



# TDDC17: Intro to Automated Planning

#### **Classical Planning**

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## **Introduction to Planning**

#### **One** way of defining planning:

Using <u>knowledge</u> about the world, including possible actions and their results, to <u>decide</u> what to do and when in order to achieve an <u>objective</u>, <u>before</u> you actually start doing it

#### You have done this before!



Using <u>knowledge</u> about the world, including possible actions and their results, to <u>decide</u> what to do and when in order to achieve an <u>objective</u>, <u>before</u> you actually start doing it



Are we done?

#### Domain-specific search guidance – too much (human) work!



Automated Planning: <u>General</u> heuristics

**Pattern Databases, Landmarks** FF,  $h^m$ , merge-and-shrink, ...  $\rightarrow$  More in **TDDD48**  Automated Planning: Entirely different search spaces

Partial Order Causal Link SAT planning, Planning Graphs, ... → More in **TDDD48** 

Need a <u>well-structured representation</u> for planners/heuristics to <u>analyse</u>!

#### **AI Planning: A Simplified View**

<u>Description</u> of specific objectives to achieve <u>Description</u> of the world + how we can <u>affect it with</u> actions

**General** planning algorithm, general heuristics

<u>Solution</u>: Plan with *actions* to perform, *ordering* constraints, Still **simplified**, because in reality:

There must be an *agent* that defines the *objectives* 

Multiple agents may plan, interact → negotiation, delegation of tasks, collaboration, competition, ...

Execution can *fail*  $\rightarrow$  feedback, monitoring, replanning, plan repair

And so on...

Execution System

Sensors

Actuators

## Modeling "Classical" Planning Problems

The basis for most extended forms of automated planning

### **Facts and States: Introduction**

Like before, we are interested in states of the world



#### Is the information above sufficient?

No – Need to **analyze** states, find **differences** compared to goal states, find **relevant** actions, ...!

#### **Facts and States: Introduction**

We need information about every state:



Efficient planning <u>depends</u> on <u>describing</u> states as collections of <u>facts</u>:

We are in a state where <u>there is dirt</u> in <u>both</u> rooms, and <u>the vacuum cleaner</u> is in the leftmost room

# Let's use a representation based on first order logic!

#### Example: Dock Worker Robots (DWR)





## **Objects 1: Object Types**







- We then specify sets of **actual objects** 
  - **robot**: { r1 }
  - location: { loc1, loc2 }
  - **crane**: { k1 }
  - **pile**: { p1, p2 }
  - container: { c1, c2, c3, pallet }



#### **Facts**



- Most planners use a <u>first-order representation</u>:
  - Every <u>fact</u> is represented as a logical <u>atom</u>: Predicate symbol + arguments
  - Properties of the world
    - raining it is raining [not in the standard DWR domain!]
  - Properties of single objects...
    - empty(crane) the crane is not holding anything
  - <u>Relations</u> between objects

r3

- attached(pile, location) the pile is in the given location
- Relations between >2 objects

pZ

can-move(robot, location, location)

 the robot can move between the two locations [Not in DWR]

Essential: Determine what is **relevant** for the **problem** and **objective**!

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## Facts / Predicates in DWR



#### - **Reference**: All predicates for DWR, and their intended meaning

"Fixed/Rigid"	adjacent	(loc1, loc2)	; can move from <i>loc1</i> directly to <i>loc2</i>
(don't	attached	(p, loc)	; pile <i>p</i> attached to <i>loc</i>
change)	belong	(k, loc)	; crane <i>k</i> belongs to <i>loc</i>
"Dynamic"	at	(r, loc)	; robot <i>r</i> is at <i>loc</i>
	occupied	(loc)	; there is a robot at <i>loc</i>
	loaded	(r, c)	; robot <i>r</i> is loaded with container <i>c</i>
	unloaded	(r)	; robot <i>r</i> is empty
(modified by actions)	holding	(k, c)	; crane <i>k</i> is holding container <i>c</i>
	empty	(k)	; crane <i>k</i> is not holding anything
	in	(c, p)	; container <i>c</i> is somewhere in pile <i>p</i>
	top	(c, p)	; container <i>c</i> is on top of pile <i>p</i>
	on	(c1, c2)	; container <i>c1</i> is on container <i>c2</i>

## States

### **States 1: State of the World**

A <u>state (of the world)</u> should specify exactly which facts (<u>ground atoms</u>) are true/false in the world at a given time

Ground = without variables



We can find all possible states!

