

# **Datorarkitektur**

**(Advanced Computer Architecture)**

**Petru Eles**  
**Institutionen för Datavetenskap (IDA)**  
**Linköpings Universitet**

**email: [petru.eles@liu.se](mailto:petru.eles@liu.se)**  
**<http://www.ida.liu.se/~petel>**  
**phone:013 281396**  
**B building, room 329:220**

# Course Information

Web page: <http://www.ida.liu.se/~TDDI03>

Examination: written, January 8th, 2024, kl. 14 - 18.

Lecture notes: available from the web page, latest 24 hours before the lecture.

Text book: William Stallings: *Computer Organization and Architecture*, Pearson/Prentice Hall, 9th edition, 2013 or 10th edition, 2016 or 11th edition, 2021.

In addition to lectures: Two big seminars!

# Preliminary Course Plan

## Lecture 1

*Introduction:* Outline, Basic computer architecture and organization, Basic functions of a computer and its main components, The von Neumann architecture.

This is to refresh your memory!!!

## Lectures 2 and 3

*The Memory System:* Memory hierarchy, Cache memories, Virtual memories, Memory management.

## Lectures 4 and 5

*Instruction Pipelining:* Organization of pipelined units, Pipeline hazards, Reducing branch penalties, Branch prediction strategies.

## Lectures 6

*RISC Architectures:* An analysis of instruction execution for code generated from high-level language programs, Compiling for RISC architectures, Main characteristics of RISC architectures, RISC-CISC trade-offs.

# Preliminary Course Plan

## Lectures 7 and 8

***Superscalar Architectures:*** Instruction level parallelism and machine parallelism, Hardware techniques for performance enhancement, Data dependencies, Policies for parallel instruction execution, Limitations of the superscalar approach.

## Lectures 9 and 10

***VLIW Architectures:*** The VLIW approach - advantages and limitations. Compiling for VLIW architectures. The Merced (Itanium) architecture.

## Lectures 11 and 12

***Architectures for Parallel Computation:*** Parallel programs, Performance of parallel computers, A classification of computer architectures, Array processors, Multiprocessors, Multicomputers, Vector processors.

**Multiprocessors on chip.**

**Multithreaded processors.**

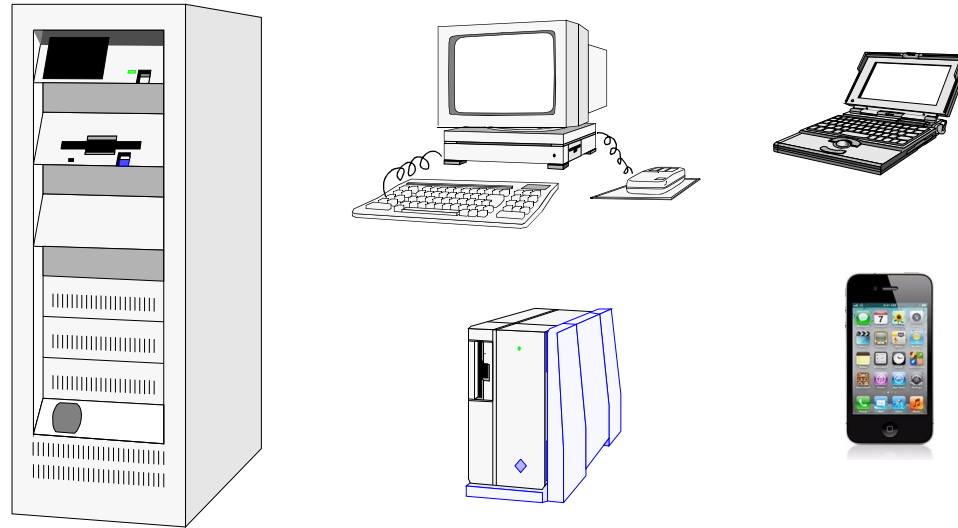
**General purpose graphic processors.**

# **COMPUTER ARCHITECTURE**

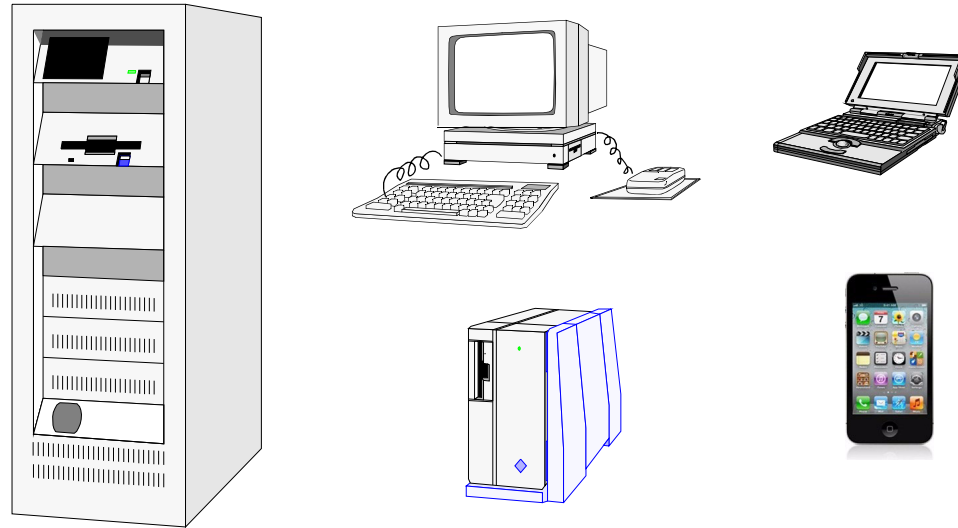
## **(BASIC ISSUES)**

- 1. What is a Computer/Computer System?**
- 2. The von Neumann Architecture**
- 3. Application Specific vs. General-Purpose**
- 4. Representation of Data and Instructions**
- 5. Instruction Execution**
- 6. The Control Unit**
- 7. The Computer System**
- 8. Main and Secondary Memory**
- 9. The Intel x86 and ARM Families**

# What Is a Computer?



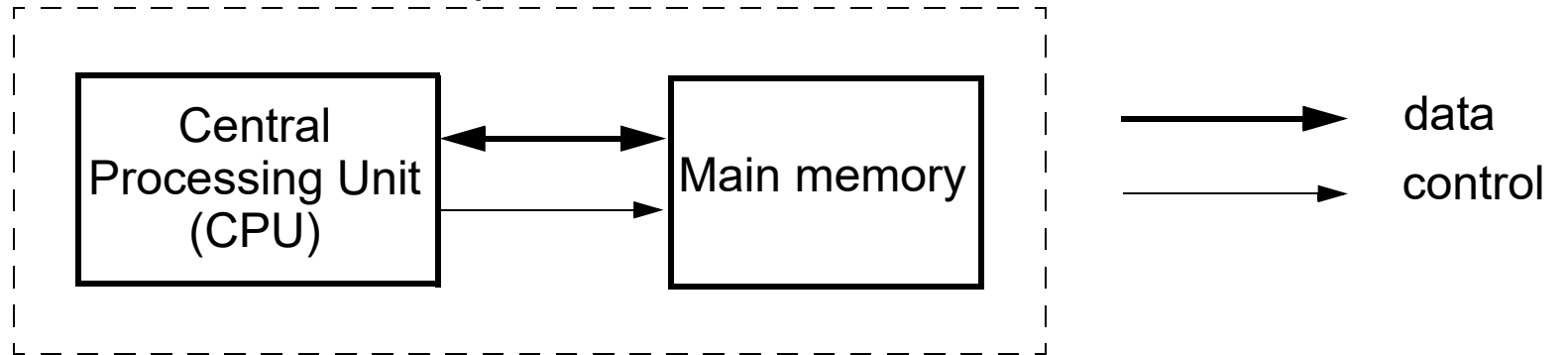
# What Is a Computer?



- A computer is a data processing machine which is operated automatically under the control of a list of instructions (called a program) stored in its main memory.

# What Is a Computer?

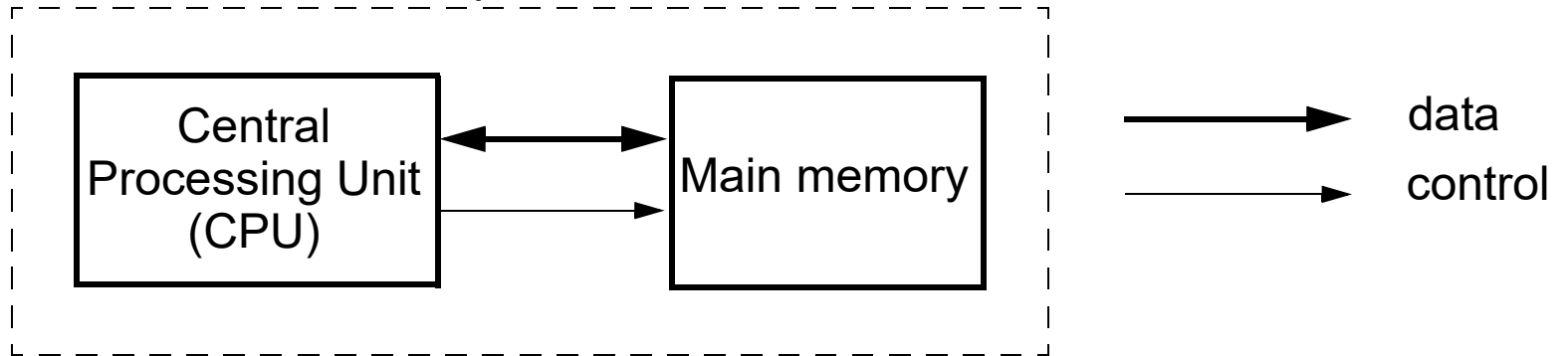
The "core" of the computer





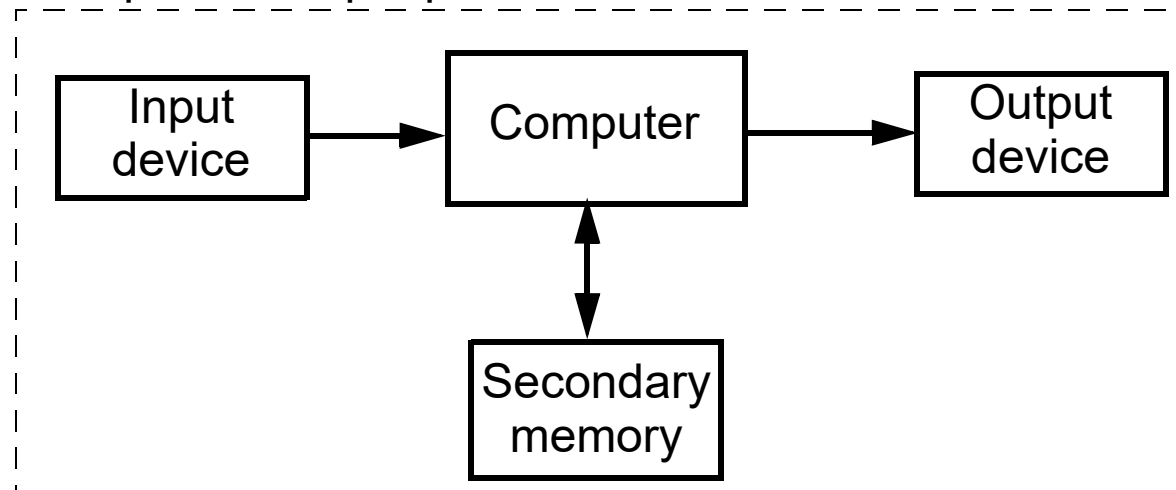
# What Is a Computer?

