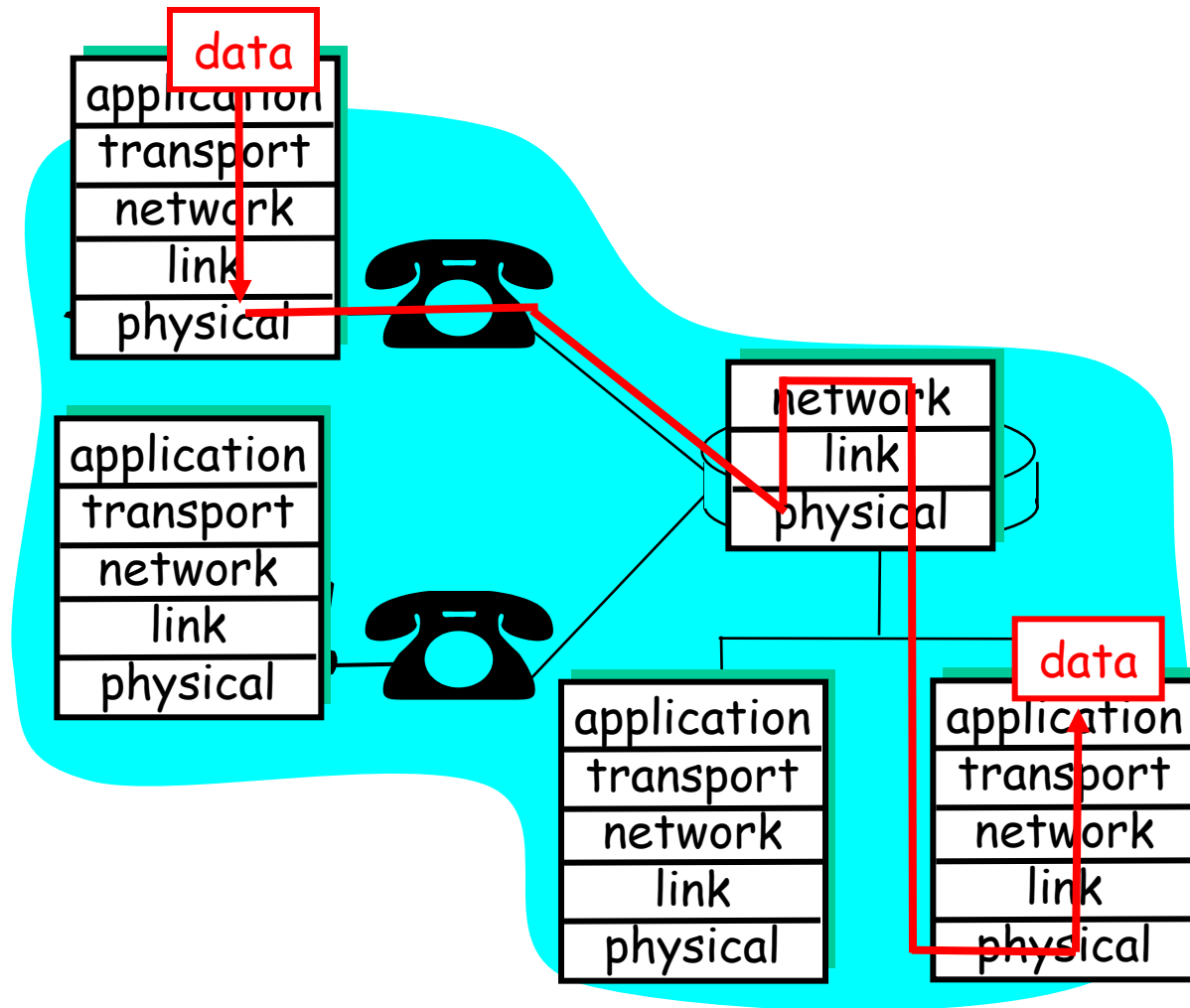
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# Layering: *physical* communication

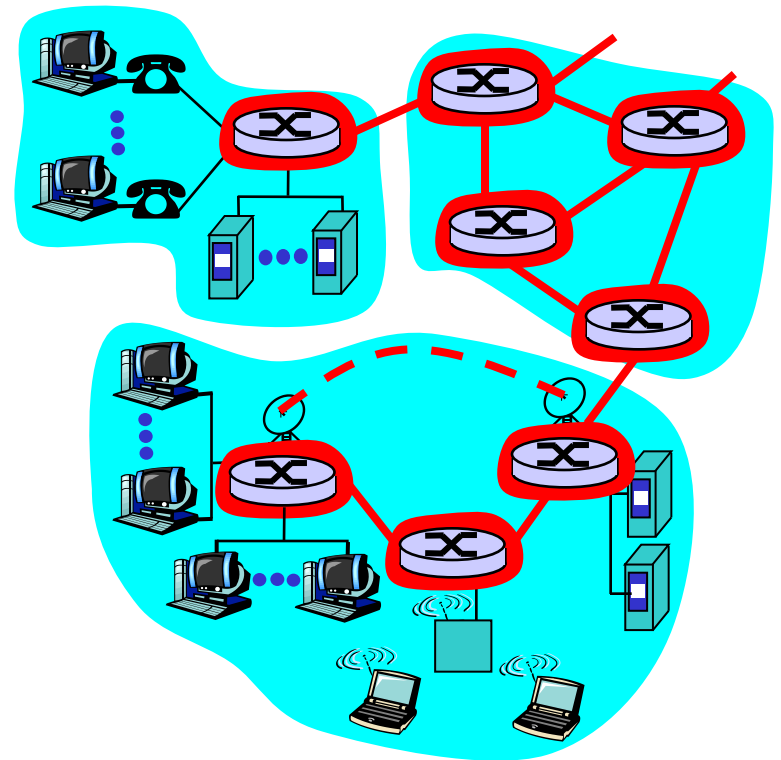




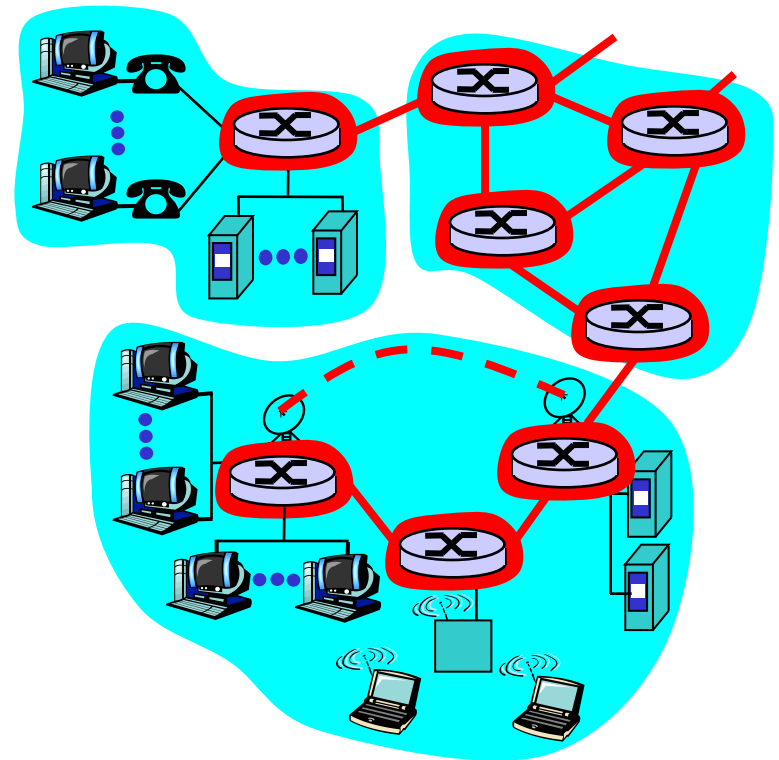
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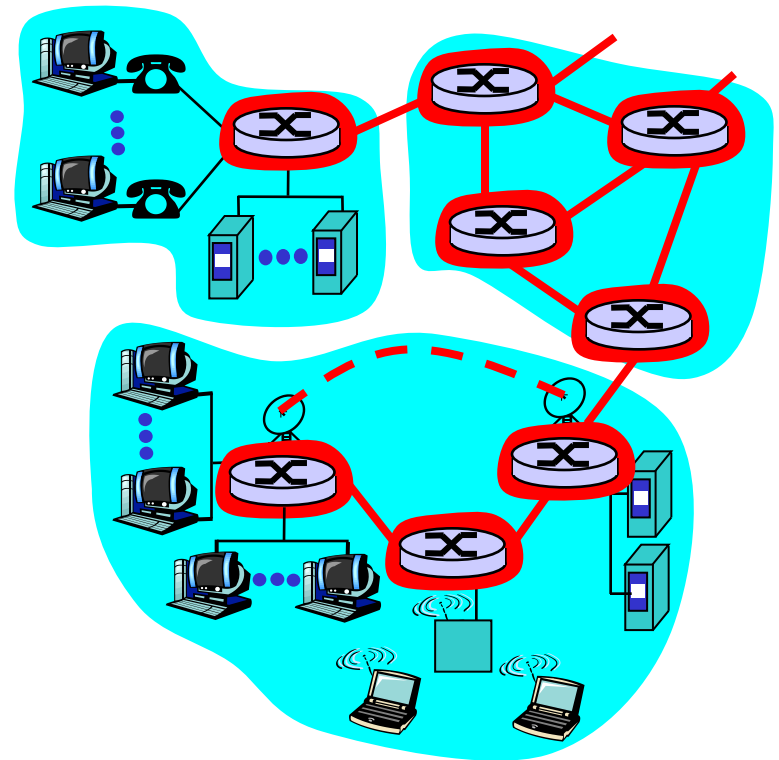
- mesh of interconnected routers