Every assignment of true/false to the ground atoms is a distinct state

Number of states: 2<sup>number of ground atoms</sup> – enormous, but finite (for classical planning!)

## States 1b: How many ground atoms, states?

• If we have r robots, l locations, k cranes, p piles, c containers:

; *rl* ground atoms

; *l* ground atoms

; *rc* ground atoms

; *r* ground atoms

- adjacent (loc1, loc2) ; l<sup>2</sup> ground atoms attached (pile, loc) ; pl ground atoms belong (crane, loc) ; kl ground atoms
  - at(rob, loc)occupied(loc)loaded(rob, cont)unloaded(rob)

holding(crane, cont); kc ground atomsempty(crane); k ground atoms

- in(cont, pile); cp ground atomstop(cont, pile); cp ground atomson(cont1, cont2); c² ground atoms
- So:
  - (l + p + k + r + 1)l + (c + 1)r + (c + 1)k + (c + 2p)c ground atoms
  - $2^{(l+p+k+r+1)l+(c+1)r+(c+1)k+(c+2p)c}$  states

#### **States 2: Efficient Representation**

#### Efficient specification and storage for a single state:

- Specify which atoms are true
  - All other atoms have to be false what else would they be?
- A <u>state of the world</u> is specified as a <u>set</u> containing all <u>ground atoms</u> that [are, were, will be] true in the world
  - s<sub>0</sub> = { on(A,B), on(B,C), in(A,2), in(B,2), in(C,2), top(A), bot(C) }



## **States 3: Initial State**



- Initial states in classical planning:
  - We assume complete information about the **initial state**  $s_0$  (before any action)

Complete **relative to the model**: We must know everything about those predicates and objects we have specified... But not whether it's *raining*!

So we can still use a set of true atoms

```
{
attached(p1, loc1), in(c1, p1), on(c1, pallet), in(c3, p1), on(c3, c1), top(c3, p1),
     attached(p2, loc1), in(c2, p2), on(c2, pallet), top(c2, p2),
     belong(crane1, loc1), empty(crane1),
     at(r1, loc2), unloaded(r1), occupied(loc2),
     adjacent(loc1, loc2), adjacent(loc2, loc1),
  }
                                                                                     ΙοςΖ
                                                ß
                                                                p2
                                                            IOC
```

#### **States 4: Goal States**

- Classical planning: Reach one of possibly many goal states
  - Can be specified as a <u>set of literals</u> that must hold

Literals = positive or negated atoms

- Example: Containers 1 and 3 should be in pile 2; container 12 should not be in pile 2
  - We don't care about their order, or any other fact





# Actions, Operators

#### **Actions 1: Intro**



- Actions in **plain search** (lectures 2-3):
  - Assumed a transition / successor function

<u>**Result**(State, Action</u>) - A description of what each action does (Transition function)



But how to <u>specify</u> it <u>succinctly</u>?

#### **Actions 2: Operators**



#### Define <u>operators</u> or <u>action schemas</u>:

- move(robot, location1, location2)
  - Precondition *tests facts*, depending on *parameters*:

at(robot, location1) ∧
adjacent(location1, location2) ∧
¬ occupied(location2)

The action is <u>applicable</u> in a state s if its precond is true in s

- The result of applying the action in state s:
  - s {negated effect facts}
  - + {positive effect facts}



Effects add or remove facts:

¬at(robot, location1),
at(robot, location2),
¬occupied(location1),
occupied(location2)

#### **Actions 3: Instances**

- The planner <u>instantiates</u> the schemas
  - Applies them to objects of the correct type
  - Example: move(r1, loc1, loc2)
    - Precondition: at(r1, loc1) ∧
       adjacent(loc1, loc2) ∧
       ¬ occupied(loc2)
    - Effects: ¬at(r1, loc1), at(r1, loc2), ¬occupied(loc1), occupied(loc2)





## Actions 4: Step by Step

Z6

In <u>classical</u> planning (the basic, limited form):



We know how states are changed by actions

→ **Deterministic**, can completely predict the state of the world after a sequence of actions!

The <u>solution</u> to the problem will be a <u>sequence</u> of actions

## Planning Domain, Problem Instance





#### Split knowledge into two parts

#### **Planning Domain**

- General properties
  - There are containers, cranes, ...
  - Each object has a location
  - Possible actions:
     Pick up container, put down container, drive to location, ...