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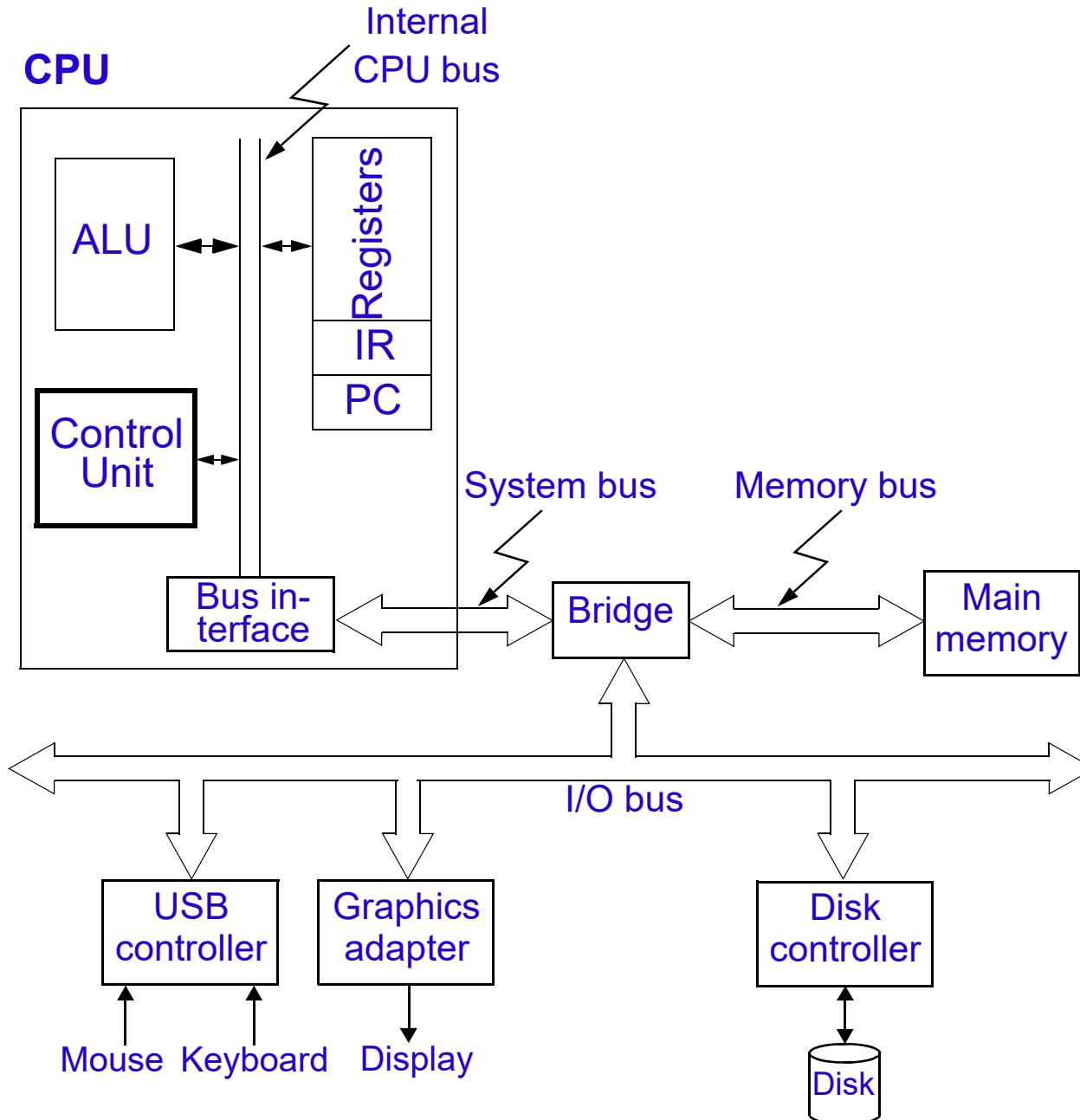


- Besides the "core" we also have the peripherals. Computer peripherals include input devices, output devices, and secondary memories. That makes the whole to a *Computer System*.

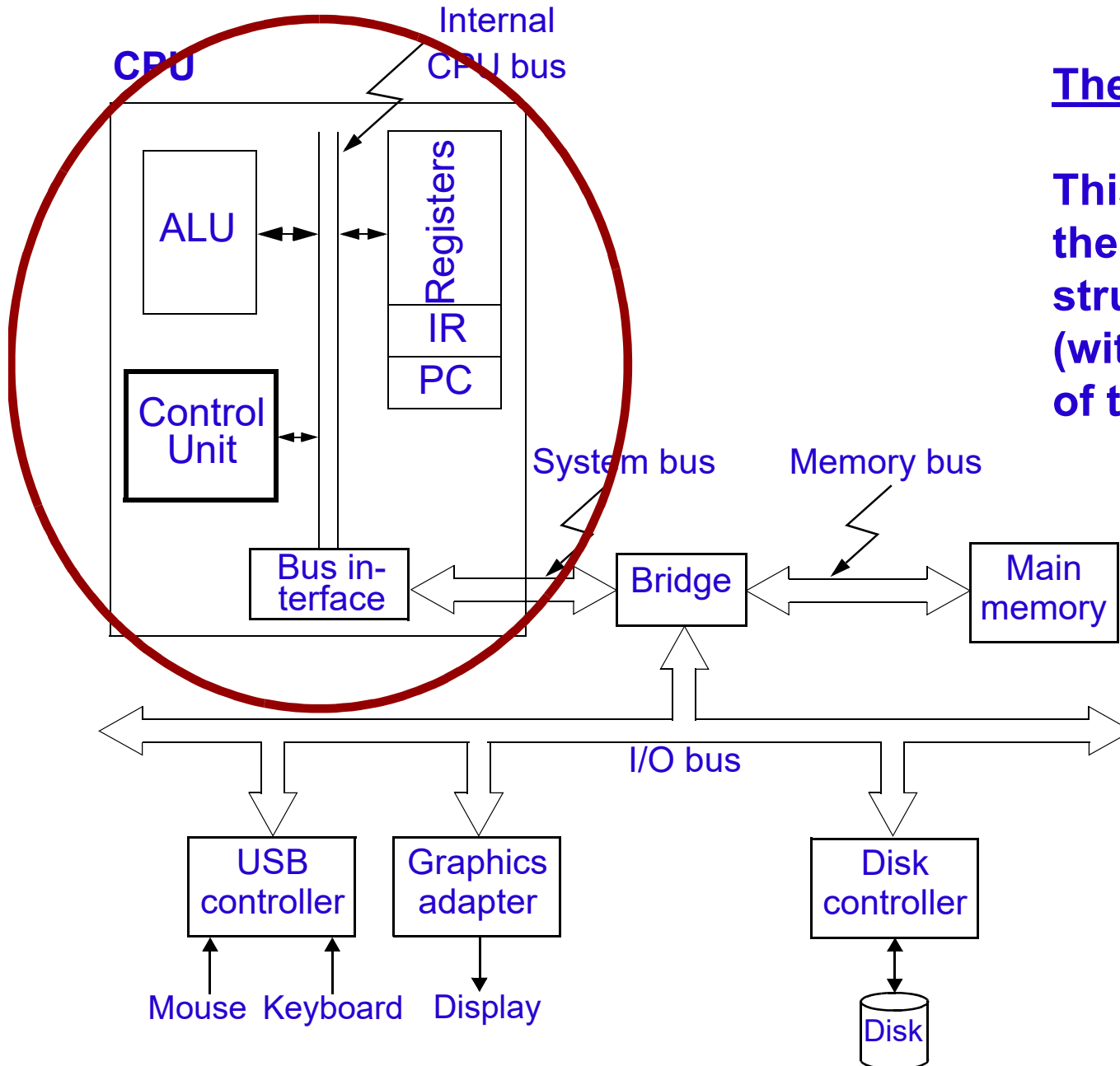
Computer with peripherals



# Computer Systems



# Computer Systems



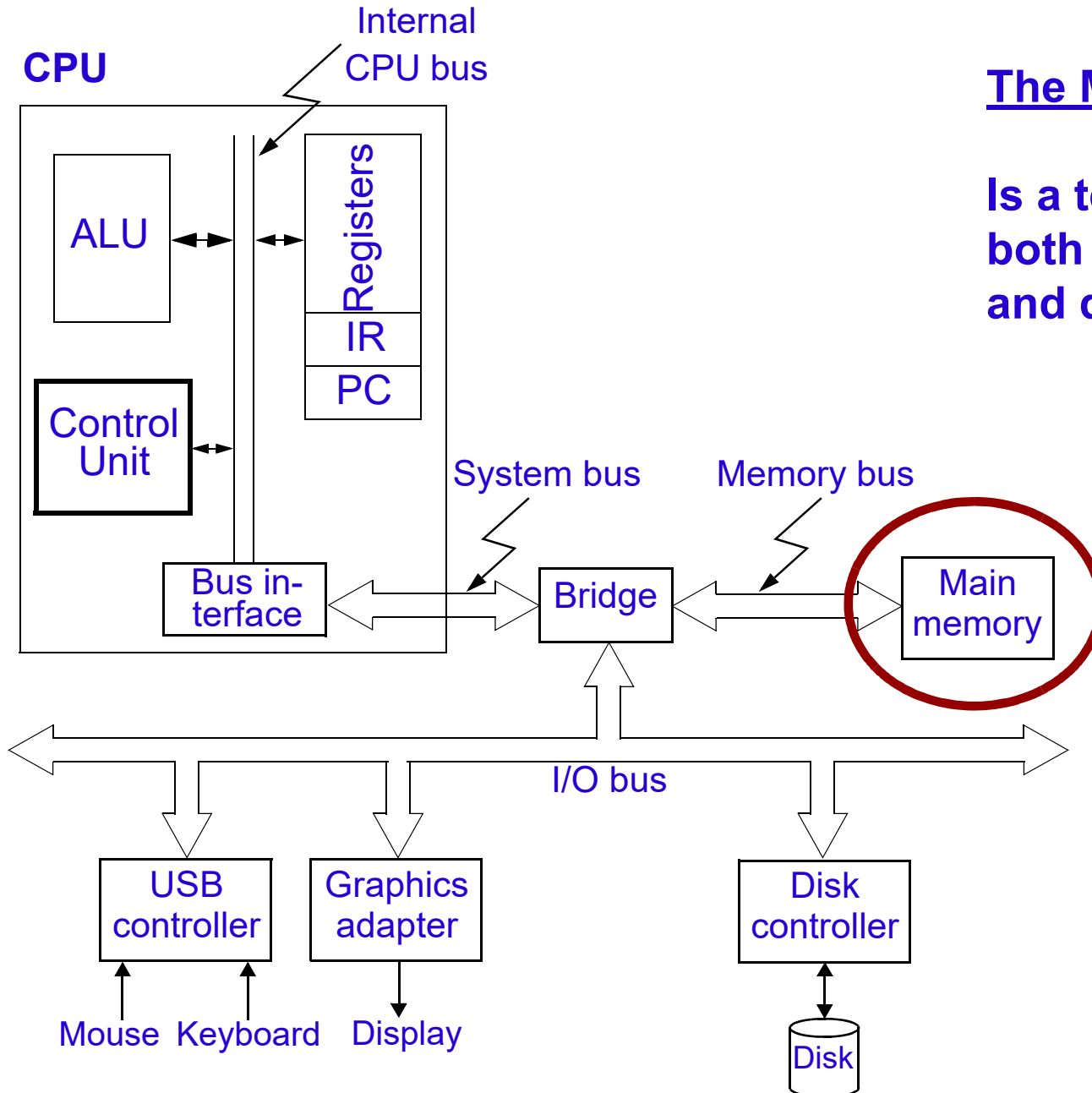
## The CPU (Central Processing Unit)

This is the hart of the system; it is the engine that interprets the instructions and executes them (with the help of other components of the computer system).