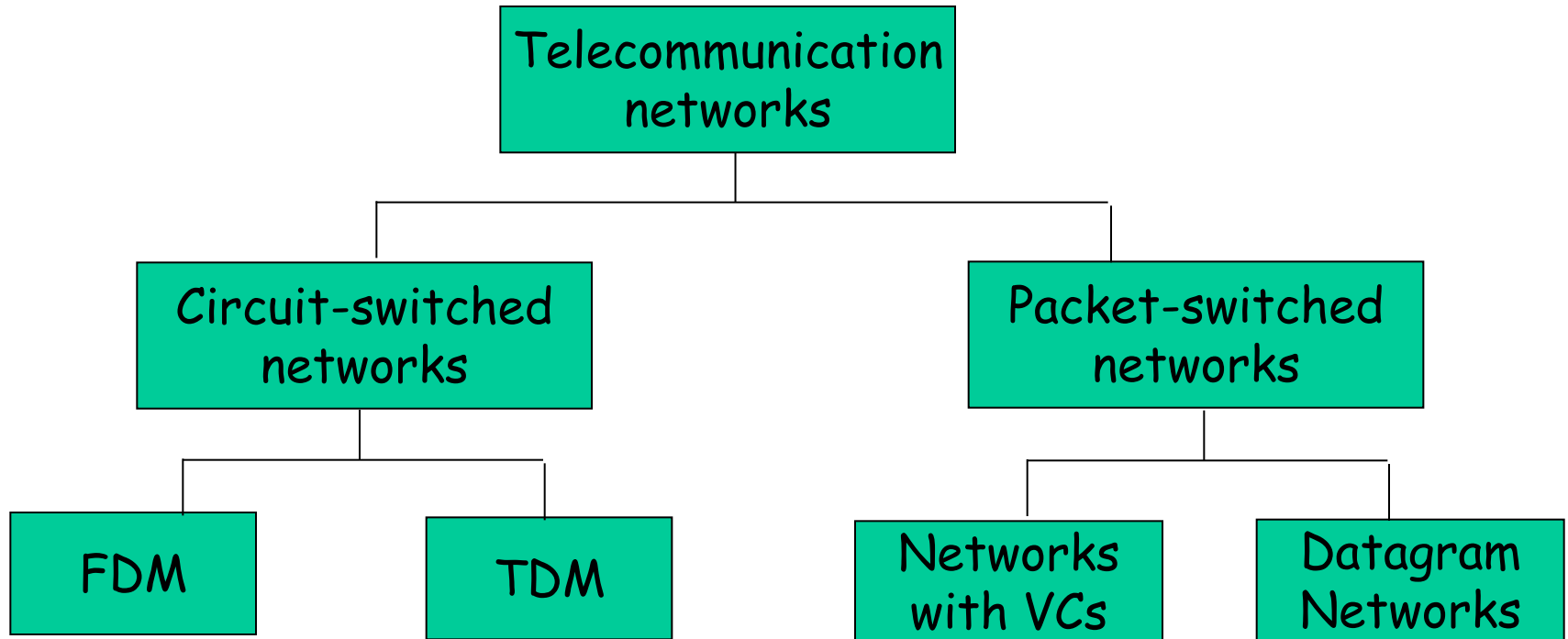
- ❑ mesh of interconnected routers
- ❑ the fundamental question: how is data transferred through net?



- ❑ mesh of interconnected routers
- ❑ the fundamental question: how is data transferred through net?
  - circuit-switching: dedicated circuit per call: telephone net
  - packet-switching: data sent thru net in discrete "chunks"



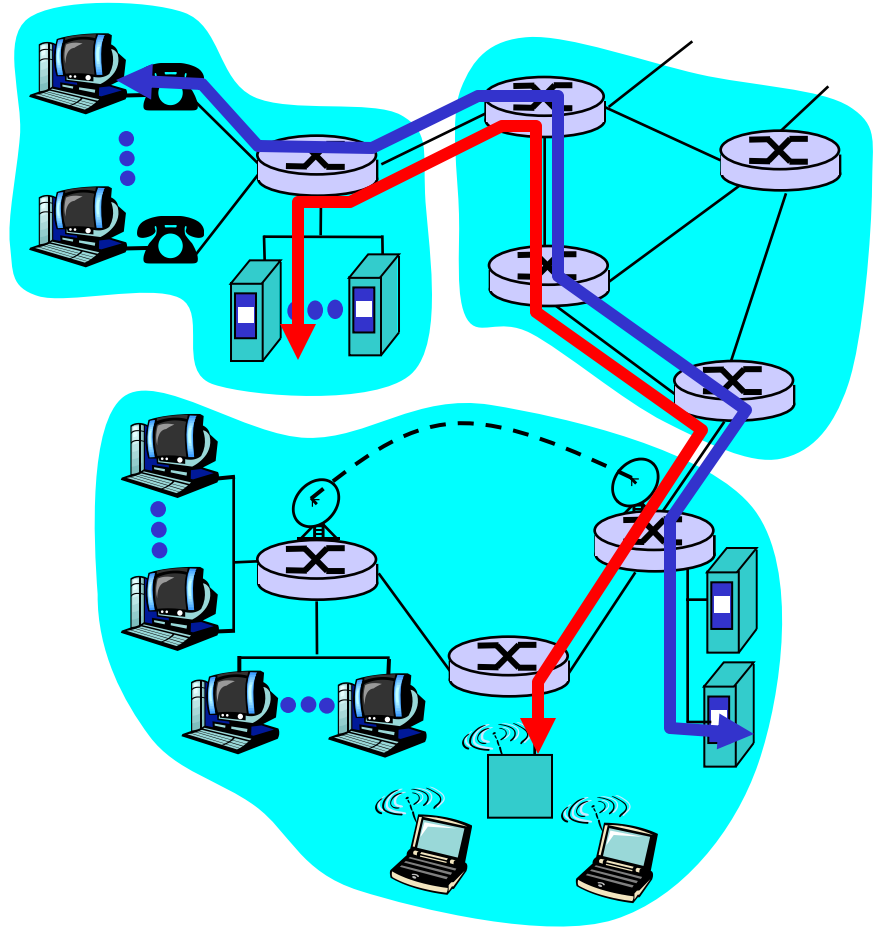
# Network Taxonomy



# Alt. 1: Circuit-Switching

End-to-end resources reserved for "call"

- ❑ Link bandwidth, switch capacity
- ❑ Dedicated resources with no sharing
- ❑ Guaranteed transmission capacity
- ❑ Call setup required
- ❑ "Blocking" may occur



# Alt. 1: Circuit-Switching

- ❑ Capacity of medium exceeds the capacity required for transmission of a single signal
  - How can we improve “efficiency”? Let's **multiplex**.
- ❑ Divide link bandwidth into “pieces”:
  - frequency division - FDMA
  - time division - TDMA
  - code division - CDMA (cellular networks)
  - wavelength division - WDM (optical)

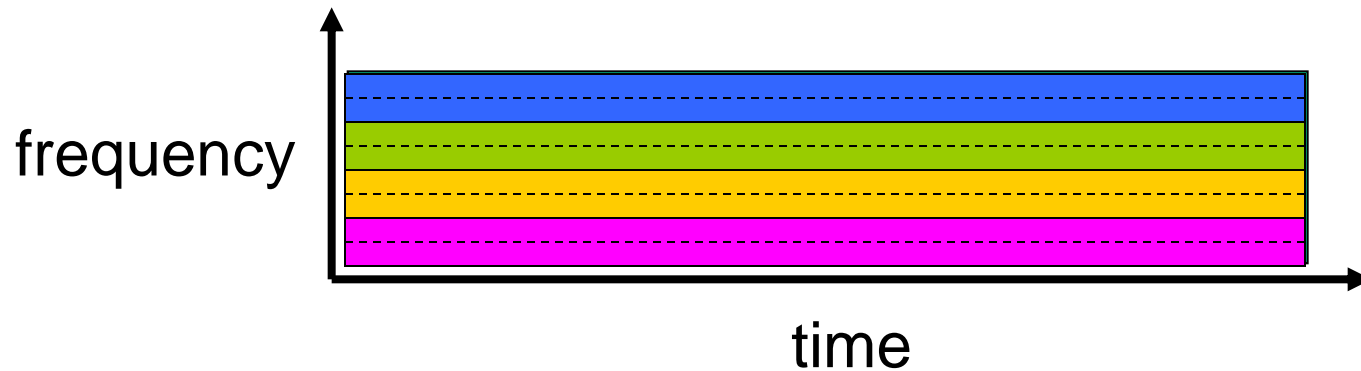
# Circuit-Switching: FDMA and TDMA

Example:

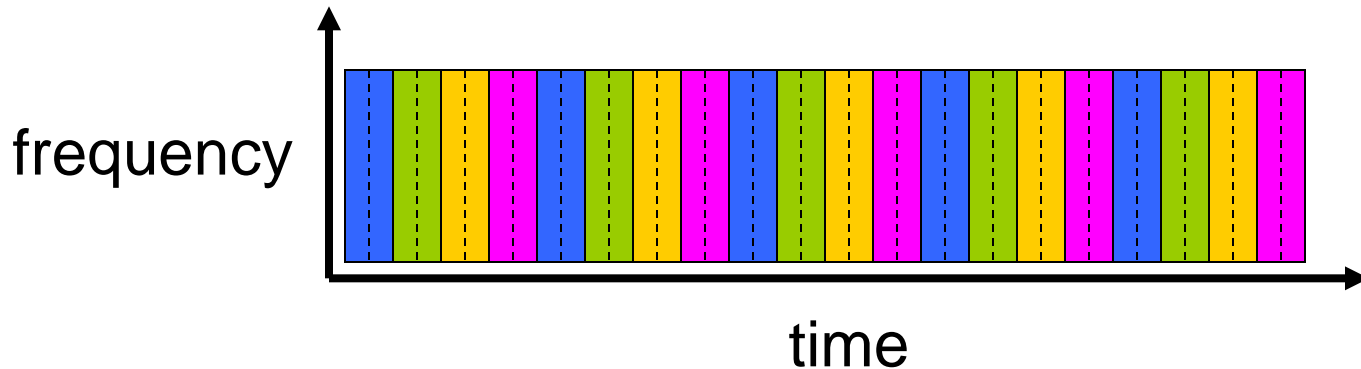
4 users



FDMA



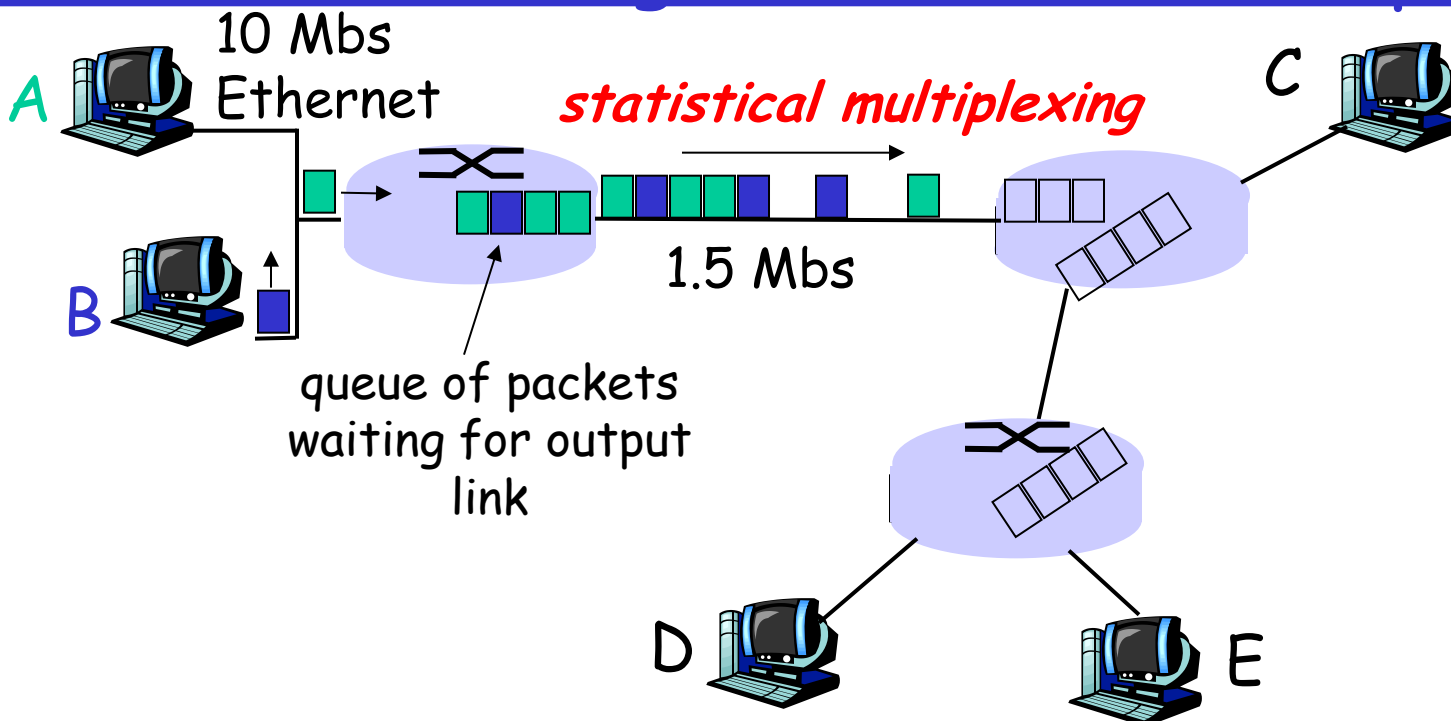
TDMA



## Alt. 2: Packet-Switching

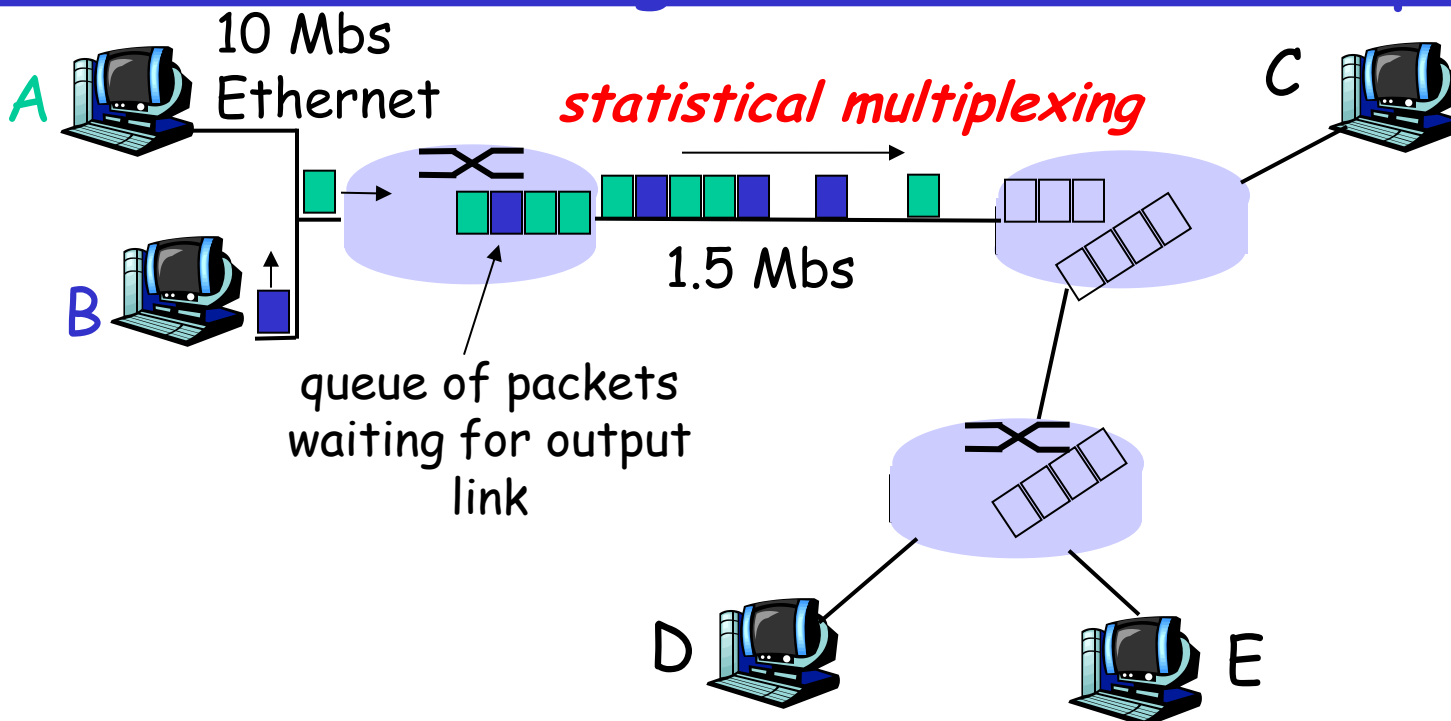
- ❑ source breaks long messages into smaller “packets”
- ❑ “store-and-forward” transmission
  - packets *share* network resources
  - each packet briefly uses full link bandwidth
- ❑ resource contention
  - aggregate resource demand can exceed amount available
  - congestion: packets queue, wait for link use
  - analogy: rush hour traffic in cities