#### **Problem Instance**

- Specific problem to solve
  - Which containers and cranes exist?
  - Where is everything?
  - Where should everything be? (More general: What should we achieve?)

## The State Space

#### **State Spaces 1: Introduction**

- Every classical planning problem has a <u>state space</u> a <u>graph</u>
  - A <u>node</u> for every world state
  - An edge for every executable action

The planning problem: Find a path (not necessarily shortest)



Example solutions: SRS, RSLS, LRLRLSSSRLRS, ...

## State Spaces 2: Intuitions?

- Now that we have a general model of <u>facts</u>:
  - <u>Every</u> combination of <u>facts</u> is a state
    - { at(robot1,loc1), at(robot1, loc2) }
    - { adjacent(loc1, loc2) } [but not adjacent(loc2, loc1)!]

Facts are like "variables" that <u>can</u> independently be true or false!

- But our **intuitions** often identify states that **we** think are:
  - "Normal"
  - "Expected"
  - "Physically possible"
- Usually:
  - The <u>initial state</u> is "normal"
    - We never specify { at(robot1,loc1), at(robot1, loc2) }
  - Preconditions/effects ensure that we can only <u>reach</u> other "normal" states
  - Mainly need to care about "normal" states... so let's focus on those!

## State Spaces 3: ToH, Actions





## State Spaces 4: Larger Example

- Larger state space interesting symmetry
  - 7 disks
  - 2187 "possible" states
  - 6558 transitions, [state, action]  $\rightarrow$  state



### State Spaces 5: Blocks World





#### **Reference: 4 operators**

- **pickup**(x) – takes x from the table •
- **putdown**(x) puts x on the table
- **<u>unstack</u>**(x, y) takes x from on top of ?y
- stack(x, y) puts x on top of y

#### **Reference: 5 predicates**

- <u>on</u>(x, y) block x is on block y
- **<u>ontable</u>**(x) x is on the table
  - <u>clear</u>(x) - we can place a block on x

- **<u>holding</u>**(x) the robot is holding block x
- **<u>handempty</u>** the robot arm is free

#### State Spaces 6: Blocks World, 3 blocks



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34

#### State Spaces 7: Blocks World, 4 blocks



#### 125 "possible" states 272 transitions



#### State Spaces 8: Blocks World, 5 blocks



#### This is tiny!
## State Spaces 9: Reachable States



Blocks	States reachable from "all on table"	Transitions (edges) in this part of the space
0	1	0
1	2	2
2	5	8
3	22	42
4	125	272
5	866	2090
6	7057	18552
7	65990	186578
8	695417	2094752
9	8145730	25951122
10	•••	
30	>197987'401295'571718'915006'598239'796 851	

Plan Generation Method 1: Forward State Space Search

## **Forward Search 1**



- Straight-forward planning: Forward search in the state space
  - Start in the initial state
  - Apply a search algorithm
    - Depth first
    - Breadth first



### Forward Search 2: Don't Precompute



The planner is <u>not</u> given a complete precomputed search graph!



Usually too large! → Generate as we go, hope we don't actually need the *entire* graph

### Forward Search 3: Initial state

- 41 Print Pri
- The <u>user</u> (robot?) <u>observes</u> the current state of the world
  - The initial state of the planning problem



- Must <u>describe</u> this using the specified <u>formal state syntax</u>...

  - ...and give it to the **planner**, which creates **one** search node
    - Here we show the path used to reach each node

[] → { clear(A), on(A,C), ontable(C), clear(B), ontable(B), clear(D), ontable(D), handempty }

### Forward Search 4: Successors

Given <u>any search node</u>...

[] → { clear(A), on(A,C), ontable(C), clear(B), ontable(B), clear(D), ontable(D), handempty }

- ...we can find <u>successors</u> by applying <u>actions</u>!
  - action pickup(D)
    - Precondition: ontable(D) ∧ clear(D) ∧ handempty
       Effects: ¬ontable(D) ∧ ¬clear(D) ∧ ¬handempty ∧ holding(D)
- This generates <u>new reachable nodes/states</u>...