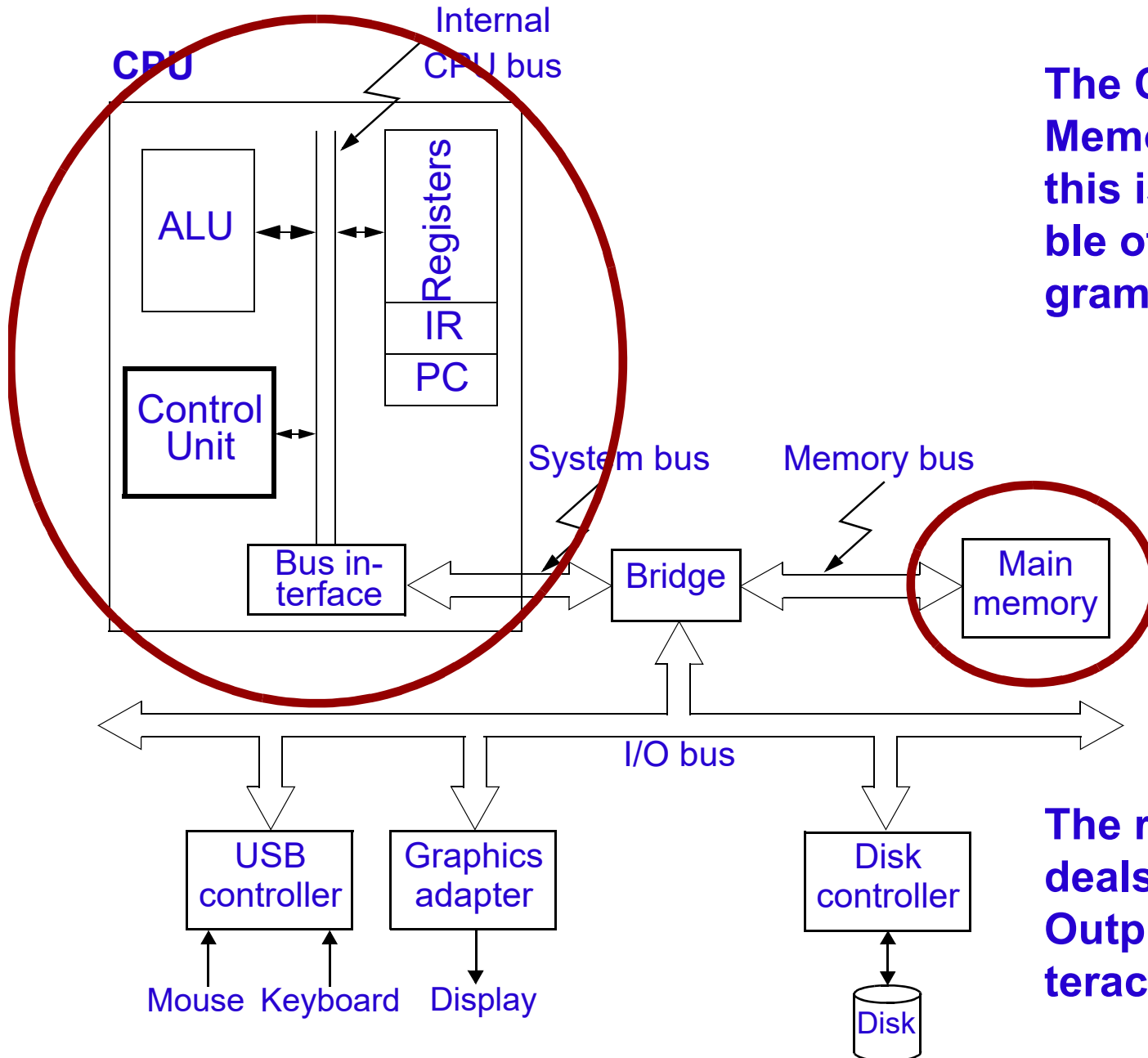
# Computer Systems

## The Main Memory

Is a temporary storage that stores both instructions (the program) and data.



# Computer Systems



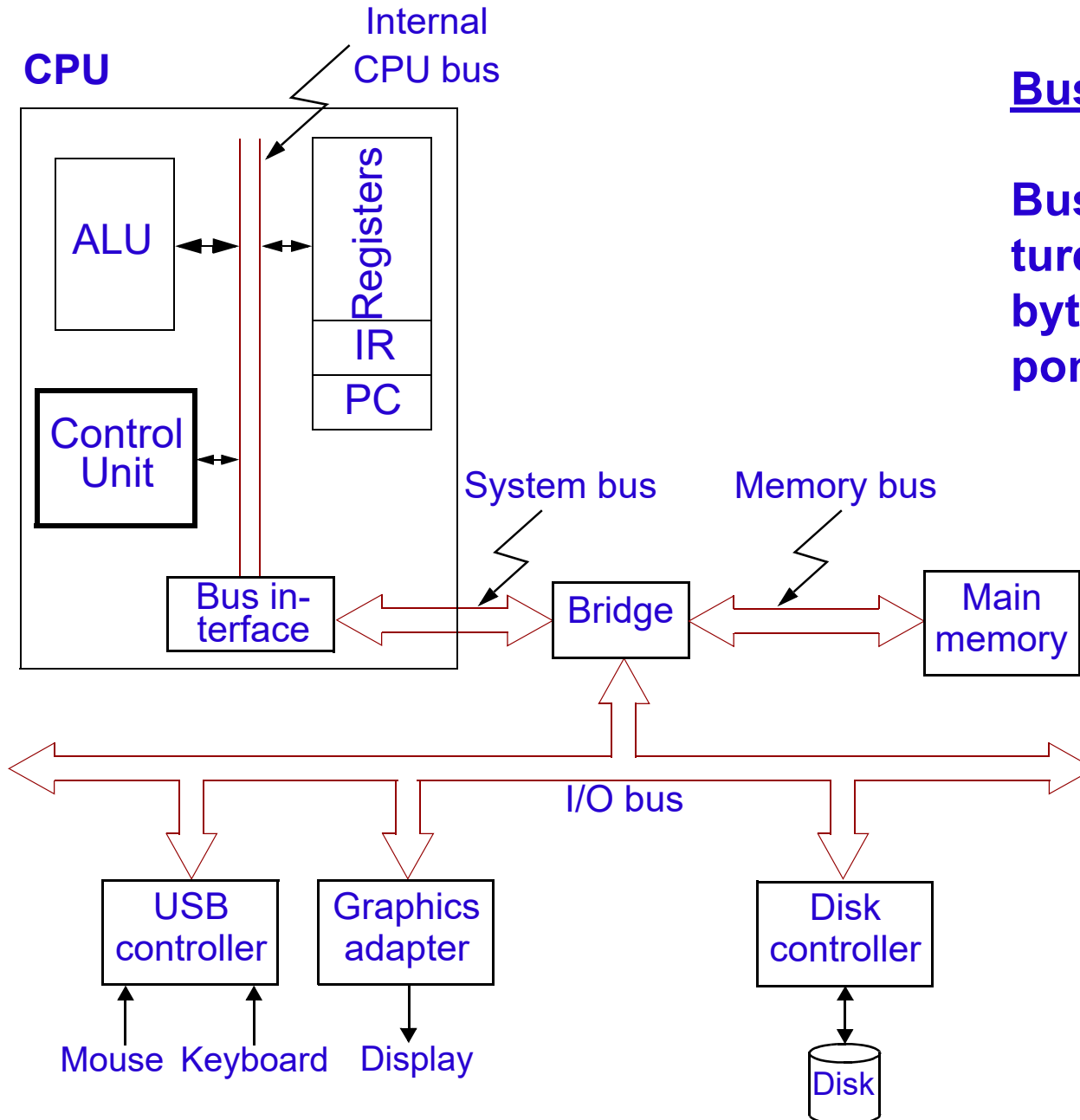
The CPU together with the Main Memory build the core computer; this is the minimal structure capable of storing and executing programs.

The rest of the computer system deals with communication, Input/Output, long term storage, and interaction with the environment.

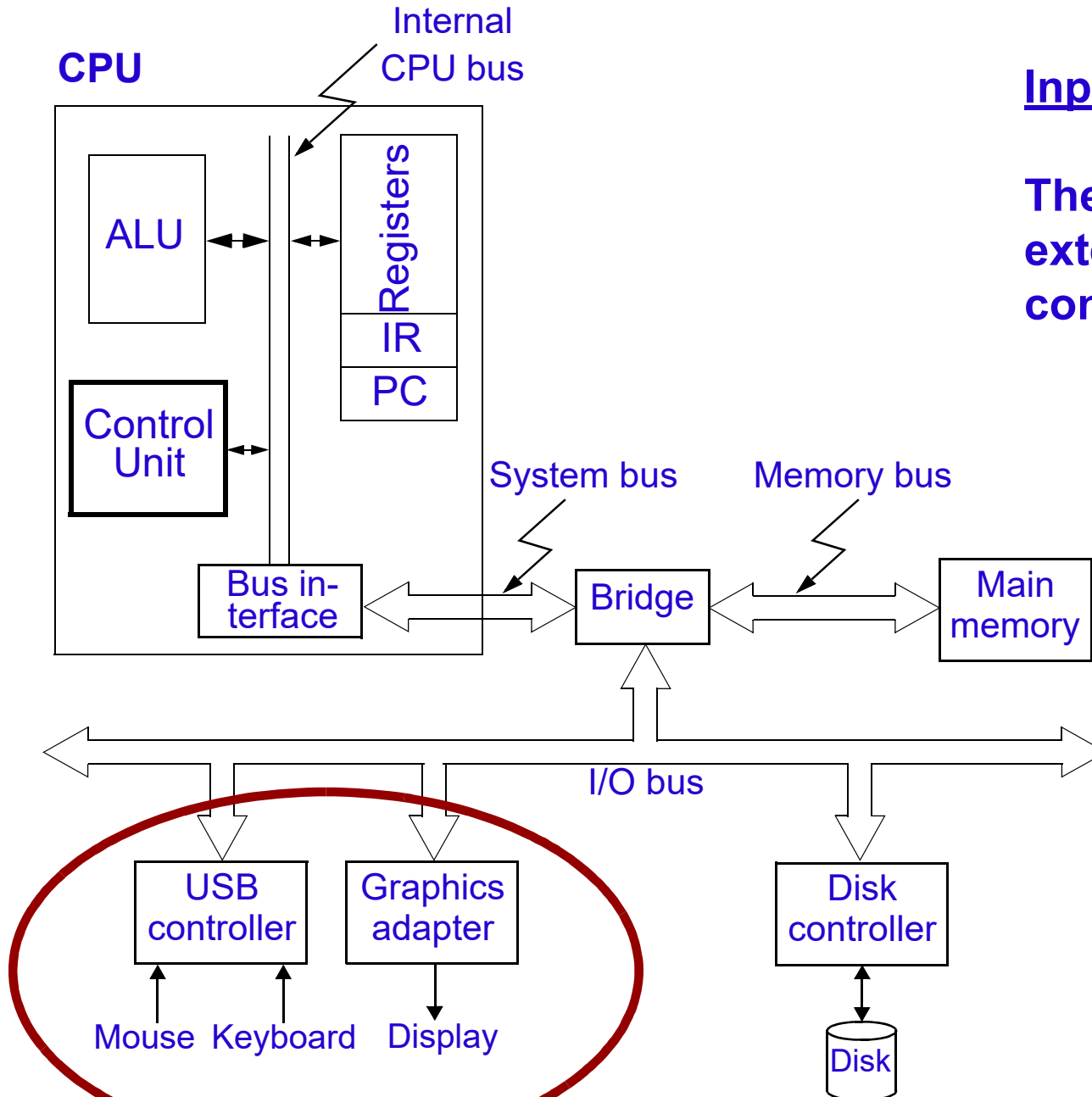
# Computer Systems

## Buses

**Buses are the physical infrastructure (electrical wiring) over which bytes are travelling between components of the computer system.**



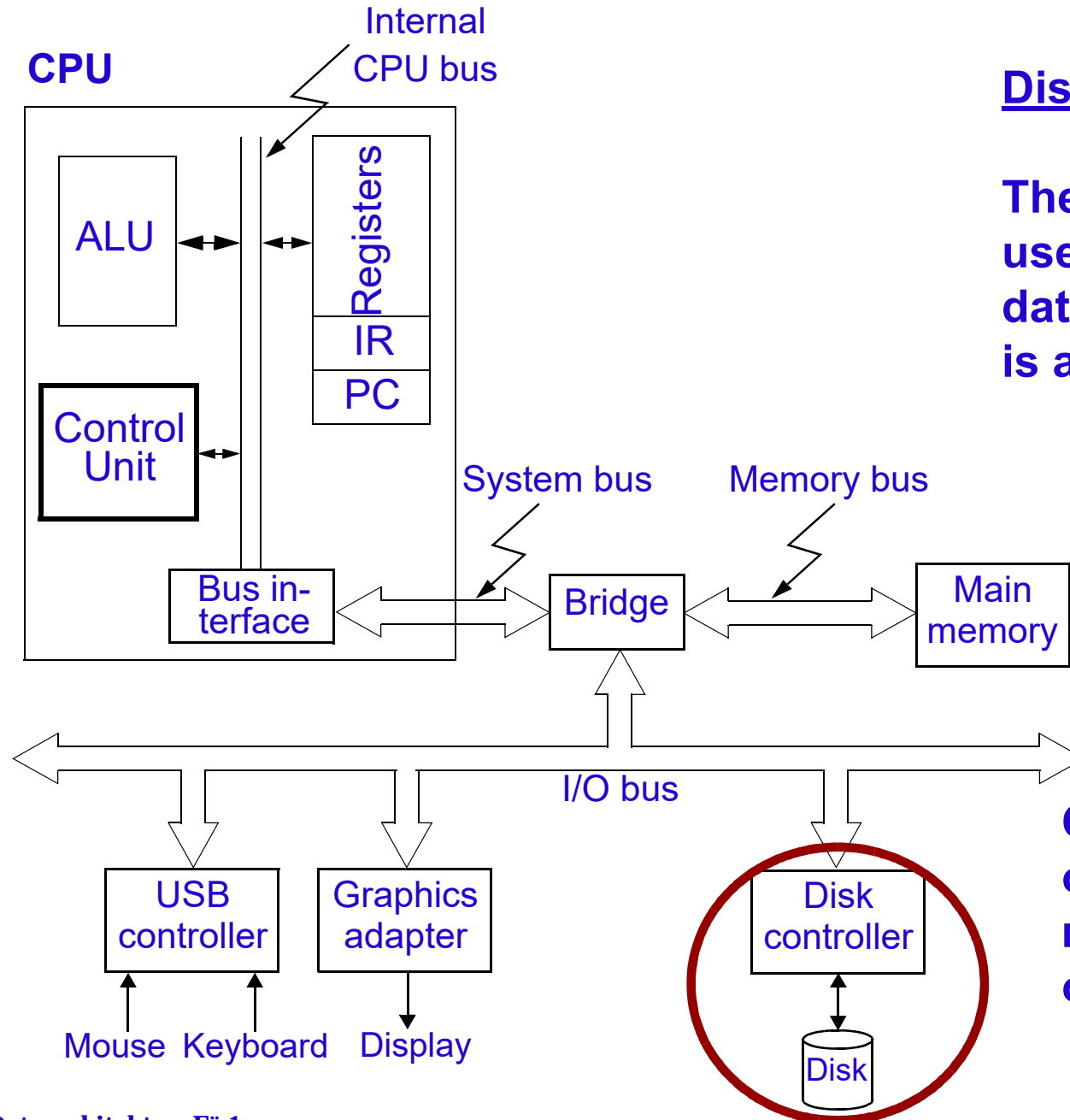
# Computer Systems



## Input/Output Devices

They connect the computer to the external world. Connection is via controllers/adaptors.

