



Data Structures for Symbol Tables

Complement to Lecture 2 (4), pages 74–78 + 80–88

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Linköpings universitet, 2008.

Symbol table



Set of symbol table items

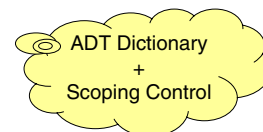
- searchable by name + scope

Data stored for each entry:

- name
- attributes
 - type (int, bool, array, ptr, function)
 - address (block, offset)
 - declared or not, used or not
 - ...

Operations

- lookup (name)
- insert (name)
- put (name, attribute, value)
- get (name, attribute)
- enter scope ()
- exit scope ()



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Data Structures for Symbol Tables



For flat symbol tables: (one block of scope)

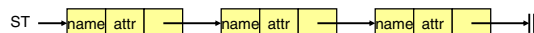
- Linear lists
 - Hash tables
 - ...
- (see data structures for ADT Dictionary)

For nested scopes:

- Trees of flat symbol tables
- Linear lists with scope control
 - Only for 1-pass-compilers
- Hash tables with scope control (see following slides)
 - Only for 1-pass-compilers

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Linear lists

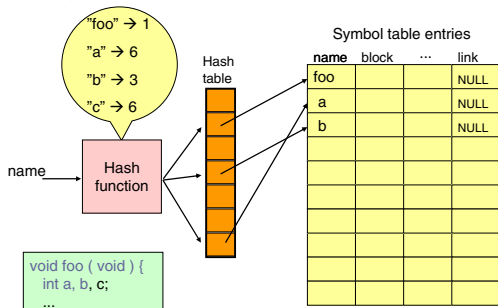


Unsorted linear lists

- Easy to implement
 - Space efficient
 - Insertion itself is fast
 - but needs lookup to check if the name was already in
 - Lookup is slow
- Inserting n identifiers and doing m lookups requires $O(n(n+m))$ string comparisons

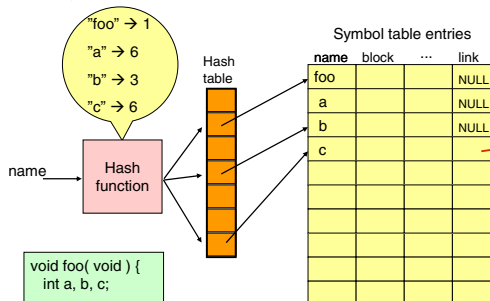
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Hash table with chaining (1)



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Hash table with chaining (2)



- Much faster lookup on average
- Degenerates towards linear list for bad hash functions

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Hierarchical Symbol Tables

For nested scope blocks

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Tree-based Symbol Table

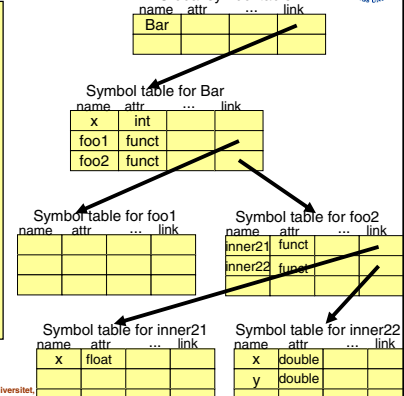
File/module scope:



```
class Bar {
  int x;
  void foo1( ... ) { ... }
  void foo2( ... ) {
    int inner21( ... ) {
      float x;
      ...
    }
    int inner22( ... ) {
      double x, y;
      ...
      foo1( x );
    }
    ...
  }
  ...
}
```

- enterscope(), exitscope()
- insert(), lookup()

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For One-Pass Compilers?