...which can also be illustrated



## Forward Search 5: Step by step

- A <u>search strategy</u> (depth first, A\*, hill climbing, ...) will:
  - <u>Choose</u> a node
  - <u>Expand</u> the node, <u>generating</u> all possible successors
    - "What actions are applicable in the current state, and where will they take me?"
    - Generates new states by applying effects



This is illustrated – the planner works with sets of facts

The blocks world is symmetric: Can always "return the same way" Not true for all domains!



<u>Uninformed</u> or <u>Informed</u> Forward State Space Search?

### **Uninformed Search**

- Can we use **uninformed** search algorithms?
  - With only 30 blocks, we have >197987401295571718915006598239796851 reachable states

#### But what if we don't need to explore all the states?

- Suppose we need to tear down a 400-block tower and build it up on another "base"
- Suppose we want <u>good</u> plans
   Juse a shortest-path algorithm such as Dijkstra's / Uniform Cost Search
- Will explore <u>all</u> plans of lower length/cost than the optimal one

#### Plans to test: More than...

16305698390789310586457967937334728775645948416347826722586241976230426399420799766425821395576658116365413711816311922048822638316916164832 04594902834106357987452326989711329392844798003040966743549740387225888734809637192406427243636291547266329397641772360103156941486368193342 17252836414001487277618002966608761037018087769490614847887418744402606226134803936935233568418055950371185351837140548515949431309313875210 82788894333711361366092831808629961795389295372200673415893327657647047564060739170102603095904030354817422127405232957963777365872245254973 845940445258650369316934041843540738326378160253394039629713918091275405325705000143444004444755554344026724320356992923177737498303751007

1.63 \* 1

102711578461258322856646764107108548826574448445631879309077796615 465441372350568748665249021991849760646988031691394386551194171193 065768422967838517772535893398611212735245298803377536493561116410 374355414584408338787093441749839774374303275575344176291224488351 906611800376194410428900071013695438359094641682253856394743335678 891705393354709843502065977868949960690415707700579763228766976414

71360002096924483494302424649061451726645947585860104976845534507479605408903828320206131072217782156434204572434616042404375211052324038225 80540571315732915984635193126556273109603937188229504400



47294834609054590571101642

44130264943230562021556885

75080732255786307776859016

75511016725485476618861912

08793077271410935265343286



### **Informed Search**



#### We need guidance!

- For example, a <u>heuristic function</u> h(n) estimating the <u>cost</u> of reaching a goal node from node n
  - Sometimes, we define cost = number of actions
  - More general: each action has a cost c(a) longer plans may be cheaper!



## **Informed Search (2)**

- Previously we **manually designed** heuristics for a problem
  - 8-puzzle  $\rightarrow$  # pieces out of place, or sum of Manhattan distances
  - Romania Travel  $\rightarrow$  straight line distance

- Rimnicu Vilcea 📩 Timisoara 211 111 Pitesti 🗖 Lugoj 70 146 🗖 Mehadia Urziceni 86 75 138 Bucharest 120 Drobeta Craiova 🗖 Giurgiu
- Now: Want to define **general** heuristic functions
  - Without knowing what planning problem is going to be solved!







### **Informed Search: Perfect Information?**

- Given a **planning problem instance** and a **current state** s:
  - $\pi^*(s)$  denotes an **optimal solution** starting in s
  - $h^*(s) = cost(\pi^*(s))$  denotes the <u>cost</u> of an optimal solution

#### • $\rightarrow h^*$ would be the "perfect heuristic"

- Admissible cannot overestimate
- Informative perfect information

#### • Great, but can we **compute** $h^*(s)$ ?

- Theoretically, yes
- Practically, as difficult as finding an optimal plan in the first place!

#### We need <u>approximations</u>

Desirable properties depend on the type of planning

## Heuristics for **Optimal** Classical Planning

## **Optimal 1: Introduction**

- In <u>optimal</u> plan generation:
  - There is a <u>quality measure</u> for plans
    - Minimal number of actions
    - Minimal sum of action costs
    - ...
  - We <u>must</u> find an optimal plan!



 Suboptimal plans (0.5% more expensive):





### Optimal 2: A\*



- Optimal Plan Generation: Often uses A\*
  - A\* focuses entirely on optimality
    - Find a guaranteed optimal plan as quickly as possible
    - But no point in trying to find a "reasonable" plan *before* the optimal one
    - Slowly expand from the initial node, systematically checking possibilities
  - A\* requires <u>admissible</u> heuristics to guarantee optimality
    - Reason: Heuristic used for *pruning* (ignoring some search nodes)
    - Non-admissible  $\rightarrow$  can ignore some nodes that would lead to optimal plans

## **Optimal 3: Relaxation?**



#### • **<u>Relaxation</u>** can be used to generate admissible heuristics...





we just generated state s, want to compute h(s)

, Relax

Relaxed problem: P' (finding a solution: fast) But how can you find relaxations that work for <u>all</u> classical planning problems?