# Computer Systems



## Disk drive

The disk drive is a special device used as a long term storage for data and programs. Such a storage is also called *Secondary Memory*.

On modern computers the secondary memory is often implemented as *solid state disk (SSD)* on *flash memory*.

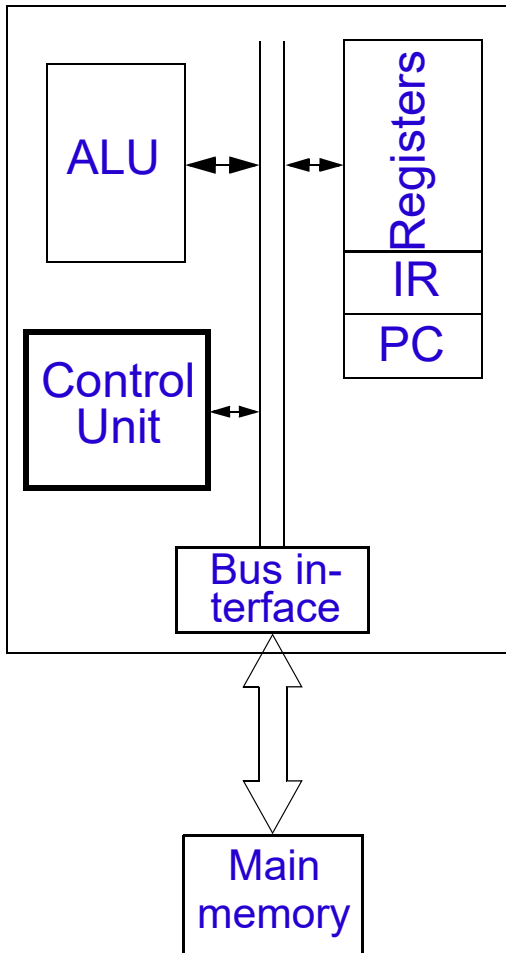


# How Does It Work? The von Neumann Architecture

- All computers in use, simple or complicated, big or small, cheap or expensive work according to the same basic concept, known as the von Neumann architecture:
  - Data and instructions are both stored in the main memory (stored program concept);
  - The content of the memory is addressable by location (without regard to what is stored in that location);
  - Instructions are executed sequentially (from one instruction to the next, in order of their location in memory) unless the order is explicitly modified.

# A Simple Computer Architecture

## CPU

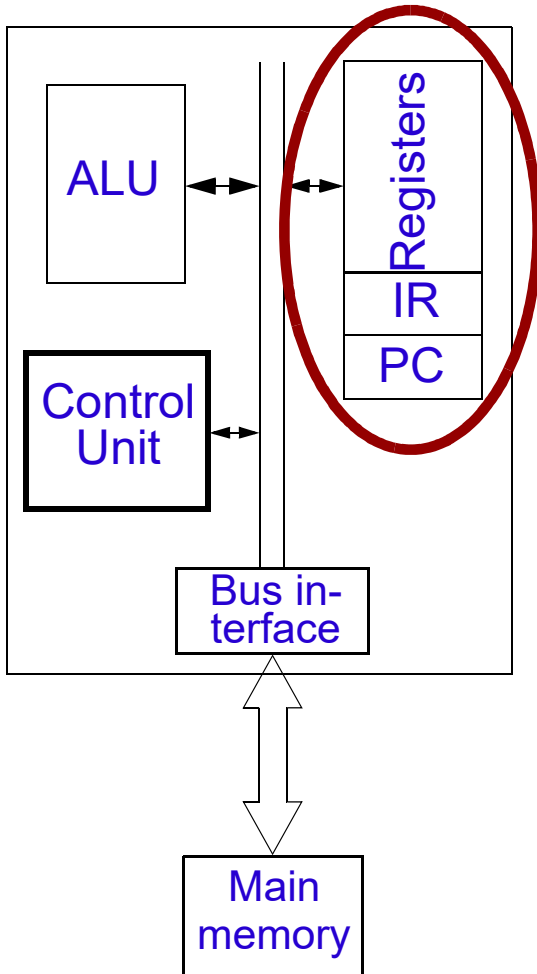


The basic von Neumann organization (architecture):

- **Central processing unit (CPU) contains:**
  - **Control unit (CU)** that coordinates the execution of instructions;
  - **Arithmetic/logic unit (ALU)** that performs arithmetic and logic operations;
  - **A set of registers.**
- **Main memory.**

# A Simple Computer Architecture

## CPU

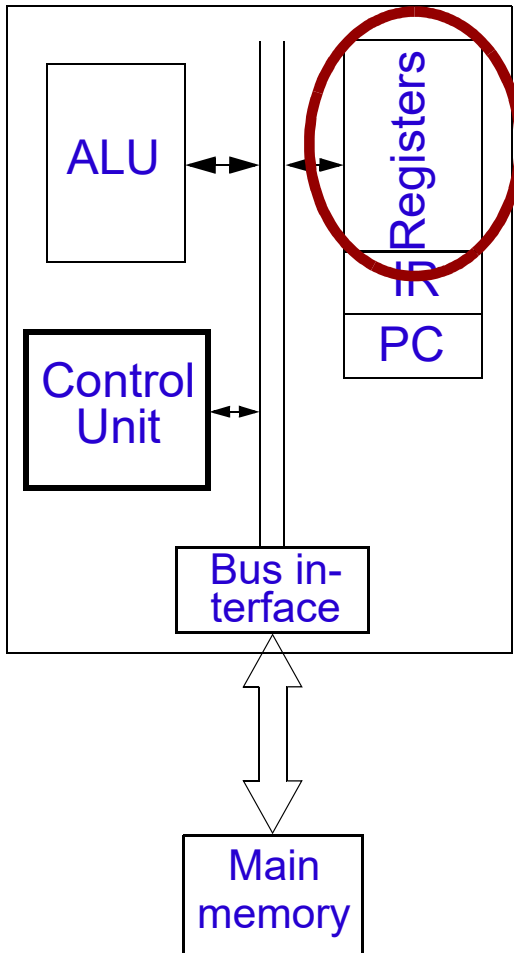


## Register Organization

- The set of registers within the CPU represents the top level of the memory hierarchy inside the computer system:
  - User visible registers: can be accessed by programs, for data storing.
  - Control and Status registers: used by the Control Unit to control the operation of the CPU; not directly accessible by the programmer.

# A Simple Computer Architecture

## CPU

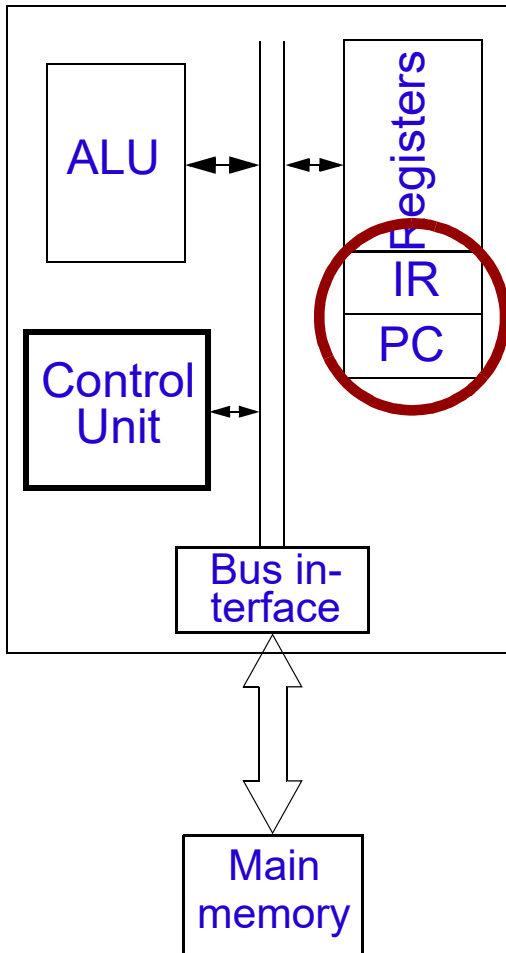


## User Visible Registers

- A set of registers which can be used without restrictions as operands for any operation and as address registers; these are so called *general-purpose registers*.

# A Simple Computer Architecture

## CPU

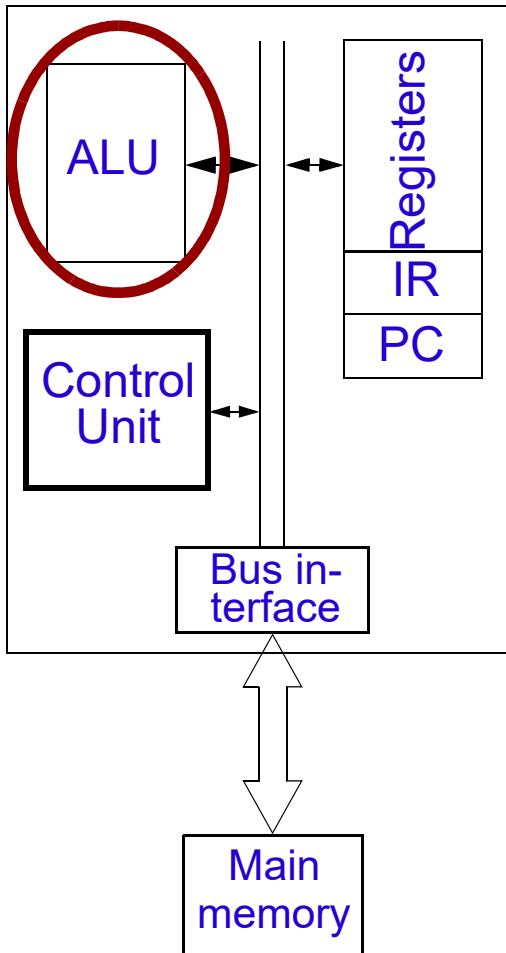


## Control and Status Registers

- ❑ Program Counter (PC): holds the address of the instruction to be fetched and executed.
- ❑ Instruction Register (IR): holds the last instruction fetched.
- ❑ Program Status Word (PSW): Condition Code Flags + other bits defining the status of the CPU.
- ❑ .....

# A Simple Computer Architecture

## CPU

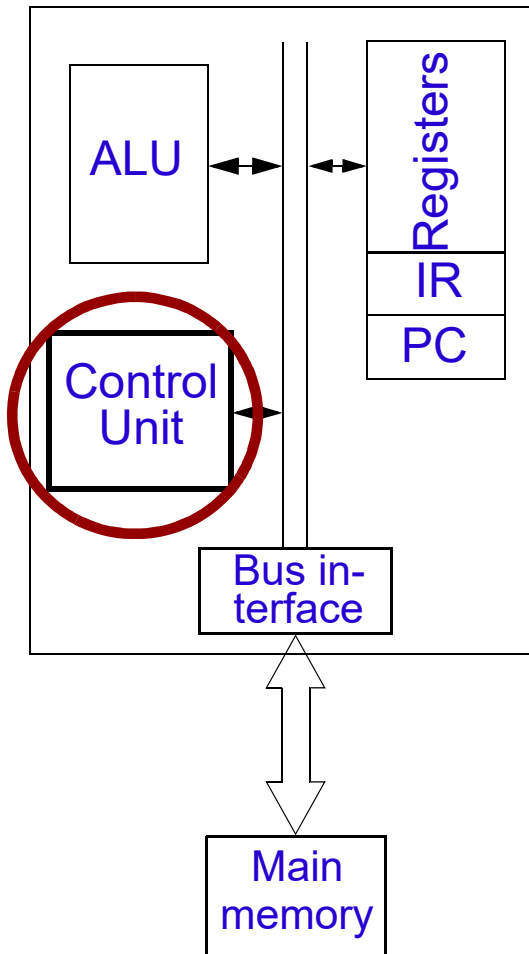


## Arithmetic Logic Unit (ALU)

- Performs arithmetic and logic operations. There might be several of them in a CPU. ALUs are different, depending on the data type they operate on: integer ALU, floating point ALU, etc.

