## EXAM

## (Tentamen)

# TDDI11 Embedded Software 

## 2020-08-19 kl: 08-12

## On-call (jour):

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## Admitted material:

- You can access your individual notes, books, and even to search the internet.
- No contacts, whether physical or virtual, are allowed during the duration of the exam with any person, whether the person is related to the course or not, except for contacting the examiner via email for questions if any.
- Any suspected breach will be systematically reported to the disciplinary board. We will use Urkund (an automatic tool to combat plagiarism).


## General instructions:

- The questions will refer to your "D1 D2" digits. These are the last two digits of the "anonymous-ID" you are assigned during the exam. You find the "anonymID" on the Lisam page you have retrieved this pdf from (i.e., the page you got the exam questions from). For instance, if your "anonymous-ID" is A-2709 then your D1 is 0 and your D2 is 9 (all in base 10). If your "anonymous-ID" is A-123 then your D1 is 2 and your D2 is 3 . Ask the examiner if this is unclear.
- The assignments are not ordered according to difficulty.
- You may answer in either English or Swedish.
- Do not take pictures or draw. We only accept PDF files obtained from a text editor or a word processor. You are free to use any text editor (examples include notepad, vim, gedit, emacs, etc) or other text-based office programs (Microsoft Word or Open/Libre Office or LaTeX). We expect you to generate and to submit a single PDF.
- Be precise in your statements.
- Clearly motivate all statements and reasoning.
- Explain calculations and solution procedures.
- If in doubt about the question, write down your interpretation and assumptions.
- Grading: U, 3, 4, 5. The preliminary grading thresholds for p points are:

| $0 \leq \mathrm{p}<20:$ | U |
| ---: | ---: | ---: |
| $20 \leq \mathrm{p} \leq 30:$ | 3 |
| $31 \leq \mathrm{p} \leq 35:$ | 4 |
| $36 \leq \mathrm{p} \leq 40:$ | 5 |

## Question 1. (10 points)

a. Explain at a high level the steps involved when sending the byte "D2" using a UART. For example, if your D2 is 3, then the question is about the steps involved in sending the byte: 0000 0011. You do not need to consider a parity bit. (2pts)
b. Give an example of an embedded system for which you would expect hard real time guarantees and an example of an embedded system for which you would expect soft-real time guarantees. Explain and justify your choices. (2pts)
c. Assume a 32-bit variable " $x$ " to which you have saved the value generated by some peripheral. Write a C function "int myCheck(int x)" that returns true exactly when the (D1 + 2)th most significant bit equals the D1:th most significant bit. (2pts)
d. Explain the difference between "concurrent" and "over-the wall" approaches to engineering projects. Which approach seems better to you? Explain. (2pts)
e. Assume a memory mapped display at address $0 x B 2000$. The sequence:

$$
\begin{aligned}
& 0 x 47,0 x 46,0 x 6 f, 0 x 73,0 x 6 f, 0 x 7 f, 0 x 64,0 x 6 c, 0 x 20,0 x 4 d \text {, } \\
& 0 x 6 \mathrm{c}, 0 \mathrm{x} 2 \mathrm{~b}, 0 \mathrm{x} 75,0 \mathrm{x} 7 \mathrm{~b}, 0 \mathrm{x} 63,0 \mathrm{x} 46,0 \mathrm{x} 6 \mathrm{~b}, 0 \mathrm{x} 42,0 \mathrm{x} 21,0 \mathrm{x} 78 \text {, } \\
& 0 \times 2 \mathrm{a}, 0 \times 68,0 \times 2 \mathrm{a}, 0 \times 1 \mathrm{~d}, 0 \times 2 \mathrm{a}, 0 \times 5 \mathrm{a}, 0 \times 2 \mathrm{a}, 0 \times 63,0 \times 2 \mathrm{a}, 0 \mathrm{x} 1 \mathrm{f} \text {, } \\
& 0 x 47,0 x 46,0 x 6 f, 0 x 73,0 x 6 f, 0 x 7 f, 0 x 64,0 x 6 c, 0 x 20,0 x 4 d \text {, } \\
& \text { 0x6c,0x2b,0x75,0x7b,0x63,0x46,0x6b,0x42,0x21,0x78, } \\
& 0 \times 2 \mathrm{a}, 0 \times 68,0 \mathrm{x} 2 \mathrm{a}, 0 \mathrm{x} 1 \mathrm{~d}, 0 \times 2 \mathrm{a}, 0 \times 5 \mathrm{a}, 0 \times 2 \mathrm{a}, 0 \mathrm{x} 63,0 \times 2 \mathrm{a}, 0 \mathrm{x} 1 \mathrm{f} \text {, } \\
& 0 \times 47,0 \times 46,0 \times 6 f, 0 x 73,0 x 6 f, 0 x 7 f, 0 x 64,0 x 6 \mathrm{c}, 0 \mathrm{x} 20,0 \times 4 \mathrm{~d} \text {, } \\
& 0 x 6 \mathrm{c}, 0 \mathrm{x} 2 \mathrm{~b}, 0 \mathrm{x} 75,0 \mathrm{x} 7 \mathrm{~b}, 0 \mathrm{x} 63,0 \mathrm{x} 46,0 \mathrm{x} 6 \mathrm{~b}, 0 \mathrm{x} 42,0 \mathrm{x} 21,0 \mathrm{x} 78 \text {, } \\
& 0 \times 2 \mathrm{a}, 0 \mathrm{x} 68,0 \mathrm{x} 2 \mathrm{a}, 0 \mathrm{x} 1 \mathrm{~d}, 0 \times 2 \mathrm{a}, 0 \times 5 \mathrm{a}, 0 \times 2 \mathrm{a}, 0 \mathrm{x} 63,0 \times 2 \mathrm{a}, 0 \mathrm{x} 1 \mathrm{f} \text {, } \\
& 0 x 47,0 x 46,0 x 6 f, 0 x 73,0 x 6 f, 0 x 7 f, 0 x 64,0 x 6 c, 0 x 20,0 x 4 d
\end{aligned}
$$