Find  $\pi^*(P')$  – <u>optimal</u> (!) plan for relaxed problem

Find cost

Solve

Compute h(s) = cost( $\pi^*(P')$ )

Since P is just 'some classical planning problem', we can't expect to find a general shortcut such as 'sum of Manhattan distances' Pattern Database Heuristics: One of <u>many</u> relaxation heuristics

## **PDB 1: Introduction**



- Main idea behind <u>relaxation</u> in <u>pattern database heuristics</u>:
  - Let's <u>ignore some facts</u> ground atoms everywhere

- Remove from *preconditions and goals*
  - Clearly makes the problem easier relaxation!
- Remove from current state and from action effects
  - No need to update these facts, when no conditions require them!



# **PDB 2: Dock Worker Robots**



- Example: Dock Worker Robots
  - Suppose we only consider container locations
    - in(container, pile), top(container, pile), on(c1,c2), ...
  - Ignore robot locations, crane locations, ...

<u>Abstract state</u> in P', information that remains after relaxation



# PDB 3: Planning in Patterns

- In P' we (pretend that we) <u>can</u> use the crane at p1 to:
  - pick up c3 (as we should)
  - place something on r1 (too far away, but that precondition disappeared...)
  - place five containers on one truck (condition "truck is free" disappeared)

A tile can move from square A to square B if A is horizontally or vertically adjacent to B and B is blank

A tile can move from square A to square B if A is adjacent to B
 A tile can move from square A to square B if B is blank
 A tile can move from square A to square B







# PDB 3b: Planning in Patterns

- In P' we still <u>can't</u>:
  - pick up c1 (preconditions about pile ordering are still there)
  - immediately place c1 below c2, ...
  - Still a planning problem P' left; need to find (cost of) optimal solution!



58

# PDB 4: Computing a Heuristic Value



#### Solve P' optimally: 4 actions

- Take c2 with the crane (it's in the way at the bottom of pile p2)
- Take c3 with the crane [relaxation not checking if the crane is busy]
- Place c3 at the bottom of pile p2
- Place c2 on top



# Let's formalize!

## PDB: Blocks World size 4



Consider physically achievable states in the blocks world, size 4:



## PDB: Blocks World size 4, facts

62 62

All ground atoms (facts) in this problem instance:

(on A A)	(on A B)	(on A C)	(on A D)
(on B A)	(on B B)	(on B C)	(on B D)
(on C A)	(on C B)	(on C C)	(on C D)
(on D A)	(on D B)	(on D C)	(on D D)

(ontable A) (ontable B) (ontable C) (ontable D)
(clear A) (clear B) (clear C) (clear D)
(holding A) (holding B) (holding C) (holding D)

(handempty)

## **PDB: Patterns, Abstract States**

- The pattern p is the set of ground facts we care about
  - A state s is represented by the **<u>abstract state</u>**  $s \cap p$
  - If  $s \cap p = s' \cap p$ , the two states are considered equivalent



A pattern generally contains few facts – for performance!

# **PDB: Ignoring Facts**



- Example: <u>only</u> consider 5 ground facts related to <u>block A</u>
  - Pattern": p={(on A B), (on A C), (on A D), (clear A), (ontable A)}



# **PDB: Transforming Actions**



- Pattern p={(on A B), (on A C), (on A D), (clear A), (ontable A)}
  - Example action: (unstack A B) an action (instance), not an operator!