# A Simple Computer Architecture

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## Control Unit

- The control unit generates the appropriate signals such that all other components of the CPU and the computer system, together, execute the current instruction.
- The current instruction to execute is stored in the instruction register (IR); it is the instruction whose memory address is stored in the program counter (PC)

# Representation of Data

- Inside a computer, data and control information (instructions) are all represented in binary format which uses only two basic symbols: "0" and "1".
- The two basic symbols are represented by electronics signals.

- Numeric data are represented using the binary system, in which the positional values are powers of 2:

$$100101 = 1 \cdot 2^0 + 0 \cdot 2^1 + 1 \cdot 2^2 + 0 \cdot 2^3 + 0 \cdot 2^4 + 1 \cdot 2^5$$

$$10110 = 0 \cdot 2^0 + 1 \cdot 2^1 + 1 \cdot 2^2 + 0 \cdot 2^3 + 1 \cdot 2^4$$

- Binary numbers are added, subtracted, multiplied and divided (by the ALU) directly; it is not needed to convert them to decimal numbers first.

$$100101 + 10110 = 111011$$

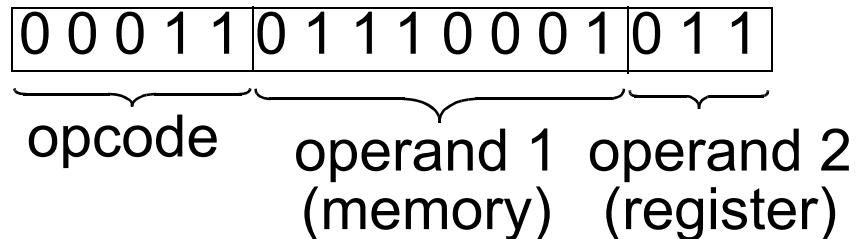


# Machine Instructions

- A CPU can only execute *machine instructions*,
- Each computer has a set of specific machine instructions which its CPU is able to recognize and execute.

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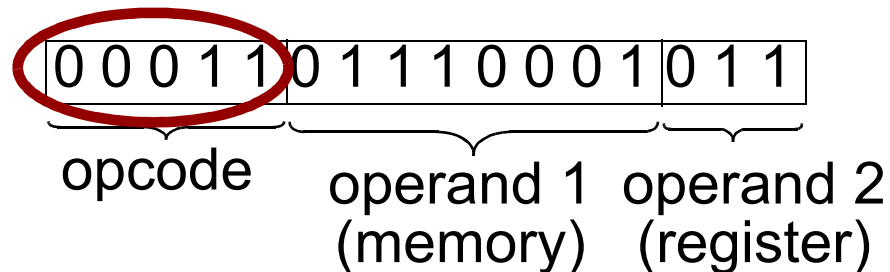
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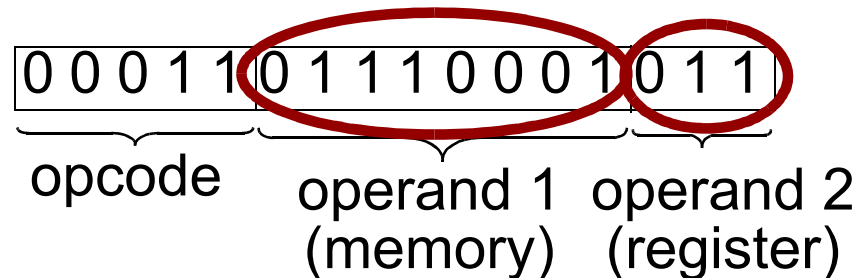
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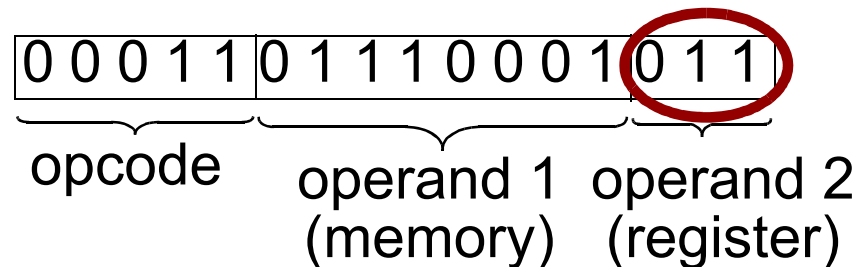
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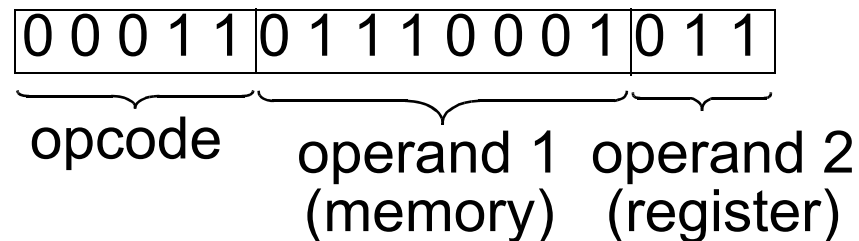
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- A machine instruction is represented as a sequence of bits (binary digits). These bits are organized into *fields* that define:
  - What has to be done (the operation code).
  - To whom the operation applies (source operands).
  - Where does the result go (destination operand); in this example CPU it is assumed that the result of the operation is stored in the same place where the second operand was stored; no additional field is needed.

# Machine Instructions

- A CPU can only execute *machine instructions*,
- Each computer has a set of specific machine instructions which its CPU is able to recognize and execute.



- Number of bits, number and length of the fields and their order is particular to each computer; this defines the *instruction format* of that computer.

# Types of Machine Instructions

- Machine instructions are of four types:
  - Data transfer between memory and CPU registers
  - Arithmetic and logic operations
  - Program control (test and branch); these are those instructions that change the flow of instruction execution by *jumping* to an instruction *different* from the instruction following the current one in memory.
  - I/O transfer

You see, there are very simple things a machine instruction does!  
But many machine instructions, together, perform the big thing!

# Instruction Execution

Let's imagine you write in a program the following instruction:

```
Z := (Y + X) * 3;
```

The instruction will be executed by the CPU as a sequence of four machine instructions!



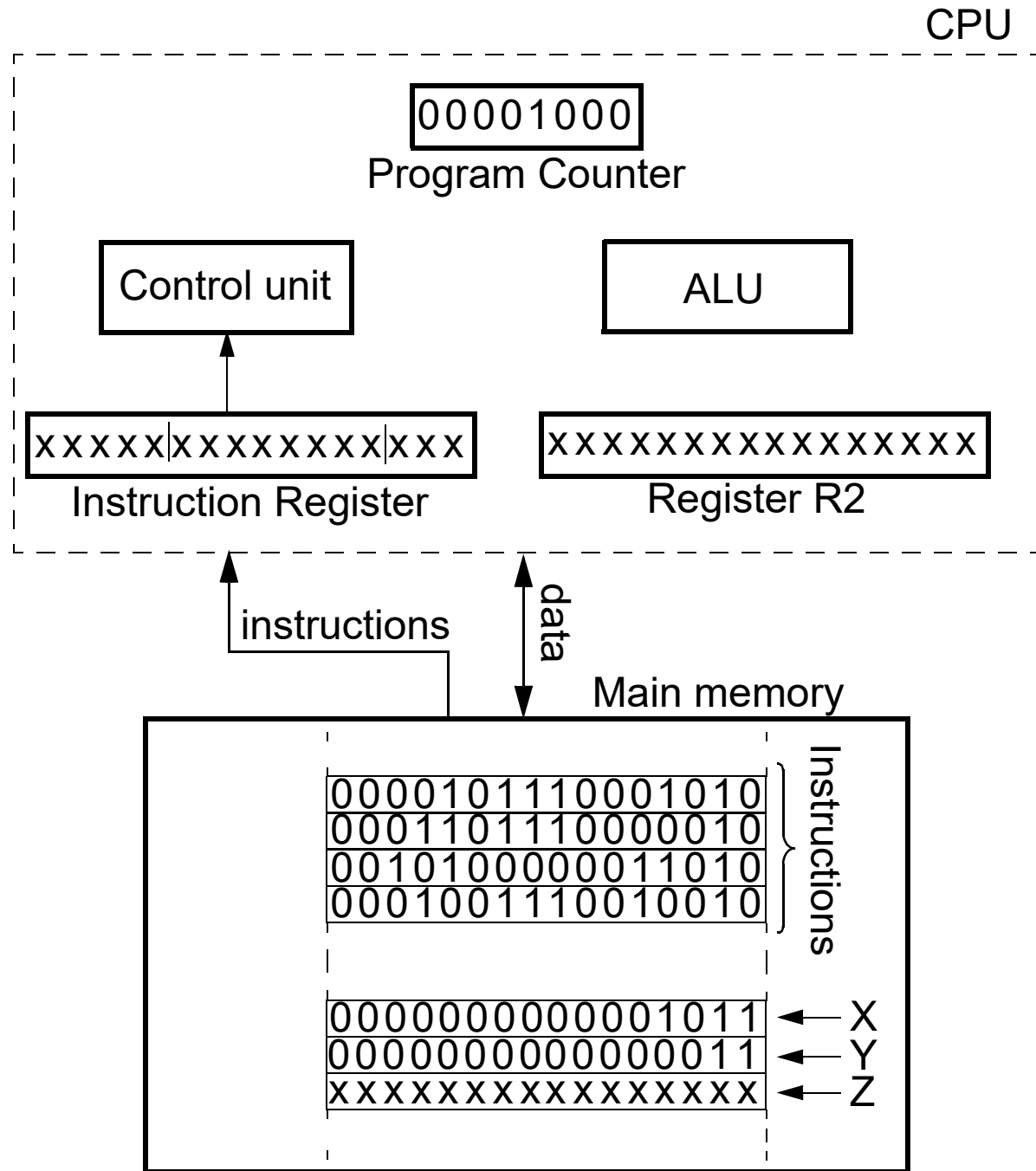
# Instruction Execution

Let's imagine you write in a program the following instruction:  $Z := (Y + X) * 3;$

	Memory address at which the instruction/data is stored	Content of the memory	
Move value of Y to Reg 2	00001000	<u>0000101110001010</u> Move addr of Y Reg 2	Instructions
Add value of X to Reg 2 (result kept in Reg 2)	00001001	<u>0001101110000010</u> Add addr of X Reg 2	
Multiply Reg 2 with 3 (result kept in Reg 2)	00001010	<u>0010100000011010</u> Mul value "3" Reg 2	
Store Reg 2 at address of Z	00001011	<u>0001001110010010</u> Move addr of Z Reg 2	
.....		.....	
Value of X: 11	01110000	0000000000001011 ← X	Data
Value of Y: 3	01110001	0000000000000011 ← Y	
Final value of Z: 42	01110010	0000000000101010 ← Z	