Consists of 100 bytes and is stored sequentially starting with $0 x 47$ at byte $0 x B 2000$ encoding the first character of the first row and $0 \times 46$ at byte $0 x B 2001$ encoding the second character of the first row. The display has 10 rows with 10 characters on each row. Each character is encoded with one byte.

1. What is the address of the byte at line D1 and at column D2?
2. Write a function char* addressOf(int row, int col) that returns the address of the byte at line row and column col. (you can assume $0 \leq$ row $\leq 9$ and $0 \leq$ col $\leq 9$ ). (2pts)


## Question 4. (6 points)

The Mealy machine described in the Table below has 4 states $\left\{s_{0}, s_{1}, s_{2}, s_{3}\right\}$ where $s_{0}$ is the initial state. The machine takes sequences of 0 s and 1 s as input. It outputs 1 exactly when the last three inputs (including overlap) build the sequence 101. For instance, the machine outputs the sequence " 00101010001 " when it reads " 10101011101 ".

|  | in:0 | in: 1 |
| :---: | :--- | :--- |
| Initial: $s_{0}$ | out: $0 /$ goto: $s_{0}$ | out: $0 /$ goto: $s_{1}$ |
| $s_{1}$ | out: $0 /$ goto: $s_{2}$ | out: $0 /$ goto: $s_{1}$ |
| $s_{2}$ | out: $0 /$ goto: $s_{0}$ | out: $1 /$ goto: $s_{3}$ |
| $s_{3}$ | out: $0 /$ goto: $s_{2}$ | out: $0 /$ goto: $s_{1}$ |

Let "d" be " $(\mathrm{D} 2 \% 3)+3$ ". If your D2 is 9 , then your d is 3 . If your D2 is 5 then your d is 5 . Give a Mealy machine that reads sequences of 0 s and 1 s as input. The machine should output 1 exactly when it reads a zero after reading a non-empty sequence of ones where the number of ones is divisible by d .

For instance, if your $d$ is 3 , then the machine should output 1 each time it reads a zero after a consecutive sequence of $n$ ones where $n$ is divisible by 3 , i.e., $n$ is in $\{3,6,9,12$, $15, \ldots\}$. It should output 0 otherwise.

Follow some examples:

- If your $\mathrm{d}=3$ and the machine reads " $0010111100011100 \ldots$... then it should output "0000000000000010..."
- If your $\mathrm{d}=4$ and the machine reads " 0010111100011100 ..." then it should output "0000000010000000 ..."
- If your $\mathrm{d}=5$ and the machine reads " $00111111111100 \ldots$... then it should output "00000000000010 ..."


## Question 5. (5 points)

As it is often the case, newly pushed stack elements get smaller addresses. A function with one argument has just been called. The stack looks as follows:

Answer the following questions:

1. The 4 bytes " $0 \times 3$ fffffe 8 , 0x3fffffe9, 0x3fffffea, $0 \times 3 f f f f f e b "$ contain an address. Explain what this address represents and why is it stored at the top of the stack? (2pts)
2. What role is usually played by the register ebp? Why is it inconvenient to only use esp? (2pts)
3. Assume a big-endian system. Suppose the argument to the function is the 32 bits integer corresponding to the sequence (in decreasing order of significance) of 4 bytes with the most significant byte containing the value 3 (written $0 x 3$ ), followed by a byte containing your value $D 1$ (if your $D 1$ is 9 then the second byte is 0 x 9 ), then a byte containing the value 4 and finally (the least significant byte) containing your value $D 2$. Give the address in the stack of each one of these 4 bytes. (3pts)
```
address
--------------
0x3fffffdc
0x3fffffdd
0x3fffffde
0x3fffffdf
0x3fffffe0
0x3fffffel
0x3fffffe2
0x3fffffe3
0x3fffffe4
0x3fffffe5
0x3fffffe6
0x3fffffe7
0x3fffffe8 <-- esp
0x3fffffe9
0x3fffffea
0x3fffffeb
0x3fffffec <-- (arg)
0x3fffffed
0x3fffffee
0x3fffffef
0x3ffffff0 <-- (caller saved register)
0x3ffffffl
0x3ffffff2
0x3ffffff3
0x3ffffff4 <-- (caller saved register)
0x3ffffff5
0x3ffffff6
0x3fffffff
0x3ffffff8 <-- ebp
0x3ffffff9
0x3ffffffa
0x3ffffffb
```