• **Example action**: (unstack C D)

#### Before transformation:

:precondition (and (handempty) (clear C) (on C D)) :effect (and (not (handempty)) (holding C) (not (clear C)) (clear D) (not (on C D)))

After transformation:

:precondition (and) :effect (and) Loses **all** preconditions and effects → never used!

transform(*a*, *p*)

## **PDB: New State Space**



Pattern p={(on A B), (on A C), (on A D), (clear A), (ontable A)}



We lose information – and the size of the search space shrinks

## **PDB: State Transition Graph**



- New reachable state transition graph:
  - Real state s: Everything on the table, hand empty, all blocks clear
    - Abstract state: s0 = { (ontable A), (clear A) }



## **PDB: Databases**



- Where did the databases go?
  - During the main search, many visited *actual states* will correspond to the <u>same</u> *abstract states* need the same value over and over again
  - Given a pattern, we **precompute** a database for all abstract states
    - Improves performance; the principle remains



## **PDB: More information**



- To make PDB heuristics more informative:
  - Calculate costs for <u>several</u> patterns
    - Suppose we only care about {clear(A), ontable(A)}
    - Suppose we only care about {on(A,B), on(C,D)}
    - Suppose we...
  - Take the <u>maximum</u> of the computed heuristic values

All are below the limit (admissible) → the largest is below the limit

- One difficulty:
  - Choosing which patterns to use...

#### bw-tower07-astar-ipdb: 7 blocks, A\* search, based on a PDB variation





- Blind A\*, h(s)=0: 43150 states calculated, 33436 visited
- A\* using iPDB: 1321 states calculated, 375 visited

#### No heuristic is perfect – visiting some additional states is fine!

## **Satisficing Planning**

# **Satisficing Planning**

- Optimal plans are <u>nice</u> but <u>often hard to find</u>
  - Larger problem instances 
     → too much time, memory (even with good heuristics)!
- Satisficing plan generation:
  - Find a plan that is sufficiently **good**, sufficiently **quickly**

#### What's sufficient?

- Usually not well-defined!
- "These strategies and heuristics seem to give pretty good results for the instances I tested..."




## **Speed: Strategies**



- One reason for speed: Other informed search strategies
  - Simple: Greedy best first search
    - Always expand the node that seems to be closest to the goal
    - Who cares if getting there was expensive? At least I might find a way to the goal!
    - (Only care about h(s), not about g(s))

#### Hill climbing

- Be stubborn:
   If one direction seems promising,
   continue in this direction
- Many others!





## **Speed: Heuristics**



• One reason for speed: Often more informative heuristics

#### Optimal planning:

- Often requires <u>admissibility</u>: "We must never overestimate, ever!"
- Result: Usually underestimates (a lot!)

#### Satisficing planning:

Extreme example – greedy best first

- Only important that the "best" successor has a low heuristic value
- For GBF, what the value is doesn't matter!



- In many cases:
  - Admissibility is not required
  - Lack of admissibility is not "only slightly harmful"
  - Lack of admissibility is irrelevant

#### 7 blocks (tiny problem)



#### Greedy, IPDB: Expand 171 states, 22 actions in the solution



Larger problem instances **>** the difference <u>increases</u>

## Example: Landmark Heuristics

#### Landmark Heuristics (1)



#### Landmark:

"a **geographic feature** used by explorers and others to **find their way** back or through an area"



## Landmarks (2)



#### Landmarks in planning:

Something you must *achieve* or *use* in *every solution* to a problem instance

Assume we are considering a state s...

#### Fact Landmark for s:

A single <u>fact</u> (ground atom) that must be true at some point

in every solution starting in s



clear(A) holding(C)

. . .

## Landmarks (3)



#### Facts, not states! Why?

Usually **many** paths lead  $S_0$ from s to goal states Few states are shared  $S_8$  $S_4$ among <u>all</u> paths Here only s0...  $S_2$  $S_5$ In S2 and S4, **numering (D)** IS under Jandmark S2 and S4, numering (D) IS under Jandmark S2 and S4, numering (D) IS under Jandmark More likely to find **facts** In S2 and S4, holding(B) is true that occur along <u>all</u> paths  $S_7$  $S_6$ Sz  $g_1$  $g_2$ 

## Landmarks (4)





## Landmarks (5): Misunderstandings



Not "we must reach (pass through) <u>the</u> landmark state"!

> Instead "we must reach some state that satisfies the landmark"

Not "A landmark fact is a state that..."

A fact is not a state. A state consists of many facts.

("A word is a sentence that...")

A landmark fact is **not** "a fact that is true in every solution"

> A solution is a <u>plan</u>. Facts are true in *states*.

A landmark fact is "a fact that is true in some state along every path from the initial state to any goal state".

# But isn't the state space graph too large to generate?

Let's try to find <u>some</u> of the landmarks, more efficiently...