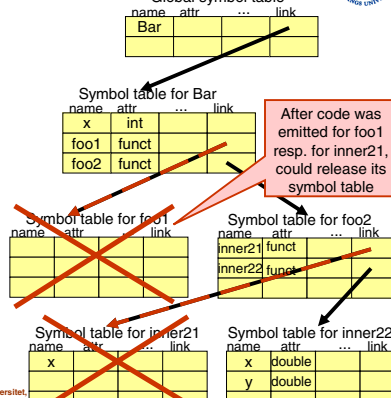
File/module scope:



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  int x;
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    int inner21( ... ) {
      float x;
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      double x, y;
      ...
      foo1( x );
    }
    ...
  }
  ...
}
```

- enterscope(), exitscope()
- insert(), lookup()

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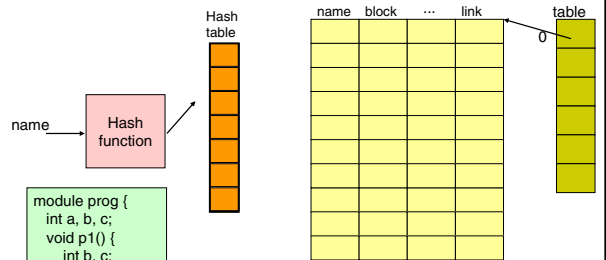


Hash tables with chaining + scoping

(For One-Pass Compilers Only)



Current scope block: 0



insert p1 and enter a new scope block (2)

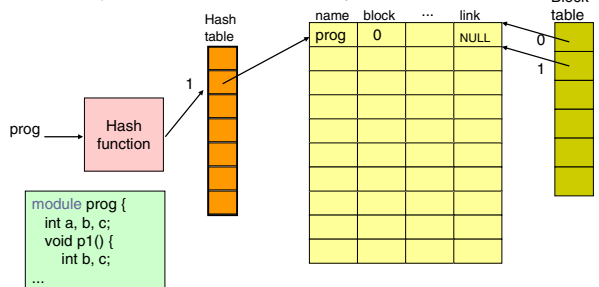
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Hash tables with chaining + scoping



Current scope block: 1



insert prog and enter a new scope block (1)

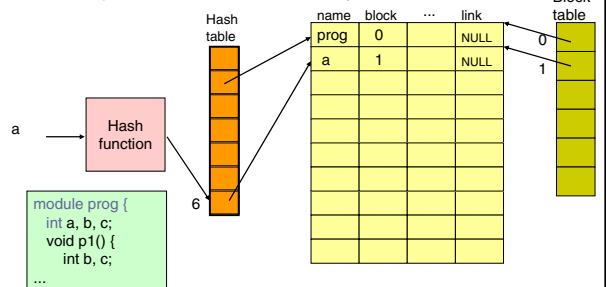
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Hash tables with chaining + scoping



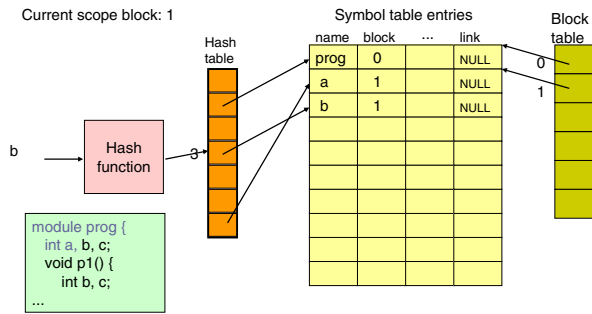
Current scope block: 1



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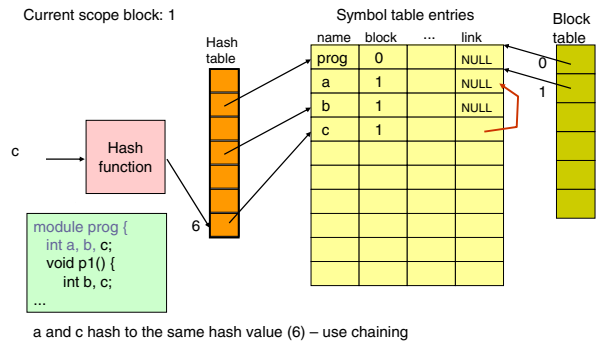
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Hash tables with chaining + scoping



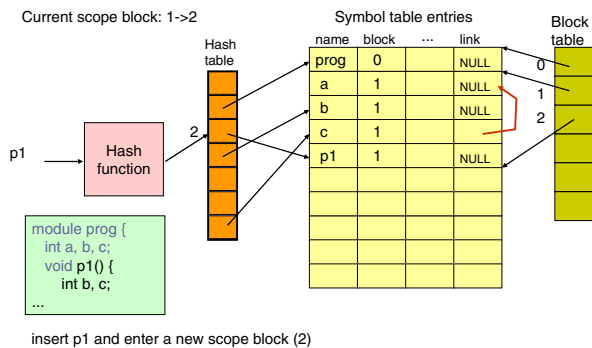
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Hash tables with chaining + scoping



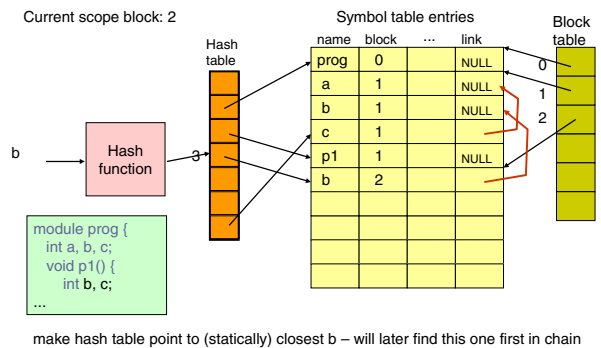
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Hash tables with chaining + scoping



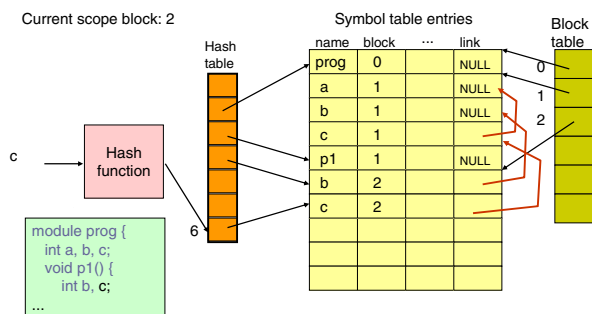