# Packet-Switching: Statistical Multiplexing



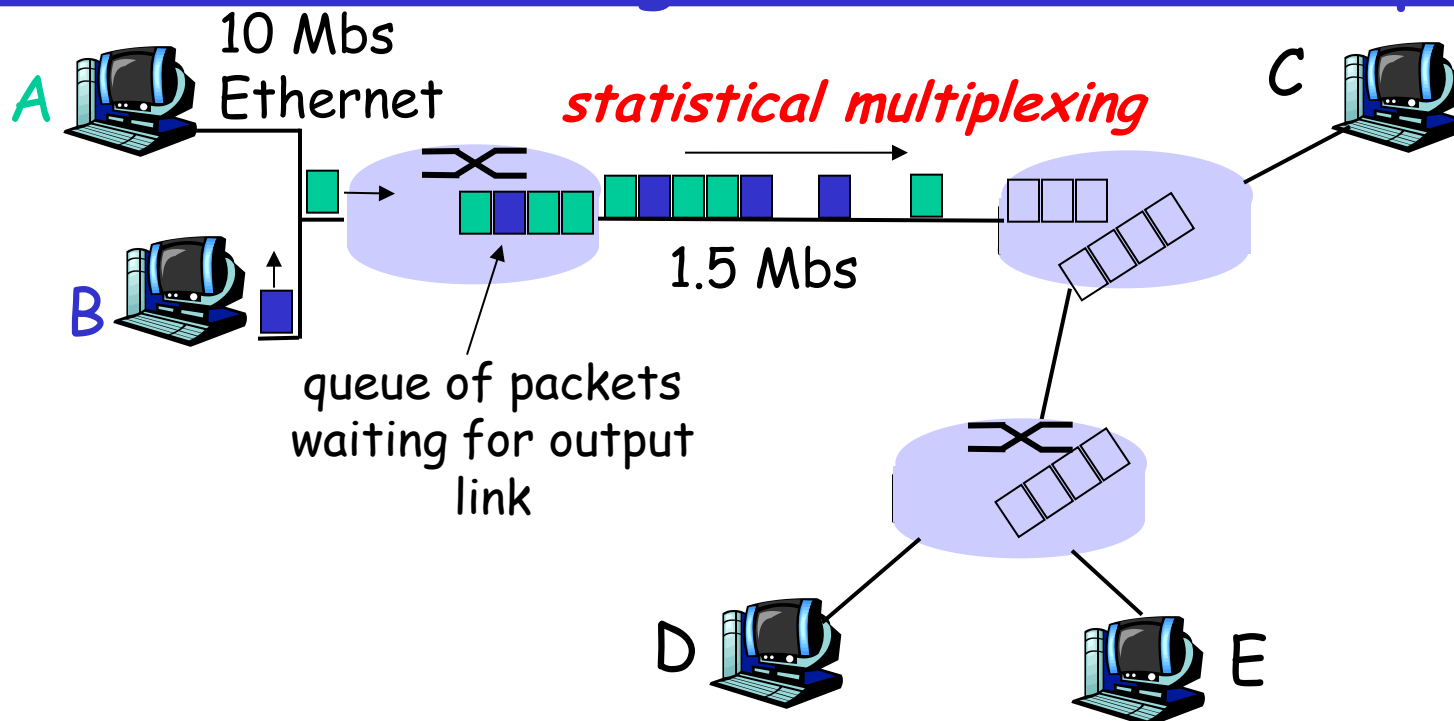
- ❑ Resource sharing great for bursty traffic
  - E.g., Sequence of A & B packets does not have fixed pattern - *statistical multiplexing*.
  - In contrast: In TDM each host gets same slot in revolving TDM frame.

# Packet-Switching: Statistical Multiplexing



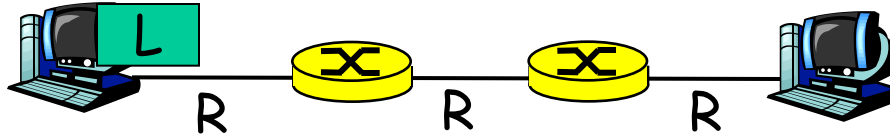
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- ❑ Is packet switching a "slam dunk" winner?
  - E.g., delay/loss and bandwidth guarantees ...

# Packet-switching: store-and-forward

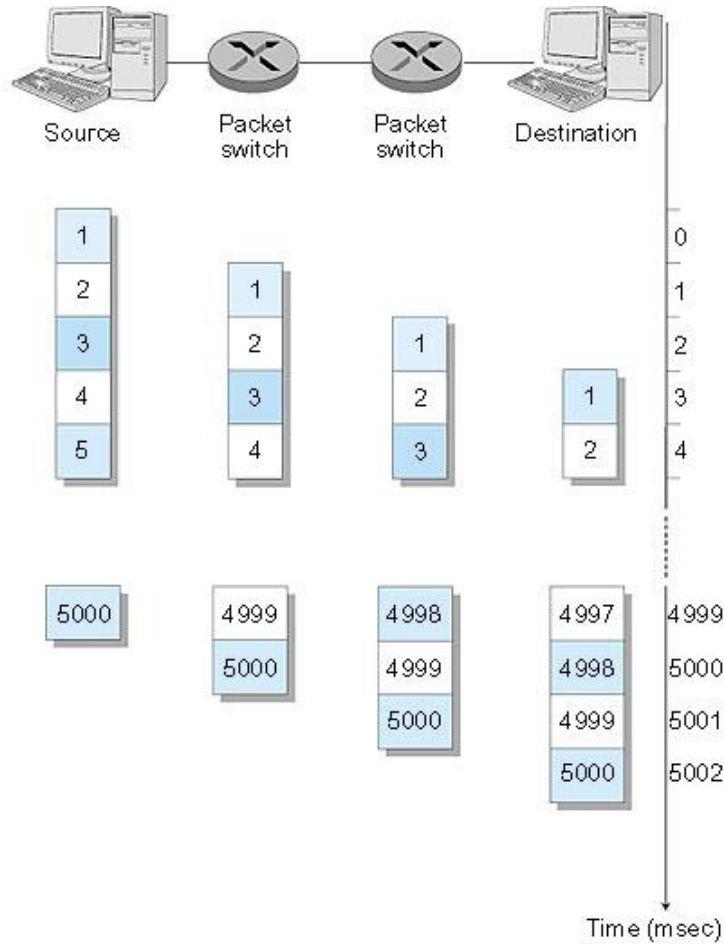


- ❑ Takes  $L/R$  seconds to transmit (push out) packet of  $L$  bits on to link or  $R$  bps
- ❑ Entire packet must arrive at router before it can be transmitted on next link: *store and forward*
- ❑ delay =  $3L/R$

## Example:

- ❑  $L = 7.5$  Mbits
- ❑  $R = 1.5$  Mbps
- ❑ delay = 15 sec

# Packet-Switching: Message Segmenting



Now break up the message into 5000 packets

- ❑ Each packet 1,500 bits
- ❑ 1 msec to transmit packet on one link
- ❑ *pipelining*: each link works in parallel
- ❑ Delay reduced from 15 sec to 5.002 sec

