

# Försättsblad till skriftlig tentamen vid Linköpings Universitet

Datum för tentamen	2009-12-22		
Sal	U11		
Tid	14-18		
Kurskod	TDDC90		
Droylod	TEN1		
ΙΙΟΥΚΟϤ			
Kursnamn/benämning	Software Security		
<b>1</b>			
Institution	IDA		
Antal uppgifter som	10		
ingår i tentamen			
Antal sidor på tentamen (inkl.	6		
försättsbladet)			
Jour/Kursansvarig	Shanai Ardi/Nahid Shahmehri		
Telefon under skrivtid	282608		
Besöker salen ca kl.	15:00, 16:30		
Kursadministratör	Madeleine Häger		
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Tillåtna hjälpmedel	Tryckt ordlista		
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LiTH, Linköpings tekniska högskola IDA, Institutionen för datavetenskap Nahid Shahmehri

## Written exam

# **TDDC90 Software Security**

# 2009-12-22

#### Permissible aids

Dictionary (printed, NOT electronic)

#### **Teacher on duty**

Shanai Ardi, 013-282608

#### Instructions

The exam is divided into two parts with a total of ten questions. You should answer all questions in all parts. In order to get the highest grade you will need sufficient points in the second part.

You may answer in Swedish or English.

#### Grading

Your grade will depend on the total points you score on the exam. The following grading scale is preliminary and might be adjusted during grading.

Grade	3	4	5
Points required	18	24	30

Important	
In order to get the highest grade you must have scored at least six points in part 2.	

## Part one

#### **Question 1: Static analysis (2 points)**

An analysis can be said to be sound. Explain what sound means in this context.

#### **Question 2: Vulnerabilities (2 points)**

Explain what a race condition vulnerability is. Give an example of code containing a race condition vulnerability.

#### **Question 3: Common Criteria (2 points)**

What is the purpose of the Common Criteria, and how can a developer benefit from it?

#### **Question 4: CMM (2 points)**

What is the capability maturity model, and how does it relate to security?

#### **Question 5: Fuzz Testing (4 points)**

The following C code calculates the line equation (in the form y = ax + b) for the line that intersects two points p1 and p2. For some inputs, the program will crash.

```
struct line { int a, b; };
struct point { int x, y; };
struct line *create_line(struct point p1, struct point p2) {
    struct line *l = malloc(sizeof(struct line));
    l.a = (p2.y - p1.y) / (p2.x - p1.x);
    l.b = p1.y - l.a * p1.x;
    return l;
}
```

Is fuzz testing using random inputs for p1 and p2 likely to detect the problem? Motivate your answer. If your conclusion is that fuzz testing is unlikely to detect the problem, then discuss how this *kind* of problem could be overcome in fuzz testing.

#### **Question 6: Best Practices (4 points)**

Name and describe four security best practices that can be used in secure software development and explain how they can improve security.

### **Question 7: Secure Design Patterns (4 points)**

Explain one secure design pattern in detail. You may choose any pattern except privilege separation (also known as PrivSep).

#### **Question 8: Threat Modeling (4 points)**

Explain what abuse/misuse cases are used for. Clarify your answer through an example.

## Part two

In order to score well on these questions you will need to show that you understand not only the technical issue or concept at hand, but also its context and its interactions with its context (e.g. processes, methods, techniques, technology, people, risks, threats, and so on). We *think* that good answers to these questions will require at least one or two handwritten pages (more or less may be required depending on how you write).

#### **Question 9: Threat Modeling (6 points)**

Explain what attack trees are. Through an example show how they are created and how they can be used.

#### **Question 10: Vulnerabilities (6 points)**

The function shown on the last page of this exam (read\_ppm) parses a portable pixmap image file. A PPM file consists of a "magic number", followed by the image dimensions, color depth, and finally the image data. Each pixel of image data is either three bytes or six bytes long, depending on the color depth.

The read\_ppm function below works quite well for normal PPM files, but contains at least two vulnerabilities that can be exploited using files with carefully chosen contents.

For each vulnerability:

- Indicate the code that contains the vulnerability.
- Explain the input that could trigger the vulnerability (you do not need to explain how to exploit it).
- Propose corrections to the code that would eliminate the vulnerability.
- Name and explain any mitigation techniques in the compiler, libraries or operating system that could prevent the vulnerabilities from being exploited.

#### **Code for question 10**

{

}

```
struct image *read ppm(FILE *fp)
   int version;
   int rows, cols, maxval;
    int pixBytes=0, rowBytes=0, rasterBytes;
   uint8 t *p;
    struct image *img;
    /* Read the magic number from the file */
    if ((fscanf(fp, " P%d ", &version) < 1) || (version != 6)) {
        return NULL;
    }
    /* Read the image dimensions and color depth from the file */
    if (fscanf(fp, " %d %d %d ", &cols, &rows, &maxval) < 3) {</pre>
        return NULL;
    }
    /* Calculate some sizes */
    pixBytes = (maxval > 255) ? 6 : 3; // Bytes per pixel
                                     // Bytes per row
    rowBytes = pixBytes * cols;
    rasterBytes = rowBytes * rows;
                                      // Bytes for the whole image
    /* Allocate the image structure and initialize its fields */
    img = malloc(sizeof(*img));
    if (img == NULL) return NULL;
    img->rows = rows;
    img->cols = cols;
    img->depth = (maxval > 255) ? 2 : 1;
    img->raster = (void*)malloc(rasterBytes);
    /* Get a pointer to the first pixel in the raster data. */
    /* It is to this pointer that all image data will be written. */
    p = img->raster;
    /* Iterate over the rows in the file */
    while (rows--) {
        /* Iterate over the columns in the file */
        cols = img->cols;
        while (cols--) {
            /* Try to read a single pixel from the file */
            if (fread(p, pixBytes, 1, fp) < 1) {
                /* If the read fails, free memory and return */
                free(img->raster);
                free(img);
                return NULL;
            }
            /* Advance the pointer to the next location to which we
               should read a single pixel. */
            p += pixBytes;
        }
    }
    /* Return the image */
   return img;
```

#### Notes on the code for those not very familiar with C

**fscanf** reads from a file to a program variable. The second argument specifies the input format. For example, fscanf(fp, " P%d ", &version) reads zero or more whitespace characters followed by an uppercase "P", followed by an integer. The value of the integer is stored in the variable named *version*. The fscanf function returns the number of items it read successfully.

**malloc** allocates memory on the heap. The parameter to malloc specifies how much memory can be allocated. Memory allocated with malloc is returned to the heap using the **free** function. When malloc fails to allocate sufficient memory, it returns NULL.

The data type **image** has four fields, rows, cols, depth and raster. When initialized, rows contains the number of rows in the image, cols the number of columns, and depth is either 1 or 2, indicating how many bytes are used to represent a single color value. The raster field contains a pointer to a memory area on the heap that holds all the pixel values. Each pixel value is a sequence of three color values.

The **fread** function reads raw bytes from a file. In this function it is used to read the image data, one pixel (i.e. 3 or 6 bytes) at a time. It returns the number of items read (in this function, the number of pixels). Specifically, the function call fread(p, pixBytes, 1, fp) will read *one* item that is *pixBytes* long from the file pointed to by *fp*, and store that item in the memory location pointed to by the variable *p*.

The **sizeof** operator returns the size of something. For example, <code>sizeof(int)</code> will return the number of byte required to store an integer, and assuming *img* is a pointer to a <code>struct image, sizeof(\*img)</code> will return the number of bytes required to store an image structure.