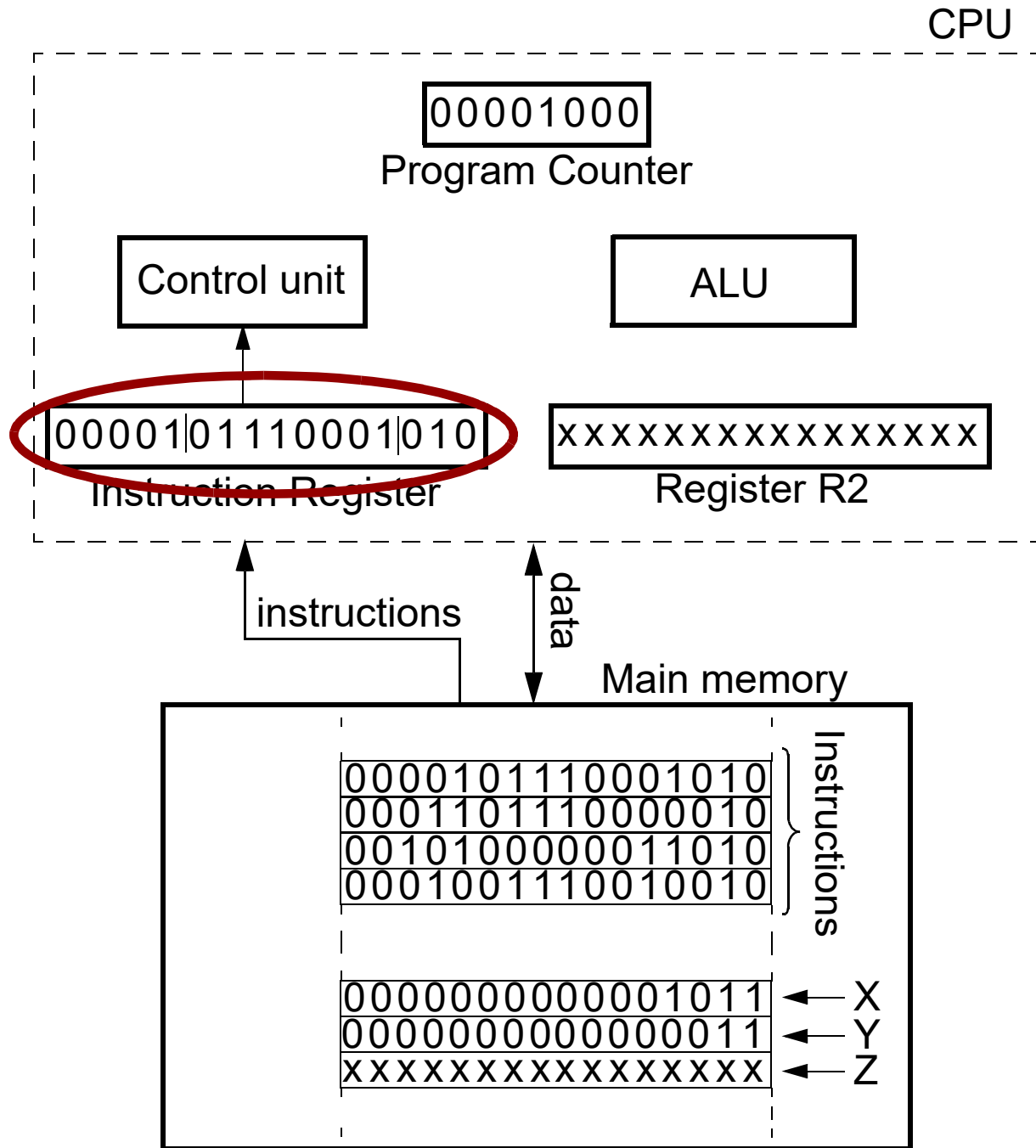
# Let's Follow the Instruction Execution

Before the first instruction 



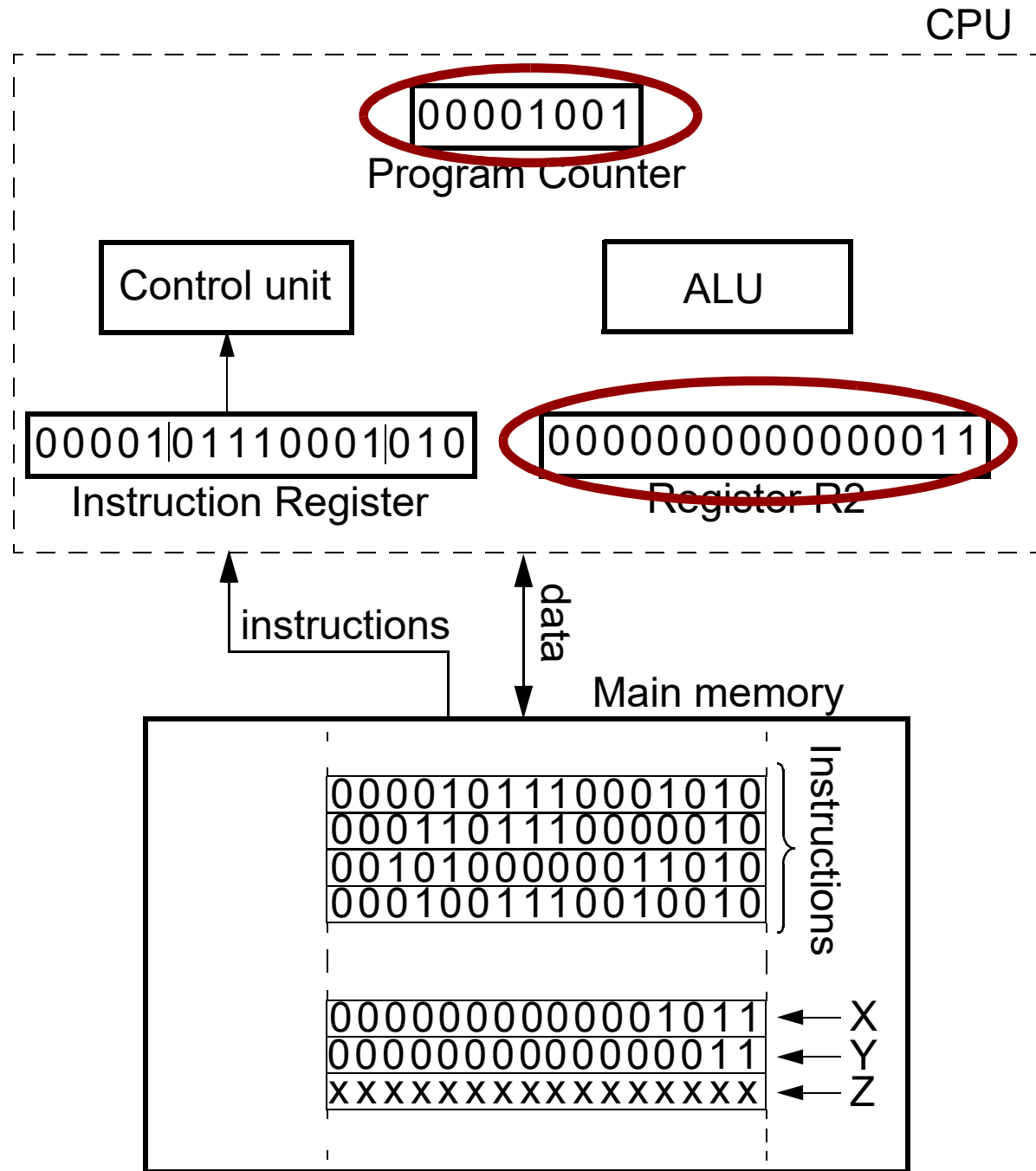
# Let's Follow the Instruction Execution

Now the first instruction is fetched



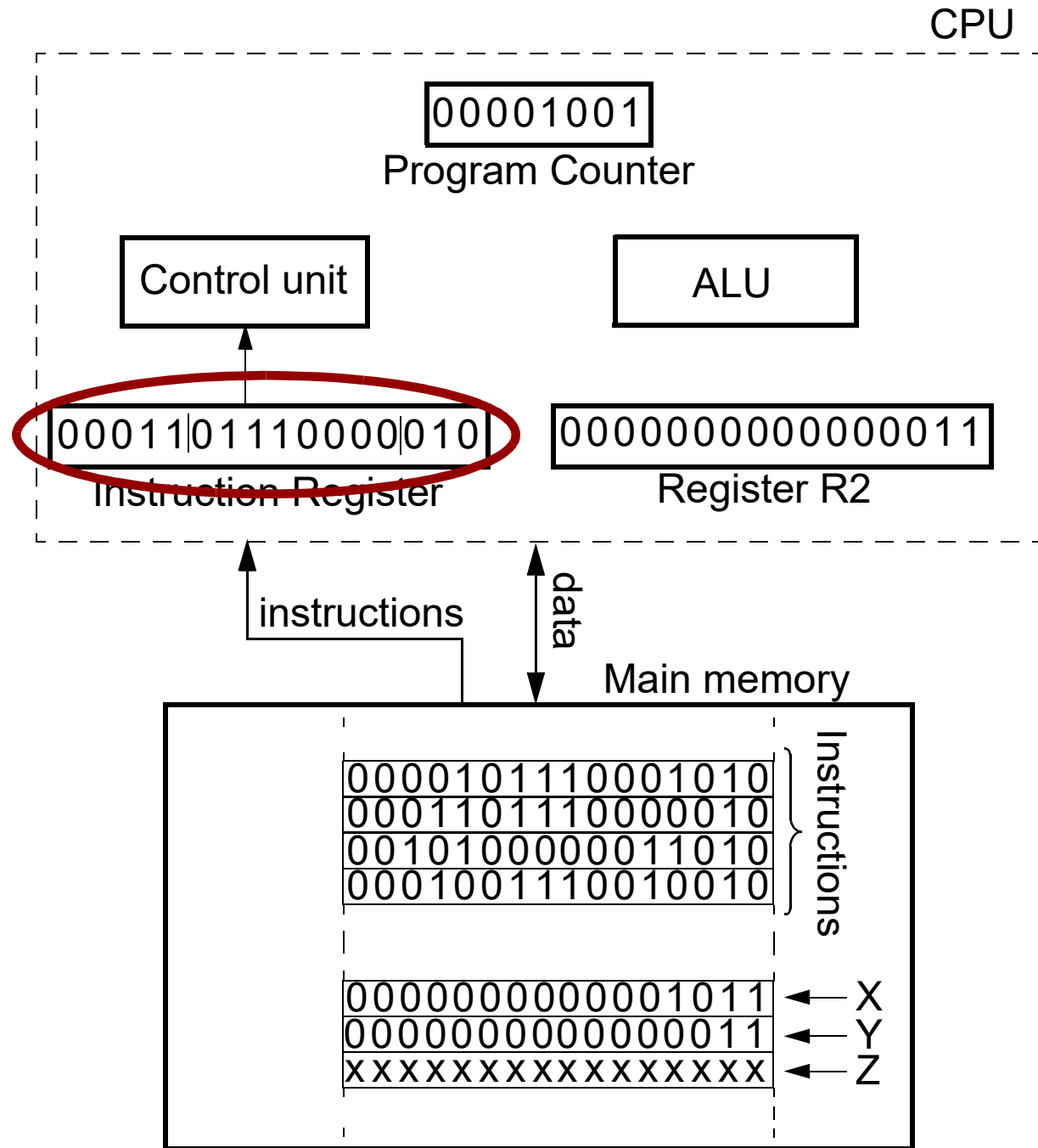
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After the first instruction 



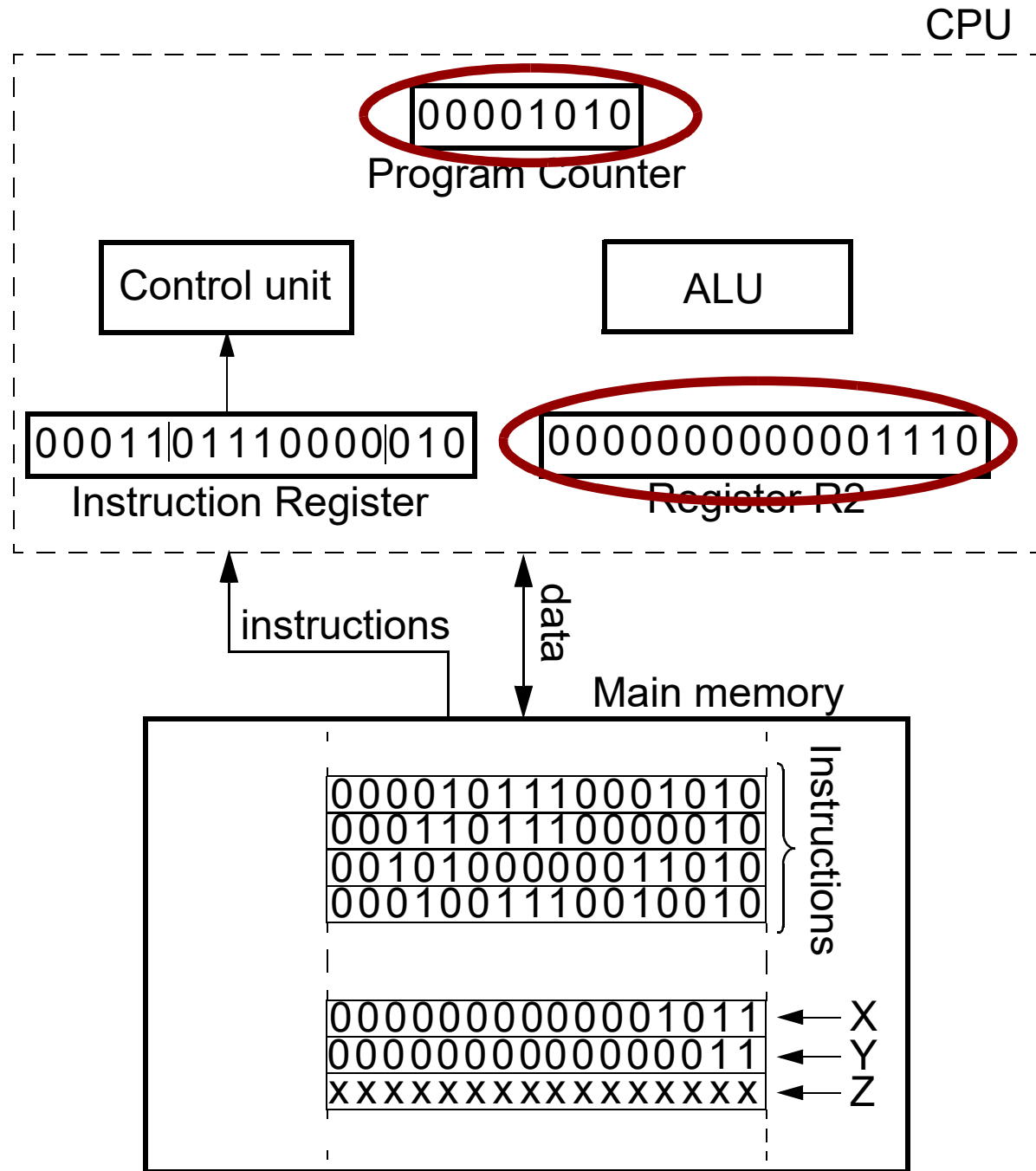
# Let's Follow the Instruction Execution