# Packet-switched networks: forwarding

## ❑ datagram network:

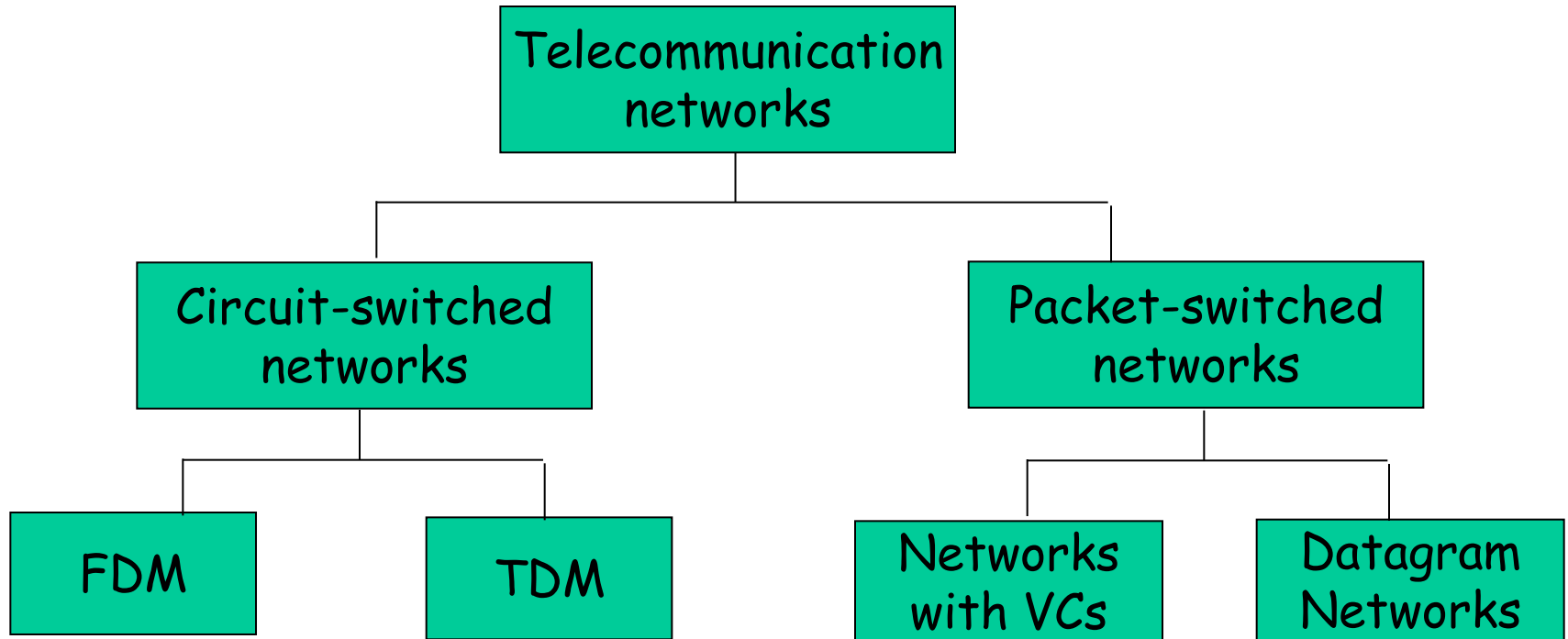
- *destination address* in packet determines next hop
- routes may change during session (flexible?)
- no “per flow” state, hence more scalable

## ❑ virtual circuit network:

- each packet carries tag (virtual circuit ID), tag determines next hop
- fixed path determined at *call setup time*
- path is **not** a dedicated path as in circuit switched (i.e., store & forward of packets)
- *routers maintain per-call state*

## ❑ datagram networks need per packet routing.

# Network Taxonomy





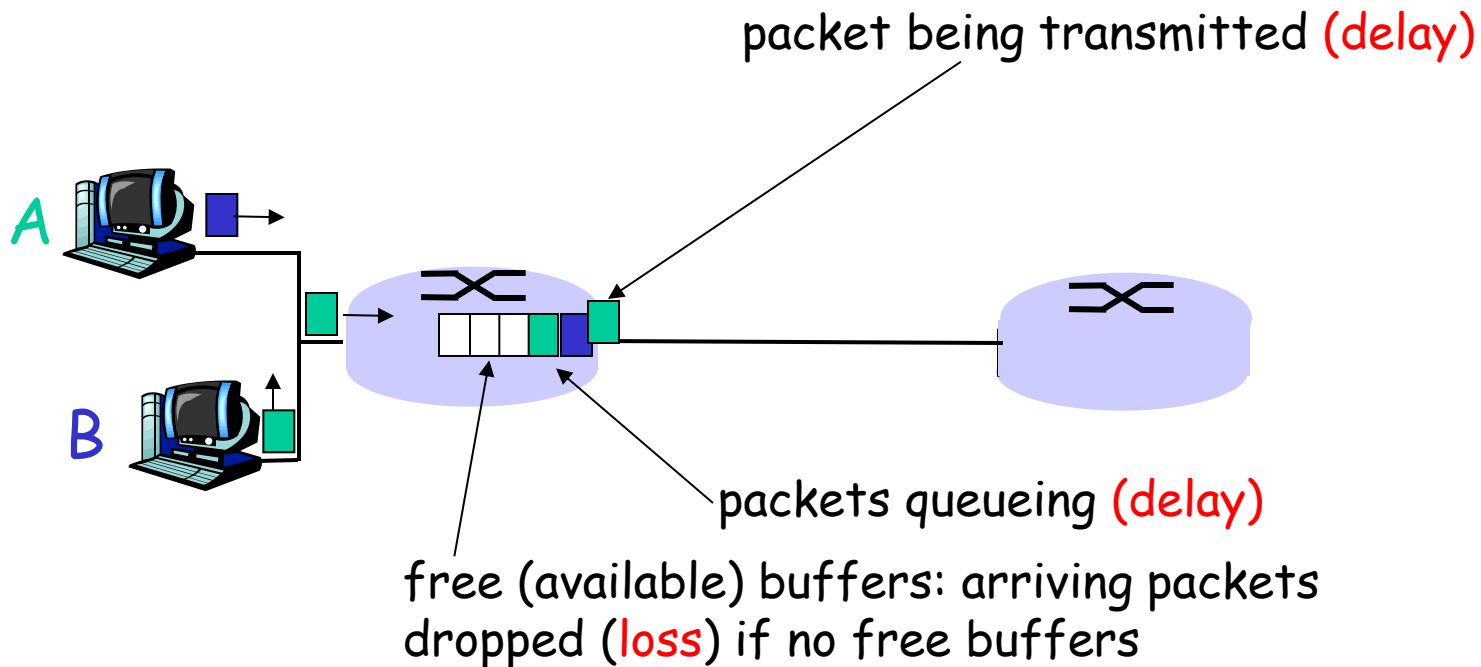
# Roadmap

- What is a Computer Network?
- Applications of Networking
- Classification of Networks
- Layered Architecture (and Protocols)
- Network Core
- Delay & Loss in Packet-switched Networks
- Structure of the Internet
- Summary

# How do loss and delay occur?

packets *queue* in router buffers

- ❑ packet arrival rate to link exceeds output link capacity
- ❑ packets queue, wait for turn
- ❑ if queue is full, arriving packets dropped (Drop-Tail)



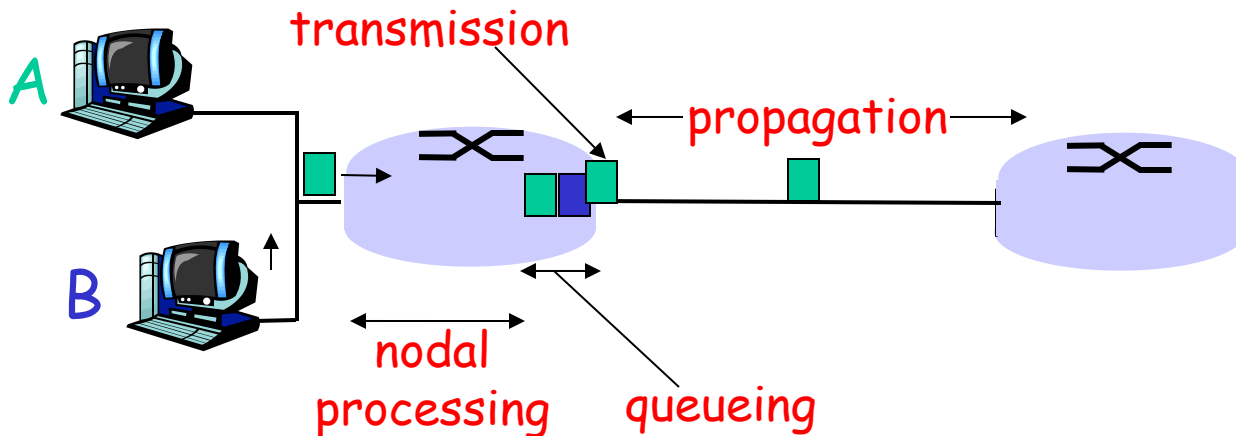
# Four sources of packet delay

## ❑ 1. Processing delay:

- check bit errors
- determine output link

## ❑ 2. Queueing delay:

- time waiting at output link for transmission
- depends on congestion level of router