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Hash tables with chaining + scoping



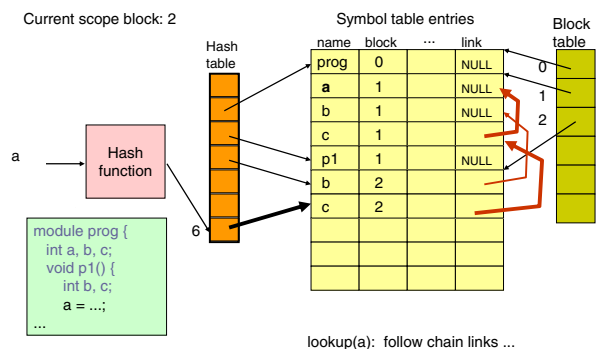
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Hash tables with chaining + scoping



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Hash tables with chaining + scoping



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More on Symbol Tables

Compiler representation of names
String space for identifiers
Arrays (see old slide set p.78, 79)

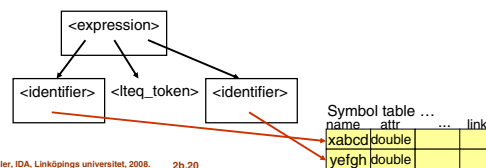
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Compiler representation of names



- A unique and compact internal representation for a name is the index (address in compiler address space) of its symbol table entry.
- Used instead of full name (string) in the internal representation of a program
- ☺ Time and space efficient

Example: Parse-tree for expression `xabcd <= yefgh;`

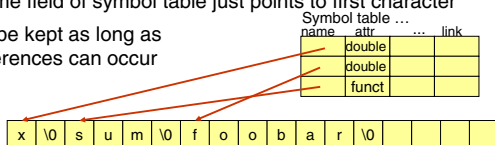


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String Space for Identifiers



- Identifiers can vary in length
- Must be stored in token table
- Name field of symbol table just points to first character
- To be kept as long as references can occur



- Usually, full names kept only during compilation
 - Exception:
Added to the program's constant pool in the .data segment if symbolic debugging or reflection should be enabled (e.g., `gcc -g file1.c` to prepare for symbolic debugging)

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