Now the second instruction is fetched



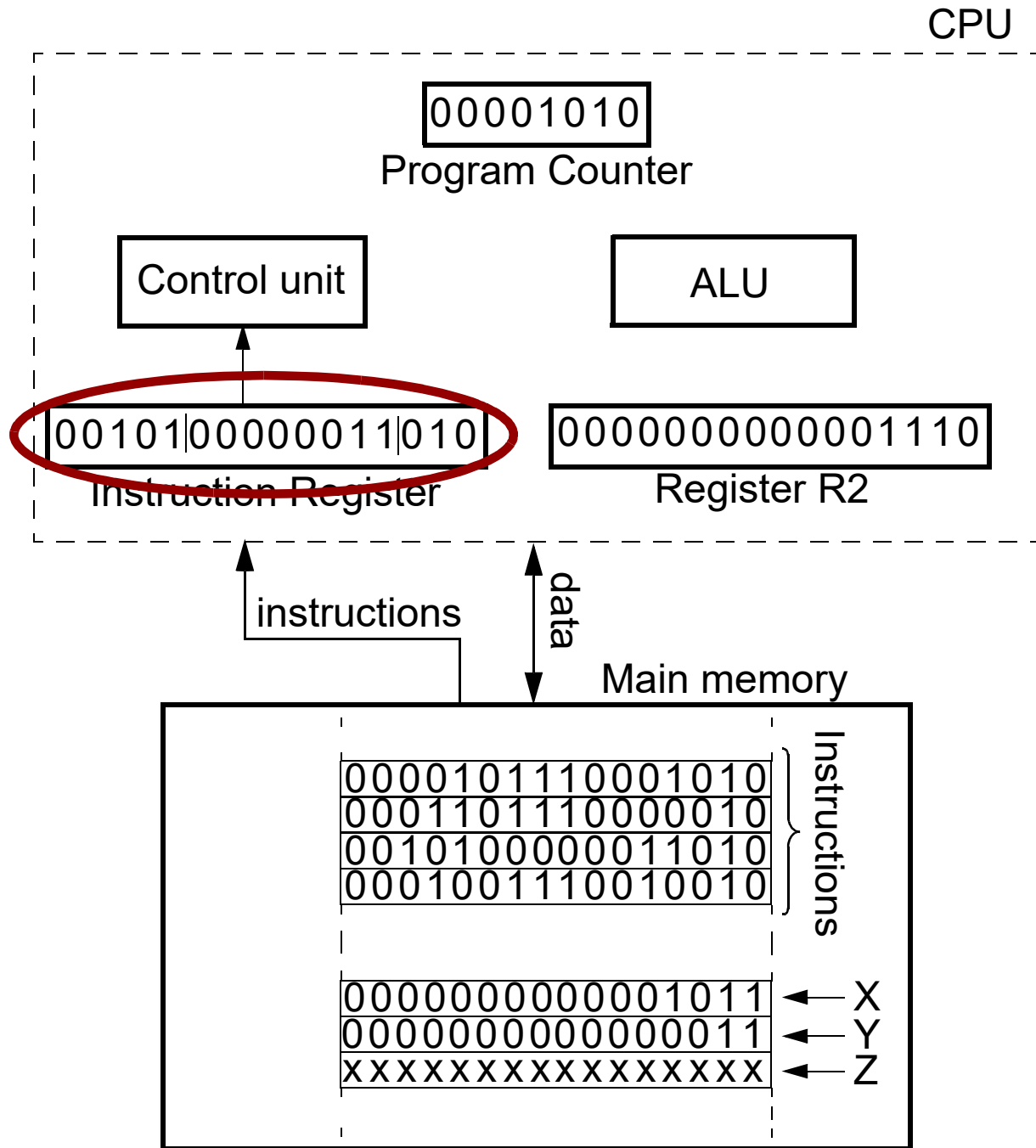
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After the second instruction 



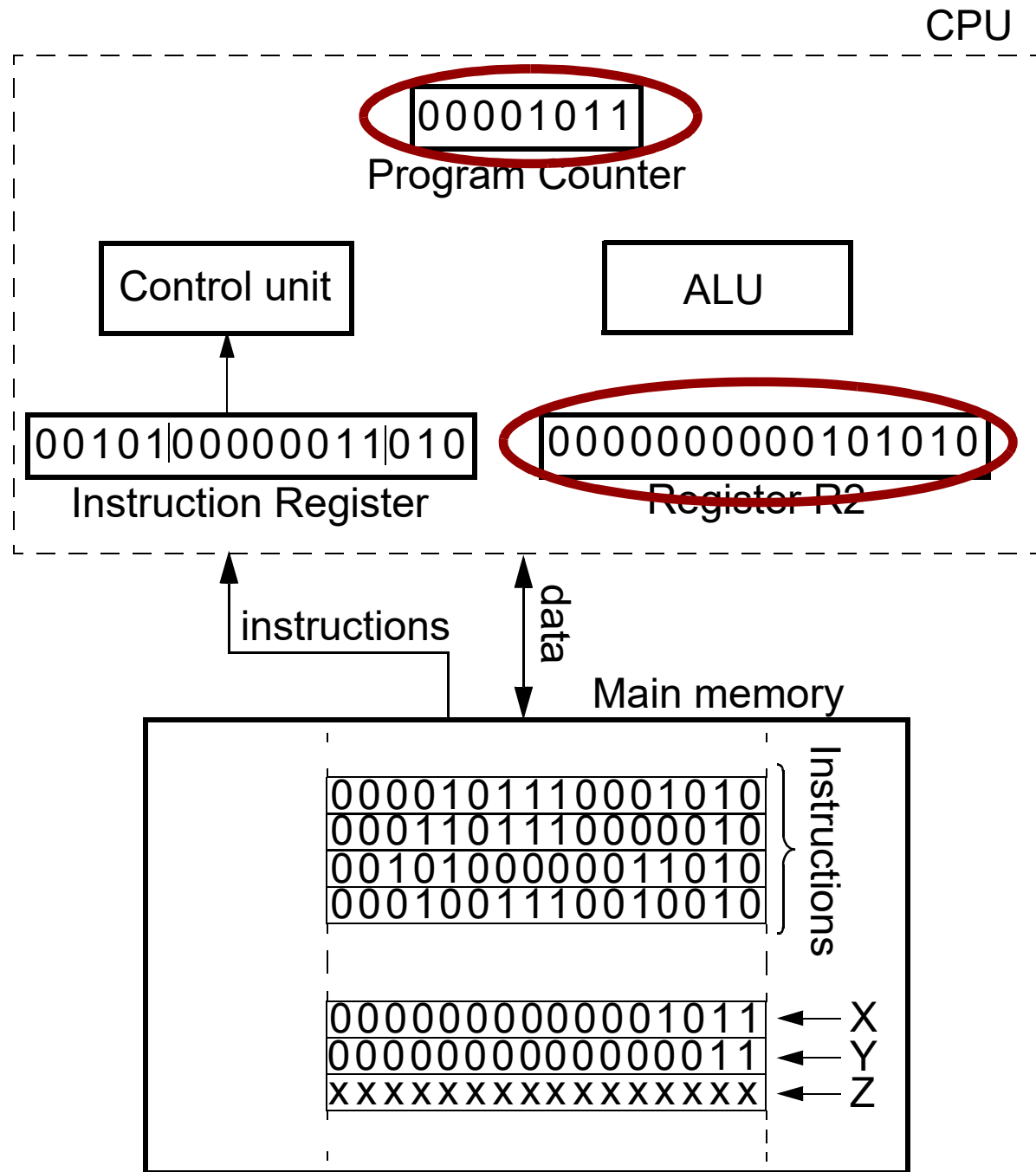
# Let's Follow the Instruction Execution