# Delay in packet-switched networks

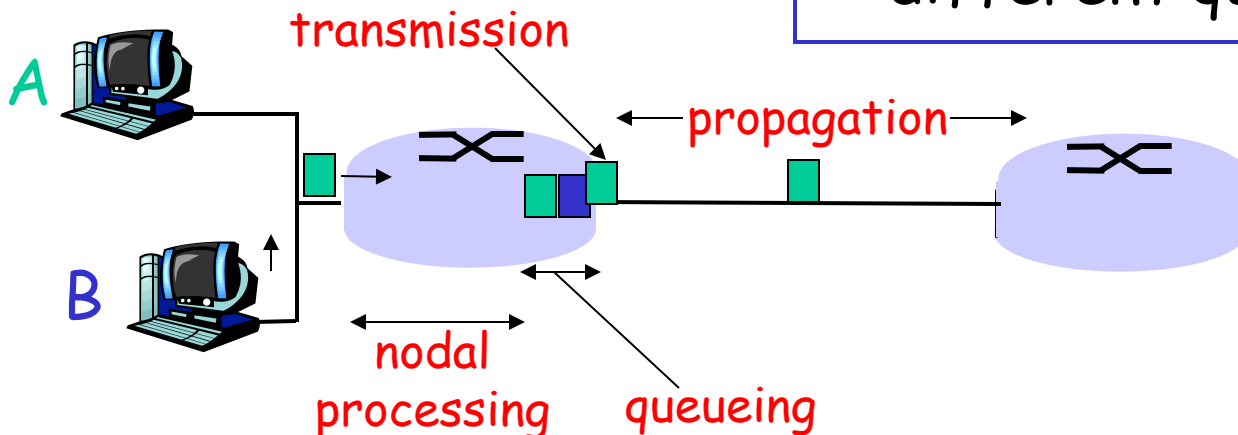
## 3. Transmission delay:

- ❑  $R$  = link bandwidth (bps)
- ❑  $L$  = packet length (bits)
- ❑ time to send bits into link =  $L/R$

## 4. Propagation delay:

- ❑  $d$  = length of physical link
- ❑  $s$  = propagation speed in medium ( $\sim 2 \times 10^8$  m/sec)
- ❑ propagation delay =  $d/s$

**Note:**  $s$  and  $R$  are very different quantities!



# Nodal processing delay

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

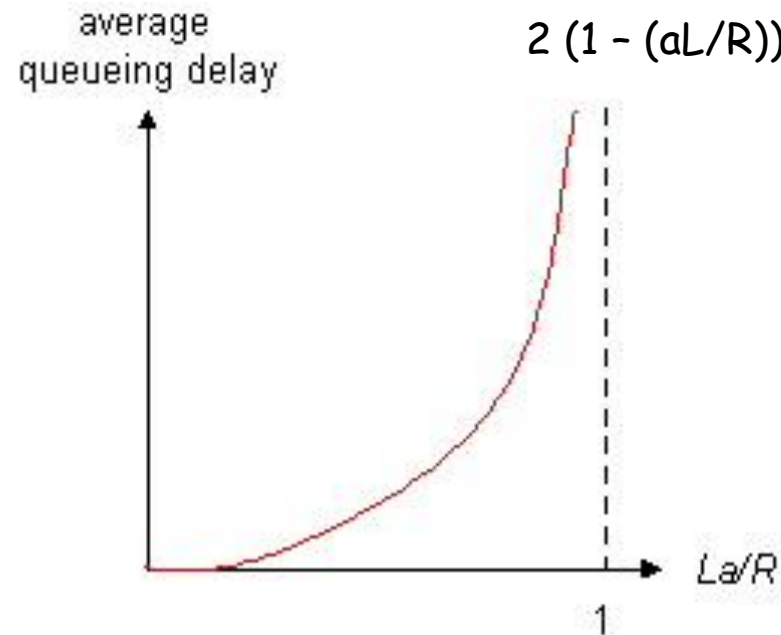
- ❑  $d_{\text{proc}}$  = processing delay
  - typically a few microsecs or less
- ❑  $d_{\text{queue}}$  = queuing delay
  - depends on congestion
- ❑  $d_{\text{trans}}$  = transmission delay
  - $= L/R$ , significant for low-speed links
- ❑  $d_{\text{prop}}$  = propagation delay
  - a few microsecs to hundreds of msecs

# Queueing delay (revisited)

$$W = \frac{L/R \ (aL/R)}{2 (1 - (aL/R))}$$

- $R$ =link bandwidth (bps)
- $L$ =packet length (bits)
- $a$ =average packet arrival rate

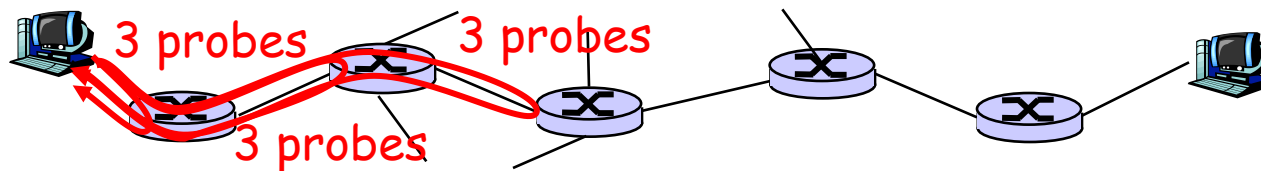
traffic intensity =  $aL/R$



- $aL/R \sim 0$ : average queueing delay small
- $aL/R \rightarrow 1$ : delays become large
- $aL/R > 1$ : more "work" arriving than can be serviced, average delay infinite!

# "Real" Internet delays and routes

- ❑ What do "real" Internet delay & loss look like?
- ❑ Traceroute program: provides delay measurement from source to router along end-to-end Internet path towards destination. For all  $i$ :
  - sends three packets that will reach router  $i$  on path towards destination
  - router  $i$  will return packets to sender
  - sender times interval between transmission and reply.