# Means-Ends Analysis Algorithm

#### Problem setup

- s = the state whose heuristic value we want, a set of true facts
  - { clear(B), ontable(B), clear(C), on(C,A),
     on(A,D), ontable(D), handempty }
- g = the goal specification, a set of desired facts
  - { clear(D), on(D,C), on(C,A), on(A,B), ontable(B) }
  - (This goal does not mention handempty!)
- One way of discovering landmarks:
   <u>means-ends analysis</u> ("backwards from the goal")
  - All facts in g must be landmarks must occur at the end of a solution path!
  - { clear(D), on(D,C), on(C,A), on(A,B), ontable(B) }
- But let's focus on the most interesting part
  - <u>4 "unachieved" landmarks, not already true in state s:</u> g - s = { clear(D), on(D,C), on(A,B), ontable(B) }



S

A

B



# Means-Ends Analysis Algorithm (2)

- Means-ends analysis, informally:
  - We start with  $g s = \{ clear(D), on(D,C), on(A,B), ontable(B) \}$
  - Now let's consider on(D,C)
    - Must be <u>achieved</u> by some action: Is a landmark, but not already true in s
  - **How** can we achieve **on(D,C)**?
    - Only using stack(D,C)



S

- What do we <u>also</u> need, in order to actually <u>execute</u> stack(D,C)?
  - All of its preconditions
  - { **holding(D), clear(C)** } must also be fact landmarks, but clear(C) is true <u>now</u>...
  - { holding(D) } another <u>unachieved</u> fact landmark!
- Updated list of 5 distinct unachieved LM found:
  - { clear(D), on(D,C), on(A,B), ontable(B), holding(D) }



g

D

B

# Means-Ends Analysis Algorithm (3)

- Some more intuitions:
  - Current list: { clear(D), on(D,C), on(A,B), ontable(B), holding(D) }
  - Let's consider <u>achieving</u> holding(D)
    - Not true now (in s), but must be *made* true
  - How can we achieve holding(D)?
    - Remember we must consider <u>all possible paths</u> to a goal state
    - → <u>All</u> actions having holding(D) as an effect
    - > { pickup(D), unstack(D,A), unstack(D,B), unstack(D,C), unstack(D,D) }
  - What do we <u>also</u> need, <u>regardless of which of these actions we use</u>?
    - The intersection of the preconditions of the 5 actions
    - Only pickup requires ontable(D); only unstack requires on(D, something)...
    - But all require { clear(D), handempty } must also be fact landmarks
    - { clear(D) } another <u>unachieved</u> fact landmark!



85

So now we have 6...

# Means-Fnds Analysis Algorithm (4)

Means-Ends Analysis Algorithm (4)	
<u>Unachieved goal facts</u> : clear(D), on(D,C), on(A,B), ontable(B)	fact-landmarks ← g – s
<ul> <li>p=on(D,C) is an unachieved fact landmark</li> <li>→ all solutions must at some point achieve on(D,C) with an action effect</li> <li>→ compute achievers = { stack(D,C) }, the only action achieving on(D,C)</li> </ul>	<pre>do {    for each p in fact-landmarks {       // Which actions could achieve p?       achievers ← {a ∈ A   p ∈ effects(a)}</pre>
<u>All</u> achievers have some common requirements / preconditions: { <i>holding(D), handempty, clear(C),</i> }	// What would <i>all</i> the achievers need? common $\leftarrow \bigcap_{a \in achievers} preconds(a)$
<pre>handempty is already true, but new = { holding(D), clear(C) } are unachieved landmarks</pre>	new ← common – s fact-landmarks ← fact-landmarks ∪ new
Maybe we can find more landmarks related to achieving <i>those</i> !	} } <b>until</b> no more fact-landmarks found
Learn to apply this algorithm! Test it on some problem instances!	

#### Landmark Counts and Costs



- One simple form of landmark heuristic: <u>Counting</u> landmarks
  - h(s) = the **<u>number</u>** of **<u>unachieved</u>** landmarks in state s

One action can actually achieve multiple landmarks at once (multiple effects) landmark count is <u>not admissible</u>

More complex (and stronger) forms of landmark heuristics also exist – pioneered by the LAMA planner

See, for example, Silvia Richter and Matthias Westphal. The LAMA planner: Guiding costbased anytime planning with landmarks. Journal of Artificial Intelligence Research, 39:127–177, 2010.