Now the third instruction is fetched



# Let's Follow the Instruction Execution

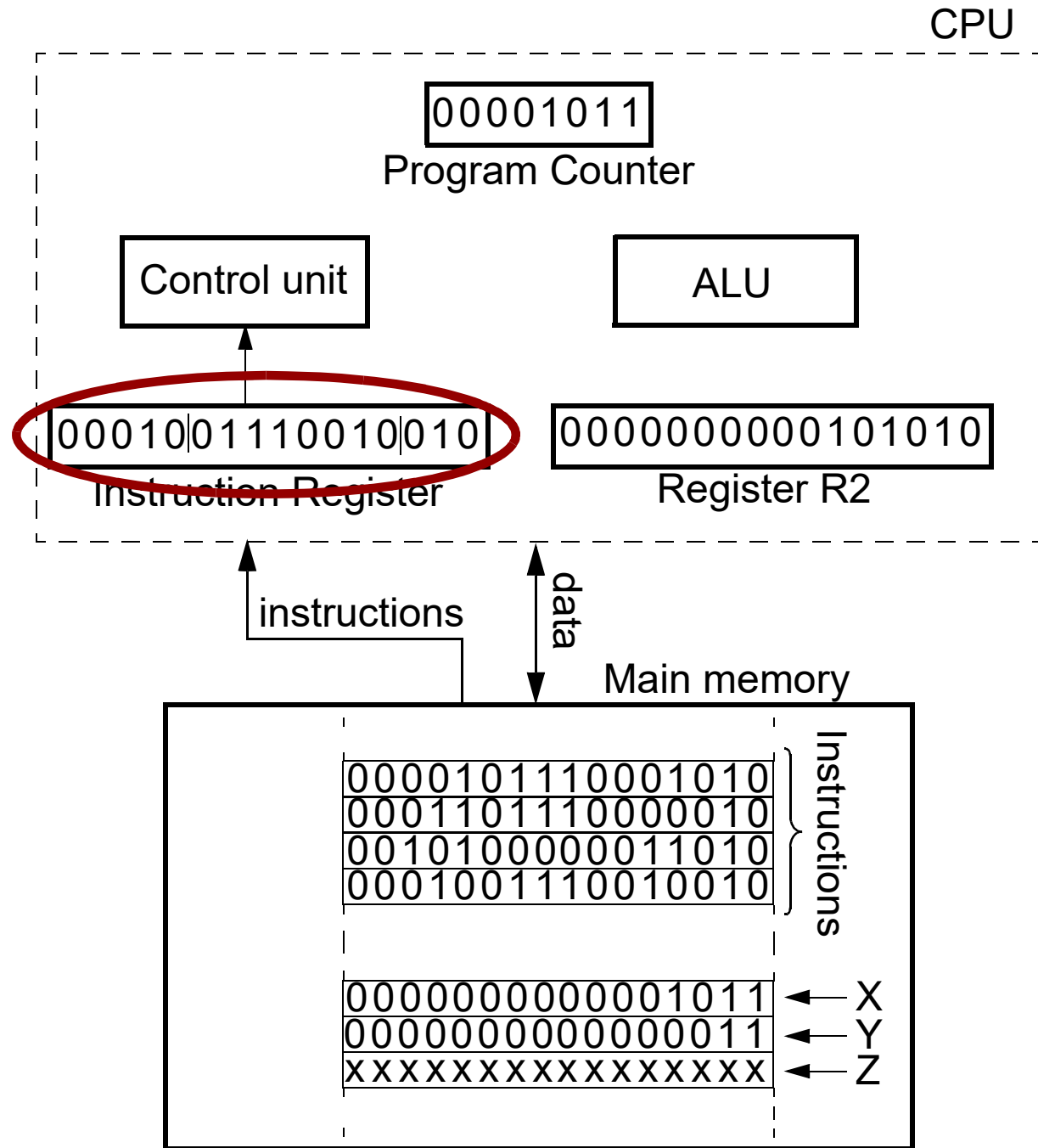
After the third instruction →





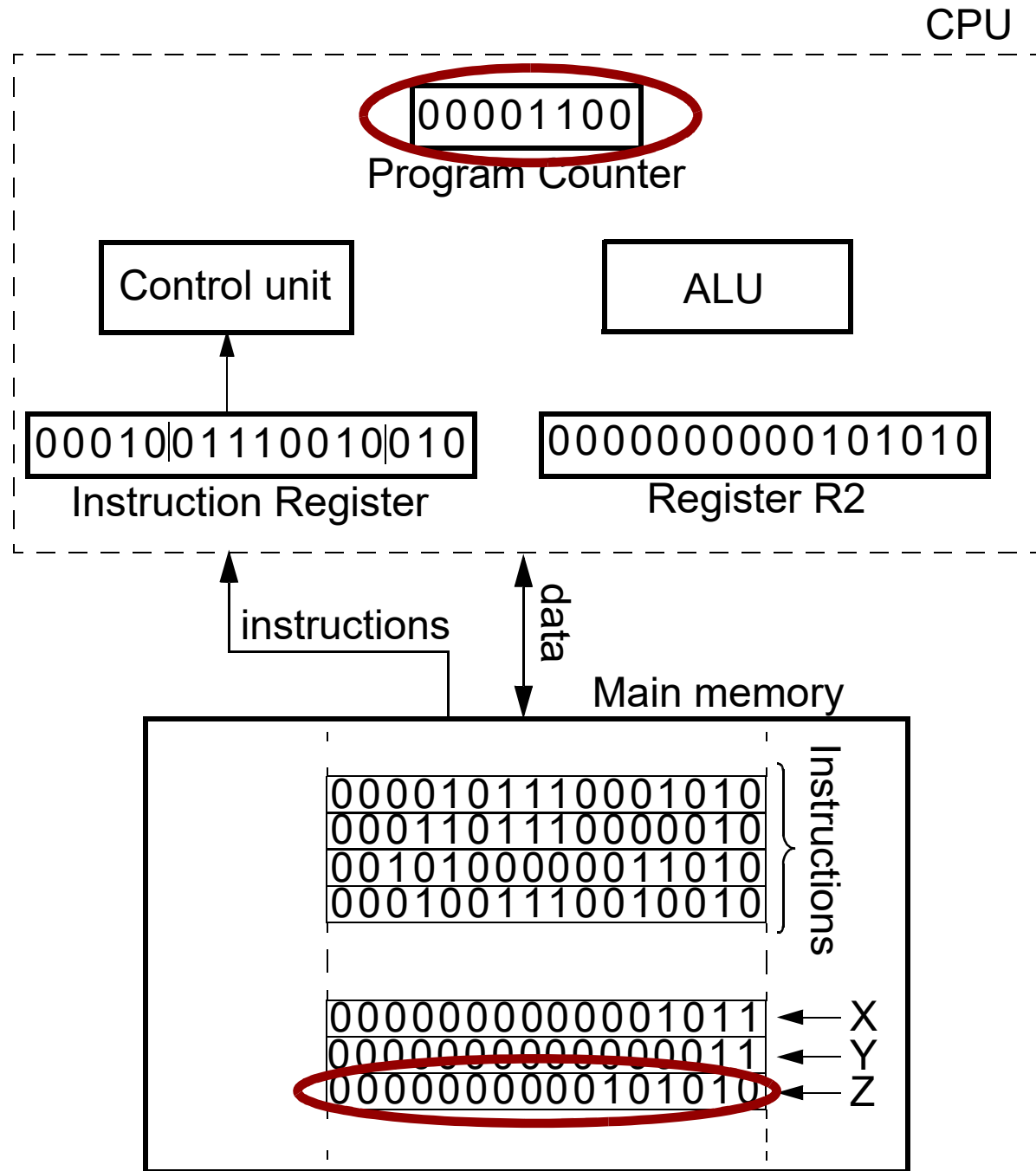
# Let's Follow the Instruction Execution

Now the fourth instruction is fetched



# Let's Follow the Instruction Execution

After the fourth and last instruction



# Compilers

We have written in our program:

`Z := (Y + X) * 3;`

High Level Language  
(e.g. C, C++, Java)

What the computer executes is:

0000101110001010
0001101110000010
0010100000011010
0001001110010010

Machine instructions  
for the particular processor that  
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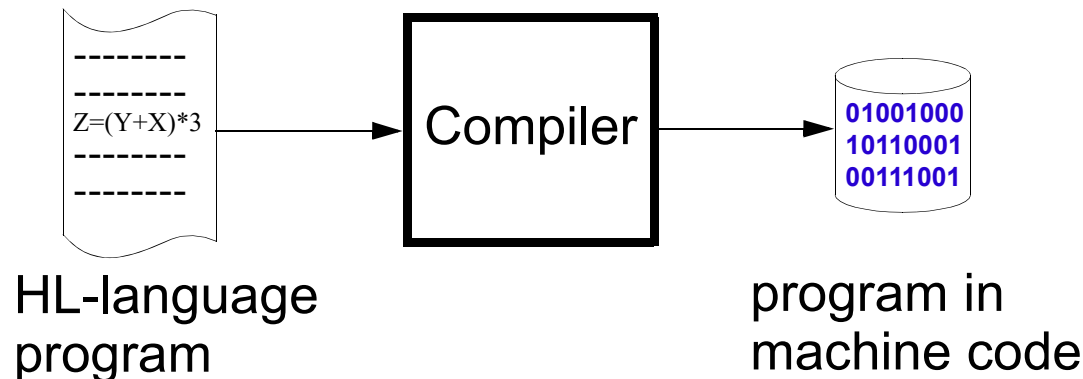
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What the computer executes is:

0	0	0	0	1	0	1	1	1	0	0	0	1	0	1	0
0	0	0	1	1	0	1	1	1	0	0	0	0	0	1	0
0	0	1	0	1	0	0	0	0	0	0	1	1	0	1	0
0	0	0	1	0	0	1	1	1	0	0	1	0	0	1	0

Machine instructions  
for the particular  
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Who brings us from our program to the machine instructions?



- A *compiler* is a program that translates programs written in a high level language into machine code to be executed on a certain processor.

# The Machine Cycle

Many things have to be done to execute a simple machine instruction:

- Fetch instruction
- Decode instruction
- Execute instruction

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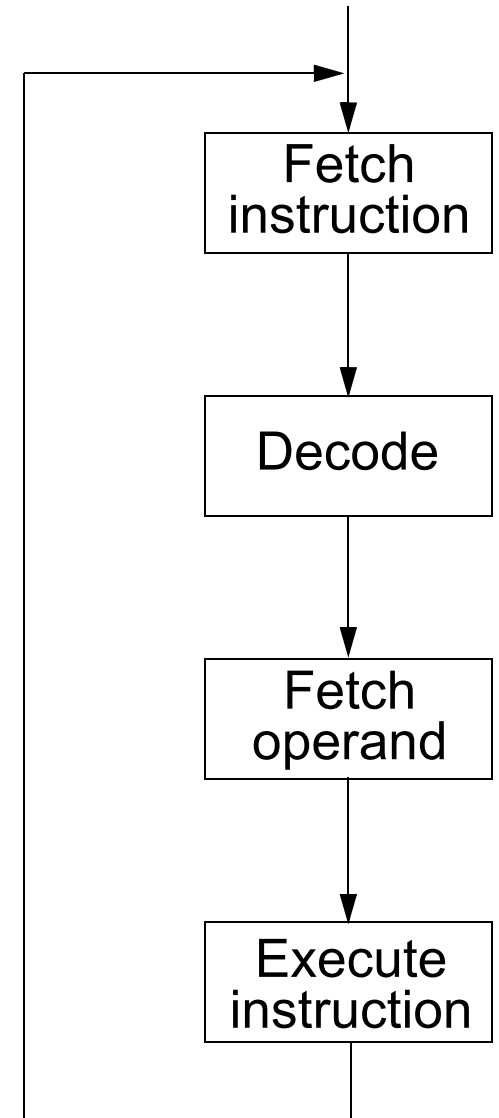
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- Fetch instruction
- Decode instruction
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## Machine Cycle

Each instruction is performed as a sequence of steps; the steps corresponding to the execution of one instruction are referred together as a *machine cycle*.

The number and nature of steps in the machine cycle differ from processor to processor.



# Case Studies

- **INTEL x86 Family**

- **The most successful example of a modern CISC (complex instruction set computer) microprocessor**

- **ARM Family**

- **The most successful RISC (reduced instruction set computer) microprocessor ever.**



# The INTEL x86

The number one microprocessor used in non-embedded systems.

First member:

- INTEL 4004 in 1971

First general purpose microprocessor:

- INTEL 8080 in 1974

# Milestones: INTEL x86

- **8080 (1974)**
  - First general purpose microprocessor;
  - 8 bits;
  - Used in first personal computer: *Altair*.
- **8086 (1978)**
  - 16 bits
  - Something like an instruction cache
  - First one used in IBM PC
- **80286 (1982)**
  - Huge increase in addressable memory (16 MByte instead of 1 MByte)
- **80386 (1985)**
  - 32 bits
- **80486 (1989)**
  - Complex cache structure and pipelining
  - Math coprocessor

# Milestones: INTEL x86

- **Pentium (1993)**
  - Introduces superscalar technology
- **Pentium Pro (1995)**
  - Advanced superscalar techniques
  - Branch prediction and speculative execution
- **Pentium II (1997)**
  - Intel MMX technology (instruction set extension for multimedia)
- **Pentium III (1999)**
  - Additional floating point instructions
  - Support for 3D graphics software
- **Pentium 4 (2000)**
  - Further improvements on the line of Pentium III
- **Core (2006)**
  - Solo: single core
  - Duo: Dual core - two processors on a chip

# Milestones: INTEL x86

- **Core 2 (2006)**
  - 64 bits
  - Dual: two processors on a chip
- **Core 2 Quad (2007)**
  - Four processors on a chip
- **Core i7 (2009/2010), mobile version 2011**
  - Six processors, Hyperthreading
- **Core i9 (2017)**
  - 10-18 processors, Hyperthreading

**Backward compatibility: newer versions can run the programs running on older versions.**

# The ARM Family

ARM processors are widely used in embedded systems (games, phones, multimedia applications, various hand-held devices, consumer products, automotive, medical equipment, wireless etc.).

- High performance, low size/cost, low power consumption

ARM Cambridge, UK, are designing single and multiprocessor architectures and *licence* them to manufacturers.

# Milestones: The ARM Family

- **ARM1 and ARM2 (1985)**
  - 32 bit RISC processor
  - 3 stage pipeline
- **ARM3 (1989)**
  - cache memory
- **ARM6 (1992)**
  - Floating point unit
- **ARM7 (1994)**
  - Most successful family.
  - First used as part of complete Systems on Chip (SoC)
- **ARM8 (1996)**
  - 5 stage pipeline
- **ARM9 (1997)**
  - separate data/instruction cache

# Milestones: The ARM Family

- **StrongArm (1996)**
  - Special version of ARM9; developed with DEC and, later, bought by Intel.
- **ARM10 (2000)**
  - 6 stage pipeline
- **ARM11 (2002)**
  - 9 stage pipeline
  - Media processing extensions
- **XScale (2002)**
  - Successor to StrongArm, by Intel
  - 7/8 stage pipeline
  - Dynamic voltage & frequency management
- **Cortex (2005)**
  - 13 stage pipeline
  - superscalar

# Milestones: The ARM Family

- **ARM11 MPCore (2005)**
  - Multicore based on ARM11; up to 4 cores
- **Cortex-A9 MPCore (2008)**
  - Multicore based on out-of-order superscalar Cortex-A9 core; up to 4 cores
- **Cortex- A15 MPCore (2011, commercial 2012)**
  - Multicore based on an out-of-order superscalar Cortex-A15 core; 2 clusters, 4 cores each.
- **ARM Cortex-A57 (2014, commercial 2016)**
  - Out-of-order superscalar, 64-bit instruction set. Used in chips by *Qualcomm, Samsung, Nvidia.*



# What is our Topic in this Course?

*We are interested in some advanced issues, typical to modern microprocessors and computer systems. These advances are at the origin of high performance achieved with today's computers.*

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  - pipelining
  - RISC architectures
  - superscalar architectures
  - VLIW architectures

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  - superscalar architectures
  - VLIW architectures
  
- **System Architectures for parallel computing**
  - performance of parallel computers
  - parallel computer architectures
  - multicore and multithreaded processors
  - general purpose graphic processors