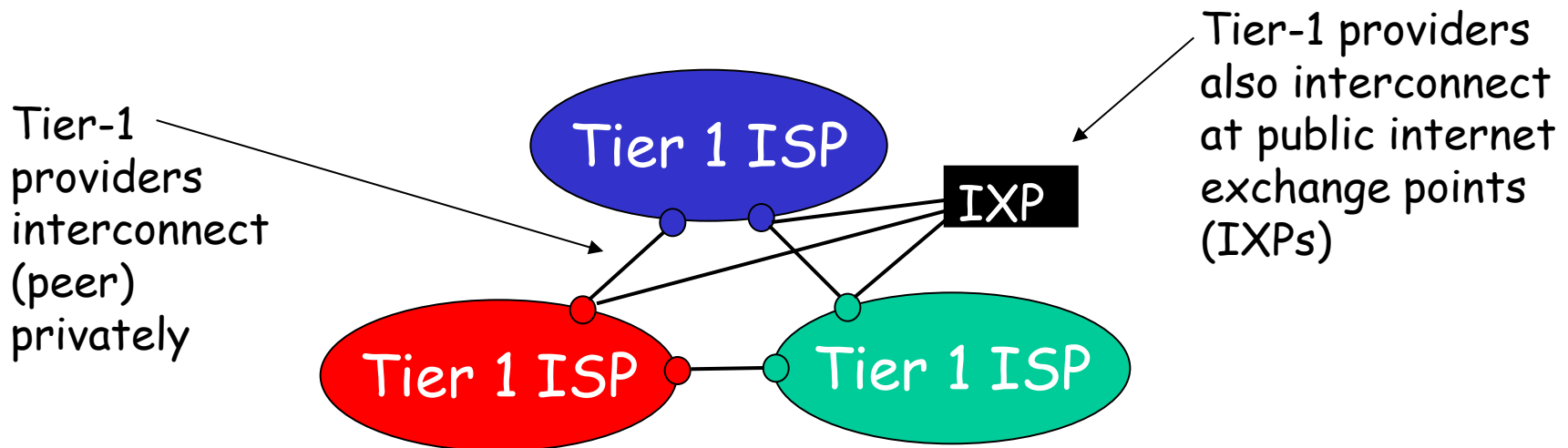


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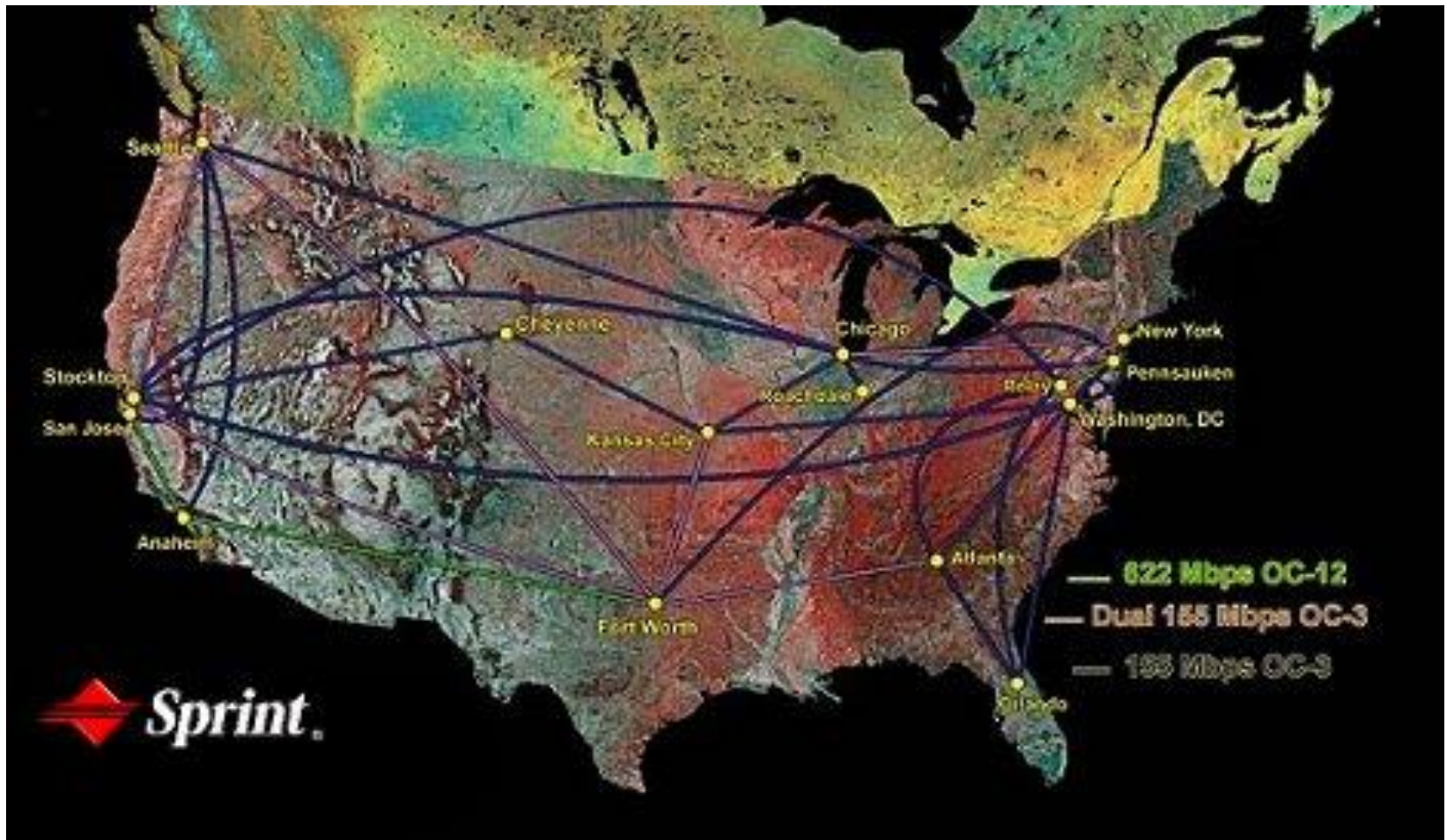
# Internet structure: network of networks

- ❑ roughly hierarchical
- ❑ **at center: "tier-1" ISPs** (e.g., UUNet, BBN/Genuity, Sprint, AT&T), national/international coverage
  - treat each other as equals



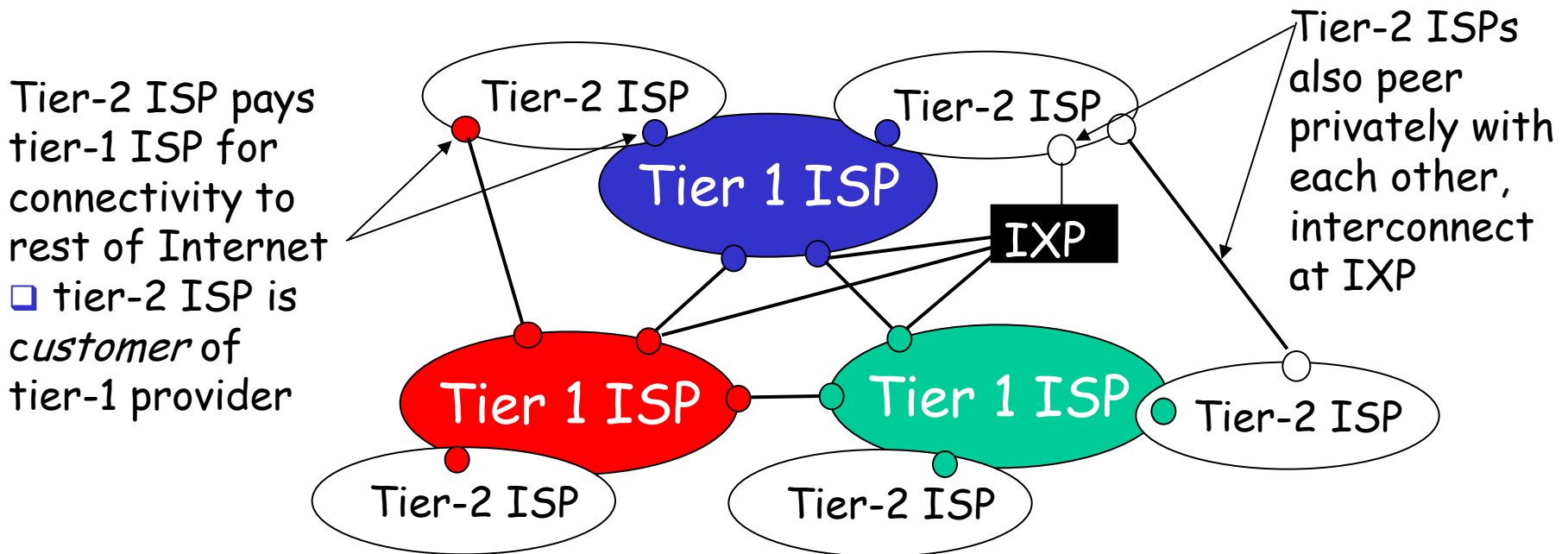
# Tier-1 ISP: e.g., Sprint

Sprint US backbone network



# Internet structure: network of networks

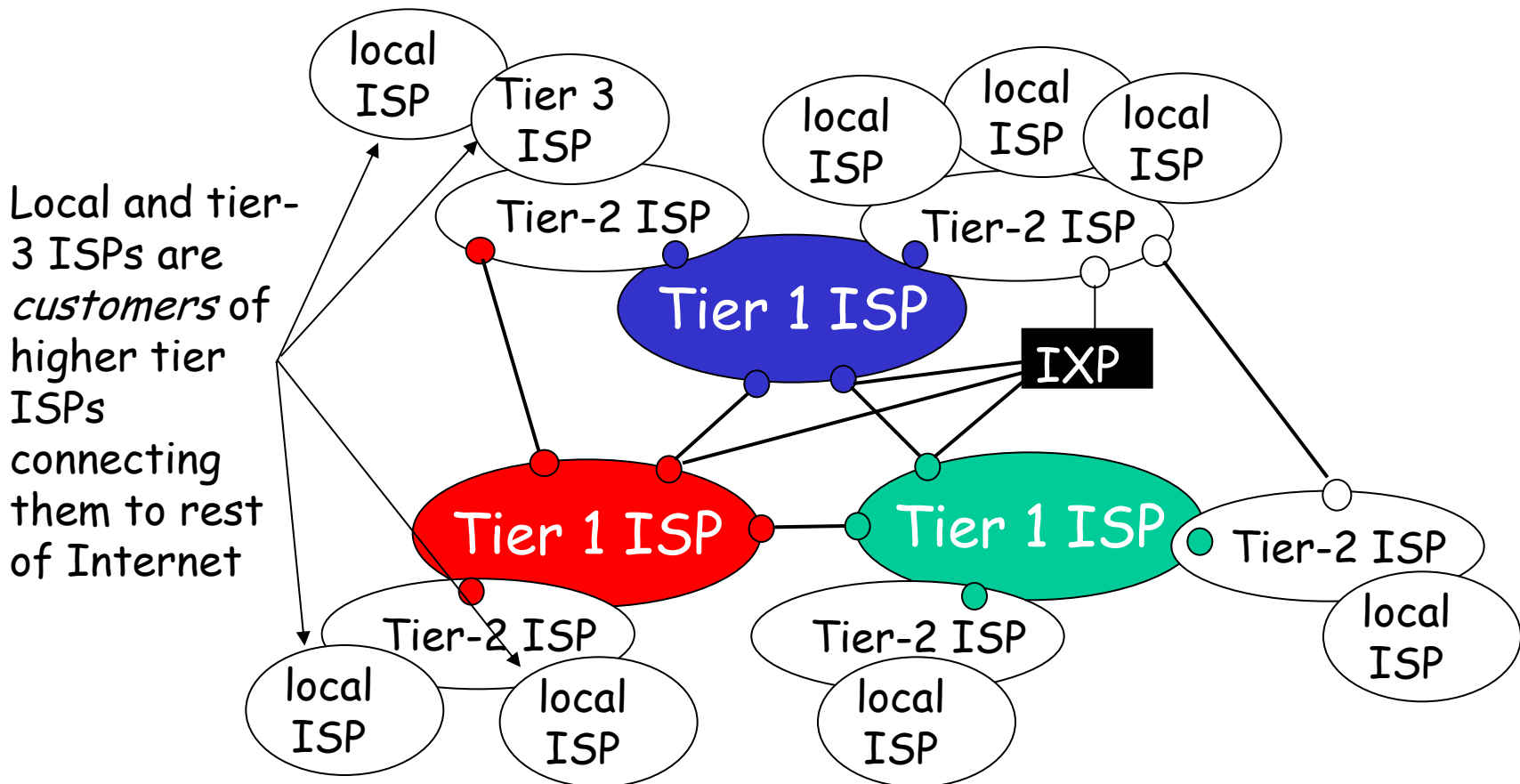
- ❑ "Tier-2" ISPs: smaller (often regional) ISPs
  - Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs



# Internet structure: network of networks

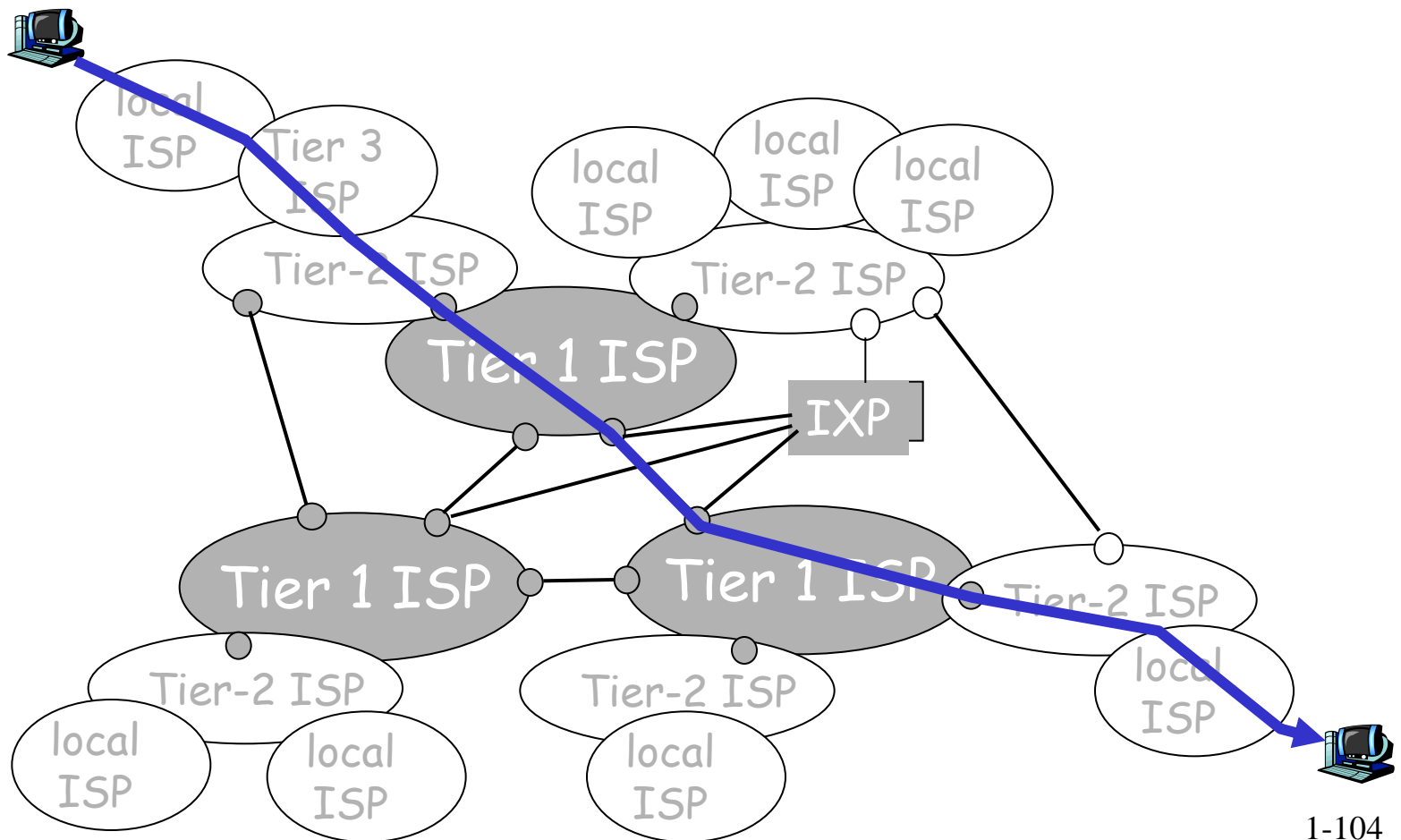
## □ "Tier-3" ISPs and local ISPs

- last hop ("access") network (closest to end systems)



# Internet structure: network of networks

- a packet passes through many networks!





# Introduction: Summary

## Covered a "ton" of material!

- ❑ Internet overview
- ❑ What's a protocol?
- ❑ Network edge, core, access network
  - packet-switching vs. circuit-switching
- ❑ Internet/ISP structure
- ❑ Performance: loss, delay
- ❑ Layering and service models
- ❑ Internet history

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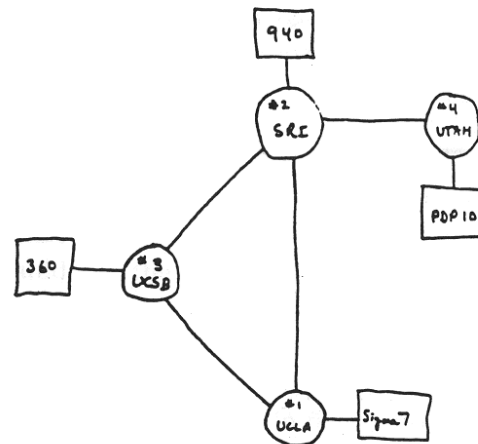
# Ohh, and the history ...

□ ...

# Internet History

## *1961-1972: Early packet-switching principles*

- ❖ **1961:** Kleinrock - queueing theory shows effectiveness of packet-switching
- ❖ **1964:** Baran - packet-switching in military nets
- ❖ **1967:** ARPAnet conceived by Advanced Research Projects Agency
- ❖ **1969:** first ARPAnet node operational
- ❖ **1972:**
  - ARPAnet public demonstration
  - NCP (Network Control Protocol) first host-host protocol
  - first e-mail program
  - ARPAnet has 15 nodes

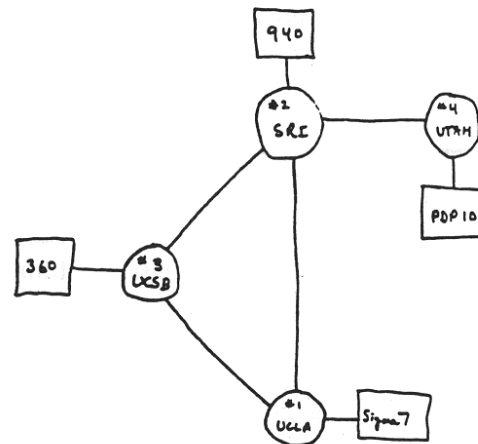


THE ARPA NETWORK

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THE ARPA NETWORK

# Internet History

## *1972-1980: Internetworking, new and proprietary nets*

- ❖ 1970: ALOHAnet satellite network in Hawaii
- ❖ 1974: Cerf and Kahn - architecture for interconnecting networks
- ❖ 1976: Ethernet at Xerox PARC
- ❖ late70's: proprietary architectures: DECnet, SNA, XNA
- ❖ late 70's: switching fixed length packets (ATM precursor)
- ❖ 1979: ARPAnet has 200 nodes

### *Cerf and Kahn's internetworking principles:*

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
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*define today's Internet  
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*1980-1990: new protocols, a proliferation of networks*

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- ❖ 1982: smtp e-mail protocol defined
- ❖ 1983: DNS defined for name-to-IP-address translation
- ❖ 1985: ftp protocol defined
- ❖ 1988: TCP congestion control
- ❖ new national networks: Cset, BITnet, NSFnet, Minitel
- ❖ 100,000 hosts connected to confederation of networks

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## *1990, 2000's: commercialization, the Web, new apps*

- ❖ early 1990's: ARPAnet decommissioned
- ❖ 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- ❖ early 1990s: Web
  - hypertext [Bush 1945, Nelson 1960's]
  - HTML, HTTP: Berners-Lee
  - 1994: Mosaic, later Netscape
  - late 1990's: commercialization of the Web

### late 1990's - 2000's:

- ❖ more killer apps: instant messaging, P2P file sharing
- ❖ network security to forefront
- ❖ est. 50 million host, 100 million+ users
- ❖ backbone links running at Gbps

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- ❖ P2P applications: BitTorrent (file sharing) Skype (VoIP), PPLive (video)
- ❖ more applications: YouTube, gaming, Twitter, facebook, ...
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